

# Power-Aware Security Protocols for the Internet of Things

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**Abstract.** le pretty abstact...

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## 1 Introduction

The Internet of Things (IoT) can be seen as web of interconnected devices that go from everyday wearable objects into fully deployed sensor networks. Despite the huge variety and characteristics of these devices, one thing that they all have in common is the constrained nature they're built upon. In order to enable the massive deployment to be expected in the near future <sup>1</sup> IoT devices must be accessible and affordable, capable of operating under lossy wireless networks while being battery powered. This poses a challenge to current Internet protocols since the assumptions regarding the devices capabilities and objectives do not hold true. To allow the IoT vision to come forward, several new protocols have been developed across the OSI layers, each addressing and tackling the challenges involved in trying to keep the quality and assurances of stronger, more expensive protocols, on constrained systems. Additionally, security is also a big topic of interest due to the fact that the interconnection of the devices around us can provide information about our choices and whereabouts, therefore reducing our privacy. This document will address these issues from a power-aware perspective, meaning the battery consumption will be of major importance.

## 2 Main Goals

Given the constraints and limitations of IoT devices described in the previous chapter and the recent protocols, the first objective of this work is to identify a working stack of protocols that takes into account these constraints and focuses on allowing a power-aware communication model.

After the analysis of the existing solutions, a baseline of power consumption will be established. Then, the focus will move towards adding confidentiality to the transmitted information by securing the channel. Once both the application level protocols and proper security solutions are defined, experiments will be performed so that the added power consumption cost of adding the security layers can be measured, profiled and documented therefore enabling the finding of the best parameters for a desired level of security.

The work will then proceed towards finding effective counter-measures against a specific group of attacks that targets the IoT devices by intensifying the use of its resources therefore draining the available power and placing the device offline. The ultimate goal is to propose an energy-efficient security mechanism that can resist these power-drain (a.k.a vampire) attacks.

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<sup>1</sup><http://blogs.wsj.com/cio/2015/06/02/internet-of-things-market-to-reach-1-7-trillion-by-2020-idc/>

## 3 Related Work

### 3.1 Protocol Analysis and Selection

There are many alternatives and some proposed standards when it comes to choosing a protocol stack for IoT communications. The decision must be based on the particularities of the devices to be used and the objective of the application itself, however a thoroughly analysis of the existing solutions is a proper way to unveil the strong and weak points of each protocol providing a good basis for an informed decision. A recent survey (January 2015) [1] of the IoT enabling technologies, protocols and applications will be the starting point for the analysis to follow. The presentation of the available protocols and solutions will follow a bottom-up approach, starting from the data link and physical layer all the way up until the application layer. In particular, the session layer will be left to the end since securing the channel is an optional feature and will be addressed after the application level protocols are properly examined.

#### 3.1.1 Data Link and Physical Layer

The first requirement for the physical layer of the IoT is the use of wireless radios. These should aim for simplicity, low-power and low-cost communications. While wireless communication are far spread and can be found from homes to airports, the type of radio commonly used, known as Wi-Fi, use a high amount of power causing concerns for battery life. In the next paragraphs, an overview of Wi-Fi(IEEE 802.11) is given with the objective of comparing it with the IEEE 802.15.4, a protocol that aims to address these issues.

##### ***IEEE 802.11***

IEEE 802.11 is a set of standards for Wireless Local Area Networks (WLAN) communications. They are the basis for the so called Wi-Fi. IEEE 802.11 is concerned with Ethernet matching speed, long ranges, message forwarding and high data throughput. These concerns directly clash with the IoT objectives and account for the added power consumption of this protocol.

