# Paraconsistent and Paracomplete Systems

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#### Abstract

Every paraconsistent or paracomplete system invalidates one of the laws of indiscernibility or one of the laws of logical explosion.

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#### Part I

# General Paraqualifications

## 1 Paracomplete Systems

A system is generally paracomplete if it invalidates one or both of the laws of indiscernibility.

#### 1.1 Indiscernibility of Identicals

$$\frac{\exists x \exists y \left[ x = y \land \exists F \left( Fx \oplus Fy \right) \right] \vdash \bot}{\exists x \exists y \left[ x = y \not\rightarrow \forall F \left( Fx \leftrightarrow Fy \right) \right] \vdash \bot} \\ \vdash \forall x \forall y \left[ x = y \rightarrow \forall F \left( Fx \leftrightarrow Fy \right) \right]}$$

#### 1.2 Identity of Indiscernibles

$$\frac{\exists x \exists y \left[ \exists F \left( Fx \leftrightarrow Fy \right) \land x \neq y \right] \vdash \bot}{\exists x \exists y \left[ \exists F \left( Fx \leftrightarrow Fy \right) \not\rightarrow x = y \right] \vdash \bot} \\ \vdash \forall x \forall y \left[ \forall F \left( Fx \leftrightarrow Fy \right) \rightarrow x = y \right]$$

### 2 Paraconsistent Systems

A system is generally paraconsistent if it invalidates some number of the laws of inconsistency.

#### 2.1 Logical Explosion

$$\frac{\exists x \exists y \left[ (x \land \neg x) \land \neg y \right] \vdash \bot}{\vdash \forall x \forall y \left[ (x \land \neg x) \rightarrow y \right]}$$

## 2.2 Logical Trivialization

$$\frac{\exists x \exists y \left[ y \wedge (\neg x \vee x) \right] \vdash \bot}{\vdash \forall x \forall y \left[ y \rightarrow (x \wedge \neg x) \right]}$$
$$\frac{\exists x \exists y \left[ y \wedge (\neg x \vee x) \right] \vdash \bot}{\vdash \forall x \forall y \left[ (\neg x \vee x) \rightarrow \neg y \right]}$$
$$\vdash \forall x \forall y \left[ \neg (x \wedge \neg x) \rightarrow \neg y \right]$$