Code Explanation:

Table of Contents

Code Explanation:	1
Mux 2x1	
4x1 Mux	
Full adder:	
Logic for Sum	
Summary	
ALU1:	3
Explanation	
ALU8:	
Instantiating ALU1 Components	
Explanation	
p-ww	

Mux 2x1

This code describes a simple 2-to-1 multiplexer:

- A multiplexer selects one of the several input signals and forwards the selected input to a single output line.
- In this 2-to-1 multiplexer:
 - There are two inputs (I0 and I1).
 - There is one select line (S).
 - There is one output (0).

The functionality is as follows:

- When the select line S is '0', the output 0 will be equal to 10.
- When the select line S is '1', the output 0 will be equal to I1.

This behavior is implemented using an if-else statement within a VHDL process, which ensures that the output 0 is updated immediately when there is any change in the inputs I0, I1, or the select line S.

4x1 Mux

This code describes a 4-to-1 multiplexer using three 2-to-1 multiplexers:

- The first level consists of two 2-to-1 multiplexers (Mux2x1_0 and Mux2x1_1). These multiplexers select between pairs of the input signals (I0, I1 and I2, I3) based on the select signal S0.
- The second level consists of one 2-to-1 multiplexer (Mux2x1_final). This multiplexer selects between the outputs of the first level (mux2x1_0_out and mux2x1_1_out) based on the select signal S1.

The 4-to-1 multiplexer works as follows:

- S0 determines which of the input pairs (I0, I1 or I2, I3) are selected in the first level.
- S1 determines which output from the first level is passed to the final output 0.

In other words:

- When S1 = 0, the final output 0 depends on the selection made by S0 between I0 and I1.
- When S1 = 1, the final output 0 depends on the selection made by S0 between I2 and I3.

By cascading two levels of 2-to-1 multiplexers, this design effectively creates a 4-to-1 multiplexer.

Full adder:

This code describes the behavior of a Full Adder:

- A Full Adder adds three bits: two significant bits A and B, and a carry-in bit Cin.
- It produces a sum bit Sum and a carry-out bit Cout.

Logic for Sum

```
Sum <= (A xor B) xor Cin;
```

The sum output Sum is calculated using the XOR operation.

- (A xor B) gives the intermediate sum of A and B without considering the carry-in.
- (A xor B) xor Cin then includes the carry-in bit to produce the final sum.

Logic for Cout

```
Cout <= (A and B) or (Cin and (A xor B));
```

he carry-out output Cout is calculated using the AND and OR operations.

• (A and B) produces a carry when both A and B are 1.

- Cin and (A xor B) produces a carry when either the intermediate sum (A xor B) is 1 and the carry-in Cin is 1.
- The OR operation combines these conditions to produce the final carry-out.

Summary

- **Inputs**: A, B, and Cin are the inputs to the Full Adder.
- **Outputs**: Sum is the sum output, and Cout is the carry-out output.
- Logic:
 - Sum is the result of the XOR operation involving A, B, and Cin.
 - Cout is the result of the OR operation combining the AND of A and B with the AND of Cin and the intermediate sum (A xor B).

This code implements a basic Full Adder, which is a fundamental building block in arithmetic logic units (ALUs) and is used in various digital circuits for binary addition.

ALU1:

The ALU1 VHDL code describes a 1-bit ALU (Arithmetic Logic Unit) named ALU1. This ALU performs basic logical and arithmetic operations based on the inputs and control signals provided. Let's break down the code to understand its structure and functionality.

Library and Use Clause:

- library IEEE; and use IEEE.STD_LOGIC_1164.ALL; include the IEEE library and the STD_LOGIC_1164 package, providing the standard logic types and operations.
- Entity ALU1:
- Defines the name of the module as ALU1.
- Ports:
 - A, B: 1-bit inputs (std_logic).
 - A_inv, B_inv: 1-bit inversion control inputs (std_logic).
 - CarryIn: 1-bit carry input (std_logic).

