## Feedback on CO661 A1 for jj333

Final score = 88 /100 Overall comment: Well done!

Task 1 - Server - Max. score 60

Criteria	Comments	Score
Tests	$round(\frac{47}{47} \times 20)$	20 / 20
Functionality		
Concurrency	This is a good solution, but I think it contains	17 / 20
and exclusion	a few bits of redundancy / odd things:	
	I don't think its necessary to create a mutual ex-	
	clusion block for getReads and isWriter; just	
	reading a value from memory is atomic, and	
	so it doesn't really matter if this interleaves	
	with other stateful (possibly atomic) interaction with these variables. Mutual exclusion is	
	really important when you are updating though,	
	so its good you have included this, for example,	
	in readUnlock().	
	I don't really understand line 199. I puzzled	
	over this line for a long time, and tried to get a	
	bug out of it, but coudln't. But I'm sure there	
	is something strange here! Closing a reader	
	releases the write semaphore, so it keeps in-	
	creasing the number of available permits for	
	the writer semaphore, even if no writer is us-	
	ing it. This means I can make more writers	
	get theourh the acquire on line 179, but then	
	this is guarded by line 178 which acquires all the reader permits so it seems okay, just odd.	
	It leads me to the question: do you even need	
	the writerSem? I'm fairly certain that this is	
	redundant because you get all the standard mu-	
	tual exclusion with the way you use your reader	
	lock.	
Race free-	Good mechanisms.	5 / 5
$\operatorname{dom}$		
Fairness	None of the sempahores are fair, although the	2 / 5
_	locks are.	
Description	Fine.	$\frac{5}{5}$
Code quality	Good, but could do with a few more local com-	4 / 5
(comments, for-	ments.	
mat, modular-		
ity/abstraction) Total	Good work.	52/60
TOTAL	GOOG WOLK.	53/60

Task 2 - Client - Max. score 20

Criteria	Comments	Score
Opening		2 / 2
Reading		2 / 2
Writing		3 / 3
Randomness		3 / 3
$Logging\ output$		3 / 3
Client/server spawning		4 / 4
Code quality		3 / 3
Total	Good!	20/20

Task 3 - Model - Max. score 20

Criteria	Comments	Score
Server and locking	Okay- This idea of making the server	3 / 6
	into basically just the semaphores and	
	having the client interact with those	
	semaphores doesn't lead to a very clear	
	model of your actualy code, but the	
	semaphores are well presented and I like	
	the use of relabelling.	
Client and top-level	Same issue as above in terms of	2 / 4
	client/server separation, but I'm not go-	
	ing to double penalise this here. One	
	strange thing though is that your client	
	reader has a branch where it needs to	
	acquire the write semaphore, and simi-	
	larly one that. That doesn't seem right,	
	and doesn't match your code.	
$Appropriate\ abstraction$	Yes the level of abstraction is about	4 / 4
	right.	
Model exhibits mutual	Yes, I could get the properties to go	6 / 6
exclusion of read and	through with appropriate positioning	
writes	of some key observable atoms (see ap-	
	pendix for my tweaked version of your	
	model)	
Total	Pretty good. Captures the right idea,	15/ <b>20</b>
	but I think it could be more clearly set	
	up.	

<sup>\*</sup> Semaphores (we have a semaphore with 1 lock for writing, and one with 3 locks for read S = acq.rel.S;

 $S2 = S \mid S;$ 

S3 = S2 | S;

st ReadWriteSem is the custom semaphore written for read and write lock management, and i

st I previously had an implementation with further handshakes for calling the functions (

- \* But the transitions provided no additional functionality, and it is simpler to underst ReadWriteSem = S[writeAcq/acq, writeRelease/rel] | S3[readAcq/acq, readRelease/rel];
- \* Client can either Client = ClientRead + ClientWrite;
- \* ClientRead has to obtain the write lock if it is available, and then one read lock.

  ClientRead = 'writeAcq.'readAcq.openRfile.ClientCloseRead + 'readAcq.openRfile.ClientCloseRead + 'readAcq.ope
- \* ClientCloseRead releases one read lock, and the write lock if obtained.
- ClientCloseRead = done.'readRelease.'writeRelease.Client + done.'readRelease.Client;
- \* ClientWrite obtains all read locks to prevent simulatenous reads, and obtains the writ ClientWrite = 'readAcq.'readAcq.'readAcq.'writeAcq.openWfile.ClientCloseWrite;
- \* ClientCloseWrite releases all locks, to enable further reads and writes.
  ClientCloseWrite = done.'writeRelease.'readRele
- \* We spawn 3 clients and 1 server (called ReadWriteSem), but in actual use we can spawn App = (Client | Client | ReadWriteSem)  $\setminus$  {readAcq, writeAcq, readRelease, writeAcq, readRele

Start = App;