# Year 3 — Number Theory

## Based on lectures by Professor Henri Johnston Notes taken by James Arthur

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These notes are not endorsed by the lecturers, and I have modified them (often significantly) after lectures. They are nowhere near accurate representations of what was actually lectured, and in particular, all errors are almost surely mine.

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## 1 Divisibility

#### 1.1 Division Algorithm

**Definition 1.1** (Well Ordering Principle). Every non-empty subset of  $\mathbb{N}_0$  contains a least element

**Theorem 1.2** (Division Algorithm). Given a  $a \in \mathbb{Z}$  and a  $b \in \mathbb{N}_1$  there exists unique integers q and r satisfying a = bq + r and  $0 \le r < b$ .

The proof splits into uniqueness and existence.

*Proof.* We shall first prove existence, define  $S := \{a - xb : x \in \mathbb{Z} \text{ and } a - xb \ge 0\}$ . We know  $S \ne 0$  since,

- if  $a \ge 0$ , then choose m = 0, then  $a mb = a \ge 0$
- if a < 0, then let a = m, so  $a mb = a ab = (-a)(b 1) \ge 0$  since -a > 0 and  $b > 0^1$

Hence S is non-empty subset of  $\mathbb{N}_0$  and so by the well ordering principle S must contain a least element  $r \geq 0$ . Since  $r \in S$ , then we have there exists a  $q \in \mathbb{Z}$  such that a - qb = r and so a = qb + r. Now it remains to check that r < b, so assume for a contradiction that  $r \geq b$ , then let there be a  $r_1 = r - b \geq 0$ . Then,

$$a = qb + r = qb + (r_1 + b) = (q + 1)b + r_1$$

and so  $a - (q+1)b = r_1 \in S$  and is smaller than r, a contradiction.

Now let us show uniqueness, assume that there exist another pair q', r' such that a = q'b + r' where  $0 \le r' < b$ . Then form a = a + qb + r = q'b + r' we have that, (q - q')b = r' - r. If q = q', then we must have r = r', suppose for a contradiction that this isn't true, then,

$$b \le |q - q'||b| = |r - r'|$$

However, since  $0 \le r, r' < b$  and so |r - r'| < b which gives a contradiction.

Here's a definition that I feel is useful that wasnt covered in the lectures.

**Definition 1.3** (Divisible). We say that some  $a \in \mathbb{Z}$  is divisible by some  $b \in \mathbb{Z}$  if and only is,

$$\exists n \in \mathbb{Z}$$
, such that  $b = na$ 

and denote it,  $a \mid b$ 

#### 1.2 Greatest Common Divisor

Let us start with a theorem.

**Theorem 1.4.** Let  $a, b \in \mathbb{Z}$ ,  $\exists d \in \mathbb{N}_0$  and non-unique  $x, y \in \mathbb{Z}$  such that,

- (i)  $d \mid a \text{ and } d \mid b$
- (ii) and if  $e \in \mathbb{Z}$ ,  $e \mid a$  and  $e \mid b$ , then  $e \mid d$
- (iii) d = ax + by

<sup>&</sup>lt;sup>1</sup>You absolute plank, there doesn't exist any numbers between 0 and 1 in  $\mathbb{Z}$ , so b>0 is the same as  $b\geq 1$ 

Proof. If a = b = 0, then d = 0Suppose that  $a \neq b \neq 0$ , then let

$$S := \{am + bn : m, n \in \mathbb{Z} \text{ and } am + bn > 0\}$$

Now  $a^2 + b^2 > 0$  so S is non-empty and a subset of  $\mathbb{N}_1$ . Hence, by the Well ordering principle then there must be some minimum element d. Then we can write d = ax + by by definition of S.

By the division Algorithm, a = qs + r for some  $q, r \in \mathbb{Z}$  with  $0 \le q < d$ . Suppose for a contradiction that  $r \ne 0$ . Then,

$$0 < r = a - qd = a - q(ax + by) = (1 - qx)a - qby$$

Hence,  $r \in S$ . But r < d, contradiciting the minimality of d in S. So we must have r = 0, i.e  $d \mid a$ . The same works for  $d \mid b$ .

