

Lecture 22 - Quantum Probability

Recap of last lecture:

↳ Mixed (qudit) state: "prob. p_i of $|Y_i\rangle \in \mathbb{C}^d$ ", $i=1\dots n$.
 ↳ encoded by density matrix $\rho \in \mathbb{C}^{d \times d}$,

$$\rho = p_1|Y_1\rangle\langle Y_1| + \dots + p_n|Y_n\rangle\langle Y_n|$$

Hermitian matrix ρ (meaning $\rho^T = \rho$) is a density matrix
 iff ① " $\rho \geq 0$ ": " ρ is positive semi-definite" ($\langle u | \rho | u \rangle \geq 0 \forall u$)
 ② " $\text{tr}(\rho) = 1$ " $\sum_{i=1}^d \rho_{ii}$

Every Hermitian matrix M acts on \mathbb{C}^d by
 "stretch by factor λ_i in direction $|v_i\rangle$ "
 (real) eigenvalues of M (orthonormal) eigenvectors of M

① $\Leftrightarrow \rho$'s eigenvals λ_i all ≥ 0 ② $\Leftrightarrow \lambda_1 + \dots + \lambda_d = 1$

\therefore any density mtx ρ equiv to mixed state
 " λ_i prob. of $|v_i\rangle$ ", $i=1\dots n$.

Actions on (mixed) states:

- Applying unitary transformation $U \in \mathbb{C}^{d \times d}$:

$$g \mapsto U g U^\dagger$$

- Measuring in basis $|u_1\rangle, \dots, |u_d\rangle$:

get outcome "i" with prob. $\langle u_i | p | u_i \rangle$

[And state collapses, changing g . You can work out formulas, but we won't bother about this now. Today, we'll never use a state again after measuring it.]

- Adjoining two states:

If g is density mtx of a d -dim. mixed state,
 σ e-dim. ,

& we decide to call the whole system one de-dim. state, its density mtx is ... $g \otimes \sigma$. (exercise/homework)

[Kronecker prod]

Linear Algebra Interlude

prop: Let $A \in \mathbb{C}^{d \times d}$, $B \in \mathbb{C}^{d' \times d}$. Then

$$\text{tr}(AB) = \text{tr}(BA).$$

proof: $\text{tr}(AB) = \sum_{i=1}^d (AB)_{ii} = \sum_{i=1}^d \sum_{j=1}^{d'} A_{ij} B_{ji}$

$$= \sum_{j=1}^{d'} \sum_{i=1}^d B_{ji} A_{ij}$$

$$\text{tr}(BA) = \sum_{j=1}^{d'} (BA)_{jj}$$
◻

Classic trick: $\langle u_i | p | u_i \rangle = \text{tr}(\underbrace{\langle u_i |}_{A} \underbrace{p | u_i \rangle}_{B})$ 1×1 matrix!

[Recall: this is prob ["i"] when measuring in $|u_1\rangle, \dots, |u_d\rangle$ basis]

$$= \text{tr}(\underbrace{p | u_i \rangle}_{\text{d} \times \text{d} \text{ "proj. onto } |u_i\rangle \text{ mtx}} \langle u_i |)$$

Also: $\text{tr}(A^\dagger B) = \sum_{ij} (A^\dagger)_{ij} B_{ji} = \sum_{ij} \underbrace{A_{ji}^*}_{i,j \text{ here}} B_{ji}$ [Could reverse i,j here]

"dot product of A, B when they're viewed as lists of #'s"

def: $\langle A, B \rangle = \dots$ "Inner product for matrices".

[Don't worry much about the dagger! Often A, B will be Hermitian, so $A = A^\dagger$.]

Today: "Quantum Probability 101"

(vs. "Classical Probability 101")

Classical

Def #1: A d -outcome prob. distrib. is a vector $p \in \mathbb{R}^d$ with $p \geq 0$, $\sum_{i=1}^d p_i = 1$.

[Outcomes called $1, 2, \dots, d$; $p_i = \Pr[i]$]

Quantum

Def #1: A d -dim. state/density mtx is a Hermitian $\rho \in \mathbb{C}^{d \times d}$ with $\rho \geq 0$, $\sum_{i=1}^d \rho_{ii} = 1$.

