

Side-Channel Attacks on NIST PQC 3rd Round Candidates

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Introduction

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NIST PQC 3rd Round Candidates

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Work on PQC SCA (and
countermeasures) may continue
for decades..

01

Introduction

Why do we care about attacks on PQC?



NIST PQC Timeline so far

- Feb 2016 – NIST PQC “Competition” announced; CFP posted.
- Dec 2017 – 1st Round begins; 69 submissions accepted (5 withdraw).
- April 2018 – 1st NIST PQC Standardization workshop (@ Florida Atlantic).
- Jan 2019 – Round 2 candidates announced (17 KEM/PKEs + 9 Signatures).
- Aug 2019 – 2nd NIST PQC Standardization workshop (@ CRYPTO).
- July 2020 – Round 3 finalists announced (4 KEMs + 3 Signatures), plus 8 alternates.
- June 2021 – 3rd NIST PQC Standardization workshop.
- October 2021 – Cut off point for any new results on candidates for NIST’s consideration.
- ~2022-2024 – Draft standards for public comment.

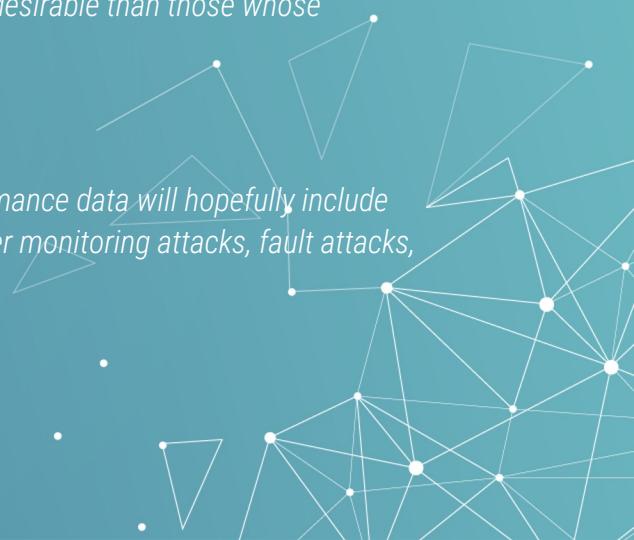


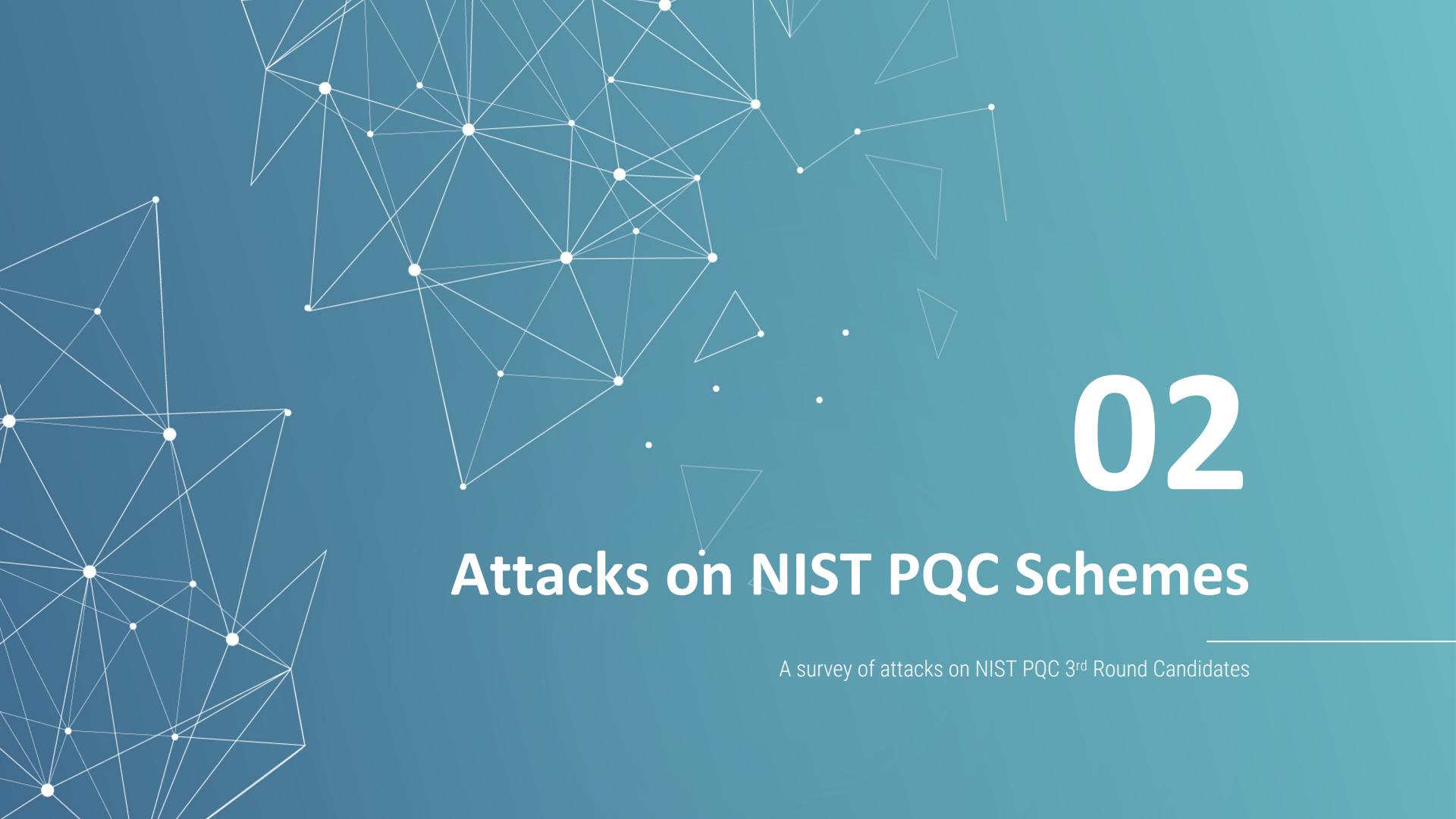
Welcome to PQC Side-Channel Attacks

This is a talk about side-channel attacks (SCA) against NIST Post-Quantum Cryptography candidates.

