# CS2040S Data Structures and Algorithms

(e-learning edition)

Welcome!

### Announcements

Lectures are now available via video (only).

Tutorials and Recitations continue in person.

Watch the lecture before recitation!

# Plan of the Day

#### Trees

- Terminology
- Traversals
- Operations

#### **Balanced Trees**

- Height-balanced binary search trees
- AVL trees
- Rotations

### Dictionary Interface

### A collection of (key, value) pairs:

interface	IDictionar	
void	insert(Key k, Value v)	insert (k,v) into table
Value	search(Key k)	get value paired with k
Key	successor(Key k)	find next key > k
Key	predecessor(Key k)	find next key < k
void	delete(Key k)	remove key k (and value)
boolean	contains(Key k)	is there a value for k?
int	size()	number of (k,v) pairs

### Dictionary

### **Implementation**

Option 1: Sorted array

- insert: add to middle of array --- O(n)
- search: binary search through array --- O(log n)

#### Option 2: Linked list

- insert: add to middle of array --- O(n)
- search : no binary search in array --- O(n)

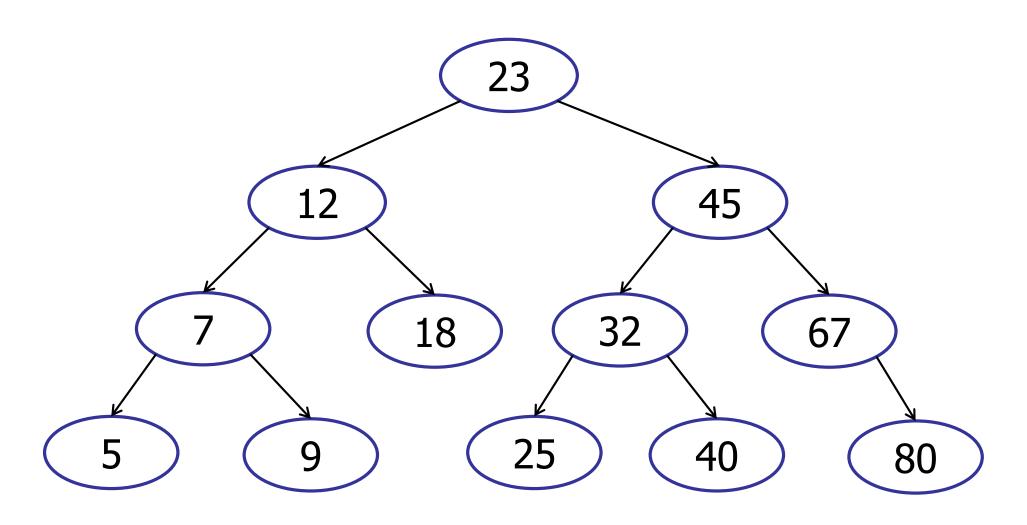
### Dictionary Implementation

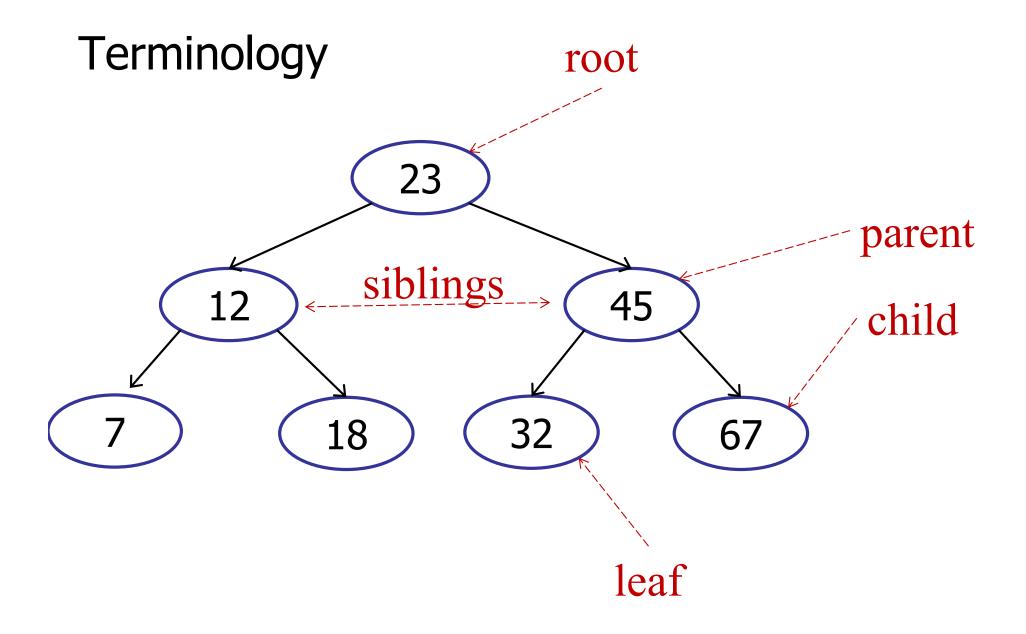
#### Possible Choices:

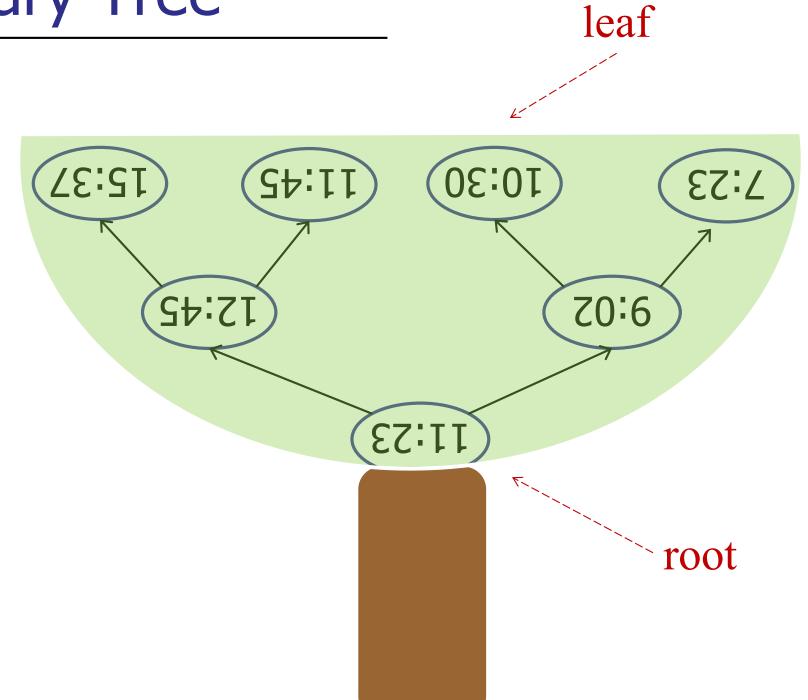
- Implement using an array (see: java.util.ArrayList).
- Implement using an array (see: java.util.Vector).
- Implement using a queue.
- Implement using a LinkedList
- ...
- Implement using a tree.

### Dictionary

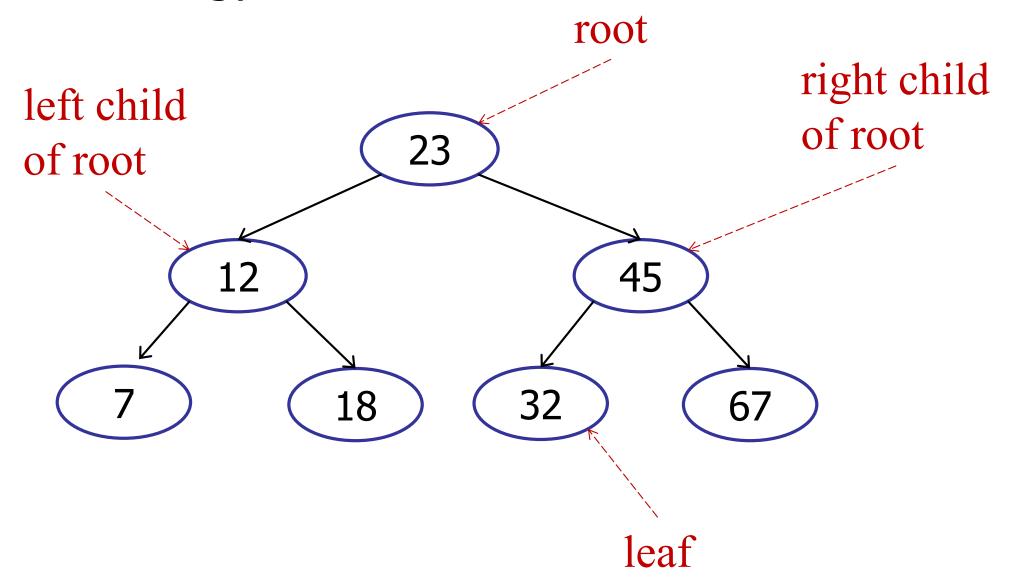
Implementation idea: Tree



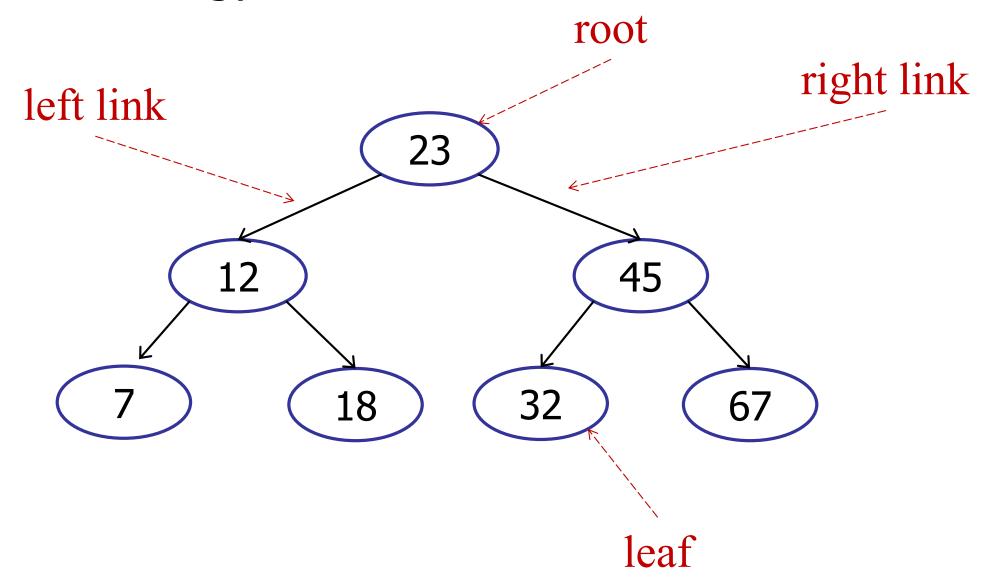


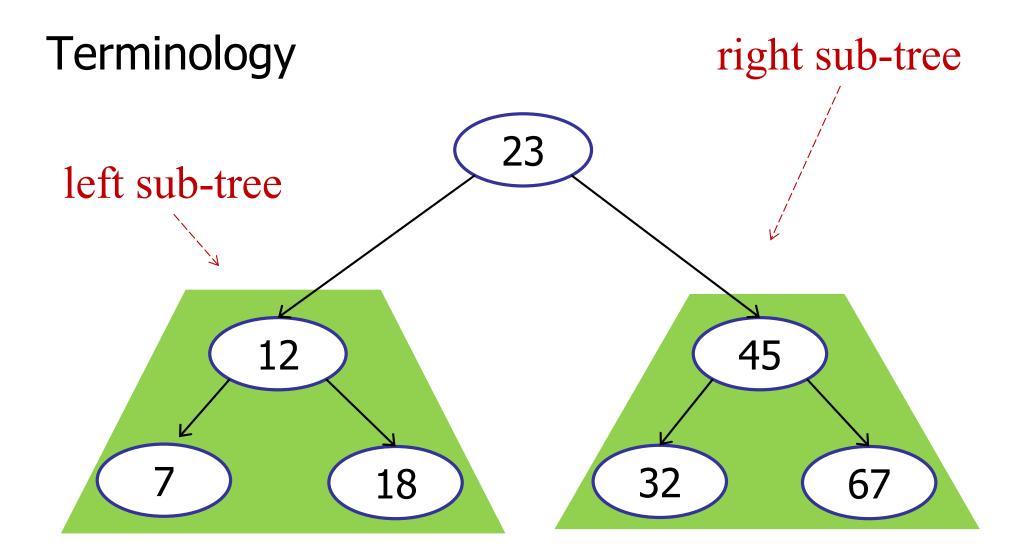


### Terminology



### Terminology





**Recursive Definition** right sub-tree 23 left sub-tree

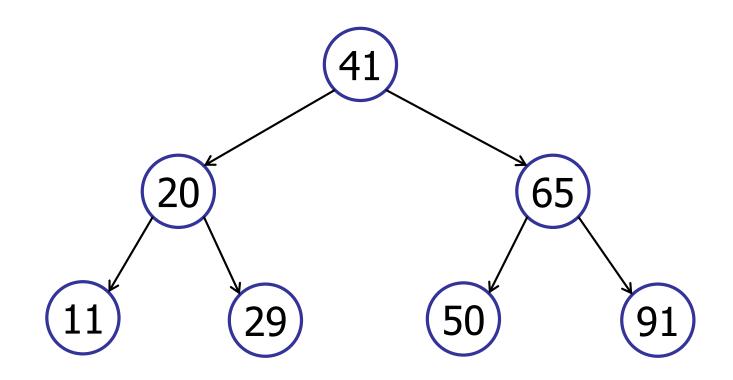
A binary tree is either:

- (a) empty
- (b) a node pointing to two binary trees

#### Java??

```
public class BinaryTree {
      private BinaryTree leftTree;
      private BinaryTree rightTree;
      private KeyType key;
      private ValueType value;
       // Remainder of binary tree implementation
```

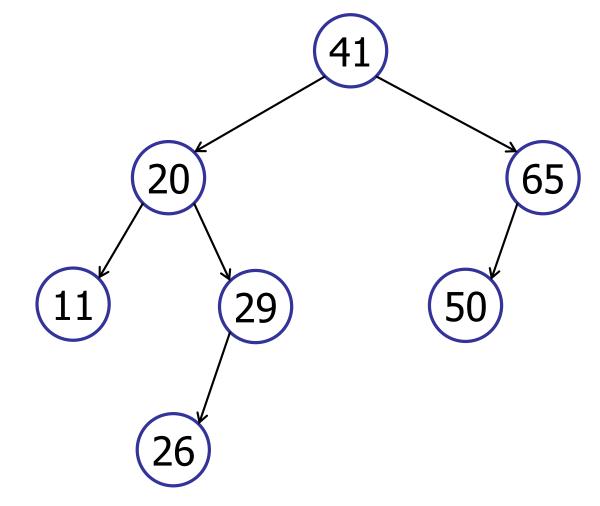
### Binary Search Trees (BST)



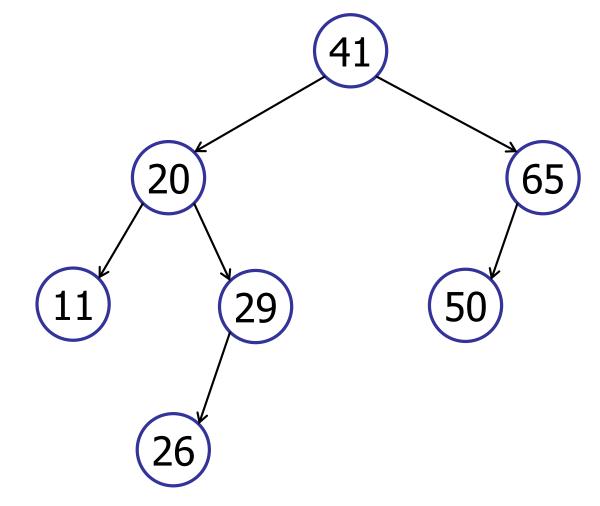
#### **BST Property:**

all in left sub-tree < key < all in right sub-right

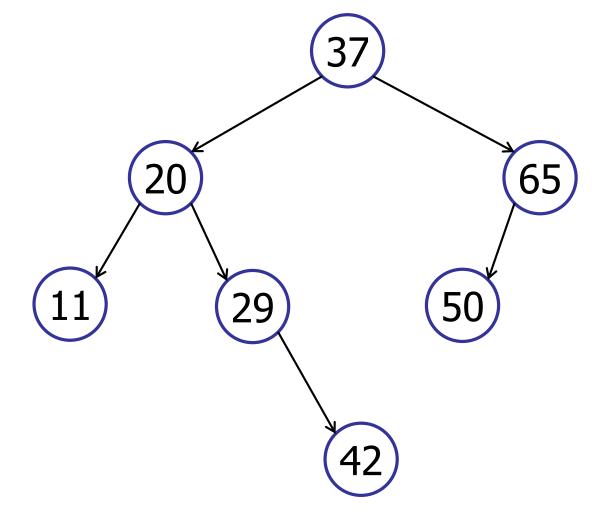
- 1. Yes
- 2. No
- 3. I don't know.



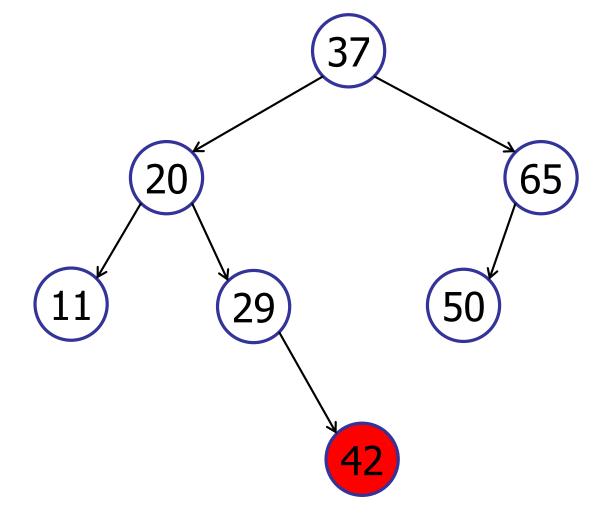
- ✓ 1. Yes
  - 2. No
  - 3. I don't know.



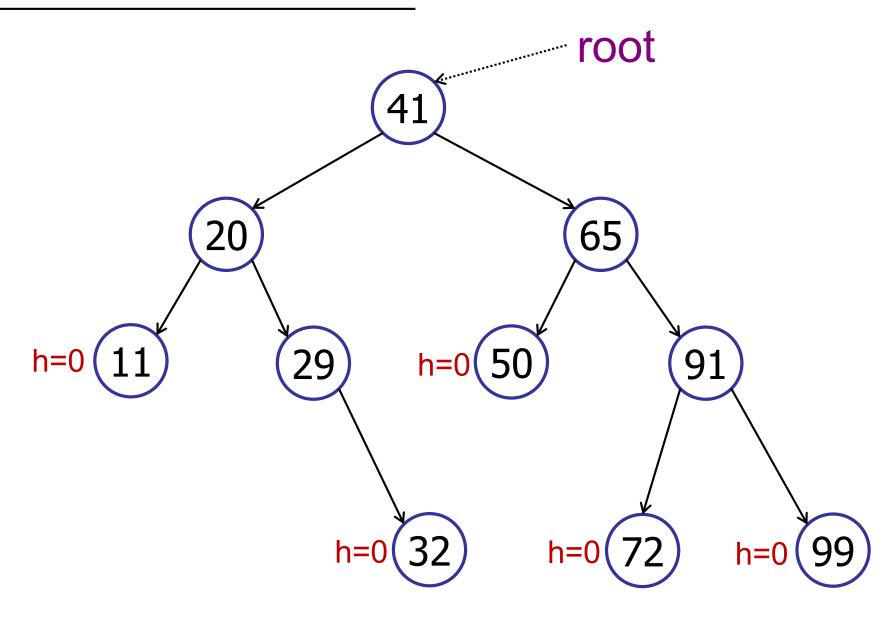
- 1. Yes
- 2. No
- 3. I don't know.

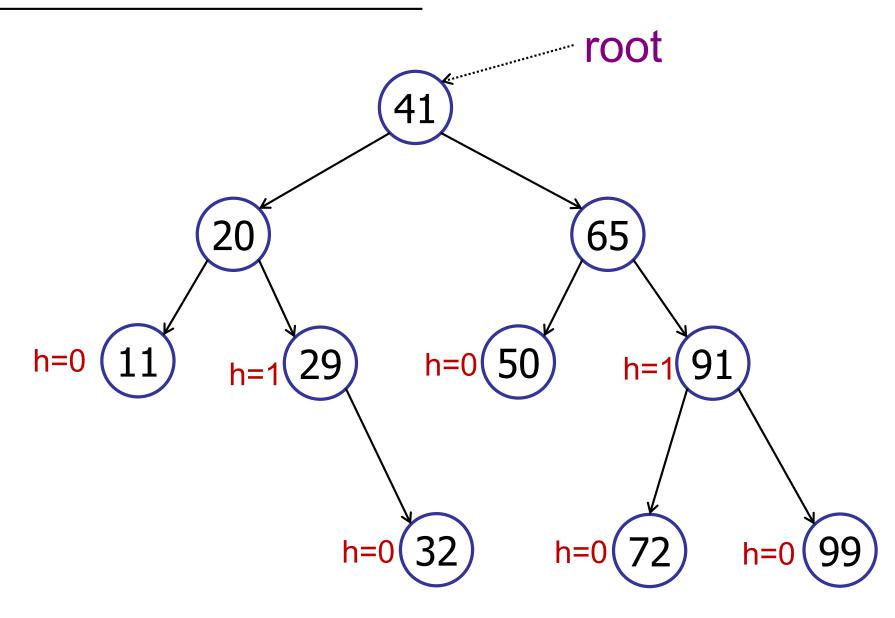


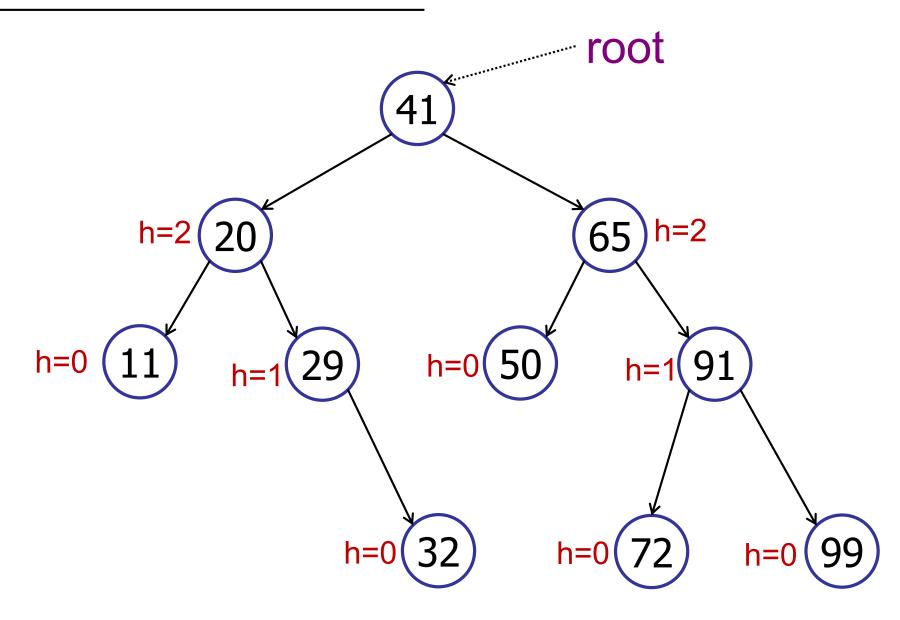
- 1. Yes
- **✓**2. No
  - 3. I don't know.