Suppose that  $e \in \mathbb{Z}$ ,  $e \mid a$  and  $e \mid b$ . Then e divides any linear combination of a and b, so  $e \mid d$ . Suppose that  $e \in \mathbb{N}_1$  also satisfies (i) and (ii). Then,  $e \mid d$  and  $d \mid e$  and so  $d = \pm e$ , but  $d, e \geq 0$  and so d = e. Thus d is unique.

Note that this is a standard trick to prove that integers divide, by just proving that r = 0 by contradiction.

Corollary 1.5. If  $a, b \in \mathbb{Z}$  then there exists a unique  $d \in \mathbb{N}_1$  such that.

- (i)  $d \mid a \text{ and } d \mid b$
- (ii) if  $e \in \mathbb{Z}$ , then  $e \mid a$  and  $e \mid b$  then  $e \mid d$

*Proof.* The existence of a d is given by the theorem. In the proof of uniqueness we only use (i) and (ii).  $\square$ 

**Definition 1.6** (Greatest Common Divisor). Let  $a, b \in \mathbb{Z}$ . Them d of the previous corollary is just the greatest common divisor of a and b, written gcd(a, b). Also sometimes seen as hcf(a, b).

If gcd(a, b) = 1, then a and b are coprime.

**Identity** (Bezouts Identity). Given  $a, b \in \mathbb{Z}$  there exist  $x, y \in \mathbb{Z}$  such that gcd(a, b) = ax + by.

**Proposition 1.7.** Let  $a, b, c \in \mathbb{Z}$ , then,

- (i) gcd(a, b) = gcd(b, a)
- (ii) gcd(a, gcd(b, c)) = gcd(gcd(a, b), c)
- (iii) gcd(ac, bc) = |c| gcd(a, b)
- (iv) gcd(1, a) = gcd(a, 1) = a
- (v) gcd(0, a) = gcd(a, 0) = |a|
- (vi)  $c \mid \gcd(a, b)$  if and only if  $c \mid a$  and  $c \mid b$
- (vii) gcd(a+cb,b) = gcd(a,b)

Then we can consider the following remark,

**Remark.** Note that gcd(a, b) = 0 if and only if, a = b = 0. Otherwise,  $gcd(a, b) \ge 1$ .

*Proof.* Checking these properties are pretty simple, for (vi) just use Bezouts.

We shall prove (iii), so let  $d = \gcd(a, b)$  and  $e = \gcd(ac, bc)$ . By (vi),  $cd \mid e = \gcd(ac, bc)$  since  $cd \mid ac$  and  $cd \mid bc$ . Then by Bezouts, there exists  $x, y \in \mathbb{Z}$  such that d = ax + by. Then,

$$cd = acx + bcy$$

and as  $e \mid ac$  and  $e \mid bc$  and so by linearity we have  $e \mid cd$ . Therefore, |e| = |cd| and so, e = |c|d.

Now, let's prove (vii), let  $e = \gcd(a + bc, b)$  and  $f = \gcd(a, b)$ . Then  $e \mid (a + bc)$  and  $e \mid b$ . Thus by linearity, we have  $e \mid a$ . Hence,  $e \mid a$  and  $e \mid b$  so by property (vi), we have  $e \mid f$ . Similarly we can get that  $f \mid a + bc$  and  $f \mid b$  and so again my (vi) we have e = f as  $f, e \geq 0$ .

**Lemma 1.8** (Euclids Lemma). Let  $a, b, c \in \mathbb{Z}$ . If  $a \mid bc$  and gcd(a, b) = 1, then  $a \mid c$ .

*Proof.* Suppose that  $a \mid bc$  and gcd(a,b) = 1. By Bezouts, we get that for some  $x,y \in \mathbb{Z}$  we get 1 + ax + by. Hence, c = acx + bcy, but  $a \mid acx$  and  $a \mid bcy$ , so  $a \mid c$  by linearity.

**Theorem 1.9** (Solubility of linear equations in  $\mathbb{Z}$ ). Let  $a, b, c \in \mathbb{Z}$ . The equation,

$$ax + by = c$$

is soluble with  $x, y \in \mathbb{Z}$  if and only if gcd(a, b) = c

*Proof.* Let  $d = \gcd(a, b)$ . Then  $d \mid a$  and  $d \mid b$  so if there exists  $x, y \in \mathbb{Z}$  such that c = ax + by then  $d \mid c$  by linearity of divisibility. Now, suppose that  $d \mid c$ . Then we can write c = qd for some  $q \in \mathbb{Z}$ . By Bezouts, there exists some  $x', y' \in \mathbb{Z}$  such that d = ax' + by'. Hence, c = qd = aqx' + bqy' and so x = qx' and y = qy' gives a suitable solution.