NIST has repeatedly stated the importance of SCA and countermeasures:

- From the original NIST PQC call for proposals in 2016:
"Schemes that can be made resistant to side-channel attacks at minimal cost are more desirable than those whose performance is severely hampered by any attempt to resist side-channel attacks."
- To the latest PQC summary document (NISTIR 8309):
"NIST hopes to see more and better data for performance in the third round. This performance data will hopefully include implementations that protect against side-channel attacks, such as timing attacks, power monitoring attacks, fault attacks, etc."



The background of the slide features a complex, abstract network graph composed of numerous white dots (nodes) and connecting lines (edges). The nodes vary in size, creating a sense of depth and connectivity. Some nodes are highlighted with a larger size, while others are smaller. The edges are thin white lines that intersect to form a dense web of triangles and other polygons across the teal-colored background.

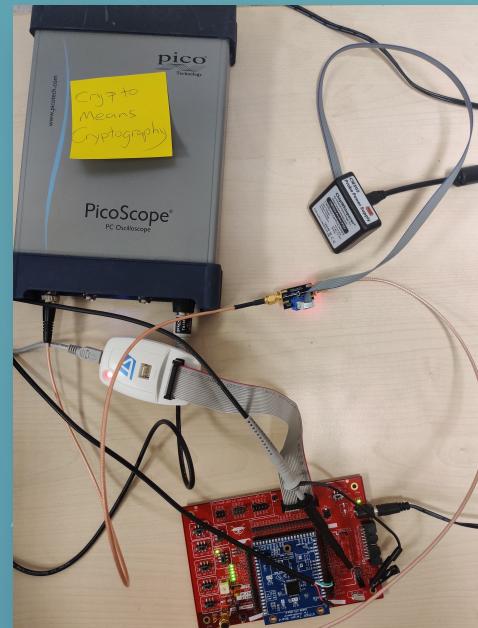
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Attacks on NIST PQC Schemes

