Thesis Notation

$March\ 28,\ 2011$

Table 1: Types

Macro	Output	Description
\type{0} \sst{0}{1}	0 0 ₁	Type macro Subscripted type
\Talpha \Tholes \Thvars \Tpvars \Taddr \Tval \Tralph \Tcalph	Σ X Θ P Addr Val Υ Ω	Alphabet Hole Alphabet Hole Variable Alphabet Propositional variables Heap addresses Values Ranked alphabet $\Sigma \cup X$
\Ttree \Tseq \Theap \Tterm \Tutree	Tree Seq Heap Term UTree	Trees Sequences Heaps Terms Uniquely-labelled trees
\Tword \Tfor	Ω^* Forest $_{\Sigma, X}$	Words Forests
\Tlist \Tlsto	Lst LStore	Lists List stores
\Tdata	D	Abstract states
\Tcont \Ttcont \Tscont \Ttercont \Tutcont	C C _{Tree} C _{Seq} C _{Term} C _{UTree}	Contexts Tree Contexts Sequence Contexts Term Contexts Uniquely-labelled tree contexts
\Tmcont \Tmtcont \Tmscont \Tmtercont	C^m C^m_{Tree} C^m_{Seq} C^m_{Term}	Multi-holed Contexts Multi-holed Tree Contexts Multi-holed Sequence Contexts Multi-holed Term Contexts

Table 1: Types

_	Macro	Output	Description
-	\TLEnv \TInterp	LEnv Interp	Logical environments Propositional interpretations
	\Tworld \Tsort \Tformula	World Sort Formula	Worlds Sorts Formulae
	\Trank	Rank	Ranks
	\DW \SW	DW SW	Duplicator-winning games Spoiler-winning games
	Mathematics	3	ı
	\Tnat \Tint \Tbool	N ℤ Bool	Natural numbers (includes 0) Integers Booleans
	Formulae	1	
	\TF \TFd \TFct \TFsd \TFsc \TFsct \TFsct \TFcd \TFcc \TFcd \TFcc \TFct \TFct \TFct \TFct	$\begin{array}{c} F \\ P \\ K \\ P_{Tree} \\ K_{Tree} \\ P^s \\ K^s \\ P^s_{Tree} \\ K^s_{Tree} \\ P^c \\ K^c \\ P^c_{Tree} \\ K^c_{Tree} \\ K^m_{Tree} \\ K^m_{Tree} \\ K^m_{Seq} \\ K^m_{Term} \end{array}$	(Propositional) formulae Data formulae Context Formulae Tree Tree context Specifically single holed With composition Multi-holed
	Programs		
	\Tprog \Tcommand	Prog Cmd	Programs Basic command
	\Tpgvar \Tscope \Tstack \Tstore \Tstate \Toutcome	Var Scope Stack Store State Outcome	Program variables Variable scope Variable stack Data store Program state Outcomes
	\Tpname	PName	Procedure names

Table 1: Types

Macro	Output	Description
\Tpdef	PDef	Procedure definition environment
\Texpr \Tbexp	Expr BExp	Expressions Boolean expressions
CAP		-
\TLVar	LVar	Logical variables
\TLState	LState	Logical state
\TSState	SState	Shared state
\TPerm	Perm	Permission
\TTok	Token	Token
\TRid	RID	Region identifier
\TAName	AName	Action name
\TAction	Action	Action
\TAMod	AMod	Action model
\TAPName	APName	Abstract predicate name
\Tassn	Assn	Assertion
\Tbassn	BAssn	
\Tiassn	lAssn	
\Tfassn	FAssn	
\Tinterp	Interp	
\Tpenv	PEnv	
\Tfenv	FEnv	

Table 2: Variables

Macro	Output	Description
\la \lb \lc	a b c	Letters of Σ
\hx \hy \hz \hxP \hyP \hzP	$egin{array}{c} x \\ y \\ z \\ \dot{x} \\ \dot{y} \\ \dot{z} \end{array}$	Holes of X
\hva \hvb \hvc	$egin{array}{c} lpha \ eta \ \gamma \end{array}$	Variables of Θ
\hpa	a	Heap addresses
\va	v	Values