##### ***IEEE 802.15.4***

IEEE 802.15.4 on the other hand was created for Low-Rate Wireless Private Area Networks (LR-WPAN) and its specifications on low power consumption, low data rate, low cost and high message throughput make it a strong candidate for IoT applications. The IEEE 802.15.4 standard supports two types of network nodes, the Full Function Device (FFD) that act as coordinator or normal nodes. And the Reduced Function Device (RFD) that are very simple, with very restricted resources and can only communicate with coordinators. The coordinators are responsible for controlling and maintaining the network. FFD are capable of storing a routing table in their memory and can implement a full Medium

Access Control (MAC). IEEE 802.15.4 supports star, peer-to-peer(mesh) and cluster-tree topologies. Regarding performance, it would be unfair to directly compare the two, since IEEE 802.11 transmission power and receiver sensitivity are much greater than 802.15.4. But if we limit both to a low power level IEEE 802.11 still outperforms IEEE 802.15.4 in terms of packet delivery ratio, throughput, latency, jitter and average energy consumption. However this comes at the cost of a far lower transmission range[2]. We can conclude that for typical LR-WPAN network requirements, IEEE 802.15.4 is better designed to address the constrained environment issues, while IEEE 802.11 would still be a suitable option if transmission range is not a problem.

### 3.1.2 Network Layer

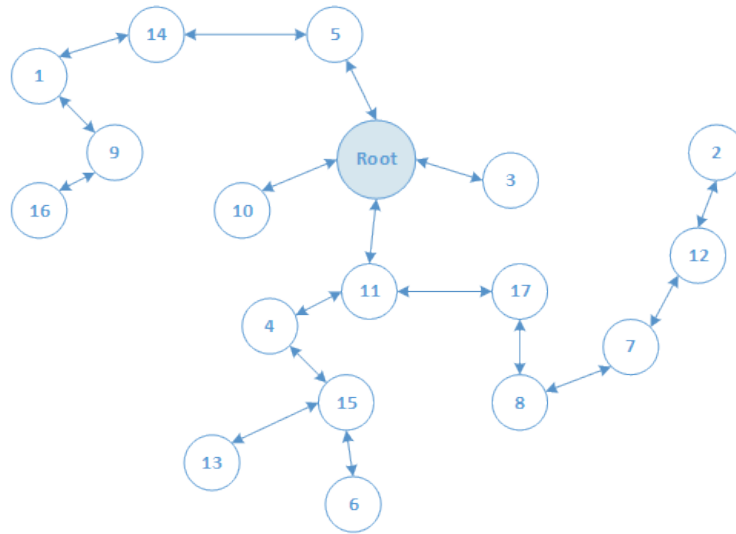
#### *6LoWPAN*

The IoT vision, as presented in the introduction, and its massive deployment can only be achieved through the use of IPv6. However, physical layers more suitable for communication over constrained networks pose some limitations to the use of the IPv6 messages. For example the limited packet size in IEEE 802.15.4 based networks. To tackle these issues, the Internet Engineering Task Force (IETF) 6LoWPAN working group developed a standard based on header compression to reduce the transmission overhead, fragmentation to meet the IPv6 Maximum Transmission Unit (MTU) requirements and forwarding to link-layer to support multi-hop delivery. [3] 6LoWPAN is able to remove a major share of IPv6 overheads, being able to compress its headers to two bytes, therefore allowing small IPv6 datagrams to be sent over IEEE 802.15.4 networks.

#### *RPL*

With the use of 6LoWPAN, upper layer routing protocols can now use the IPv6 addressing scheme. Given the possible frequent topology changes associated with the radio-link instability, successful solutions must take these requirements into account on their specification. RPL can support a wide variety of link-layers and is prepared for devices with very limited resources. It is able to build up network routes, distribute routing knowledge among nodes and adapt the topology in a very efficient way. More in depth, RPL creates a Destination Oriented Directed Acyclic Graph (DODAG) between the 6LoWPAN network nodes (Figure 1) that supports unidirectional traffic towards the DODAG root and bidirectional traffic between devices. Each node has a rank that indicates its position relative to other nodes and with respect to the root. This rank is used to create optimized network paths. In order to allow packets to propagate downwards the topology, either source routing or stateful routing tables are used. (More Information on this two types of routing are given in sections 3.2.1 and 3.2.2). For both modes, the DODAG root always maintains a complete list of the network nodes. RPL provides a set of control messages in order to exchange routing graph

