Chess AI

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1 Introduction

In this report, I will discuss my approach to creating a Chess AI program that will be able to select good moves against both AI and human opponents. I implemented my ai move selection using the provided Chesspresso library, and in this assignment I have included 3 separate chess AI classes. The first of which, MiniMax.java will use minimax search to select moves against its opponent. The second, AlphaBeta.java, will similarly use minimax search, but with the addition of Alpha Beta pruning to search more efficiently by only searching through nodes in the game tree that are viable. Finally, the last chess AI class is AlphaBetaTransposition, which does a minimax search with both alpha beta pruning as well as the addition of a transposition table to further eliminate redundancies in the search algorithm by keeping track of a visited set of equivalent positions, or "transpositions". The transposition table utilizes the hascode function already implemented within Chesspresso's position class for hashing.

2 Utility and Evaluate Function

In order to conduct a minimax search, it is necessary to have a utility function that will be able to return the value of leaf nodes in the game tree. This will allow the value of terminal states to be passed up the tree to the starting node so it can make a move decision. I implemented a method called utility() in my classes that would take a position and return its utility. It is shown below.

```
public int utility(Position pos) {
       int value = 0;
5
       //checkmate
       if (pos.isMate()) {
         //ai checkmated opponent
9
         if (pos.getToPlay() != color) {
10
           value = BESTCASE;
11
         //ai got checkmated
         if (pos.getToPlay() == color) {
14
           value = WORSTCASE;
17
18
       else if (pos.isStaleMate())
19
         value = 0;
20
21
       else {
22
         value = evaluate(pos);
23
24
       return value;
25
26
```

This utility function will test if the current is a terminal node, indicated by a checkmate for either side or a stalemate, and then return a large positive or large negative value depending on whether the state is a win or loss for the AI. For a stalemate, the function returns a value of zero. For non-terminal states, the value will call an evaluate function to determine the utility value of the position. The evaluate() function is shown below.

```
public int evaluate(Position position) {

int value = 0;
```

```
//If the turn is the AI's, then the getMaterial score is correct
       if (position.getToPlay()=color) {
6
         //getMaterial gets material score
         //getDomination weighs occupation of center squares more, not as important as material score
         value += position.getMaterial();
         value += 0.5 * position . getDomination();
11
       //If the turn is the other person's, then negate score to get AI's score
13
14
       if (position.getToPlay()!=color) {
         value -= position.getMaterial();
16
17
         value -= 0.5 * position.getDomination();
18
      return value;
20
```

Here, the evaluate function will evaluate non-terminal positions by considering the material score of the position. Because the getMaterial() method in Chesspresso takes into account whose turn it is, we have to negate the score when it's not the AI's turn, because we always want to evaluate the material score of the position from the AI's standpoint. Additionally, Chesspresso has a method called getDomination(), which will increase scores for positions that occupy spaces near the center, as that signifies a dominant position for a player during a chess game, as the center of the board gives the player many options to attack from. However, this domination score is weighted by 0.5 because it is of secondary importance to the material score.

3 Cutoff Test

In addition to the utility and evaluate functions, we need a cutoff test that tells minimax search when to stop. This is because minimax search essentially does a depth first search through the game tree and requires a base case to stop the search. The two main base cases we look at are whether the search has exceeded a predetermined depth and whether the search has encountered a terminal node. The cutoff function is shown below.

```
public boolean cutOff(Position pos, int depth) {
       if (depth >= maxD){
         return true;
       if (pos.isMate()){
         return true;
9
       if (pos.isStaleMate()) {
10
         return true:
11
       if (pos.isTerminal()){
13
         return true;
14
15
       return false;
16
     }
```