A survey of attacks on NIST PQC 3rd Round Candidates

Types of Attacks Considered

- **Classical cryptanalysis** mathematically analyzes a cryptosystem.
- **Timing analysis** exploits variable runtime of an algorithm.
- **Fault attacks** are semi-invasive methods to intentionally induce faults to reveal cryptographic internal states.
- **Simple, differential, correlation power analysis** non-invasively exploits variations in power consumption of a cryptographic algorithm.
- **Electromagnetic attacks** exploit radiation from a cryptographic algorithm.
- **Template attacks** profiles a sensitive device to gain access to the secret.
- **Cold-boot attacks** exploit memory remanence to read data out of a computer's memory after the computer has been powered off.
- **Countermeasures** protect/hinder attacks via hiding or masking methods.



A Disclaimer to the Survey

The purpose of this talk is to motivate more attacks.

By showing gaps in the state-of-the-art.

We try to remain as unbiased/neutral as possible:

- Presenting papers that directly attack a candidate.
- Thus, we only focus on KEMs and signatures.
- We try to assume implementations are correct.
- Oracle timings attacks [GJN20], etc., *should* be fixed by now.



A Summary of the Attacks on NIST PQC Candidates

Table1: Attacks on NIST PQC third round candidates categorized by classical cryptanalysis (**CC**), static timing analysis (**STA**), fault attacks (**FA**), simple power analysis (**SPA**), advanced (correlation/differential) power analysis (**APA**), electromagnetic (**EM**) attacks, template attacks (**TA**), cold-boot attacks (**CB**), and countermeasures (**CM**).

General Observations on Attacks on PQC

Many candidates could have used inspections for secure coding practices.

- Lots of these were posted on the NIST PQC forum¹.

Things to be cautious of in code-based cryptography:

- Decoding *must* be implemented in constant-time.
- Be cautious of decoding methods that are insecure (e.g. Reed-Solomon).

What about hash-based signatures?

- Hash-based signatures are particularly sensitive to fault attacks.
- Stateful hash-based signatures need care in state management.
- SPHINCS+ has not had a lot of focus compared to others.



1. <https://groups.google.com/a/list.nist.gov/g/pqc-forum>

General Observations on Attacks on PQC

Things to be cautious of in lattice-based cryptography:

- Plenty of attacks found on schemes designed before the competition.
- No major cryptanalysis attacks on lattice hardness assumptions.
- Decryption failures need to be infeasible, even one is too many [DRV20].
- In general, masking schemes perform quite well.

What about isogeny-based cryptography?

- Unclear how effective attacks/countermeasures transfer from ECC.
- SIKE is the only NIST candidate, promising schemes designed since.
- CSIDH had a few damaging quantum attacks [Pei20, CSCDJRH20].
- New signature schemes proposed Wave [DST19] and SQISign [DFKL+20].



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Recent Highlights

Some selected attacks relevant to the
NIST PQC standardization project.



New Classical Cryptanalysis of Rainbow and GeMSS in Round 3

Multivariate schemes in the 3rd Round:

- Rainbow (Finalist Signature)
- GeMSS (Alternate Signature)

Both substantially attacked near the start of the 3rd Round by new MinRank-style attacks [Beu20, TPD20]

- (Was not part of end-of-2nd-Round decision-making)

Rainbow loses ~16 bits of security at Level 1, and ~55 bits of security at Level 5

- (Rainbow team proposes new security analysis accounting for memory costs – correct?)