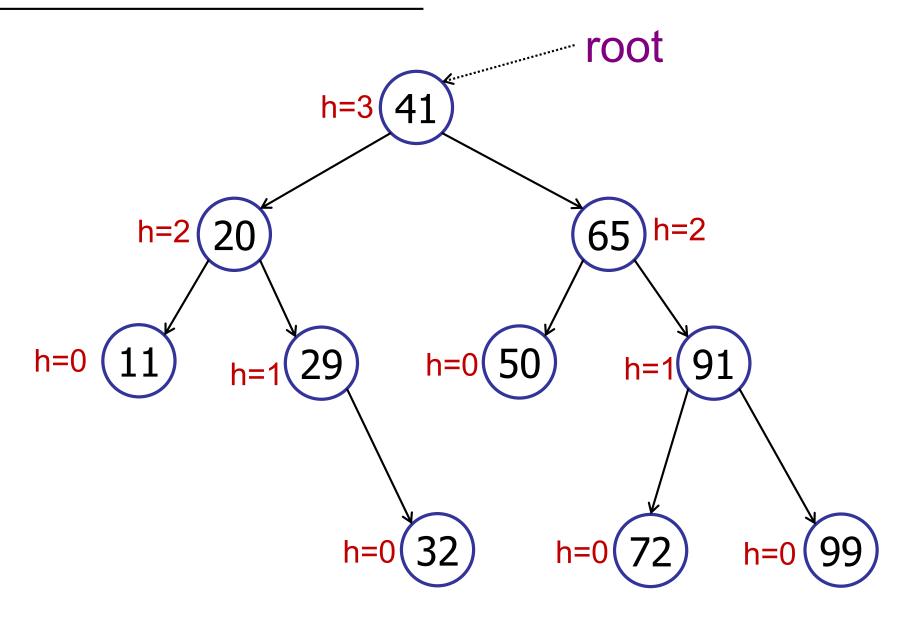


- 1. Terminology and Definitions
- 2. Basic operations:
  - height
  - search, insert
  - searchMin, searchMax
- 3. Traversals
  - in-order, pre-order, post-order
- 4. Other operations



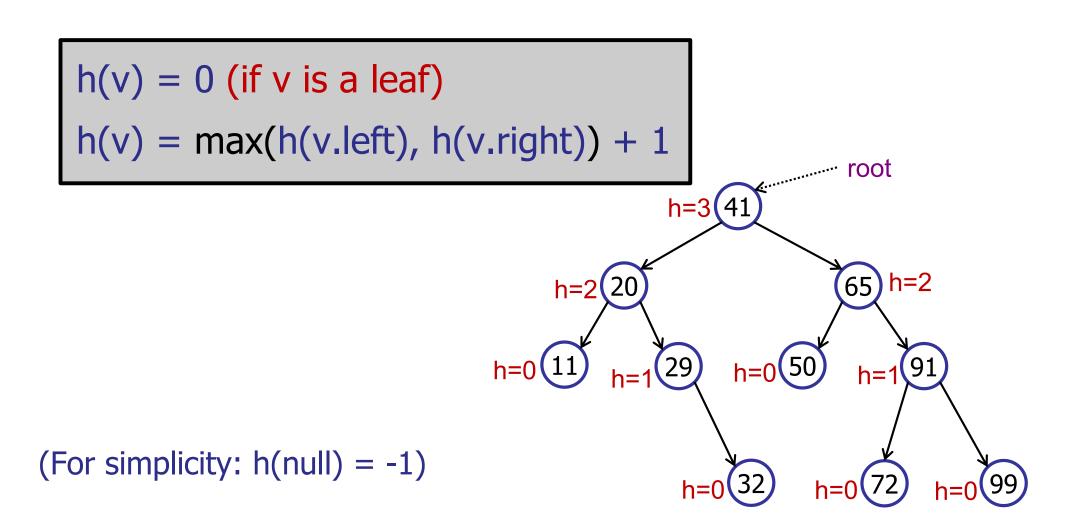






#### Height:

Number of edges on longest path from root to leaf.



#### Calculating the heights

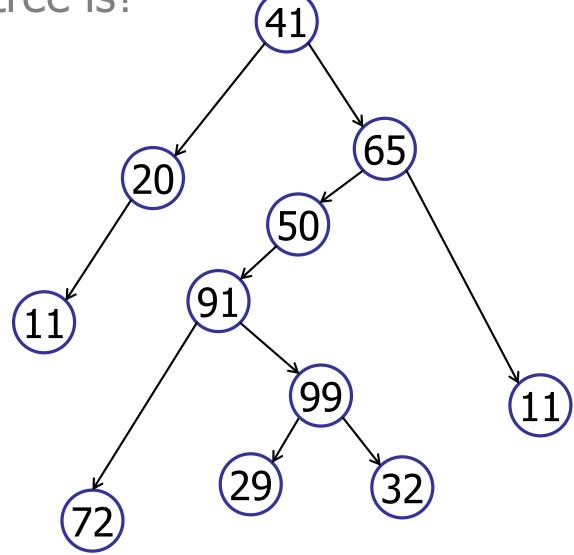
#### check for null

```
public int height() {
      int leftHeight = -1;
       int rightHeight = -1;
       if (m leftTree != null)
             leftHeight = m leftTree.height();
       if (m rightTree != null)
             rightHeight = m rightTree.height();
       return max(leftHeight, rightHeight) + 1;
```

max of subtrees

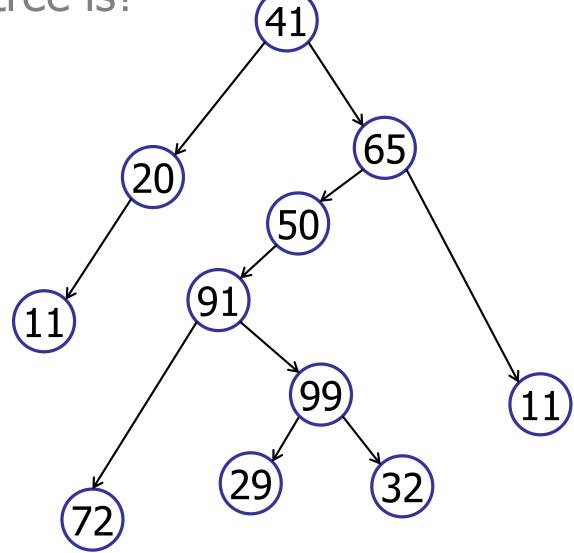
The height of this tree is?

- 1. 2
- 2. 4
- 3. 5
- 4. 6
- 5. 7
- 6. 42



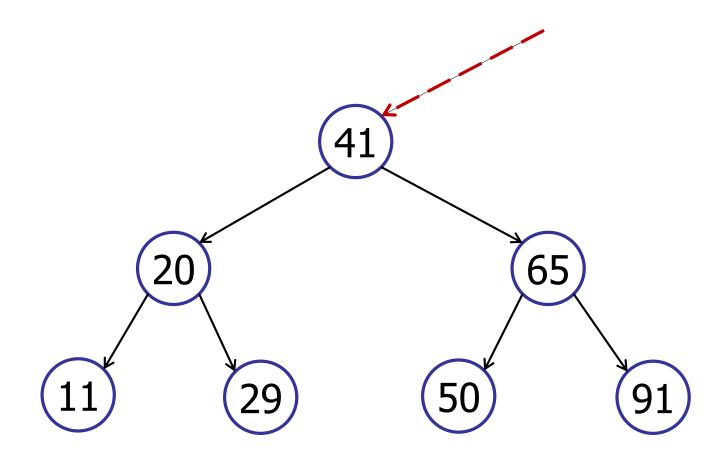
The height of this tree is?

- 1. 2
- 2. 4
- **√**3. 5
  - 4. 6
  - 5. 7
  - 6. 42

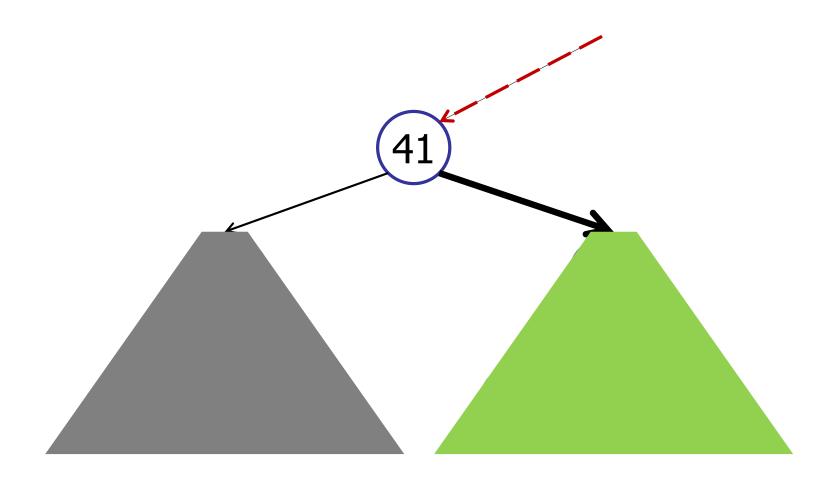


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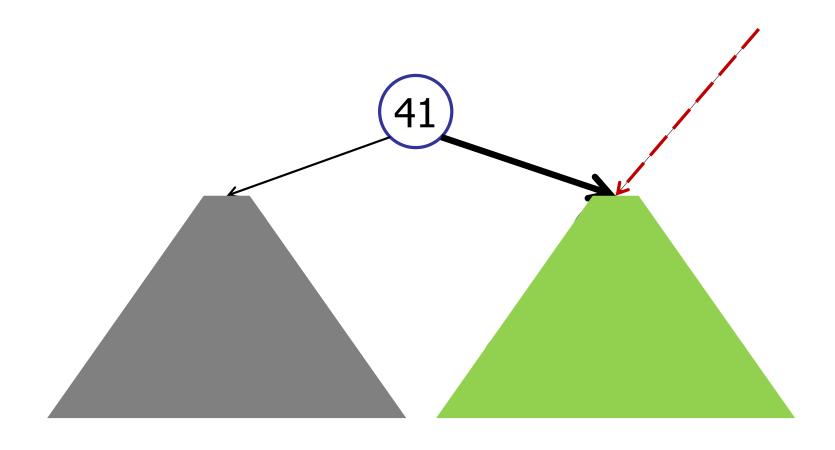
Search for the maximum key:



Search for the maximum key:



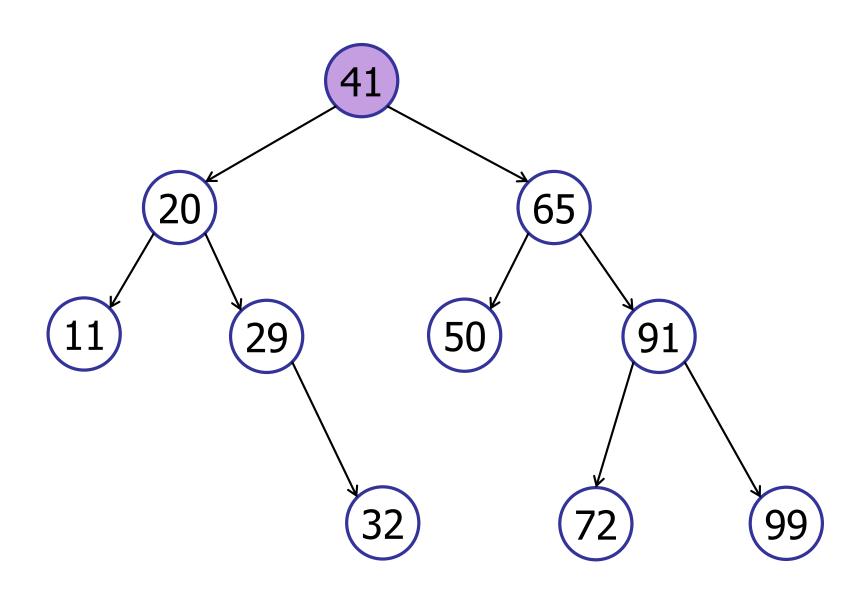
Search for maximum key



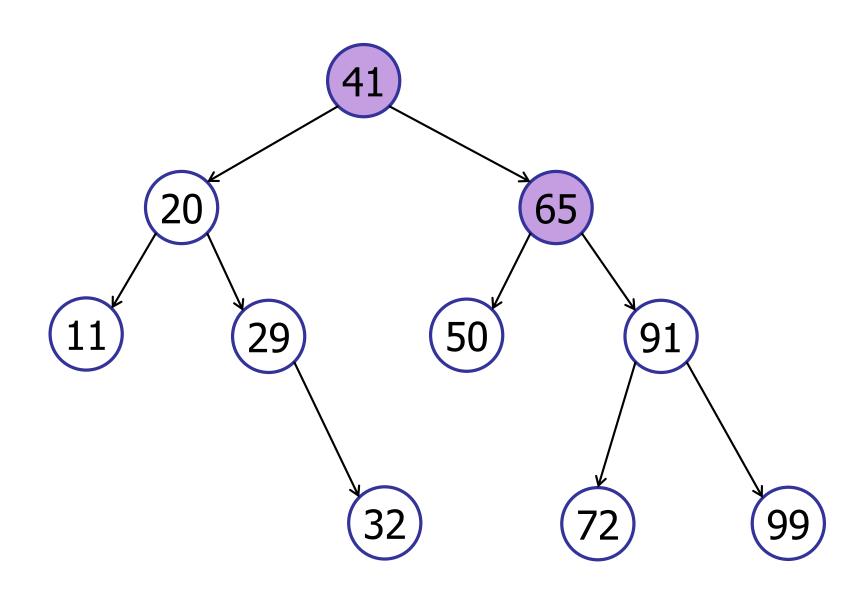
#### Searching for the maximum key

```
public BinaryTree searchMax() {
    if (rightTree != null) {
        return rightTree.searchMax();
    }
    else return this; // Key is here!
}
```