4 Minimax Search

Finally, we have all the pieces required to conduct a Minimax search. This is shown in the miniMax(), mini(), and max() methods shown below.

```
public short miniMax(Position pos, int startDepth) {

//initialize variables for minimax, assuming maximizer
short bestMove = 0;
int bestValue = WORSTCASE;
int currentValue;

//get set of possible moves and try each one out
for (short move : pos.getAllMoves()) {

try {
```

```
pos.doMove(move);
13
           currentValue = mini(pos, 0);
14
           //if value is better, then update bestValue
16
17
           if (currentValue > bestValue) {
18
             bestValue = currentValue;
             bestMove = move;
19
20
21
           //undo move to try more moves
           pos.undoMove();
23
           catch (IllegalMoveException e) {
24
25
       }
26
       return bestMove;
27
28
29
     //returns value of min player's move choice
30
     public int mini(Position pos, int depth) {
31
32
       nodes++:
33
34
       int currentDepth = depth+1;
35
36
       //Base case: pass cutoff
       if (cutOff(pos, currentDepth)){
37
         return utility(pos);
38
39
       //initially, minimizer's bestValue is +infinity
40
       int bestValue = BESTCASE;
41
42
       //get moves and try each one out
43
       for (short move : pos.getAllMoves()) {
44
         try {
45
           pos.doMove(move);
46
           //if move is better, then bestValue is updated
47
           bestValue = Math.min(bestValue, max(pos, currentDepth));
48
           pos.undoMove();
49
           catch (IllegalMoveException e) {
50
51
52
       return bestValue;
53
54
55
     //returns value of max player's move choice
56
     public int max(Position pos, int depth) {
57
58
59
       nodes++:
       int currentDepth = depth+1;
60
61
       //Base Case: Pass Cutoff
62
       if (cutOff(pos, currentDepth)){
63
         return utility (pos);
64
65
66
       int bestValue = WORSTCASE;
67
68
       for (short move : pos.getAllMoves()) {
         try {
69
           pos.doMove(move);
70
           bestValue = Math.max(bestValue, mini(pos, currentDepth));
71
72
           pos.undoMove();
           catch (IllegalMoveException e) {
73
74
       }
       return bestValue;
76
```

The miniMax method will take a position as well as a starting depth. The position is the initial game state that the search will start from and the starting depth is what depth the search starts at. To clarify, being passed a starting depth of 3 does not mean that the search will search 3 layers into the tree. Rather, it means the search will "start" at depth 3 and search until the max depth is reached, so in total the Minimax search will search the (maximum depth - starting depth) layers. In this class, the maximum depth is stored in the maxD instance variable in MiniMax.java.

The miniMax method will start the search assuming the AI is in the maximizer position, so it will call min on its possible moves, which will in turn call max on its moves, etc. until a base case is reached. Here, the base case is a position passes

a cutoff test - if a position passes it, then the search will immediately return the utility of the current position up the tree and stop searching downward. Additionally, both max and min will keep track of the value of their best possible move in bestValue, and that will be the value they return up the tree if they are not at a leaf node. Thus, leaf nodes will return their utility up the tree after hitting the base case, and tree nodes will return the best possible value amongst the value of all their possible moves. Additionally, I'm keeping track of the node count with an instance variable called node to later compare the efficiency of the various searches.

5 Iterative Deepening MiniMax

After implementing depth-limited Minimax, I used a for loop and an instance variable called bestMove to keep track of the results from running my Minimax search up to varying depths. The method idmm() will do the work of iteratively calling the depth-limited minimax function, and getMove() will simply return the bestMove chosen by idmm(). Again, because Minimax takes into account the starting depth, and not the max depth as its parameter, the for loop will have to decrement with each loop, first starting at max depth -1, so it will search 1 layer to reach max depth. The final loop will start at depth 0 and search up to max depth, so it will in total search max depth layers.

```
public short getMove(Position pos) {
2
       Position position = new Position (pos);
       short move = idmm(position);
5
       System.out.println("Minimax nodes searched: " + nodes);
6
       return move;
8
     public short idmm(Position pos) {
11
       //This loop will run minimax search for depth 1 up to depth maxD
12
       for (int i = maxD-1; i >= 0; i--) {
13
14
         //save results from each loop in bestMove instance variable
15
         bestMove = miniMax(pos, i);
17
         //if found checkmate, no need to go deeper
18
19
           pos.doMove(bestMove);
20
           if ( utility ( pos )==BESTCASE) {
21
             return bestMove;
22
23
           pos.undoMove();
24
25
         catch (IllegalMoveException e){
27
       return bestMove;
29
30
```