- OPO, OP1: 1-bit operation select inputs (std_logic).
- CarryOut: 1-bit carry output (std_logic).
- Result: 1-bit result output (std logic).

Architecture Behavioral:

- Defines the implementation details of the ALU1 entity.
- The architecture is named Behavioral, indicating it describes the behavior of the ALU.
- Component Declarations:
- MUX2to1: A 2-to-1 multiplexer.
- MUX4x1: A 4-to-1 multiplexer.
- FullAdder: A full adder.
- Signal Declarations:
- not_A, not_B: Signals to hold the inverted values of A and B.
- mux2x1_A, mux2x1_B: Signals to hold the outputs of the first level of multiplexers.
- and_w, or_w, adder_w: Signals to hold intermediate results for AND, OR, and ADD operations respectively.

Inversion Logic:

- not_A <= not A;: Inverts A and assigns it to not_A.
- not_B <= not B;: Inverts B and assigns it to not_B.

2. First Level Multiplexers:

- Mux1: A 2-to-1 multiplexer that selects between A and not_A based on A_inv, outputting to mux2x1_A.
- Mux2: A 2-to-1 multiplexer that selects between B and not_B based on B_inv, outputting to mux2x1_B.

3. Logical Operations:

- and_w <= mux2x1_A and mux2x1_B;: Performs bitwise AND on mux2x1_A and mux2x1_B.
- or_w <= mux2x1_A or mux2x1_B;: Performs bitwise OR on mux2x1_A and mux2x1_B.

4. Full Adder:

• Adder: A full adder that sums mux2x1_A, mux2x1_B, and CarryIn. The sum is assigned to adder_w and the carry-out is assigned to CarryOut.

5. Final Multiplexer:

• mux4_1: A 4-to-1 multiplexer that selects between and_w, or_w, adder_w, and 0 based on OP0 and OP1, outputting to Result.

Explanation

This 1-bit ALU performs the following operations based on the control signals A_inv, B_inv, OPO, and OP1:

• Inversion Control:

- A_inv: When set to 1, inverts A before further processing.
- B_inv: When set to 1, inverts B before further processing.

• Operations:

- AND (OP1 = 0, OP0 = 0): Result = mux2x1_A AND mux2x1_B.
- OR (OP1 = 0, OP0 = 1): Result = mux2x1_A OR mux2x1_B.
- ADD (0P1 = 1, 0P0 = 0): Result = mux2x1_A + mux2x1_B + CarryIn.
- Zero (OP1 = 1, OP0 = 1): Result = 0.

This ALU module can be used as a building block for more complex multi-bit ALUs by cascading multiple instances of this 1-bit ALU, each handling one bit of the multi-bit operation.

ALU8:

The provided VHDL code describes an 8-bit ALU (Arithmetic Logic Unit) named ALU8. This ALU is built by instantiating eight 1-bit ALU components (ALU1) to handle 8-bit operations.

Library and Use Clause:

- library IEEE; and use IEEE.STD_LOGIC_1164.ALL; include the IEEE library and the STD_LOGIC_1164 package, providing the standard logic types and operations.
- Entity ALU8:
- Defines the name of the module as ALU8.
- Ports:
 - A, B: 8-bit inputs (std_logic_vector(7 downto 0)).
 - A in, B inv: 1-bit control inputs (std logic) for inverting A and B.
 - CarryIn: 1-bit carry input (std_logic).
 - OP: 2-bit operation select input (std_logic_vector (1 downto 0)).
 - CarryOut: 1-bit carry output (std_logic).
 - Result: 8-bit result output (std_logic_vector(7 downto 0)).

Architecture Behavioral:

- Defines the implementation details of the ALU8 entity.
- The architecture is named Behavioral, indicating it describes the behavior of the ALU.
- Component Declaration:
- **ALU1**: A 1-bit ALU component which is instantiated eight times, one for each bit of the 8-bit ALU.
- Signal Declaration:

- carry: A 9-bit signal (std_logic_vector(8 downto 0)) to hold intermediate carry signals between each 1-bit ALU.
- Initial Carry Input:
- carry(0) <= CarryIn;: Sets the initial carry input to the value of CarryIn.