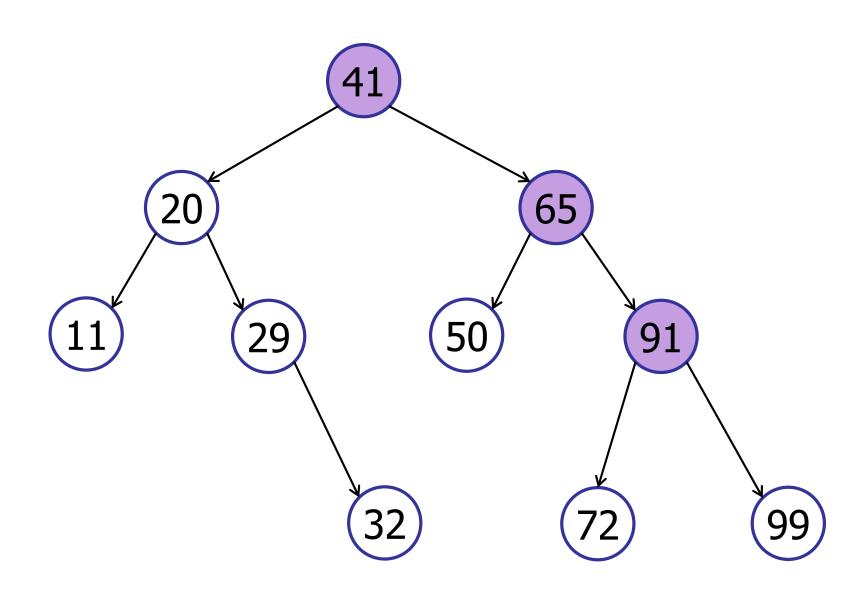
searchMax()



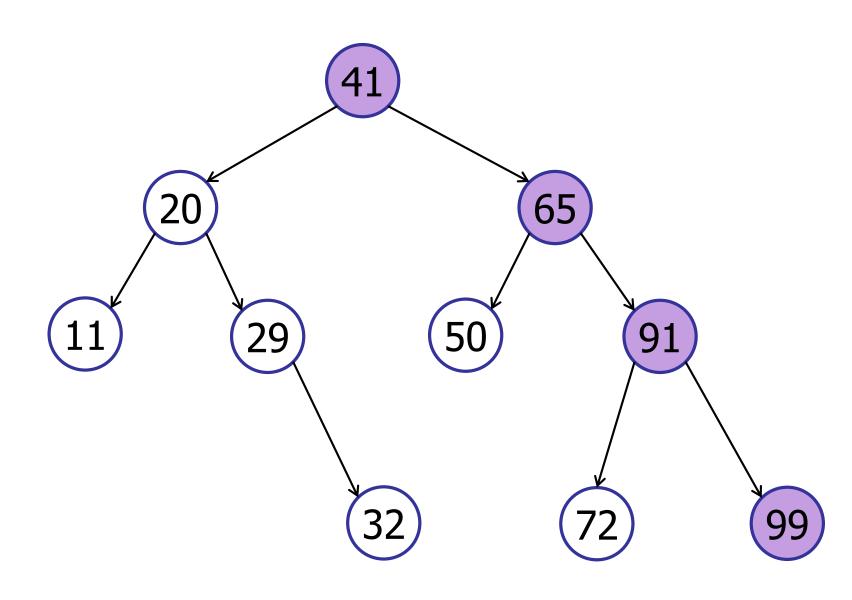
searchMax()



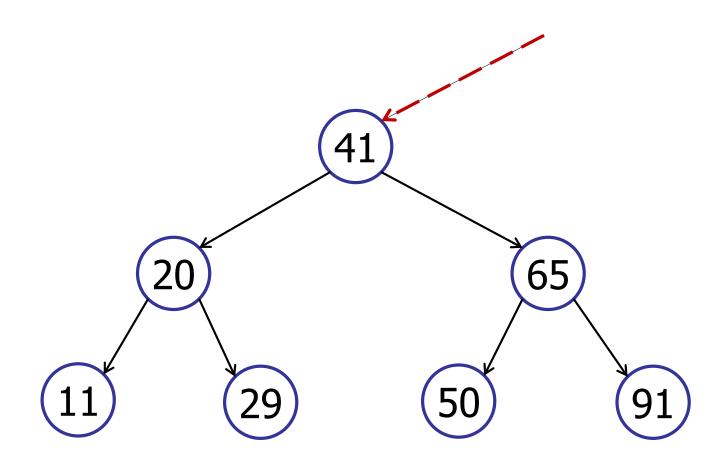
searchMax()



searchMax()



Search for the minimum key:

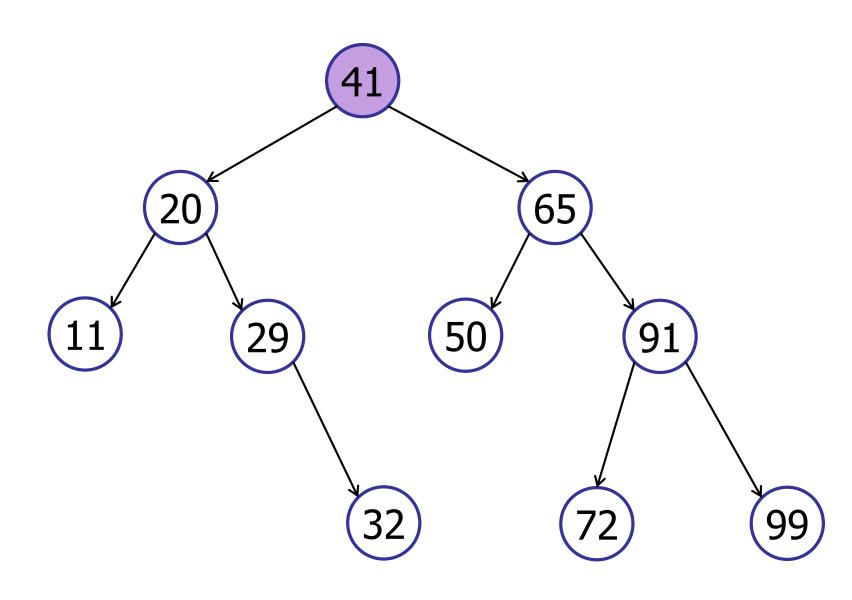


### Binary Tree

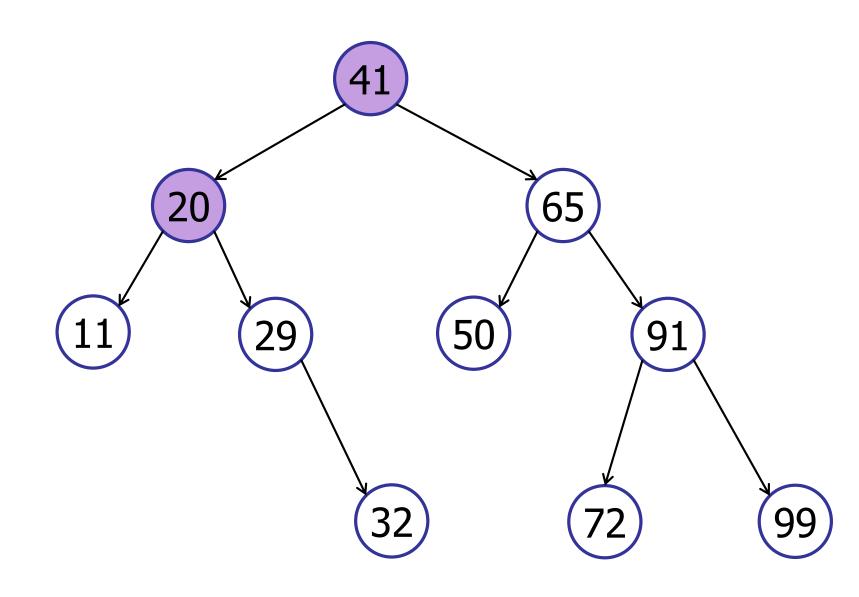
#### Searching for the minimum key

```
public BinaryTree searchMin() {
    if (m_leftTree != null) {
        return leftTree.searchMin();
    }
    else return this; // Key is here!
}
```

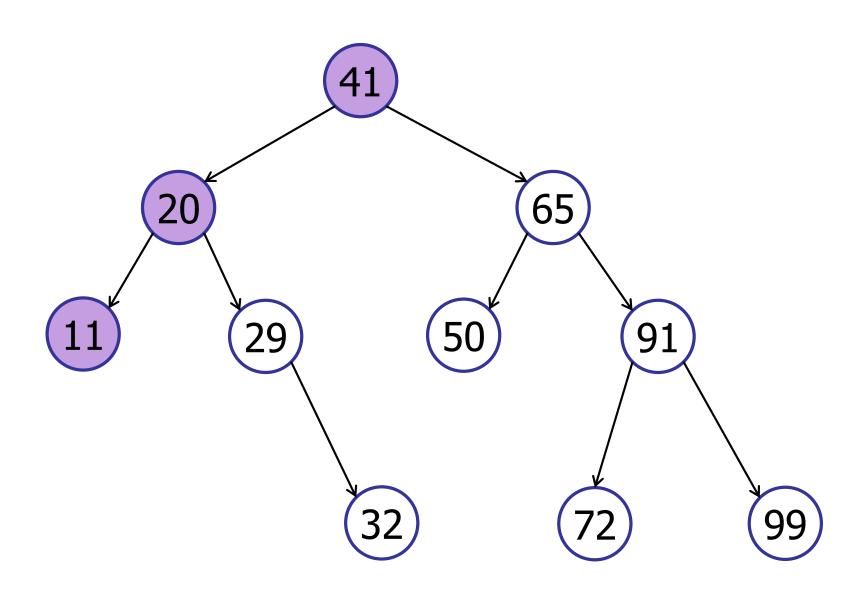
searchMin()



searchMin()

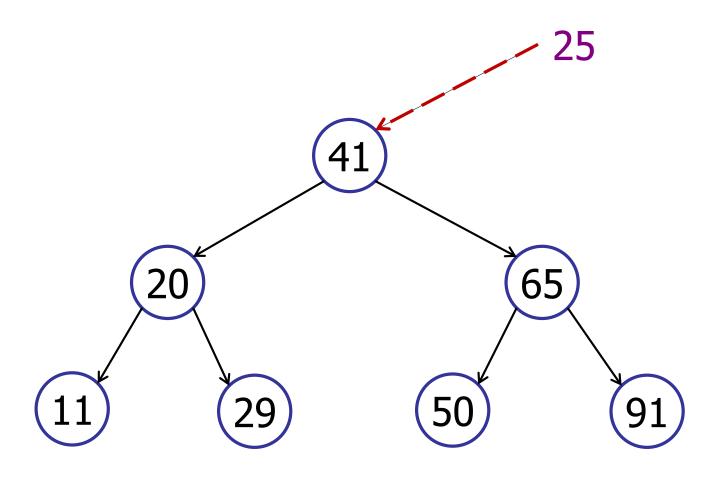


searchMin()