#### 1.3 Euclids Algorithm

**Theorem 1.10** (Euclids Algorithm). Let  $a, b \in \mathbb{N}_1$  with a > b > 0 and  $b \nmid a$ . Let  $r_0 = a$ ,  $r_1 = b$  and apply the division Algorithm repeatedly to obtain a sequence of remainders defined successively,

$$r_0 = r_1q_1 + r_2$$
  $0 < r_2 < r_1$   $0 < r_3 < r_2$   $\vdots$   $0 < r_n < r_{n-1}$   $0 < r_n < r_{n-1}$ 

Then the last non-zero remainder,  $r_n$  is the gcd(a, b).

*Proof.* There is a stage at which  $r_{n+1} = 0$  because the  $r_i$  are strictly decreasing non-negative integers. We have,

$$\gcd(r_i, r_{i+1}) = \gcd(r_{i+1}q_{i+1} + r_{i+2}r_{i+1})$$
$$= \gcd(r_{i+2}r_{i+1})$$
$$= \gcd(r_{i+1}, r_{i+2})$$

Applying this result repeatedly,

$$\gcd(a, b) = \gcd(r_0, r_1)$$

$$= \gcd(r_2, r_3)$$

$$= \dots$$

$$= \gcd(r_{n-1}, r_n)$$

$$= r_n$$

Where the last equality is because  $r_n \mid r_{n-1}$ 

**Remark.** One can also use Euclids Algorithm to find the  $x, y \in \mathbb{Z}$  Bezouts Identity state to exist by working backwards. These aren't unique.

#### 1.4 Extended Euclidean Algorithm

Instead of doing Euclids, and working backwards we can compute our bezouts x, y during euclids. This is the extended Euclids Algorithm. This time we are going to define sequences of integers  $x_i$  and  $y_i$ , such that  $r_i = ax_i + by_i$ . Recall that  $r_n$  is the last non-zero remainder and that  $r_n = \gcd(a, b)$ . Therefore  $\gcd(a, b) = r_n = ax_n + by_n$  and so  $(x, y) := (x_n, y_n)$ .

We have that  $r_0 = a$  and  $r_1 = b$ . Hence, we see  $r_0 = 1 \times a + 0 \times b$  and  $r_1 = 0 \times a + 1 \times b$ , and so we set  $(x_0, y_0) := (1, 0)$  and  $(x_1, y_1) := (0, 1)$ . So, now we consider for  $i \ge 2$  we have a pair  $(x_j, y_j)$  for j < i. Then  $r_{i-2} = r_{i-1}q_{i-1} + r_i$  and so,

$$\begin{aligned} r_i &= r_{i-2} - r_{i-1}q_{i-1} \\ &= (ax_{i-2} + by_{i-2}) + (ax_{i-1} + by_{i-1})q_{i-1} \\ &= a(x_{i-2} - x_{i-1}q_{i-1}) + b(y_{i-2} - y_{i-1}q_{i-1}) \end{aligned}$$

Thus we set  $x_i := x_{i-2} - x_{i-1}q_{i-1}$  and  $y_i := y_{i-2} - y_{i-1}q_{i-1}$ . These can be defined recursively this way.

$$(x_i, y_i) := (x_{i-2}, y_{i-2}) - q_{i-1}(x_{i-1}, y_{i-1})$$

**Example.** We compute gcd(841, 160) use Extended Euclidean Algorithm.

i	$r_{i-2}$		$r_{i-1}$		$q_{i-1}$		$r_i$	$x_i$	$y_i$
0							841	1	0
1							160	0	1
2	841	=	160	×	5	+	41	1	-5
3	160	=	41	×	3	+	37	-3	16
4	41	=	37	×	1	+	4	4	-21
5	37	=	4	×	9	+	1	-39	205
6	4	=	1	×	4	+	0		

Therefore,  $gcd(841, 160) = 1 = 841 \times (-39) + 160 \times 205$ .