6 Alpha Beta Pruning

In addition to Minimax search, I also implemented Alpha Beta pruning on my Minimax search - this would prune off certain branches of the tree where it was not necessary to search any deeper. It would go about this by keeping track of two variables, alpha and beta, which represented the maximizer's worst case scenario and the minimizer's worst case scenario, respectively. If maximizer has a possible position in which the value of that position is greater than beta, then we prune off the rest of that tree and don't explore that position's sucessors. The reason for this is that the maxizer will always pick something of higher or equal value to the node we looked at, but this value is higher than minimizer's worst case scenario. Ultimately, this means that beta will not allow maximizer to even get to this node because its worst case scenario is better than letting maximizer get to choose that position. The same logic, applies to minimizer - if it reaches a node that potentially has a value less than alpha, then we prune off that branch because maximizer will not let it get to a value that's lower than its current worst case scenario. The three methods for Minimax search with Alpha Beta pruning are shown below.

```
//alpha beta MiniMax search
public short abMiniMax(Position pos, int startDepth) {
short bestMove = 0;
int bestValue = WORSTCASE;
```

```
int currentValue;
       int alpha = WORSTCASE;
8
       int beta = BESTCASE;
9
10
11
       for (short move : pos.getAllMoves()) {
12
13
         try {
14
           pos.doMove(move);
           currentValue = mini(pos, 0, alpha, beta);
15
           if (currentValue > bestValue) {
17
             bestValue = currentValue;
18
19
             bestMove = move;
           }
20
21
           pos.undoMove();
22
          catch (IllegalMoveException e) {
23
24
25
26
       return bestMove;
    }
27
28
     //returns mini's best move value
29
30
     public int mini(Position pos, int depth, int alpha, int beta) {
31
       nodes++:
32
       int currentDepth = depth+1;
33
34
       //Base case: pass cutoff
35
       if (cutOff(pos, currentDepth)){
36
        return utility (pos);
37
38
39
       int bestValue = BESTCASE;
40
       for (short move : pos.getAllMoves()) {
41
42
         try {
43
           pos.doMove(move);
           bestValue = Math.min(bestValue, max(pos, currentDepth, alpha, beta));
44
           pos.undoMove();
46
           //Pruning - if value is less than alpha (max's worst case), no need to continue
47
           if (bestValue <= alpha) {
48
             return bestValue;
49
50
51
           //update beta if bestValue is better
52
           beta = Math.min(beta, bestValue);
53
54
          catch (IllegalMoveException e) {
55
56
57
       return bestValue;
58
59
60
     //returns max's best move value
61
62
     public int max(Position pos, int depth, int alpha, int beta) {
63
64
       int currentDepth = depth+1;
65
66
67
       if (cutOff(pos, currentDepth)){
         return utility(pos);
68
69
70
       int bestValue = WORSTCASE;
71
72
       for (short move : pos.getAllMoves()) {
73
           pos.doMove(move);
74
           bestValue = Math.max(bestValue, mini(pos, currentDepth, alpha, beta));
75
76
           pos.undoMove();
77
           //Pruning - if value is greater than beta (mini's worst case), no need to continue
78
           if (bestValue>=beta) {
             return bestValue;
80
81
82
```

```
//update alpha if bestValue is better alpha=Math.max(alpha, bestValue);

// catch (IllegalMoveException e) {

// catch (Il
```

The rest of the methods, including the iterative deepening component, are the same between Minimax and Minimax with Alpha Beta Pruning.

7 Minimax Search with Alpha Beta and Transposition Table

Next, I used a transposition table to keep track of visited equivalent states in a hashtable in order to further reduce redundancies in the search algorithm. This was implemented by using a hashtable that would store a the hashcode that results from calling getHashCode() on a given position (implemented in Chesspresso already), and using that as a key that mapped to the corresponding position's value and depth, stored as an integer array. For each move that is tried, the resulting position is hashed and we check if the position's hashcode is already inside the hashtable.

If it's already inside, then we check whether the depth that we're at currently is less than the depth associated with the previously found position. If it's less than the previous depth, then we update the depth associated with the position in the hashtable. Additionally, we don't need to keep searching down the tree because we already know the best value from this point, as it is already stored in the transposition table, so we can just return that value. If it's not less than the previous depth, then we keep searching down the tree.