information. DODAG Information Objects (DIO) are used to advertise information needed to build the DODAG. Destination Advertisement Object (DAO) are used to advertise information so that downwards traffic can go through the nodes towards the leafs. Nodes may also resort to DODAG Information Solicitation (DIS) messages to request graph information from neighbour nodes. Finally, RPL has a built in topology repair mechanism that acts in the case of a routing topology failure, link failure or node failure. In the case the topology needs to be rebuilt, a link layer metric is used to calculate the new route. The new path is considered fit for work if the link layer acknowledgements are received on it.



**Fig. 1.** A Sample RPL DODAG

### 3.1.3 Application Layer

#### *Hypertext Transfer Protocol (HTTP)*

HTTP is an application level protocol that works in the request-response model and is the foundation of data communication on the World Wide Web (WWW). It is primarily designed to run over Transmission Control Protocol (TCP) which is a problem in lossy and constrained environments due to the delivery assurances and congestion control algorithms it employs. Besides, HTTP is verbose, text-based, and not suited for compact message exchanges. Moreover, the header size required for a message exchange can leave too few payload space in constrained networks like the IEEE 802.15.4-based networks where the MTU size of the protocol is 127 bytes. These protocol specifications would not raise any

issues in standard WWW communications, but when it comes to constrained environments it is clear that the protocol is not adequate to the necessities of IoT devices and networks.

### ***Constrained Application Protocol (CoAP)***

CoAP is a document transfer protocol based on REpresentational State Transfer (REST) on top of HTTP functionalities. CoAP objective is to enable tiny constrained devices to use RESTful interactions, where clients and servers expose and consume web services using Universal Resource Identifiers (URIs) together with HTTP get, post, put and delete methods. Unlike REST, CoAP runs over User Datagram Protocol (UDP) instead of TCP which makes it suitable for full IP networking in small micro-controllers. Retries and reordering are implemented at the application stack using a messaging sub-layer that detects duplicated messages and provides reliable communication using different types of messages. Confirmable messages must be acknowledged by the receiver, non-confirmable follow the fire and forget model. While being a lightweight protocol, CoAP still provides important features:

- Resource Observation - CoAP can extend the HTTP request model with the ability to observe a resource therefore monitoring resources of interest using a publish/subscribe mechanism.
- Resource Discovery - CoAP servers provide a list of resources using well-known URIs that allow clients to discover what resources are provided and their types.
- Interoperability - since CoAP is based on the REST architecture, a simple proxy enables CoAP to easily interoperate with HTTP.

A study that compared CoAP and HTTP using mobile networks concluded that there is no situation where CoAP would consume more resources than HTTP [4]

### ***Message Queue Telemetry Transport (MQTT)***

MQTT is a publish/subscribe messaging protocol designed for lightweight Machine to Machine (M2M) communications. It employs a client/server model and consists of three components, the publisher, the subscriber and a broker. Subscribers register their interest for a specific topic and then get informed by the broker when a publisher generates data regarding that topic. Every message is a discrete chunk of data, opaque to the broker. The broker, on its side, checks authorization of the publishers and subscribers. MQTT supports three Application Level Quality of Service (QoS) levels:

- At Most Once (Fire and Forget): A message will not be acknowledged by the receiver or stored and redelivered by the sender.

- At Least Once: It is guaranteed that the message will be delivered to the receiver, but more than one can reach the destination due to message resending. The sender stores the message until it gets an acknowledgment from the receiver.
- Exactly Once: A four-way handshake mechanism is used to guarantee that the message will be received exactly once by the counterpart.