## 2 Primes and Congurences

We start by defining primes and composite numbers,

**Definition 2.1** (Prime). A number  $p \in \mathbb{N}_1$  with p > 1 is prime if and only if it's only divisors are 1 and p, i.e.

$$n \mid p \implies n = 1 \text{ or } n = p$$

**Definition 2.2** (Composite Numbers). A number  $n \in \mathbb{N}_1$  with n > 1 is composite if and only if it is not prime, i.e.

$$n = ab$$
  $1 < a, b \in \mathbb{N}$ 

One is neither composite nor prime.

**Proposition 2.3.** If  $n \in \mathbb{N}_1$  with n > 1, then n has a prime factor.

*Proof.* Use strong induction, so assume for 1 < m < n where  $m \in \mathbb{N}_1$  that m has a prime factor.

Case (i): If n is prime, then n is a prime factor of n.

Case (ii): If n is composite, then n = ab where a, b > 1 and so, 1 < a < n. By the induction hypothesis, there is a prime p such that  $p \mid a$ . Hence,  $p \mid a$  and  $a \mid n$  so, by transitivity  $p \mid n$ .

**Proposition 2.4.** If  $1 < n \in \mathbb{N}_1$ , then we can write  $n = p_1 p_2 \dots p_k$  where  $k \in \mathbb{N}_1$  and  $p_i$  are primes.

*Proof.* If n is prime, then the result is clear. So suppose that n is composite. Then n must have a prime factor, so  $n = p_1 n_1$  where  $1 < n_1 \in \mathbb{N}_1$ . If  $n_1$  is prime, we are done. If  $n_1$  is composite, then we can write  $n_1 = p_2 n_2$  and so on... This process terminates as  $n > n_1 > n_2 > \cdots > 1$ . Hence after at least n steps we obtain a prime factorisation of n.

Example.

$$666 = 3 \times 222 = 3 \times 2 \times 111 = 3 \times 2 \times 3 \times 37$$

**Theorem 2.5.** There are infinitely many primes

Euclid's Proof. For a contradition, assume there are finitely many primes,  $\{p_1, p_2, p_3, \ldots, p_n\}$  and that is a complete list. Consider  $N := p_1 p_2 \ldots p_n + 1 \in \mathbb{N}$ . Then N > 1 so by the first proposition, N has a prime factor p. However, every prime is one of the elements of the list, so  $p = p_i$ . Hence,  $p_i \mid (p_1 p_2 \ldots p_n)$  so  $p \mid (N-1)$ . However,  $p \mid N$  and we can write 1 = N - (N-1), so  $p \mid 1$ , which is a contradition.

#### 2.1 Fundemental Theorem of Arithmetic

**Lemma 2.6.** Let  $n \in \mathbb{Z}$ , then if  $p \nmid n$  then gcd(p, n) = 1

*Proof.* Let  $d = \gcd(p, n)$ . Then  $d \mid p$  so by definition of prime either d = 1 or d = p. But  $d \mid n$  so  $d \neq p$  because  $p \nmid n$ . Hence, d = 1.

**Theorem 2.7** (Euclid's Lemma for Primes). Let  $a, b \in \mathbb{Z}$  and p be a prime. If  $p \mid ab$ , then  $p \mid a$  or  $p \mid b$ .

*Proof.* Assume  $p \mid ab$  and that  $p \nmid a$ . We shall prove  $p \mid b$ . By Lemma, gcd(p, a) = 1, so by Euclids lemma,  $p \mid b$ .

Remark. Euclids Lemma for primes immediately generalises to several factors.

**Definition 2.8.** Let  $n \in \mathbb{N}_1$  and p be a prime. Then,

$$v_p(n) := \max\{k \in \mathbb{N} \cup \{0\} : p^k \mid n\}$$

In other words, k is the unique non-negative integer such that  $p^k \mid n$  but  $p^{k+1} \mid n$ . Equivalently,  $v_p(n) = k$  if and only if  $n = p^k n'$  where  $n' \in \mathbb{N}$  and  $p \nmid n'$ .