If the position was not already inside, then we need to continue down the tree to get its best value. After we get its best value, we can record its information in the transposition table for later comparison. The code for the transposition table implementation is shown below. Note: the depth of each position is (max depth - start depth) because that tells you how many layers deep it has searched so far.

```
//alpha beta transposition table mini max search
     public short abttMiniMax(Position pos, int startDepth) {
3
       short bestMove = 0;
       int bestValue = WORSTCASE;
       int current Value;
       int alpha = WORSTCASE;
       int beta = BESTCASE;
9
11
       //try all moves and get their utility
       for (short move : pos.getAllMoves()) {
12
13
14
           pos.doMove(move);
16
           //if already in transposition table - uses position's getHashCode
17
           if (tt.containsKey(pos.getHashCode())) {
18
19
             //info array stores position's value and index
20
21
             int info[] = tt.get(pos.getHashCode());
22
             //if search depth is less than previous search depth
23
             if (maxD-startDepth < info[1]) {
24
               //update position's info with this position
26
               currentValue = info[0];
27
               int[] updateInfo = {currentValue, maxD-startDepth};
28
               tt.put(pos.getHashCode(), updateInfo);
29
30
             //otherwise, keep searching down the tree
31
32
               currentValue = mini(pos, startDepth, alpha, beta);
33
34
           }
36
           //if position not already in table, keep searching down tree
37
38
           //then add position into table with its info
```

```
else {
39
              currentValue = mini(pos, startDepth, alpha, beta);
40
              int[] updateInfo = {currentValue, maxD-startDepth};
41
              tt.put(pos.getHashCode(), updateInfo);
42
43
44
            if (currentValue > bestValue) {
45
46
              bestValue = currentValue;
              bestMove = move;
47
49
50
            pos.undoMove();
51
           catch (IllegalMoveException e) {
53
       return bestMove;
54
55
56
57
58
     public int mini(Position pos, int depth, int alpha, int beta) {
59
       nodes++;
60
       int currentDepth = depth+1;
61
62
63
        //Base case: pass cutoff
        if (cutOff(pos, currentDepth)){
64
         return utility(pos);
65
66
67
       int bestValue = BESTCASE;
68
        for (short move : pos.getAllMoves()) {
69
         try {
70
           pos.doMove(move);
71
72
73
            //if already in transposition table
            if (tt.containsKey(pos.getHashCode()))
74
              //get the position's value and depth info
75
              int info[] = tt.get(pos.getHashCode());
76
77
              //if position's search depth (not current depth)  cpreviously found
78
              //then update table with position's info
79
80
              if (\max D-current Depth < info[1]) {
                bestValue \, = \, info \, [\, 0 \, ] \, ;
81
                int[] updateInfo = {bestValue, maxD-currentDepth};
                tt.put(pos.getHashCode(), updateInfo);
83
84
85
              //not less than previously found - continue down tree
86
87
              else {
                bestValue = Math.min(bestValue, max(pos, currentDepth, alpha, beta));
88
90
            //not in table - continue down tree, then update table with position's info
91
92
            else {
93
94
              bestValue = Math.min(bestValue, max(pos, currentDepth, alpha, beta));
              int[] updateInfo = {bestValue, maxD-currentDepth};
95
              tt.put(pos.getHashCode(), updateInfo);
96
            }
97
            pos.undoMove();
98
99
            //alpha beta pruning
            if (bestValue <= alpha) {
              return bestValue;
104
            beta = Math.min(beta, bestValue);
106
           catch (IllegalMoveException e) {
108
       }
       return bestValue;
112
     public int max(Position pos, int depth, int alpha, int beta) {
113
114
```

```
nodes++;
       int currentDepth = depth+1;
116
117
        if (cutOff(pos, currentDepth)){
118
119
         return utility(pos);
120
121
       int bestValue = WORSTCASE;
        //try all moves
124
        for (short move : pos.getAllMoves()) {
125
         try {
126
            pos.doMove(move);
127
128
            //if in table already
129
            if (tt.containsKey(pos.getHashCode())) {
130
131
              //get the position's value and depth info
              int info[] = tt.get(pos.getHashCode());
134
              //if position's search depth (not current depth) is less than previously found
               then update table with position's info
136
              if (maxD-currentDepth < info[1]) {</pre>
138
                bestValue = info [0];
                int[] updateInfo = {bestValue, maxD-currentDepth};
139
                tt.put(pos.getHashCode(), updateInfo);
140
141
142
              //if position's search depth >= previously found - continue down tree
143
              else {
144
                bestValue = Math.max(bestValue, mini(pos, currentDepth, alpha, beta));
145
146
           }
147
148
149
            //if not already in table - continue down tree, and then update
150
              bestValue = Math.max(bestValue, mini(pos, currentDepth, alpha, beta));
              int[] updateInfo = {bestValue, maxD-currentDepth};
153
              tt.put(pos.getHashCode(), updateInfo);
            pos.undoMove();
156
            //alpha beta pruning
            if (bestValue>=beta) {
              return bestValue;
160
161
            alpha=Math.max(alpha, bestValue);
163
            catch (IllegalMoveException e) {
165
166
        return bestValue;
167
168
```