[Think of ρ as a "source of quantum randomness"]

Classical \rightarrow quantum : replace "vector" by
"Hermitian matrix"

quantum \rightarrow classical : take every Hermitian
matrix to be diagonal

Example:

Recall: if state ρ measured in $\{|u_1\rangle, \dots, |u_d\rangle\}$ basis, $\Pr["i"] = \langle u_i | \rho | u_i \rangle = \text{tr}(\rho |u_i\rangle \langle u_i|)$

[Needn't worry about dagger[†], $\because \rho^\dagger = \rho$] $\rightarrow = \langle \rho, E_i \rangle, E_i := |u_i\rangle \langle u_i|$
 ("proj. onto $|u_i\rangle$ " mtx)

Remark: $E_i \geq 0$ & $E_1 + \dots + E_d = I_{d \times d} = \mathbb{I}_{d \times d}$

[why?]

[It is a pure-state dens. mtx]

[Matrix that stretches by 1 in every direction: all eigenvalues are 1.]

[Two notations for the $d \times d$ identity matrix.]

Compare: Given prob. dist $p \in \mathbb{R}^d$, let

$$e_i = (0, \dots, 0, \underset{i\text{th position.}}{1}, 0, \dots, 0)$$

[Indicator of "ith outcome"]

$$p_i = \Pr_p["i"] = \langle p, e_i \rangle.$$

$$\text{And } e_i \geq 0, e_1 + \dots + e_d = (1, 1, \dots, 1) = \vec{\mathbb{I}}$$

Probability 101 def. #2: Events

[[My descriptions here will be a bit weird,
to accommodate the subsequent quantum analogy.]]

Given prob. dist $p \in \mathbb{R}^d$, let's define some
"mutually exclusive, collectively exhaustive" events
[[meaning always exactly one happens]]
 $\text{Event}_1, \dots, \text{Event}_m$.

Naivest: $m=d$, $\text{Event}_i = "i \text{ is drawn from } p"$
[[like in preceding example]]

Lumping: Some outcomes grouped together.

E.g. $\text{Event}_1 = "the draw from } p \text{ is odd"$
 $\text{Event}_2 = " \text{ even"}$

Identify with "indicator vectors":

$$e_1 = (1, 0, 1, 0, \dots)$$

$$e_2 = (0, 1, 0, 1, \dots)$$

$$\Pr_p[\text{Event}_1] = p_1 + p_3 + p_5 + \dots = \langle p, e_1 \rangle; \quad \Pr_p[\text{Event}_2] = \langle p, e_2 \rangle$$

Note: e_i is nonneg. & sum to $\vec{1} = (1, 1, \dots, 1)$.

[[This ensures $\langle p, e_i \rangle \geq 0$ & $\sum_i \langle p, e_i \rangle = 1$; exactly
one Event occurs.]]

[[In "lumping", you have $\leq d$ events. Now we'll describe a way to have $>d$ events. Involves injecting additional randomness.]]

Using additional randomness:

(in quotes b/c need
not be q/I)

Suppose e_1, \dots, e_m are any "indicator" vectors in \mathbb{R}^d with $e_i \geq 0$ & $e_1 + \dots + e_m = \vec{1}$.
[entrywise]

E.g.: $e_1 = (0, .2, .3, 0, \dots)$

$m < d$ here, but
could easily
have $m > d$

$e_2 = (1, .7, .2, 0, \dots)$

$e_3 = (0, .1, .5, 1, \dots)$

Event₁, Event₂, Event₃ ??

Think: Draw from p.

If outcome 1 : Event₂ occurs.

If outcome 2 : { Event₁ occurs w.p. $\frac{.2}{.7}$
Event₂ $\frac{.7}{.7}$
Event₃ $\frac{.1}{.7}$

If outcome 3 : { Event₁ occurs w.p. $\frac{.3}{.7}$
etc.

[[Yes, you need additional randomness to implement. But still makes sense to say "exactly one event always occurs", and...]]

$$\Pr[\text{Event}_j] = \sum_{i=1}^d p_i (e_j)_i = \langle p, e_j \rangle.$$

Quantum Generalization

How can you measure [and thereby get classical outcomes] a qudit state $|g\rangle$?

- in orthonormal basis $\rightarrow d$ outcomes
- "lumping": Say $d=2^q$, $|g\rangle$ has q qubits.
Can do a "partial measurement" of r qubits: yields $2^r < d$ outcomes
- "additional randomness": Can add additional ["ancilla"] s qubits, do a 2^{r+s} -dim unitary, partially/fully measure: up to $2^{r+s} > d$ outcomes

→ [Can also do non-powers-of-2 things]

Most general: Think of whole procedure as a "BODM": Big Ol' Device for Measuring.