Instantiating ALU1 Components

The ALU8 is constructed by instantiating eight ALU1 components, each handling one bit of the 8-bit inputs.

```
-- Instantiate ALU1 components for each bit
ALU1 0: ALU1 port map (
  A
        => A(0),
  A_{inv} => A_{in}
       => B(0),
  B_{inv} => B_{inv}
  CarryIn => carry(0),
       => OP(0),
  OP0
  OP1
         => OP(1),
  CarryOut => carry(1),
  Result => Result(0)
);
ALU1_1: ALU1 port map (
  A => A(1),
  A_{inv} => A_{in}
  B => B(1),
  B_{inv} => B_{inv}
  CarryIn => carry(1),
  OP0
        => OP(0),
         => OP(1),
  OP1
  CarryOut => carry(2),
  Result => Result(1)
);
ALU1_2: ALU1 port map (
  A => A(2),
  A_{inv} => A_{in}
        => B(2),
  В
  B_{inv} => B_{inv}
  CarryIn => carry(2),
  OP0 => OP(0),
         => OP(1),
  OP1
  CarryOut => carry(3),
  Result => Result(2)
);
ALU1_3: ALU1 port map (
        => A(3),
  A_{inv} => A_{in}
       => B(3),
```

```
B_{inv} => B_{inv}
  CarryIn => carry(3),
  OP0
         => OP(0),
  OP1
         => OP(1),
  CarryOut => carry(4),
  Result => Result(3)
);
ALU1_4: ALU1 port map (
  A => A(4),
  A_{inv} => A_{in}
        => B(4),
  B_{inv} => B_{inv}
  CarryIn => carry(4),
  OP0
         => OP(0),
  OP1
         => OP(1),
  CarryOut => carry(5),
  Result => Result(4)
);
ALU1_5: ALU1 port map (
        => A(5),
  A_{inv} => A_{in}
        => B(5),
  B_{inv} => B_{inv}
  CarryIn => carry(5),
  OP0
         => OP(0),
  OP1
         => OP(1),
  CarryOut => carry(6),
  Result => Result(5)
);
ALU1_6: ALU1 port map (
        => A(6),
  Α
  A_{inv} => A_{in}
       => B(6),
  B_inv => B_inv,
  CarryIn => carry(6),
  OP0
         => OP(0),
  OP1
         => OP(1),
  CarryOut => carry(7),
  Result => Result(6)
);
ALU1_7: ALU1 port map (
  Α
       => A(7),
  A_{inv} => A_{in}
        => B(7),
  B_inv => B_inv,
  CarryIn => carry(7),
  OP0
         => OP(0),
  OP1
         => OP(1),
```

```
CarryOut => carry(8),
Result => Result(7)
);
```

Each instance of ALU1 is connected to a corresponding bit of the 8-bit inputs A and B. The inversion control signals A_in and B_inv, and the operation select signals OP(0) and OP(1) are common to all instances. The carry-out of each ALU1 instance is connected to the carry-in of the next ALU1 instance in sequence.

Final Carry Output

```
-- Final carry output
CarryOut <= carry(8);</pre>
```

The final carry output of the ALU8 is taken from carry(8).

Explanation

1. Initialization:

• The initial carry input is provided by CarryIn.

2. Bit-wise ALU Operations:

- Each ALU1 instance performs operations on one bit of A and B, with carry propagation handled by the carry signal.
- The results of these operations are collected in the Result signal.

3. Carry Propagation:

• The carry generated by each ALU1 instance is propagated to the next bit, ensuring correct arithmetic operations across the 8 bits.

4. Final Carry Out:

• The final carry output after all 8 bits have been processed is provided as CarryOut.

This modular approach allows for easy scalability and clear separation of logic for each bit, facilitating debugging and extension to larger bit-width ALUs if needed.