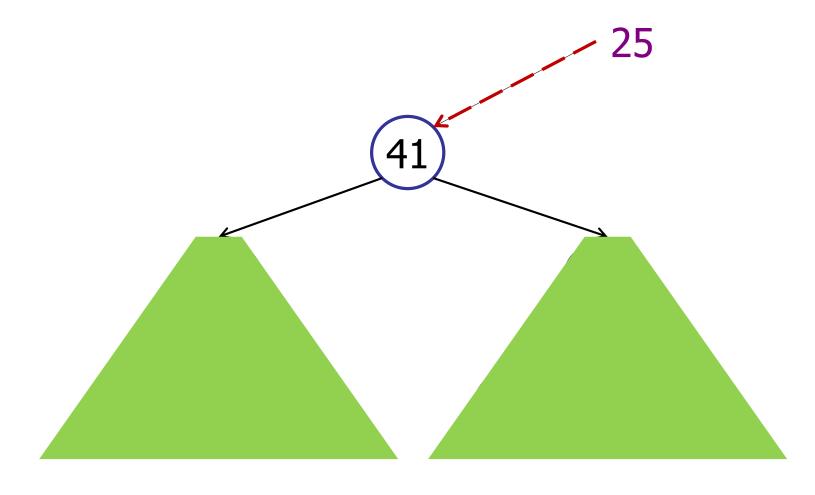


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  - search, insert
- 3. Traversals
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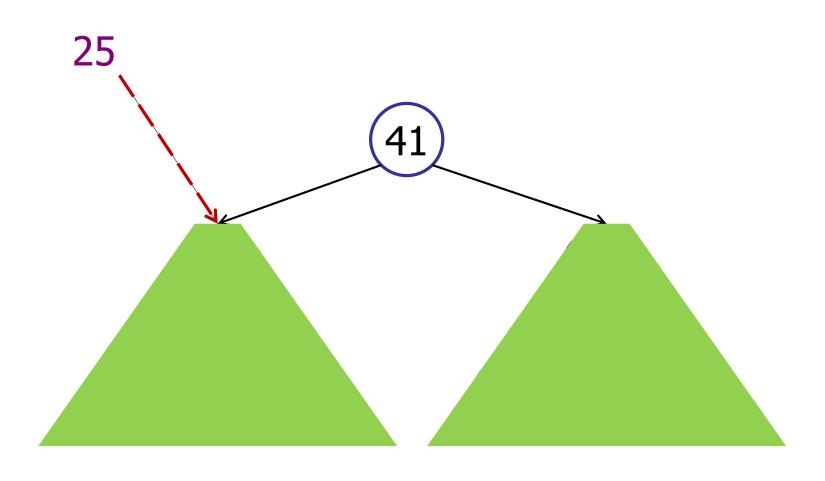
Search for a key:



Search for a key: 25 < 41



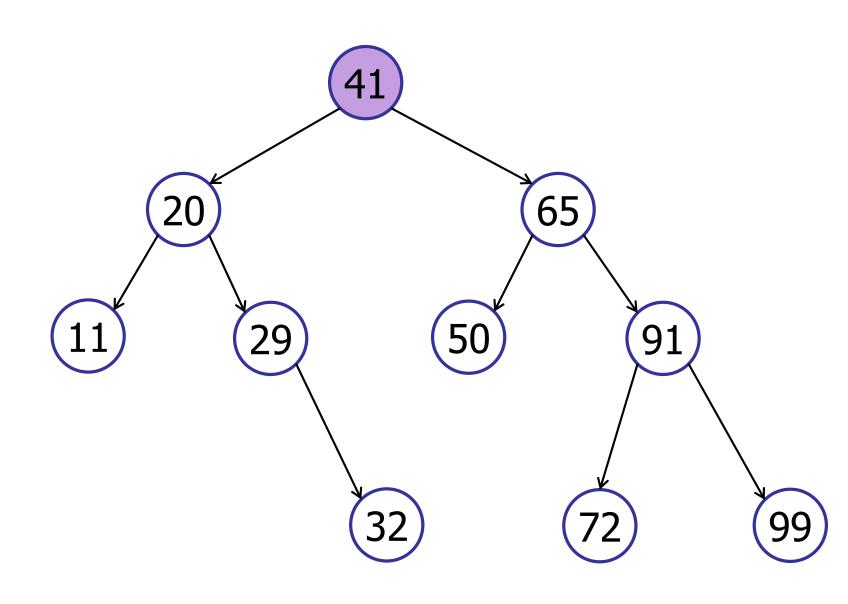
#### Search for a key:

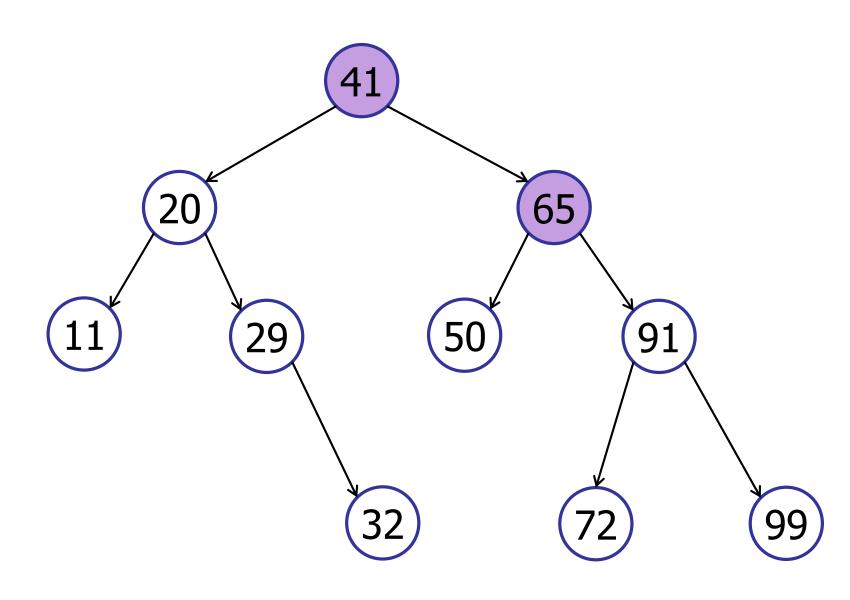


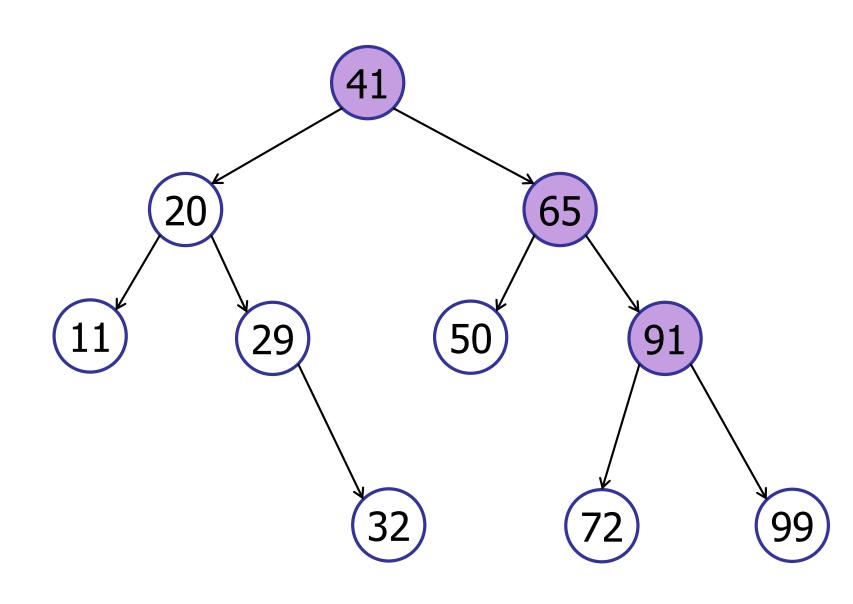
### Binary Tree

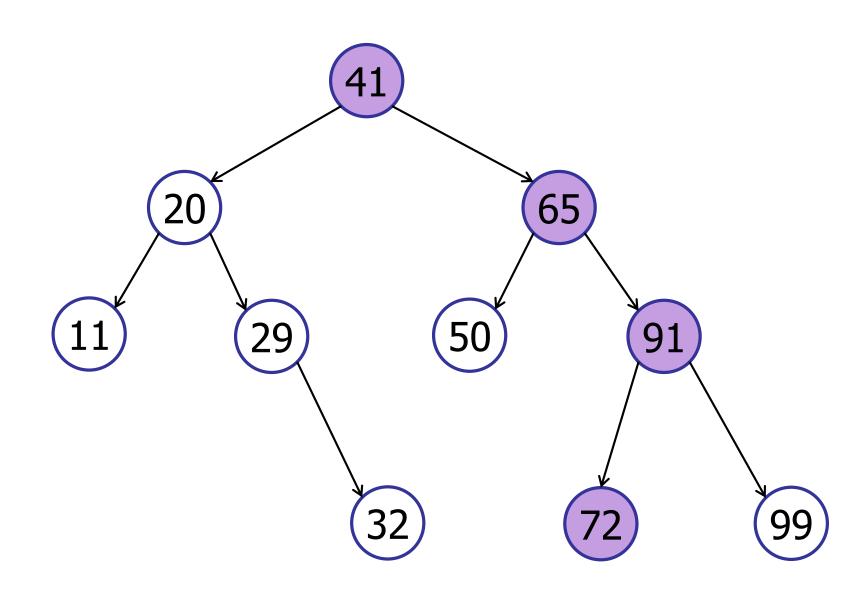
#### Inserting a new key

```
public BinaryTree search(KeyType queryKey) {
       if (queryKey.compareTo(key) < 0) {</pre>
              if (leftTree != null)
                     return leftTree.search(key);
              else return null;
       else if (queryKey.compareTo(key) > 0) {
              if (rightTree != null)
                     return rightTree.search(key);
              else return null:
       else return this; // Key is here!
```



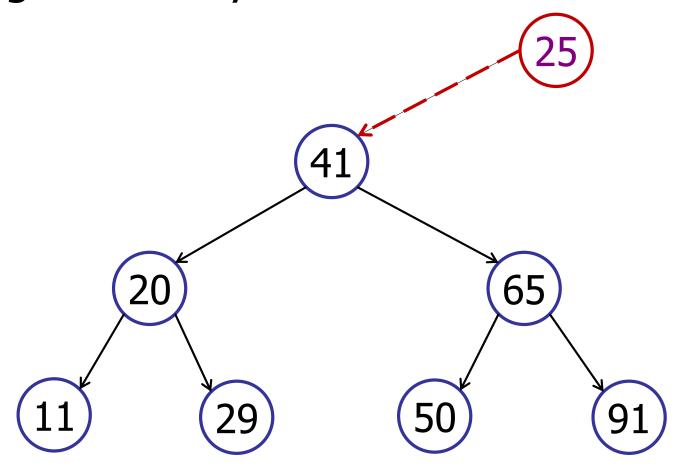






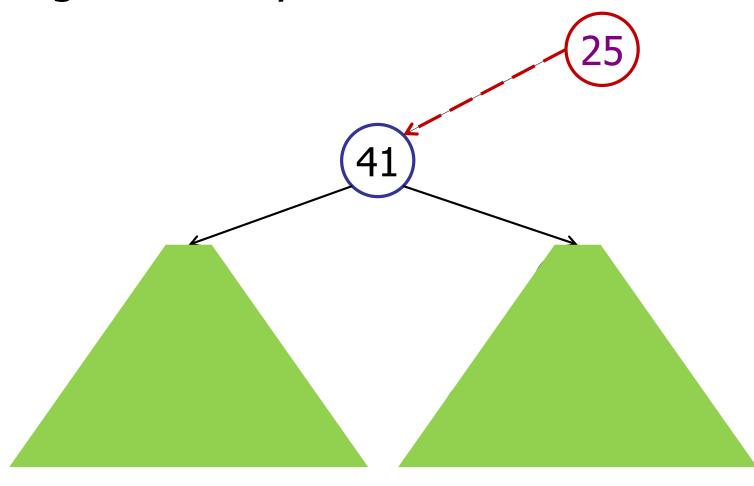
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Inserting a new key:

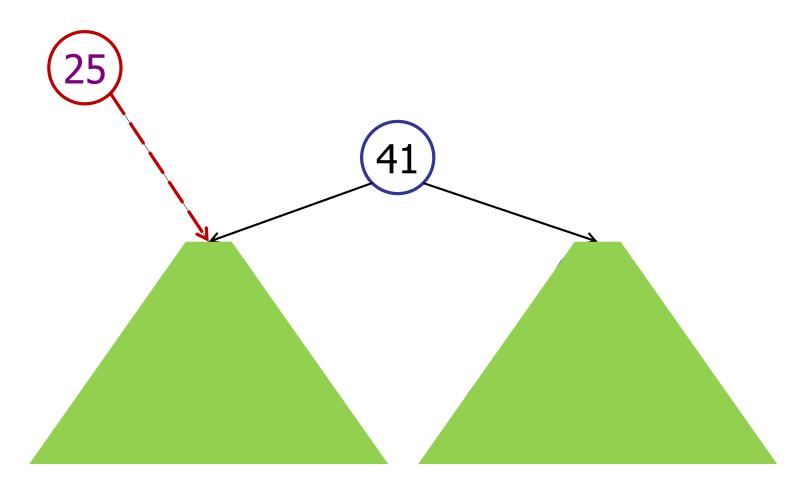


25 < 41

Inserting a new key:



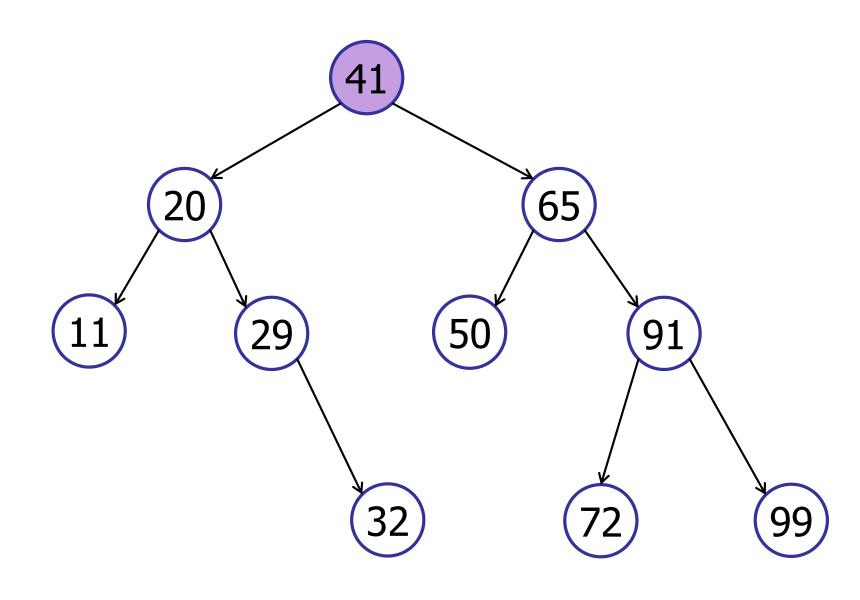
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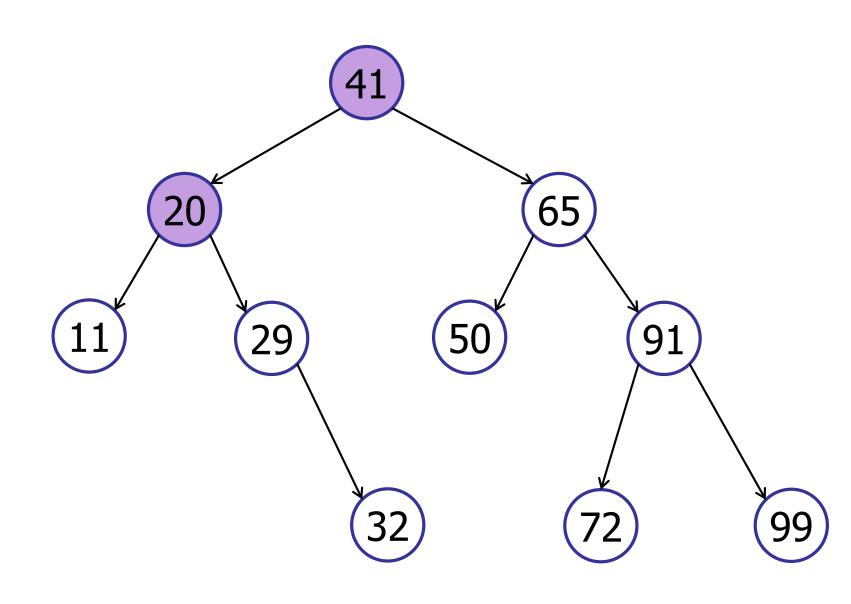


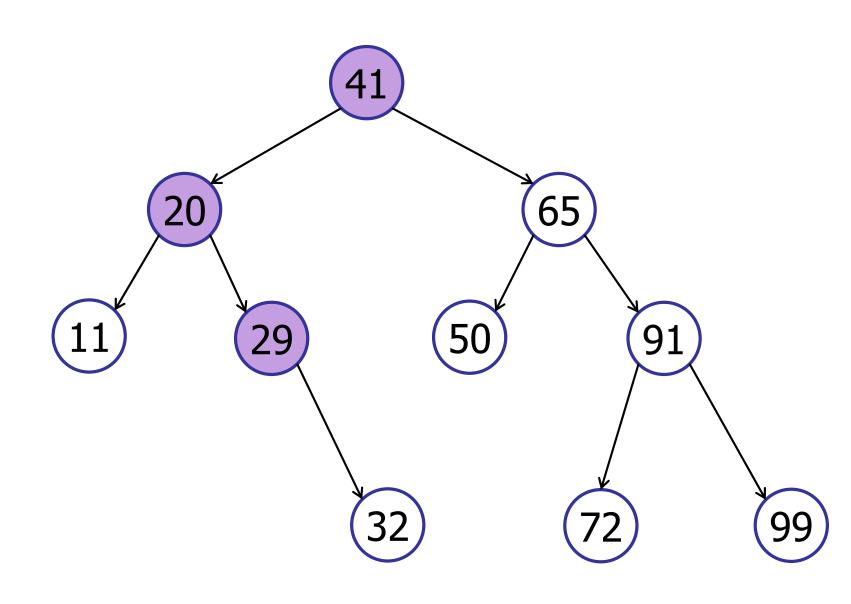
#### Binary Tree

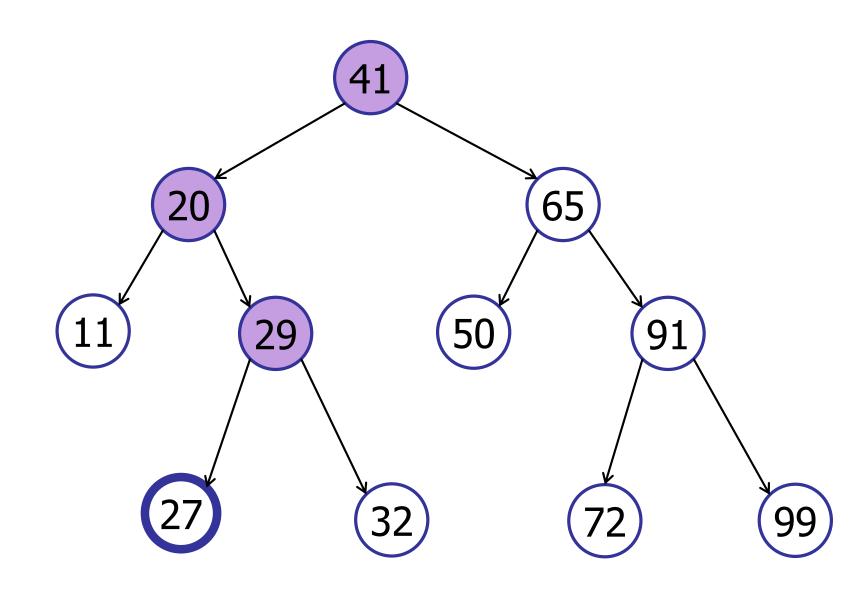
#### Inserting a new key

```
public void insert(Key insKey) {
       if (insKey.compareTo(key) < 0) {
             if (leftTree != null)
                    leftTree.insert(insKey);
             else leftTree = new BinaryTree(insKey);
      else if (insKey.compareTo(key) > 0) {
             if (rightTree != null)
                    rightTree.insert(insKey);
             else rightTree = new BinaryTree(insKey);
      else return; // Key is already in the tree!
```









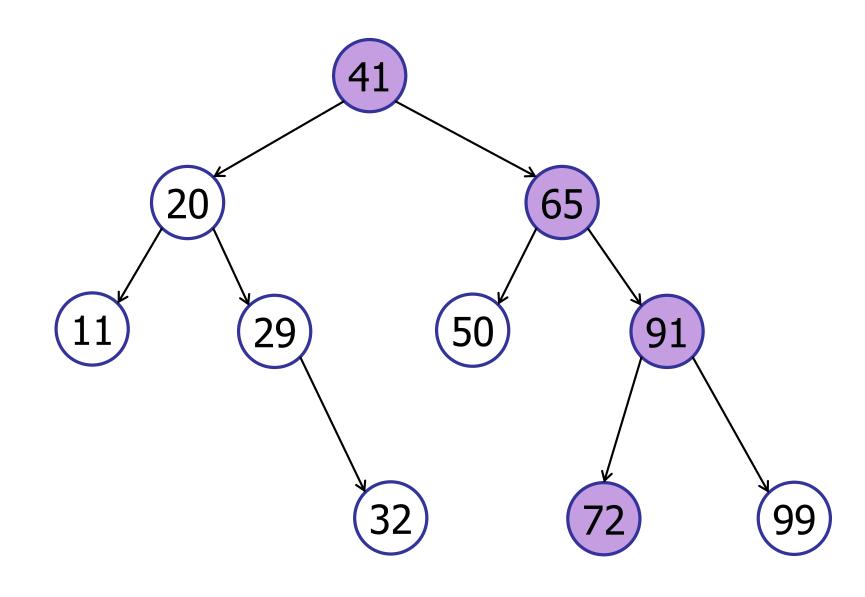
What is the worst-case running time of search in a BST?

- 1. O(1)
- 2. O(log n)
- 3. O(n)
- 4.  $O(n^2)$
- 5.  $O(n^3)$
- 6.  $O(2^n)$

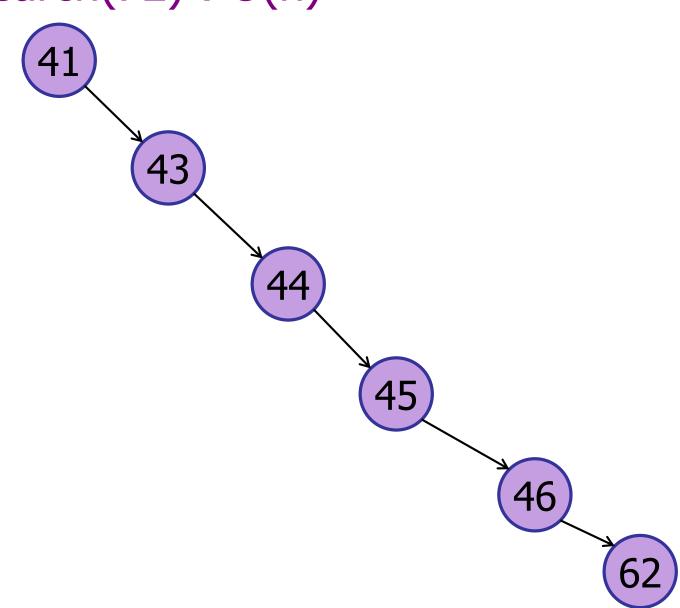
What is the worst-case running time of search in a BST?