Table 2: Variables

Macro	Output	Description
\vb	w	
\vc	u	
,		D ::: 1 : 11
/pvp	p	Propositional variables
\pvq	q	
\tr	t	Trees
\sq	s	Sequences
\hp	h	Heaps
\tm	r	Terms
\lst	l	Lists
\lstc	lc	
\ls	ls	
\lsc	lsc	
\word	w	Words (confusing with worlds?)
\state	s	Abstract state
\cont	c	Contexts
\contd	d	
\cidc	i	Element of I
\lenv	σ	Logical Environments of LEnv
\interp	ι	Propositional Interpretation of Interp
\ca	\mathcal{A}	Context Algebra
\sor	ς	Sort
\rank	r	Rank
\rn	m	Non-adjunct component
\rs	s	Adjunct component
\rL	$\stackrel{\circ}{L}$	Label component
\rV	V	Variable component
\± v	,	
\world	$w_{\underline{c}}$	World
\swor	w^{S}	Spoiler's world
\dwor	w^{D}	Duplicator's world
\opvar	*	Operator
\tv	b	Truth value
Programs	S	
\cmd	\mathbb{C}	Program
\bcmd	φ	Basic command
\pdef	μ	Procedure definition
\penv	γ	Procedure definition environment
	ı '	

Table 2: Variables

Macro \stack	Output	
	Output	Description
\	S	Stack
\scope	ρ	Scope
\stor	χ	Store
\expr	E	Expression
\bexp	B	Boolean expression
\oc	o	Outcome
\pse	Г	Procedure specification environment
Logical f	ormulae	
\Fp	F	Propositional formula
\Fdv	P	Data formula
\Fdva	Q	
\Fdvb	R	
\Fcv	K	Context formula
\Fo	P	Arbitrary formula
\FSd	S_{P}	Specific data formula
\FSc	S_{K}	1
\FS	S	
Abstract	modules	
\ama	A	Abstract module
\amb	B	Tibstract module
	עוו	
CAP	I	I
\lstate	l	Logical state
\sstate	s	Shared state
\rid	r	Region identifier
\aname	γ	Action name
\perm	ϕ	Permission
\amod	ζ	Action model
\tok	t	Token
\action	a	Action
\apname	α	Abstract predicate name
	I	Interference assertion
\ias		
\ias \ridexp	R	Region expression
	$R \atop \pi$	Region expression Permission expression

Table 3: Constants

Macro	Output	Description
\nullref	nil	null value
\empt \hole	Ø -	Empty tree Context hole
\truth \falsity	T F	Truth

Table 4: Operators

Macro	Output	Description
\tpar \tlab{0}{1} \scat \cell{0}{1} \hsep \hemp \lcell{0}{1}		Tree parallel Tree label Sequence concatenation Heap cell Heap combination Empty heap List cell
\app \comp \mcc{0} \subs{0}{1}	• © [0/1]	Context application Context composition Multi-holed composition Substitution
\cid \czero	I 0	Identity of a context algebra Zero of a context algebra
\holes \dom \range \freevars	holes dom range fv	Holes function
\fresh	#	Fresh
Mathematics		
\pfun \fpfun	⊸ ⊸ _{fin}	Partial function Finite partial function
\fupd{0}{1}{2}	$0[1\mapsto 2]$	(Partial) function update
<pre>\powset{0} \fpowset{0} \Set{0} \Setb{0}{1} \union \intersect \Card{0}</pre>	$ \begin{array}{c} \mathcal{P}(0) \\ \mathcal{P}_{\mathrm{fin}}(0) \\ \{0\} \\ \{0 \mid 1\} \\ \cup \\ \cap \\ 0 \end{array} $	Powerset Set Set builder Union Intersection Cardinality

Table 4: Operators

Macro	Output	Description
\disjunion	₩	Disjoint union
\Family{0}{1}{2} \family{0}{1}{2}	$ \begin{cases} {0_2} \\ {2 \in 1} \\ {0} \\ {1 \in 2} \end{cases} $	Family Alternate macro
\randthen	9	Relation composition
Logical formulae		
\Ftzero \Ftpar \Ftlab{0}{1} \Fscat	0 ⊗ 0[1]	Tree zero Tree parallel Tree label Sequence concatenation
\Fapp \Fappr \Frapp \FapprE \FrappE	•- •-∃ ∃	Application Application adjoint (tree formula) Application adjoint (context formula) Existential versions of adjoints
\Fid	I	Context identity
\Fcomp \Fcompr \Frcomp \FcomprE \FrcompE	° ~ ~ ~ ~	Composition Composition adjoints Existential versions of adjoints
\Fimp \Fand \For \Fnot \Fiff	→	Material implication Conjunction Disjunction Negation Material equivalence
\Fdtt \Fdff \Fctt \Fcff \Ftt \Fff	true false True False T ⊥	Data true Data false Context true Context false Propositional true Propositional false
\new	И	Fresh quantification
\Fcell{0}{1} \Femp \Fsep \Fwand \FwandE	$0 \mapsto 1$ emp $*$ $-*$ $-*$ $-*$	Heap cell Separating conjunction Magic wand
\Fstarprod{0}	\prod_0^*	Star Product
\bigoast	*	