MQTT has support for persistent messages stored on the broker, where the most recent message will be sent to a client that subscribes that topic. Clients can register a custom message to be sent to the broker on disconnect enabling other subscribers to know when a device disconnects. MQTT runs on TCP which in some cases causes drawbacks in performance. A performance evaluation of MQTT and CoAP [5] provides comparisons on several protocol facets:

- Influence of Packet Loss on Delay: With low values of packet loss, MQTT experienced lower delays, but as the packet loss increased CoAP performed better. This is due to the greater TCP overheads involved in the retransmissions of messages when compared to UDP.
- Influence of Packet Loss on Data Transfer: CoAP generated less data for each packet loss versus all the MQTT QoS levels.
- Overheads for Message Sizes: When packet loss rate is low, CoAP generates less overhead than MQTT for all message sizes, but as message size grows, the reverse is true. This happens because when the message size is large, the probability that UDP loses the message is higher than TCP which causes CoAP to retransmit the whole message more often than MQTT.

In order to address the drawbacks on constrained devices, Message Queue Telemetry Transport for Sensor Networks (MQTT-SN) protocol[6] was created. Among the improvements and new features, MQTT-SN runs on UDP, adds broker support for indexing topic names, provides a discovery procedure to help clients without a pre-configured server address and supports devices in sleep state. With this approach, an extra gateway is necessary to convert from MQTT-SN to MQTT so the communications can be understood by the broker.

#### 3.1.4 Session Layer

So far security issues have not been addressed in any of the previous layers, this is because security is an expensive, optional feature. The application layer protocols rely on underneath layers to achieve secure communications, and network layer protocols assume that if security is necessary then it has already been handled in upper layer protocols. In fact, the session layer is where the security mechanisms are implemented and provides an abstraction layer to application layer protocols. These mechanisms work on top of the transport layer and aim to provide authentication, confidentiality and message integrity.



### ***Transport Layer Security (TLS)***

TLS is a well-known security protocol that is used to provide secure transport layer for TCP communications, allowing the upper layer protocols to be left untouched. TLS operation consists of two phases: the handshake and then the data encryption. During the handshake, both parties negotiate which algorithms will be used during the session, authenticate themselves, and prepare the shared secret for the data encryption. Both HTTP and MQTT work over TCP and use TLS as the adopted security protocol.

### ***Datagram Transport Layer Security (DTLS)***

DTLS aims to be the equivalent of TLS over UDP transport layer. DTLS works over datagrams that can be lost, duplicated, or received in the wrong order, therefore needing some extra mechanisms(application layer protocols QoS) to cope with that. Although both CoAP and MQTT-SN work over UDP and use DTLS as the adopted security, some authors argue that DTLS is not a suitable option [7] and defend the need of a new integrated security solution. Some of the presented drawbacks are:

- There is no multicast support, which is a key feature in IoT (topology discovery and update for example).
- Handshake phase is prone to exhaustion attacks on the device resources.
- The loss of a message in-flight requires the retransmission of all the messages in-flight.

A final overview of the analysed protocols and security solutions is given in Table 1. And a comparison of the protocol stack is shown in Table 2.

**Table 1.** IoT Application Protocols Comparison

Application Protocol	RESTful	Request/ Response	Publish/ Subscribe	Adjustable QoS	Transport	Security
HTTP	✓	✓	✗	✗	TCP	TLS
CoAP	✓	✓	✓	✓	UDP	DTLS
MQTT	✗	✗	✓	✓	TCP	TLS
MQTT-SN	✗	✗	✓	✓	UDP	DTLS

**Table 2.** Protocol Stack Comparison Overview

Layer	Web	IoT
Application	HTTP	CoAP
Session	TLS	DTLS
Transport	TCP	UDP
Network	IPv6	6LoWPAN
Data-Link/Phy	802.11	802.15.4