- 1. O(1)
- 2. O(log n)
- **✓**3. O(n)
  - 4.  $O(n^2)$
  - 5.  $O(n^3)$
  - 6.  $O(2^n)$

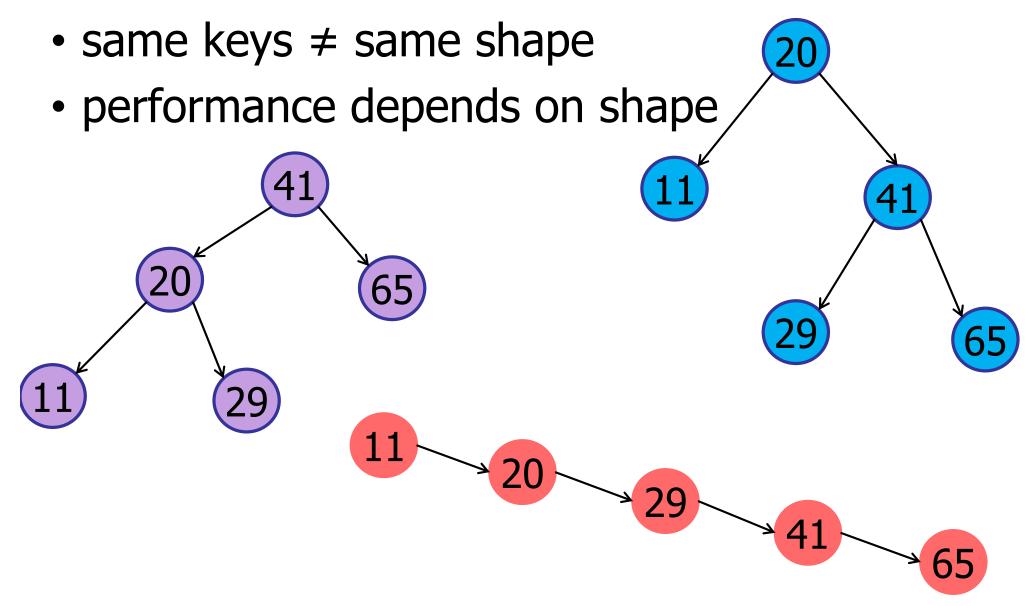
search(72) : O(h)



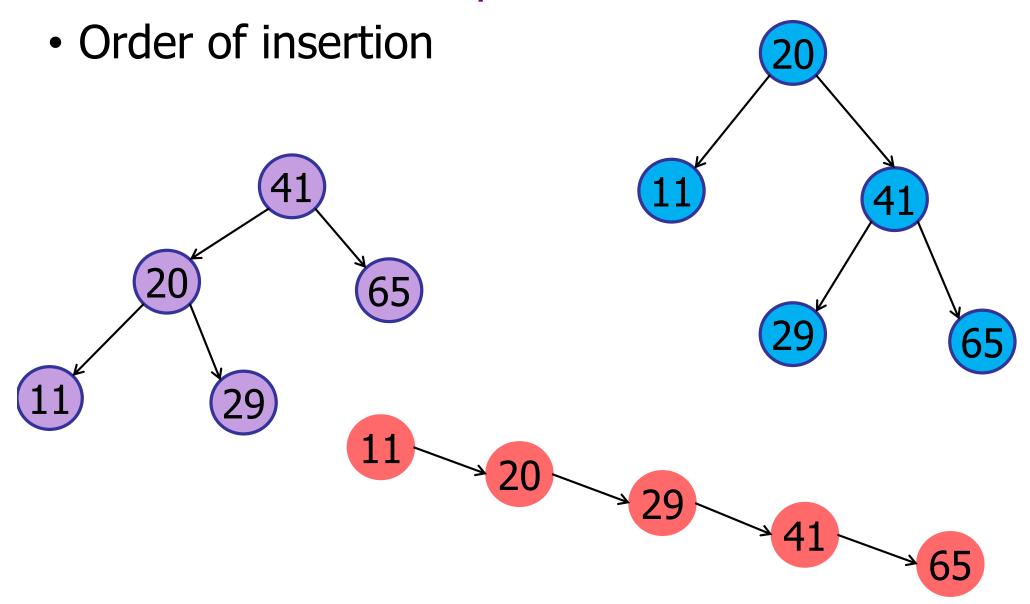
search(72) : O(h)



#### Trees come in many shapes



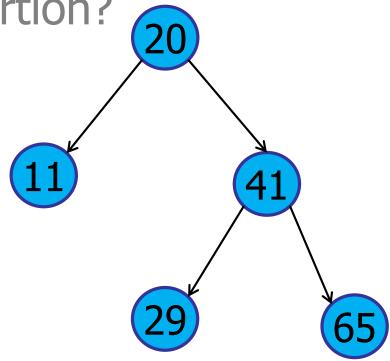
#### What determines shape?



What was the order of insertion?



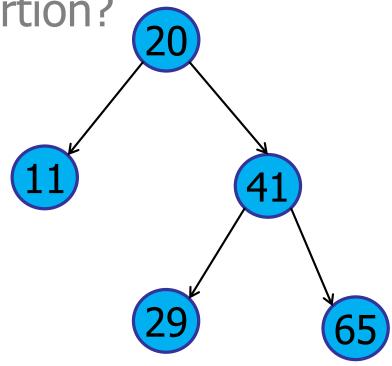
- 2. 20, 11, 41, 29, 65
- 3. 11, 20, 41, 29, 65
- 4. 65, 41, 29, 20, 11
- 5. Impossible to tell.



What was the order of insertion?

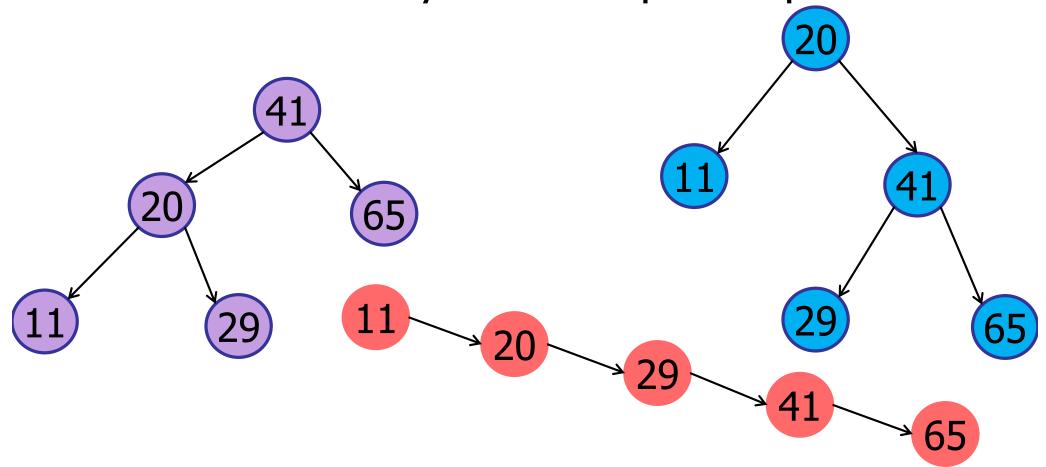


- **✓** 2. 20, 11, 41, 29, 65
  - 3. 11, 20, 41, 29, 65
  - 4. 65, 41, 29, 20, 11
  - 5. Impossible to tell.



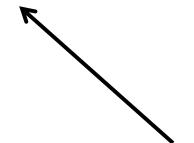
#### What determines shape?

- Order of insertion
- Does each order yield a unique shape?



#### What determines shape?

- Order of insertion
- Does each order yield a unique shape? NO
  - # ways to order insertions: n!
  - − # shapes of a binary tree? ~4<sup>n</sup>



Catalan Numbers

#### Catalan Numbers

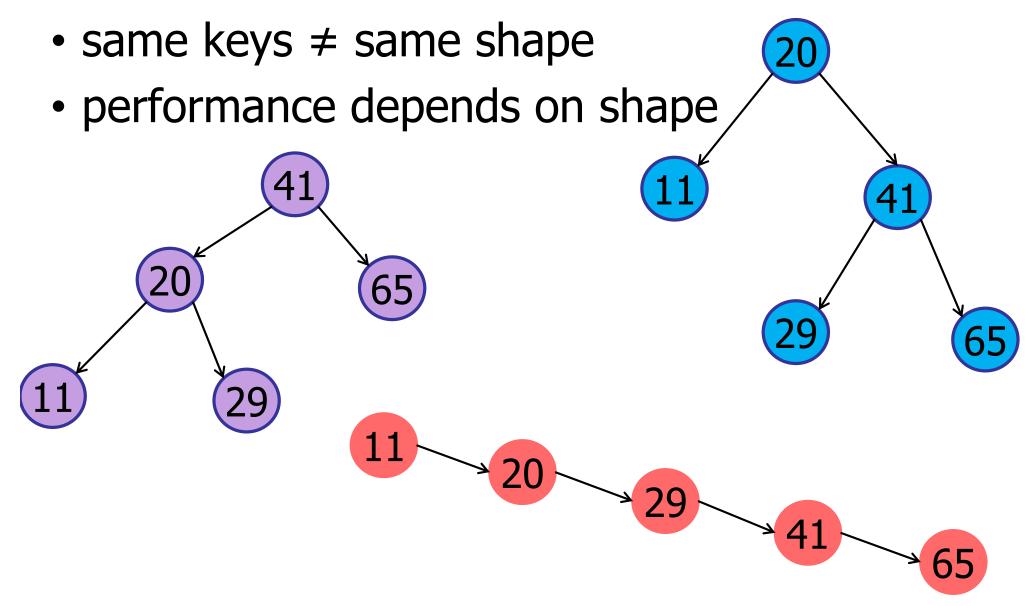
 $C_n = \#$  of trees with (n+1) leaves

C<sub>n</sub> = # expressions with n pairs of matched parentheses

((())) ()(()) (()()) (()())

Why are these the same?

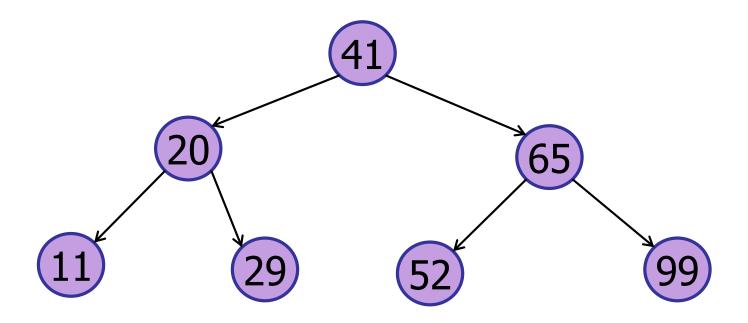
#### Trees come in many shapes



# Tree Shape

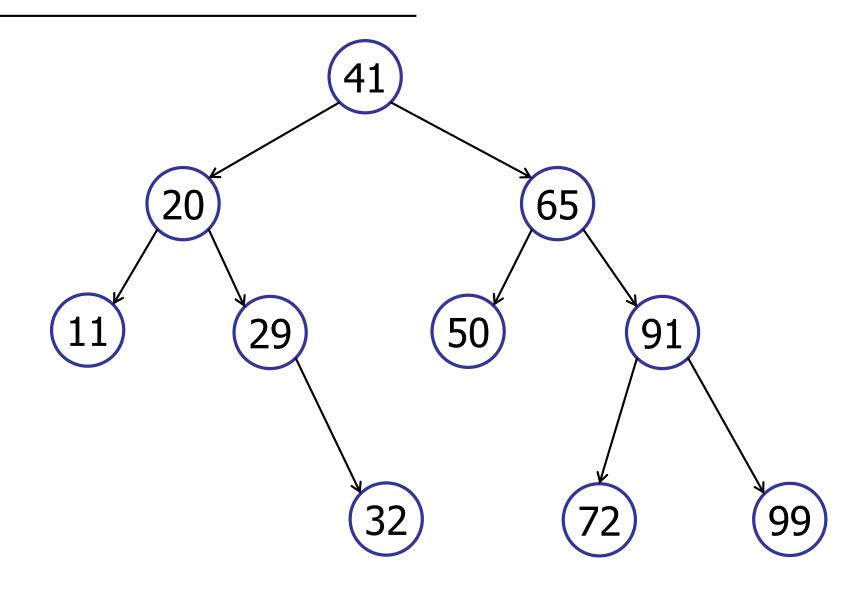
#### Trees come in many shapes

- same keys ≠ same shape
- performance depends on shape
- insert keys in a random order ⇒ balanced