GeMSS loses up to 71 bits of security at Level 1, and up to 196 bits of security at Level 5

- Notably, there seems to be very little security gain between their Level 1 and Level 5 parameter sets now



LWE With Side Information

"When a side-channel attack fails, what can you still do with it?" – Léo Ducas.

- Lattice reduction (e.g. [ACD+18]) is a common method for deriving security levels.
- [DDGR20] propose a tool¹ to integrate "hints" from side-channels to use in lattice reduction.
- e.g. if you discover via side-channels $\text{HW}(s_0) = 2 \rightarrow s_0 \in \{3,5\}$, for $s_i \in \{-5, \dots, 5\}$.
- There are four types of hints:

Perfect:

$$\langle \mathbf{s}, \mathbf{v} \rangle = l$$

Modular:

$$\langle \mathbf{s}, \mathbf{v} \rangle = l \bmod k$$

Approximate:

$$\langle \mathbf{s}, \mathbf{v} \rangle = l + \epsilon_\sigma$$

Short Vector Hints:

$$\mathbf{v} \in \Lambda$$

- These hints can reduce the BKZ block size, making attacks easier.



1. <https://github.com/lducas/leaky-LWE-Estimator>

LWE With Side Information

"When a side-channel attack fails, what can you still do with it?" – Léo Ducas.

This work may have impacts in the future:

- This could potentially affect certifications of cryptographic modules.
 - Especially certifiers for Common Criteria, perhaps even FIPS 140-3 and more.
 - Some certifiers (at least for symmetric) may require 2^{80} , or 2^{100} remaining keys **after** SCA.
 - Thus, if this is unsatisfied, the implementation can ‘fail’.
-
- **But** having this a priori knowledge could help certifications.
 - One may now set cryptographic parameters with side channels in mind.



Masking Lattice-Based KEMs

Masked Ring-LWE (à la NewHope) [OSPG18] has large performance differences:

- Overheads in sampling and A2B conversion due to prime modulus.
- Overall, the performance for 1st order protection is 5.7x slower than unprotected.
- Using Masked Comparison [BPO+20] this could be reduced further, but was broken.

The choice between Kyber and Saber may come down to side-channels and masking:

- Differences between these schemes is small, but their moduli differ.
- Masked Saber [Ver19,BDK+20] is 2.5x slower for 1st order protection.
- Recent work [LS19,CHK+20] shows Saber and NTRU can benefit from using NTTs.



Masking Lattice-Based KEMs

We have seen attacks on Saber's 1st order masking scheme:

- In [NDJ21], session keys and secret keys are retrieved using deep neural networks (created at a profiling stage), but require 2.5k and 62k traces.
- They target multiple points in the masked logical shifting on arithmetic shares function, poly_A2A(), and its shuffled variant, poly_A2A_shuffled().
- In [NDGJ21] the session key and long-term secret-key are recovered from 16 traces using deep learning power analysis, without needing to profiling the scheme with masks deactivated.
- They target similar points-of-interest, poly_A2A() and POL2MSG(), for message recovery, and for the secret-key using maps from error-correcting codes using single and multiple traces.
- The target device is once more a very leaky Cortex M4, running at a lower frequency compared to [SKB21].



Masking Lattice-Based KEMs

Kyber have also designed a masked scheme [BGRSV21], for first-order as well as high-order protection. This is a relatively new publication so attacks have not been seen yet.

In [XPROYZ21] the unmasked Kyber is targeted to first find the secret key using EM, requiring 4 traces, on the reference implementation. The same attack is not viable for the assembly-optimised version, thus they target message recovery by also attacking the decoding function which requires profiling to find POIs.

More countermeasures are designed in [HP21] to protect against DPA and fault attacks. The DPA protection is realised in a RNR-protected NTT multiplier and achieves a relatively low overhead of 1.13x compared to unprotected.