**Example.** We can see that,

- $-v_2(720) = 4 \text{ as } 2^4 \mid 720 \text{ but } 2^5 \nmid 720$
- $-v_3(720) = 2 \text{ as } 3^2 \mid 720 \text{ but } 3^3 \nmid 720$
- $-v_5(720) = 1 \text{ as } 5^1 \mid 720 \text{ but } 5^2 \nmid 720$
- if  $p \ge 7$ , then  $v_p(720) = 0$  as  $p \nmid 720$ .

**Lemma 2.9.** Let  $n, m \in \mathbb{N}_1$  and p be a prime. Then  $v_p(mn) = v_p(m) + v_p(n)$ 

Proof. Let  $k = v_p(m)$  and  $\ell = v_p(n)$ . Then we write  $m = p^k m'$  where  $p \nmid m'$  and  $n = p^\ell n'$  where  $p \nmid n'$ . Then  $nm = p^{k+\ell}m'n'$  and so by euclids lemma  $p \nmid m'n'$  as if it did then  $p \mid n'$  or  $p \mid m'$  but it doesn't. So  $v_p(mn) = v_p(m) + v_p(n)$ .

**Theorem 2.10** (Fundemental Theorem of Arithmetic). Let  $1 < n \in \mathbb{N}_1$ . Then,

- (i) (Existence) The number n can be written as a product of primes.
- (ii) (Uniqueness) Suppose that,

$$n = p_1 \dots p_r = q_1 \dots q_s$$

where each  $p_i$  and  $q_j$  are prime. Assume further that,

$$p_1 \le p_2 \le \dots \le p_r$$
 and  $q_1 \le q_2 \le \dots \le q_s$ 

Then r = s and  $p_i = q_i$  for all i

Remark. If 1 is a prime, then the Uniqueness here is broken, as,

$$6 = 3 \times 2 = 3 \times 2 \times 1 = \dots$$

**Remark.** A consequence of the FTA is that the integral domain  $\mathbb{Z}$  is in fact a UFD.

*Proof.* The existence is something we have done before. The harder part is uniqueness. Let  $\ell$  be any prime. Then we have,

$$v_e ll(n) = v_\ell(p_1 \dots p_r)$$
  
=  $v_\ell(p_1) + \dots + v_\ell(p_r)$ 

However,

$$v_{\ell}(p_i) = \begin{cases} 1 & \text{if } \ell = p_i \\ 0 & \text{if } \ell \neq p_i \end{cases}$$

Therefore,

$$v_{\ell}(n) = \#$$
 of  $i$  for which  $\ell = p_i$   
=  $\#$  of times  $\ell$  appears in the factorisation  $n = p_1 \dots p_r$ 

Similarly,

$$v_{\ell}(n) = \#$$
 of times  $\ell$  appears in the factorisation  $n = q_1 \dots q_s$ 

Thus every prime  $\ell$  appears the same number of times in each factorisation, giving the desired result.  $\Box$ 

**Remark.** Another way of interpreting this result is to say that for  $n \in \mathbb{N}_1$ ,

$$n = p_1^{v_{p_1}(n)} p_2^{v_{p_2}(n)} \dots p_r^{v_{p_r}(n)}$$

where  $p_1, \ldots, p_r$  are the distinct prime factors of n. Note that we take the empty product to be 1, which covers the case for n = 1.

**Lemma 2.11.** Let  $n = \prod_{i=1}^r p_i^{a_i}$  where each  $a_i \in \mathbb{N}_0$  and the  $p_i$ 's are distinct primes. The set of positive divisors of n is the set of numbers of the form  $\prod_{i=1}^r p_i^{c_i}$  where  $0 \le c_i \le a_i$  for  $i = 1, \ldots, r$ .

Proof. Exercise 
$$\Box$$

### 2.2 Congurences

**Definition 2.12.** Suppose  $a, b \in \mathbb{Z}$  and  $n \in \mathbb{N}_1$ . We write  $a \equiv b \mod n$ , and say 'a is congruent to b mod n', if and only if  $n \mid (a-b)$ . If  $n \nmid (a-b)$  we say that a and b are incongruent mod n.