### 3.2 Attack Analysis, Detection and Mitigation

Exploitation of existing solutions in the forms of malicious attacks can be found at all the studied OSI layers. They can go from the well-known Denial of Service (DoS) at the application layer to a physical intruder replacing some node on a sensor field. However, given the characteristics of the devices and networks used in IoT combined with the power consumption focus of this work, a specific kind of attacks performed at the network layer is of special interest and importance:

#### *Battery Depletion Attacks aka Vampire Attacks*

Battery Depletion Attacks aim at draining the battery, "life", of the network devices, working over time to entirely disable a network, hence being called Vampire Attacks. These attacks do not focus on flooding the network with many packages, instead they drain the node's life by delaying the packets transmission. Many of the existing attacks are not protocol specific [8], while others target specific protocols and implementations [9]. The following attacks aim at giving an overview of the existing attack possibilities on different routing solutions as well as existing mitigation strategies. Additionally, a range of attacks that target the RPL routing protocol is also analysed. Since RPL is the selected protocol of our energy efficient stack, it is of special importance to consider and assure the mitigation of attacks that would drain the devices batteries by exploiting this light weight protocol inner workings.

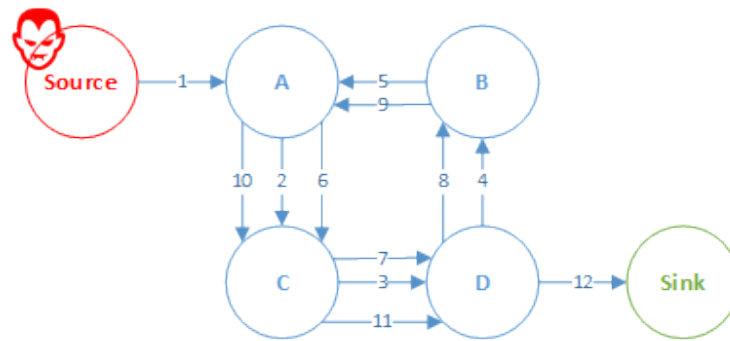
#### 3.2.1 Stateless Protocols

In systems that use this type of routing protocols, the source node specifies the entire route to the destination in the packet header. This means that intermediaries do not make decisions regarding the next hop, they only forward to the next node as specified in the original path therefore reducing the amount of computation performed and used energy. However, the source node must ensure that the route is valid at the time of sending and that the neighbour relations among the devices allow the specified forwarding path. Using this transmission

scheme, a malicious device can specify paths through the network that are far from optimal, wasting energy at the intermediate nodes who follow the included malicious source route. A couple examples of these attacks are the Carousel and Stretch Attacks.

### ***Carousel Attack***

The objective of this attack is to send a packet along a route composed as a series of loops. This way a single node may forward the malicious packet several times increasing the total energy consumption by a factor of the number of loops the attacker has introduced on the packet header path. It targets source routing protocols by exploiting the limited verification of the packets headers at the intermediary nodes. Figure 2 shows an example where a vampire node created a path composed of circles around the network when it could exit after the first hop through the D node.

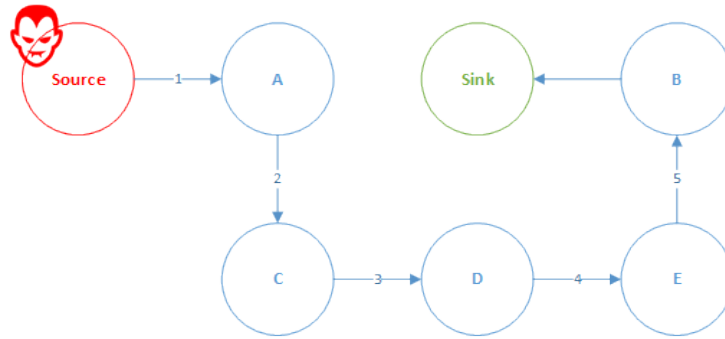


**Fig. 2.** Carousel Attack

Existing mitigations strategies rely on checking the source route for loops on intermediary nodes, either selecting an appropriate route for the packet or simply dropping it.