Masking Lattice-Based Signatures

This issue is also seen in lattice-based signature schemes:

- Masked Dilithium [MGTF20] 7-9x faster using power-of-two modulus.
- Dilithium 1st order masking 5.6x slower compared to unprotected.
- Masking qTESLA [GR19] also gained efficiencies by changing modulus.

Masking Falcon will be slightly different:

- But now “isochronous” with constant runtime of Gaussian sampling [HPRR20].
- Its use of floating-point arithmetic makes masking an open problem.
- But has been considered before in MPC [ABZS13, GHK+20].
- Recent result called Mitaka: a simpler, parallelisable, maskable variant of Falcon.



Countermeasures are no Guarantee

But in general, we've only just begun protecting these schemes.

- There's many attack vectors to find and countermeasures need to be tested.

So far, we only have masked designs for Saber, Kyber and Dilithium.

- For first-order [MGT+19] and higher-order [BDK+19,BGR+21].

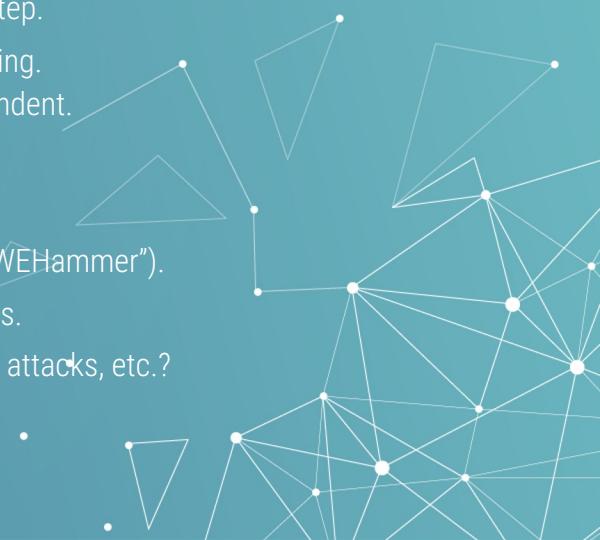
And obviously, countermeasures also not a guarantee.

- The masking scheme for the code-based scheme QcBits [RHH+17] was attacked [SKC+19].
- Masked comparison [BPO+20] used in lattice-based masked KEMs was attacked recently [BDH+21].
- As already discussed, Saber's masking has been broken by appropriate profiling and power analysis.



Active Side-Channel Attacks: QuantumHammer (LWEHammer?)

- QuantumHammer is a Rowhammer-style fault attack [MIS20] against LUOV β .
- Didn't have a large impact *on the process*, because of a concurrent work -- computational attack on LUOV [DDSVZ20].
- Intuition:
 1. Profile a victim device (e.g. in the cloud, like AWS) in an offline, pre-processing step.
 2. Online physical manipulation of the device while LUOV signature scheme is running.
This causes the device to output malformed signatures that are secret key dependent.
 3. Computationally recover the secret key given enough malformed signatures.
- Is this more broadly applicable? (Upcoming work, joint NIST PQC + Univ of AR team – “LWEHammer”).
- Rowhammer is an “old attack” (2014); modern cloud architectures *ought* to have defenses.
- Are the Rowhammer defenses enough in practice? Can Rowhammer be replaced by fault attacks, etc.?



Fault Attacks Against PQC

Determinism is generally considered preferable from a security perspective.

Fault attacks have exploited determinism in SPHINCS+ [CMP18] and Dilithium [BP18,RJH+19].

- Many schemes counter this by adding random salt as an option or as standard.

Hedging [AOT+20] is an interesting alternative to mitigating these fault attacks (and randomness failures) is by deriving the per-signature randomness from a combination of the secret-key, message, and a nonce.

This is formalized for Fiat-Shamir signatures and apply the results to hedged versions of XEdDSA, a variant of EdDSA used in the Signal messaging protocol, and to Picnic2, and show hedging mitigates many of the possible fault attacks.



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Takeaways for the Future

What can we learn from these attacks?