**Remark.** In particular,  $a \equiv 0 \mod n$  if and only if  $m \mid a$ 

**Example.** Here are some examples:

- $-4 \equiv 30 \mod 13 \text{ since } 13 \mid (4-30) = -26$
- $-17 \not\equiv -17 \mod 4 \text{ since } 17 (-17) = 34 \text{ but } 4 \nmid 34.$
- -n is even if and only if  $n \equiv 0 \mod 2$
- n is odd if and only if  $n \equiv 1 \mod 2$
- $-a \equiv b \mod 1 \text{ for all } a, b \in \mathbb{Z}$

**Proposition 2.13.** Let  $n \in \mathbb{N}_1$  being congruent mod n is an equivience relation, so,

- (i) Reflexive:  $\forall a \in \mathbb{Z}, a \equiv a \mod n$
- (ii) Symmetric:  $\forall a, b \in \mathbb{Z}, a \equiv b \mod n \implies b \equiv a \mod n$
- (iii) Transitive:  $\forall a, b \in \mathbb{Z}, a \equiv b \mod n \text{ and } b \equiv c \mod n \implies a \equiv c \mod n.$

*Proof.* The proof follows from,

- (i)  $n \mid 0$ .
- (ii) If  $n \mid (a-b)$  then  $n \mid (b-a)$
- (iii) If  $n \mid (a b) + (b c) = (a c)$

**Proposition 2.14.** Congruences respect addition, subtraction and multiplication. Then let  $a, b, \alpha, \beta \in \mathbb{Z}$ . Suppose that  $a \equiv \alpha \mod n$  and  $b \equiv \beta \mod n$ . Then,

- (i)  $a + b \equiv \alpha + \beta \mod n$
- (ii)  $a b \equiv \alpha \beta \mod n$
- (iii)  $ab \equiv \alpha\beta \mod n$

Moreover, if  $f(x) \in \mathbb{Z}[x]$  then  $f(a) \equiv f(\alpha) \mod n$ 

*Proof.* Check that  $ab \equiv \alpha\beta \mod n$ . Since,  $a \equiv \alpha \mod n$  and so,  $n \mid (a - \alpha)$  and so  $a = \alpha + ns$  for some  $s \in \mathbb{Z}$ , Similarly  $b = \beta + nt$ . Hence,

$$ab = (\alpha + ns)(\beta + nt) = \alpha\beta + n(s\beta + t\alpha + nst)$$

and so  $n \mid (ab - \alpha\beta)$ . Therefore,  $ab \equiv \alpha\beta \mod n$ , as required.

**Example.** Let  $n \in \mathbb{N}_1$  and write n in decimal notation,

$$n = \sum_{i=0}^{k} a_i \times 10^i \qquad 0 \le a_i \le 9$$

Then, define f(x) by,

$$f(x) = \sum_{i=0}^{k} a_i x^i$$

Then, since  $10 \equiv -1 \mod 11$ , we see that  $n = f(10) \equiv f(-1) \mod 11$ , whence,

$$11 \mid n \iff 11 \mid f(-1) \iff 11 \mid (a_0 - a_1 + a_2 - a_3 + \dots + (-1)^k a_k)$$

This is an easy way to test for divisibility by 11.

**Example.** Does  $x^2 - 3y^2 = 2$  have a solution with  $x, y \in \mathbb{Z}$ . Let  $x, y \in \mathbb{Z}$ . Note that  $x^2 - 3y^2 \equiv x^2 \mod 3$ . Now,  $x \equiv 0, 1, 2 \mod 3$ , so  $x^2 \equiv 0, 1, 4 \mod 3 \equiv 0, 1 \mod 3$ . Hence,  $x^2 - 3y^2 \equiv x^2 \not\equiv 2 \mod 3$  and so  $x^2 - 3y^2 \not\equiv 2$ .

**Remark.** Suppose we have  $f \in \mathbb{Z}[x_1, \ldots, x_m]$  if we have  $a_1, \ldots, a_m \in \mathbb{Z}$  such that  $f(a_1, \ldots, a_m) = 0$  then  $f(a_1, \ldots, a_m) \equiv 0 \mod n$  for every  $n \in \mathbb{N}$ . Therefore if there exist an  $n \in \mathbb{N}_1$  such that  $f(x_1, \ldots, x_m) \equiv 0 \mod n$  has no solution, there there cannot exist  $a_1, \ldots, a_m \in \mathbb{Z}$  such that  $f(a_1, \ldots, a_n) = 0$ .