Table 4: Operators

Macro	Output	Description
\Fstarquant{0}	※ 0.	Star Quantification
\Fnoh{0} \Fahole	∅	No hole Just a hole
\Fswsc	•	Strongest weaker stable
\Fswsb \Fwssb	҈ ҈ ℝ	Weakest stronger stable
\Fgdiab	\$	
\Fgdia \Fsws \Fwss	© © ® R	
\shared{0} \Fshare{0}{1}{2} \shareds{0} \Fshares{0}{1}{2}	$\begin{bmatrix} 0 \\ 0 \\ 2 \\ 0 \\ 0 \end{bmatrix}_2$	
\actquant \actto \actap{0}{1}{2}{3}{4} \acta{0}{1}{2}{3} \actan{0}{1}{2}{3} \actan{0}{1}{2}	$\begin{array}{c} \exists \\ \leadsto \\ 0(1): \ \exists 2. (3 \leadsto 4) \\ 0(1): \ 2 \leadsto 3 \\ 0: \ 1 \leadsto 2 \\ 0: \ \exists 1. (2 \leadsto 3) \end{array}$	
\Fall \Fperm{0}{1}{2}{3} \Fpermn{0}{1}{2}	$\begin{array}{c} \operatorname{all} \\ \left[0(1)\right]_{3}^{2} \\ \left[0\right]_{2}^{1} \end{array}$	
\upimp \Fupimp{0}{1}	$\overset{\Longrightarrow}{\Longrightarrow}_{\{0\}\{1\}}$	
\Fup{0}{1}	$\left[rac{0}{1} ight>$	
\charac{0}{1}	D_1^0	Rank-r characteristic of a world

Table 5: Relations

Macro	Output	Description
\mods	⊨ ⊭	Models Doesn't model
\entails \lequiv \nlequiv	⊆ ≡ ≢	Entailment? Equivalence

Table 5: Relations

Macro	Output	Description
\rankwfr	⊲	Rank well-founded relation
$\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ $	\mathfrak{R}_0	Rank constructor relation
$move{0}$	\mathfrak{M}_0	Game move relation
\spo	\leq	Simulation preorder

Table 6: Misc

Macro	Output	Description
\Gdef \Gbar	::=	BNF notation
\eqdef	def =	'is defined to be'
\vect{0}	$\overrightarrow{0}$	
$\sm(0)$	[[0]]	Semantics
Meta syntax	ı	ı
\IFF \IMPLIES \OR \AND \NOT \EXISTS{0} \EXIST{0} \FORALL{0}	$\iff \Longrightarrow$ or and not there exists 0 s.t. there axist 0 s.t. for all 0,	If and only if Implies Or And Not There exists There exist For all
\undef	undefined	Undefined
\blank	(\cdot)	Blank
Logic notation	ns	ı
\CLs \CLc \CLm	CL CL^s CL^c CL^m	
Text macros	ı	ı
\ie \etc \etal \eg	$i.e.$ $etc.$ $et\ al.$ $e.g.$	
Games stuff	I	I
\Spoiler \Duplicator	Spoiler Duplicator	

Table 6: Misc

Macro	Output	Description
Sorts		
\mksort{0}	0	
\Sdata	d	
\Scont	С	

Table 7: Automata

Macro	Output	Description
\Lang \Langi{0}	$\mathcal{L} \ \mathcal{L}_0$	Language
\nil	ε	Epsilon
\1b1	a	Label
\1bm	b	
\aut	\mathcal{A}	Automaton
\States	Q	State set
\Accs	A	Accepting states
\st	q	State
\inist	e	Initial state
\tra{0}	f^0	Transition relation
\nilclose	arepsilon-closure	ε -closure
\auti{0}	\mathcal{A}_0	Automaton
\Statesi{0}	Q_0	State set
\Accsi{0}	A_0	Accepting states
\sti{0}	q_0	State
\inisti{0}	e_0	Initial state
\trai{0}{1}	f_1^0	Transition relation
$\nilclosei{0}$	ε -closure $_0$	ε -closure
\hataut	$\hat{\mathcal{A}}$	Partial automata
\hatauti{0}	$\hat{\mathcal{A}}_0$	
\aint{0}	$\llbracket 0 \rrbracket_{\mathcal{A}}$	Induced mappings
$\aintw{0}{1}$	$[0]_1$	
\afint{0}	$(0)_{\mathcal{A}}$	
$\afintw{0}{1}$	$(0)_1$	
\acompl{0}	$\overline{0}$	Complementation
\nlsub	0	Non-deterministic linear substitution
\nllef	⊘ ∃	
\nlrig	_⊘∃	
\usub	\odot	Uniform substitution
\ulef	∃	
\urig	_⊙∃	

