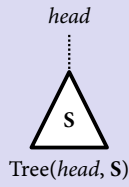


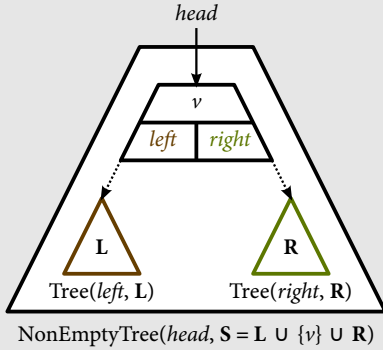
Precondition: *head* is a tree representing some set *S*. We are inserting *value*.



Is *head* null?

head \neq null

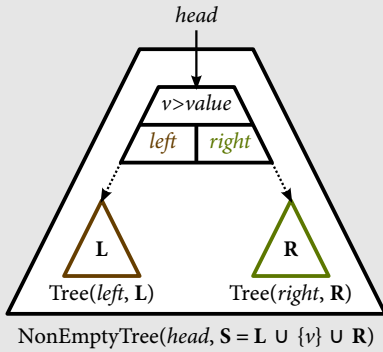
head points to a non-empty tree with some *v* at the root and two, possibly empty, subtrees.



compare(*value*, *v*)

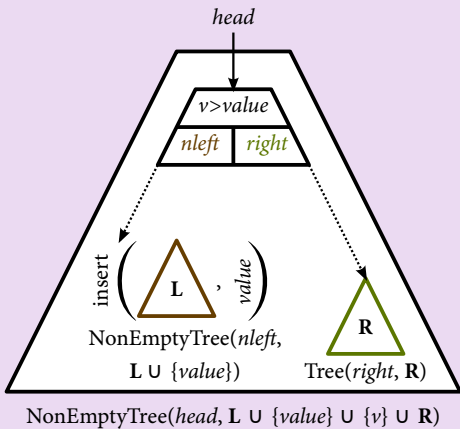
value < *v*

We cannot insert into the *right* tree, because that would break $\forall r \in R. v < r$.

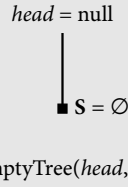


Recursively call insert(*left*, *value*), yielding *nleft*; set *left* field to *nleft*.

Tree represents $L \cup \{value\} \cup \{v\} \cup R$, which $= S \cup \{value\}$. Return *head*.

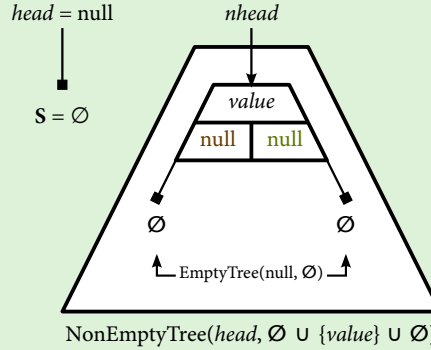


The tree at *head* is empty, and represents the set \emptyset . We can return a one-node tree.



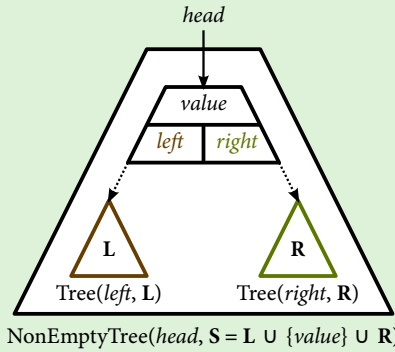
Construct one-node tree containing *value*, pointed at by new variable *nhead*.

We require $\emptyset \cup \{value\} = \{value\}$; new tree is $\emptyset \cup \{value\} \cup \emptyset = \{value\}$. Return *nhead*.



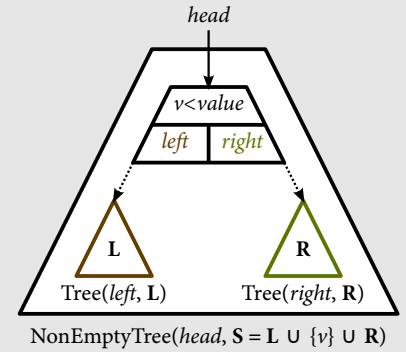
value = *v*

value $\in \{v\}$, so *value* $\in L \cup \{v\} \cup R$, so *value* $\in S$, so $S \cup \{value\} = S$. Return *head*.



value > *v*

We cannot insert into the *left* tree, because that would break $\forall l \in L. v > l$.



Recursively call insert(*right*, *value*), yielding *nright*; set *right* field to *nright*.

Tree represents $L \cup \{v\} \cup R \cup \{value\}$, which $= S \cup \{value\}$. Return *head*.