What can we learn from these attacks?

Implementation complexity will significantly increase with these standards.

- NIST set the focus on ARM Cortex M4 and Xilinx Artix-7.
- Will attacks or countermeasures we've seen be as effective on other devices?
- Will we see more attack vectors? Will we need more countermeasures?

Complexity also increased by the large number of fragile/sensitive operations.

- Past attacks highlight many sensitive/fragile operations that can break schemes.
- Even constant-time; e.g. how will *ctgrind* work with rejection sampling?

In order to learn the relevance of these attacks you should consider your use case.

- Some candidates will not fit on these devices, let alone smart cards.



Selecting the Correct Device

In order to learn the relevance of these attacks you should consider your use case.

Many of the attacks shown are on the ARM Cortex M4, a device NIST chose for benchmarking.

One recent research result looked at benchmarking Falcon and Dilithium on ARM Cortex M7.

- Similar to M4, except the M7 has a full 64-bit floating-point unit, needed by Falcon.

We wanted to benchmark, profile, and analyse Falcon on this embedded target using the FPU.

- Falcon signing ran more than 6x faster with the FPU, Dilithium had no improvements.
- Profiling Falcon we saw a 15-17x speed-up in many operations inside key gen and sign.

However, we saw some irregularities on the constant time performance of Falcon...



Selecting the Correct Device

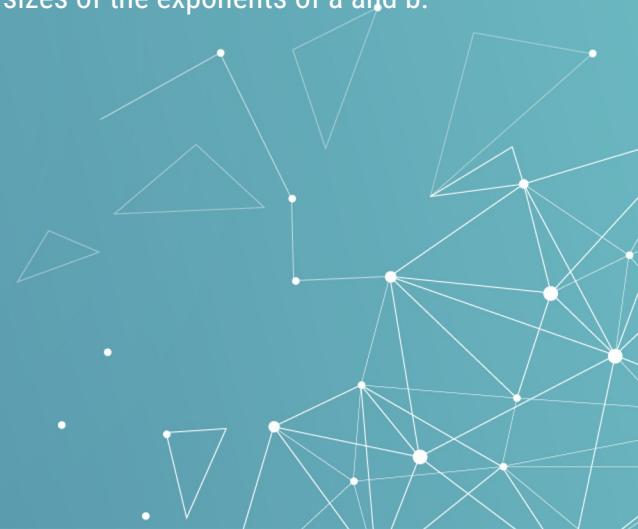
We found a few subtle constant time issues with Falcon on Cortex M7 on STM32 development boards and ARMv7.

Firstly, 64-bit floating point operations are not fully constant-time on four STM32 development boards we tested that had ARM Cortex M7 CPUs.

For instance, $a \times b$ is much faster if $b = 0$. Worse, the timing of $a + b$ depends on relative sizes of the exponents of a and b .

We found constant-time issues with almost every operation on the four boards.

We did not investigate how to exploit these issues.



Selecting the Correct Device

Secondly, we discovered is an issue with Falcon on ARMv7.

When casting a double type to an int64_t, C prescribes to round towards zero.

There is no native instruction to do such a double→int64 truncation on ARMv7.

Instead, the compiler calls the runtime symbol `_fixdfi`, aka `__aeabi_d2lz`.

This might or might not be implemented in constant time.

In LLVM it is not – and it leaks the sign.

This is also the case for the Raspberry Pi 3, which they target in [6].

We reported this issue to the Falcon team and proposed a fix.



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Conclusions



In Summary

We hope this survey will motivate further evaluations of NIST PQC candidates.

NIST will take these into consideration in their final decisions for standardization.

NIST PQC will add implementation complexities.

- We have seen many novel attacks due to the fragility of some operations.
- Many attacks were also enabled by implementation errors.

We discuss some open questions and implications of the attacks found.

Start your PQC transition now and consider the attacks that will affect your use cases.





Thanks

Does anyone have any questions?

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