Table 7: Automata

Macro	Output	Description
\nusub	0	
\nurig	-⊚∃	
\reach	reachable	Reachability

Table 8: Programs

Macro	Output	Description
\psyntax{0}	0	Program syntax
\pvs{0}	0	Program variable
$\protect\pro$	0	Procedure name
Expressions		ı
\eP	+	
\eM	_	
\eT	*	
\eLT	<	
\eEQ	=	
\eIMP	=>	
\eFALSE	false	
Programs		
\pskip	skip	Skip
\pas{0}{1}	0 := 1	Assign
\seq	:	Sequential composition
\pifelse{0}{1}{2}	if 0 then 1 else 2	If-then-else
(P110100003 (13 (13)	if 0 then	
	1	
$\left(0\right)_{1}_{2}$	else	
	2	
	if 0 then	
\pifv{0}{1}	1	
\pwhile{0}{1}	while 0 do 1	While-do
\h-il(0)(1)	while 0 do	
<page-header></page-header>	1	
\procdef{0}{1}{2}{3}	$0 := 1(2) \{3\}$	Procdef
_	$0 := 1(2)$ {	
\procdefv{0}{1}{2}{3}	3	
•	}	
	$0 := 1(2)$ {	
\procdefvb{0}{1}{2}{3}	3	
<u> </u>	}	
	0(1) {	
\procdefvc{0}{1}{2}	$\frac{1}{2}$	
.r	}	
)	I

Table 8: Programs

Macro	Output	Description
\procs{0}{1}	procs 0 in 1 procs 0	Procedures
\procsv{0}{1}	in 1	
\pcal1{0}{1}{2}	call 0 := 1(2)	Procedure call
\pcallo{0}{1} \plocal{0}{1}	$\begin{array}{c c} \operatorname{call} 0(1) \\ \operatorname{local} 0 \text{ in } 1 \end{array}$	Local variable declaration
_	local 0 in	Local variable declaration
\plocalv{0}{1}	1	
$\alloc{0}{1}$	$\mathit{0} := \mathtt{alloc}(1)$	Allocate
$\displaystyle \dispose\{0\}\{1\}$	$\mathtt{dispose}(0,1)$	Dispose
\sto{0}{1}	[0] := 1	Store
$\left\{0\right\}\left\{1\right\}$	0 := [1]	Fetch
Scope	I	I
\sceq	\Rightarrow	Scope value
\scval{0}{1}	$0 \Rightarrow 1$	
\scc	*	Scope combination
\scemp	Ø	Empty scope
Semantics	ı	
\tosco{0}	[0]	Stack o Scope
$\ensuremath{\tt (0)}$	$\mathcal{E} \llbracket 0 rbracket$	$Expr \to (Scope \to Val)$
$\protect\$	$\mathcal{P} \llbracket 0 rbracket$	$BExp \to \mathcal{P}(Scope)$
$\begin{tabular}{l} \begin{tabular}{l} tabu$	$\mathcal{B} \llbracket 0 rbracket$	$ BExp \to (Scope \to 2)$
\csem{0}	$\mathcal{C} \llbracket 0 rbracket$	$Cmd \to (Scope \to \mathcal{P}(Scope) \cup \{ \not \xi \})$
\bigto	~→	Big step
\nbigto	→	
\fault	\$	Fault
\writeVar	writeVar	
\writeVars	writeVars	
\lookup	lookup	
\axioms{0}	Ax [[0]]	
Abstract modules	I	
$\am\{0\}$	¥	
\admH	H	Heap module
$\arraycolsep=0.00000000000000000000000000000000000$	\mathbb{T}	Tree module
\amL	\mathbb{L}	List module
Module translations		
\trl	au	Translation