### ***Stretch Attack***

The objective of this attack is to create a longer source route around the network than the one who would be required to transverse the network from the source to the sink. The number of elements in the path would be greater than the optimal path, therefore increasing the total energy consumption by a factor of the number of additional hops. It's success rests on intermediary nodes not checking for better paths. Figure 3 shows an example where a vampire node created a path that goes through a greater number of nodes than required to reach the sink.



**Fig. 3.** Stretch Attack

A limited way of mitigating this attack would be to ensure that path routes have less than the total number of devices on the network. Vasserman and Hoper proposed a property called "no-backtracking" that assures the packet is always moving closer to the sink on every hop. [8]

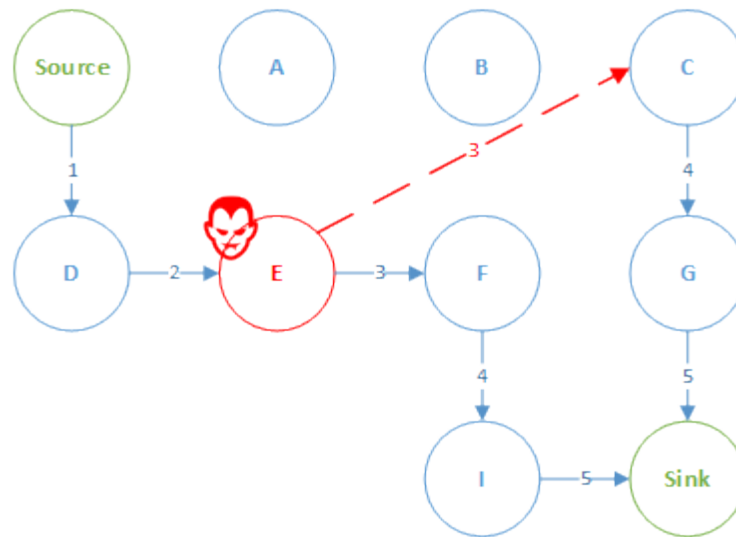
### 3.2.2 Stateful Protocols

In systems that use this type of routing protocols, network nodes are aware of the network topology and its state, being able to make local decisions on the node to whom they will forward the packet. The effect of the Vampires on this type of routing is limited since the route is built dynamically from many independent forwarding decisions. However, attackers can still cause damage by forcing packet forwarding through nodes that would not be on the optimal path, for example by forwarding the packet back to the source. A couple examples of these attacks are the Directional Antenna and Wormhole Attacks.

#### *Directional Antenna Attack*

In this attack, the attacker takes the role of an intermediary and not the source of a packet. If the attacker has the resources to use a directional antenna, it can deposit a packet on arbitrary parts of the network while also forwarding the packet locally. This causes nodes that were not on the optimal path to also consume energy by forwarding a packet they would not normally receive, therefore increasing the total energy consumption by a factor of the directions the attacker can position the antenna and the distance between the receiver and the sink. Figure 4 shows an example where a vampire intermediary deposited a node on a distant location of the network, causing the packet to follow 2 different routes towards its destination

A mitigation strategy could be to analyse the route paths of a given packet that reached the sink more than one time. The last node identifier to appear duplicated before the path started to diverge would be one who then directed the packet to multiple regions, the attacker.



**Fig. 4.** Directional Antenna Attack

### *Wormhole Attack*

### 3.2.3 RPL Specific Attacks

#### *Hello Flooding*

*Selective Forwarding* não sei onde meter, mas tem de aparecer aqui o clone attack...

### 3.3 Secure Bootstrapping

dizer o que é, de que forma se relaciona, que problemas resolve e um paragrafo com titulo a bold por cada forma que pode resolver o problema

## 4 Proposed Solution

## 5 Work Evaluation

## 6 Work Planning

## 7 Conclusion

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