8 Profiling

One way that I profiled my AI was by counting the number of nodes that it searched through. This was done by keeping track of an instance variable called nodes and incrementing it every time max or min was called. I then compared the number of nodes searched for, which was essentially a proxy for number of function calls. I tested this by making the AI black and the human player White. I then did the same opening move (pawn from d2 to d4) and tested how many nodes the AI searched before making its move. The output is shown below.

```
Minimax: making move 5835 rnbqkbnr/ppppppppppppp/8/8/3P4/8/PPP1PPPP/RNBQKBNR b KQkq d3 0 1 Minimax nodes searched: 1499220 making move 6834 rnbqkbnr/pp1ppppp/2p5/8/3P4/8/PPP1PPPP/RNBQKBNR w KQkq - 0 2
```

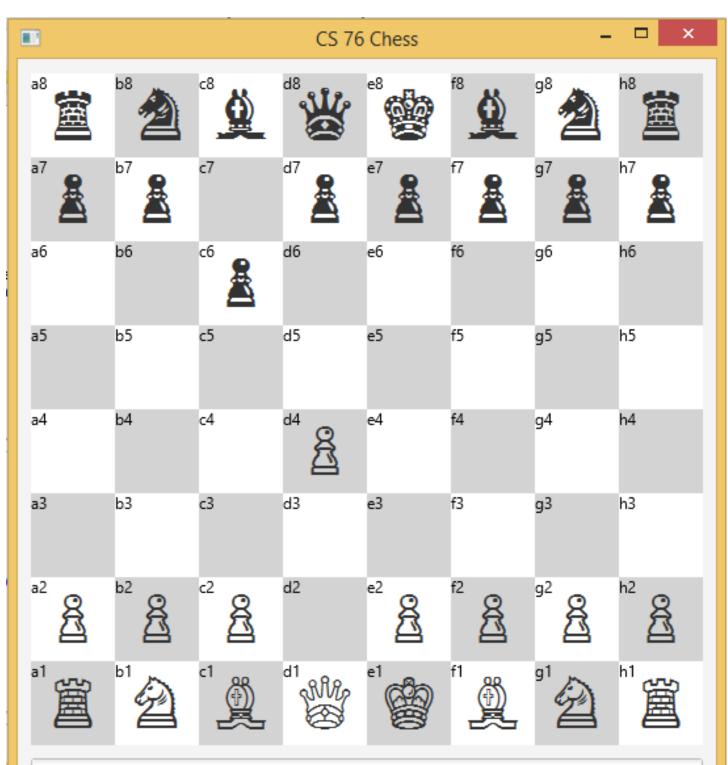
Minimax with Alpha Beta: making move 5835 rnbqkbnr/ppppppppp/8/8/3P4/8/PPP1PPPP/RNBQKBNR b KQkq d3 0 1 Alpha Beta nodes searched: 246832 making move 6834 rnbqkbnr/pp1ppppp/2p5/8/3P4/8/PPP1PPPP/RNBQKBNR w KQkq - 0

Alpha Beta and Transposition Table Minimax: making move 5835 rnbqkbnr/ppppppppp/8/8/3P4/8/PPP1PPPP/RNBQKBNR b KQkq d3 0 1 Alpha Beta W/ Transposition Table nodes searched: 66272 making move 6834 rnbqkbnr/pp1ppppp/2p5/8/3P4/8/PPP1PPPP/RNBQKBNR w KQkq - 0 2

Minimax searched around 1.5 million nodes, and with the addition of alpha beta that number was cut down to around 250,000, and with the transposition table added, it went down to around 66,000. This clearly shows that the addition of more clever search algorithms reduced the number of function calls and as a result increased the speed of the search significantly.

9 Testing

In order to test my algorithms, I made sure that they would always pick the same moves given the same opening state. I made myself the white player and manually chose the move pawn from b2 to b4. The following was the result from all three AIs, showing that they weren't pruning off branches that were important.



Welcome to CS 76 chess. Moves can be made using algebraic notation; for example the command c2c3 would move the piece at c2 to c3.