Table 8: Programs

Macro	Output	Description
\ar	α	Abstraction relation
$\displaystyle \min\{0\}$	[0]	Implementation
\prtr{0}	[0]	Predicate translation
\TIF	\mathcal{I}	Interface set
\TIFi	$\mathcal{I}_{ ext{in}}$	
\TIFo	$\mathcal{I}_{ ext{out}}$	
\intf	I	Interface
\intfi	in	
\intfo	out	
\TCP	\mathcal{F}	Crust parameter set
\crup	F	Crust parameter
\drep{0}{1}	$\langle\langle 0 \rangle\rangle^1$	Data rep
\crep{0}{1}{2}	$\langle\langle 0\rangle\rangle_2^1$	Context rep
\crust{0}{1}	\mathbb{n}_0^{1}	Crust
\rsem{0}	(0)	The Description
\idpt{0}{1}	$(0)^1$	Intermediate Data Translation
$\left(1\right) \left(1\right) \left(2\right) $	$(0)_2^1$	Intermediate Context Translation
Predicates	1	
\vsafe	vsafe	
\bsafe	bsafe	
Hoare logic	,	'
\triple{0}{1}{2}	{0} 1 {2}	Hoare triple
\vtriple{0}{1}{2}	{0} 1 {2}	
\pspec{0}{1}{2}	$0:1\rightarrowtail 2$	
Hoare logic rules		
\mkrule{0}	0	
\raxiom	AXIOM	Axiom
\rframe	Frame	Frame
\rcons	Cons	Consequence
\rdisj	Disj	Disjunction
\rskip	SKIP	Skip
\rseq	SEQ	Sequencing
\rif	IF	lf o
\rwhile	WHILE	While
\rassgn	Assgn	Assignment
\rlocal	Local	Local variable
\rpdef	PDEF	Procedure definition
\rpcall	PCALL	Procedure call
\rpwk	PWK	Procedure weakening

Table 8: Programs

Macro	Output	Description
\rconj	Conj	Conjunction

Table 9: Cap

		~··r
Macro	Output	Description
\wf	wf	Well-formedness
\stable{0}	stable(0)	Stability
\locp{0} \shap{0} \hpp{0} \pmp{0} \stap{0} \actp{0}	$\begin{array}{c} 0_{\rm L} \\ 0_{\rm S} \\ 0_{\rm H} \\ 0_{\rm P} \\ 0_{\rm R} \\ 0_{\rm A} \end{array}$	Local portion Shared portion Heap portion Permission portion State portion Action portion
\amplus	Ц	Action model combination
\psep \zperm \lssep \wsep \zworld	$egin{array}{c} egin{array}{c} \egin{array}{c} \egin{array}{c} \egin{array}{c} \egin{array}{c} \egin{array}{c} \egin{array}$	Permission combination Zero permission Logical state combination World combination Zero world
\act{0} \abp{0}	0	Action name Abstract predicate
\lcol{0} \hcol{0}	[0] [0] _H	LState collapse Heap collapse
\Guar \Guarc \Guarr \Guars	G G_c \overline{G} \widehat{G}	Guarantee relation Construction operation Repartitioning guarantee Step-close guarantee
\Rely \Relyc	$\begin{array}{c} \mathrm{R} \\ \mathrm{R}_c \end{array}$	Rely relation Construction
\PRED \PREDq \PREDr \pred \predq	$egin{array}{c} P \ Q \ R \ p \ q \end{array}$	Predicates
\spredq	Q	
\pde	Δ	Abstract predicate axioims
\len	len	Vector length

Table 9: Cap

Macro	Output	Description
\fpex{0}{1}{2}{3}	$0 \uparrow 1, 2, 3$	Fixed-point Existence
Programming langua	ge	ı
\atomic{0} \ppar \letin{0}{1} \pstmts	$ \begin{vmatrix} \langle 0 \rangle \\ \parallel \\ let \ 0 \ in \ 1 \\ c \end{vmatrix} $	
$\sim {0}{1}$	$safe_0(1)$	
Proof rules		
\ratomic	Atomic	Atomic
\rprim	Prim	Primitive
\rpar	Par	Parallel
\rguar1	Guar-L	Guarantee left
\rguarr	Guar-R	Guarantee right
\rpredi	Pred-I	Predicate introduction
\rprede	Pred-E	Predicate elimination
\rlet	LET	Let
\rcall	CALL	Call
\rloop	Loop	Loop
\rchoice	Сноісе	Choice
\rexist	EXIST	Existential