We are going to prove the following theorem,

**Theorem 2.15.** There are infinitely many primes p with  $p \equiv 3 \mod 4$ 

*Proof.* Suppose that p is a prime. Then  $p \equiv 0, 1, 2, 3 \mod 4$ , but  $p \not\equiv 0 \mod 4$  because  $4 \nmid p$ . If  $p \equiv 2 \mod 4$  then p = 4k + 2 for some  $k \in \mathbb{Z}$ , so  $2 \mid p$  so in fact p = 2. Therefore there are three types of primes,

- (i) p = 2
- (ii)  $p \equiv 1 \mod 4$
- (iii)  $p \equiv 3 \mod 4$

Let  $N \in \mathbb{N}$  it suffices to show that there exist a type (iii) prime with p > N. Let 4(N!) - 1 and so  $M \ge 3$  and so by the existence of FTA we can write  $M = p_1 \dots p_k$ . If  $p \le N$ , then  $M \equiv -1 \mod p$  so  $p \nmid M$ . Hence,  $p_j > N$  for all j. Moreover  $p_j \ne 2$  for all j because M is odd. Therefore for each j we have  $p_j \equiv 1, 3 \mod 4$ . If  $p_j \equiv 3 \mod 4$  for any j then we are done. If this is not the case, then  $p_j \equiv 1 \mod 4$  for all j, and so,  $M \equiv 1 \times 1 \times \cdots \times 1 \mod 4 \equiv 1 \mod 4$ ; but by defintion of M we have  $M \equiv -1 \equiv 3 \mod 4$  contradition!

**Remark.** Congruences do not repsect division,  $4 \equiv 14 \mod 10$  but  $2 \not\equiv 7 \mod 10$ 

**Proposition 2.16.** Let  $a, b, s \in \mathbb{Z}$  and  $d, n \in \mathbb{N}_1$ .

- (i) If  $a \mid b \mod n$  and  $d \mid n$  them  $a \mid b \mod d$
- (ii) Suppose  $s \neq 0$ . Then  $a \equiv b \mod n$  if and only if  $as \equiv bs \mod ns$

*Proof.* (i) follows from transitivity of divisibility;

(ii) follows from multiplication and cancellation properties.

**Theorem 2.17** (Cancellation law for Congruences). Let  $a, b, c \in \mathbb{Z}$  and  $n \in \mathbb{N}_1$ . Let  $d = \gcd(c, n)$ . Then  $ac \mid bc \mod n \iff a \equiv b \mod \frac{n}{d}$ . In particular, if n and c are coprime, then  $ac \equiv bc \mod n \iff a \equiv b \mod n$ .

*Proof.* Since,  $d = \gcd(c, n)$ , we may write n = dn' and c = dc' where  $n', c' \in \mathbb{Z}$ . Suppose  $ac \equiv bc \mod n$ . Then  $n \mid c(a - b)$  and so  $n' \mid c'(a - b)$ . However,  $\gcd(n', c') = 1$  and so  $n' \mid (a - b)$  by Euclids Lemma. Thus,  $a \equiv b \mod n'$ .

Suppose conversely  $a \equiv b \mod n'$  and so,  $n' \mid (a-b)$  and so  $n \mid d(a-b)$ . But  $d \mid c$  and so  $d(a-b) \mid c(a-b)$  and thus  $n \mid c(a-b)$  by the transitivity of divisibility. Thus  $ac \equiv bc \mod n$ .

**Proposition 2.18.** Let  $a, m, n \in \mathbb{Z}$ . If m and n are coprime and if  $m \mid a$  and  $n \mid a$  then  $nm \mid a$ .

*Proof.* Since  $m \mid a$  we can write a = mc for some  $c \in \mathbb{Z}$ . Now  $n \mid a = mc$  and  $\gcd(m, n) = 1$  and so by Euclids Lemma,  $n \mid c$ . Hence,  $mn \mid mc = a$ .

Corollary 2.19. Let  $m, n \in \mathbb{N}$  be coprime and let  $a, b \in \mathbb{Z}$ . If  $a \equiv b \mod m$  and  $a \mid b \mod n$  then  $a \equiv b \mod mn$ .

*Proof.* We have  $n \mid (a-b)$  and  $m \mid (a-b)$ . Since m and n are coprime we therefore have  $mn \mid (a-b)$ .  $\square$