Call readouts/outcomes " 1 ", " 2 ", ..., " m "

[m may be $< d$, $= d$, $> d$]

Turns out: [if you do the math, which we won't]

BODM yields m Hermitian $E_1, \dots, E_m \in \mathbb{C}^{d \times d}$

satisfying: (i) " $E_i \geq 0$ " $\forall i$ (positive semidefinite)
 $\langle w | E_i | w \rangle \geq 0 \quad \forall i, w$

(ii) $E_1 + \dots + E_m = \mathbf{I}_{d \times d}$ (identity matrix)

And: $\Pr[\text{BODM applied to } g \text{ reads out } 'i'] = \langle g, E_i \rangle.$

ex: check (i) $\Rightarrow \langle g, E_i \rangle \geq 0 \quad \forall i$ (uses $g \geq 0$)

(ii) $\Rightarrow \langle g, E_1 \rangle + \dots + \langle g, E_m \rangle = 1$ (uses $\sum_i E_i = I$)

Conversely: given Hermitian E_1, \dots, E_m satisfying
(i), (ii), can in principle build a
physical BODM with associated
behavior.

Real name of BODM is "POVM".

(You'll see "POVM" a lot. Stands for

"Positive Operator-Valued Measure" for some
abstruse math reason)

[Anyway, POVM $\{E_1, \dots, E_m\}$ is quantum generalization of M.E.C.E. events.]

Question: If POVM registers outcome "i", how does state ρ collapse?

Answer: Depends on how it's implemented!

Can't tell just from E_1, \dots, E_m .

[If you want to know, need to know the circuit implementing the measurement. Then, once you know this, can work out the formula. For posterity, if (nonuniquely)

$E_i = M_i^+ M_i$ for $d \times d$ M_i , then you can implement it in such a way that upon reading out "i", ρ collapses to

$$\frac{M_i \rho M_i^+}{\langle \rho, M_i M_i^+ \rangle}$$

But generally, when people discuss POVMs, they usually only care

about measurement outcome probabilities, and don't plan to continue "using" the state.]

↑
see homework

Probability 101 def #3: Random Variables

Given prob. distribution $p \in \mathbb{R}^d$, a random variable is just a real number x_i for each outcome $1 \leq i \leq d$.

So it's any d -vector $x \in \mathbb{R}^d$.

Expected value is $E_p[x] = \sum_{i=1}^d p_i x_i = \langle p, x \rangle$.

Quantum generalization: "observable"
for d -dim. states $|f\rangle$:

Any old Hermitian matrix $X \in \mathbb{C}^{d \times d}$.

recall: can express $X = \sum_{i=1}^d x_i |u_i\rangle\langle u_i|$

\downarrow

\nearrow

real eigenvalues

projection onto eigenvectors

$|u_i\rangle$

$$X = \sum_{i=1}^d x_i |u_i; X u_i|$$

↑ projection
onto eigenvect.
 $|u_i\rangle$

real eigenvalues

Physically, could build an instrument that:

- measured in basis $|u_1\rangle, \dots, |u_d\rangle$
- on outcome " i ", read out real # x_i

If you applied this instrument to ρ , what's expected value of readout?

$$\sum_{i=1}^d \underbrace{\langle \rho, |u_i; X u_i| \rangle}_{\text{Pr(measure "i")}} x_i = \langle \rho, \sum_{i=1}^d x_i |u_i; X u_i| \rangle = \langle \rho, X \rangle \quad \text{😊}$$

[Matches classical situation.]

Notation: $E_\rho[X]$ ← [Good notation: linearity of expectation holds.]

Rem: X^2 operator stretches by x_i^2 factor in direction $|u_i\rangle$. I.e., it's $\sum_{i=1}^d x_i^2 |u_i; X u_i|$

$$\therefore \text{expected (readout)}^2 = E_\rho[X^2].$$

[This is not tautological! We had to check it. But further reassures that notation is good.]

Quantum probability theory:

sources of randomness: ρ ✓

events: POVMs ✓

random: observables ✓

Done 😊 [Really, what else is there in
Probability 101? 😊]

Warning: If X, Y are observables
(Hermitian matrices)

$XY \neq YX$ necessarily.

Indeed, "quantum probability" often
called "noncommutative probability".

[Holds iff XY is Hermitian; i.e., iff
 XY is itself an "observable".]

[[Actually, of course, there's much much more. From quantum probability, can develop quantum... • statistics

- information theory
- learning theory
-

We'll see a little statistics on homework, including... the "Uncertainty Principle",
 $\text{stddev}_p[X] \cdot \text{stddev}_p[Y] \geq |E_p\left[\frac{i}{\hbar}(XY - YX)\right]|.$

Some info theory in next lecture.

My research is about learning & statistics of quantum states; e.g., the "tomography = state learning" problem:

How many copies, n , of state ρ are needed to approximately learn it; by applying POVMs to $\rho \otimes \rho \otimes \dots \otimes \rho$?]]
(n times)