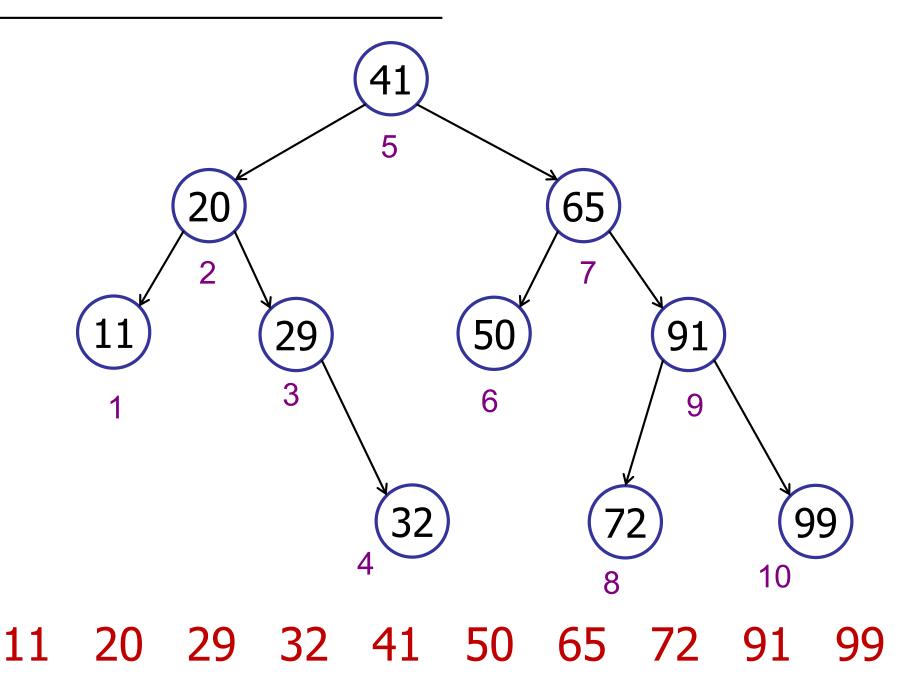


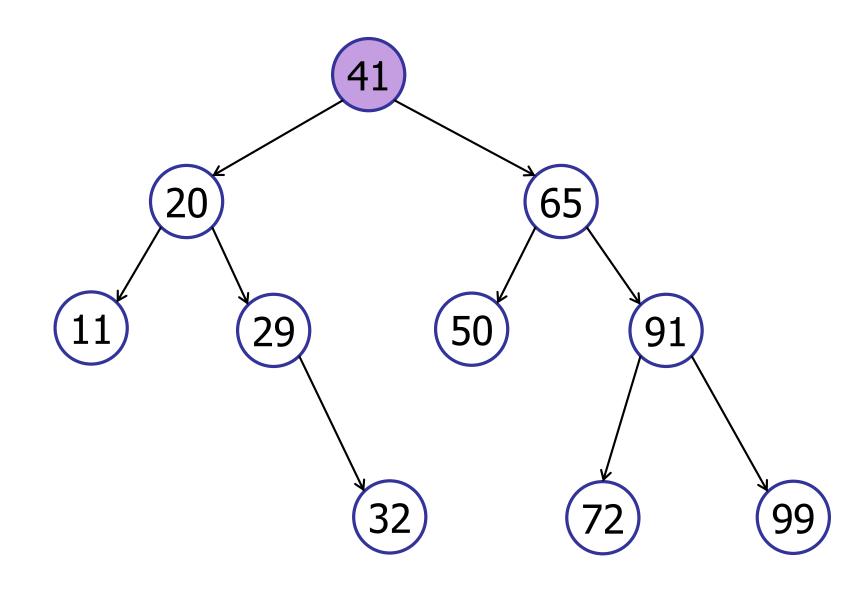
# Binary Search Trees

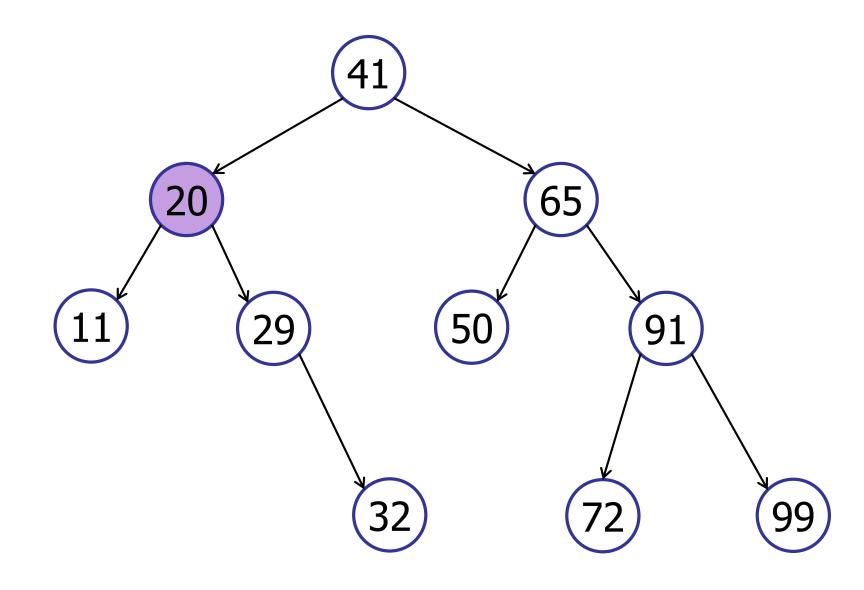
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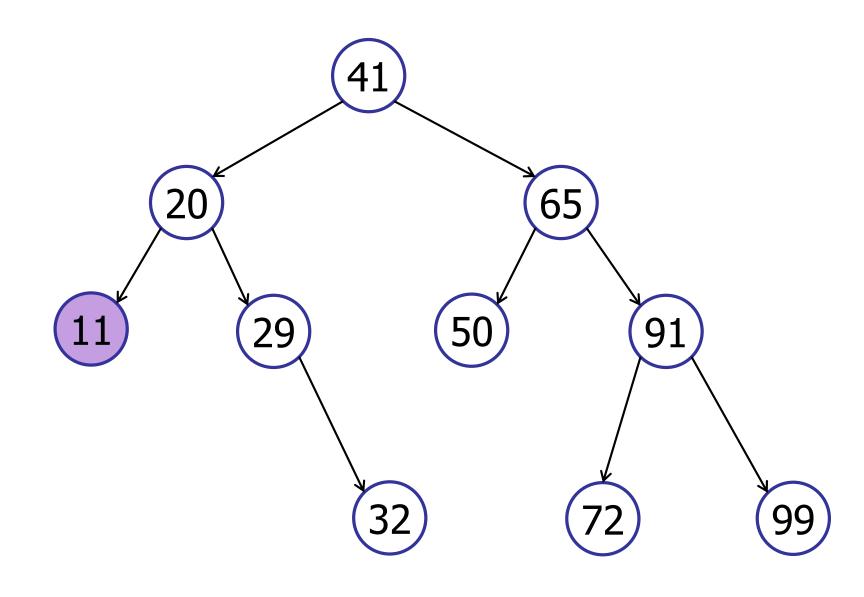


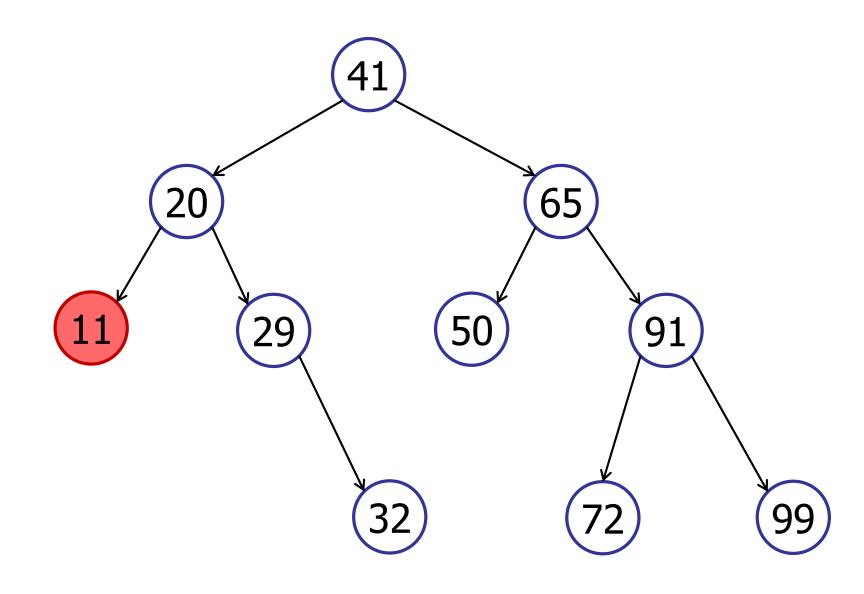
11 20 29 32 41 50 65 72 91 99

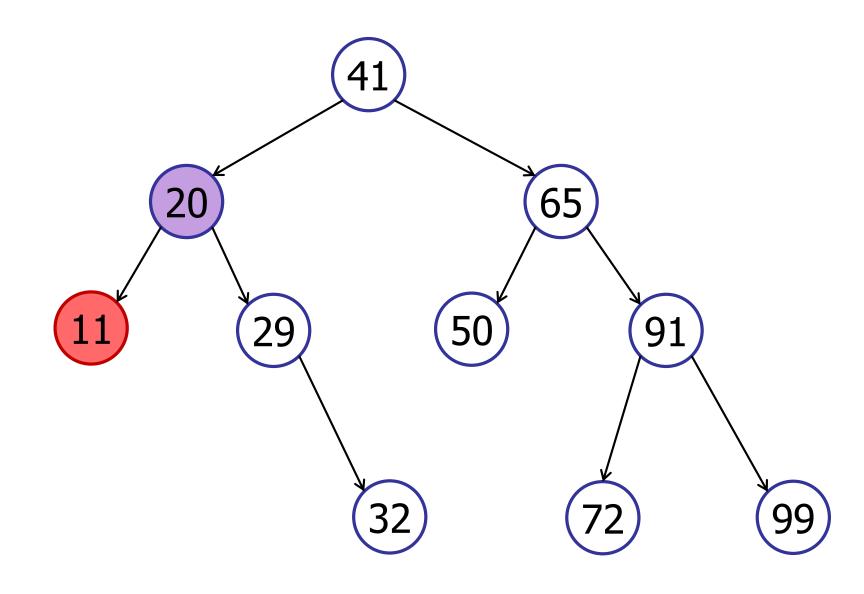


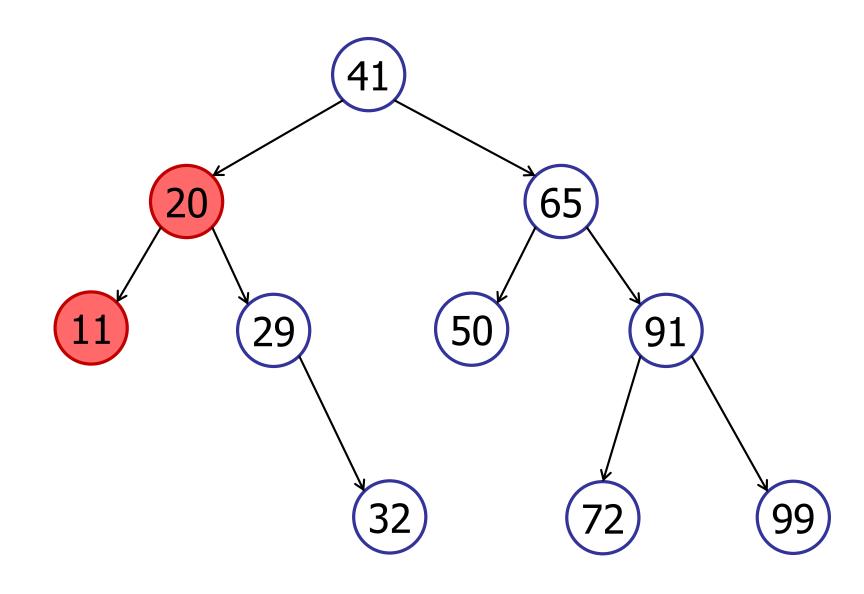


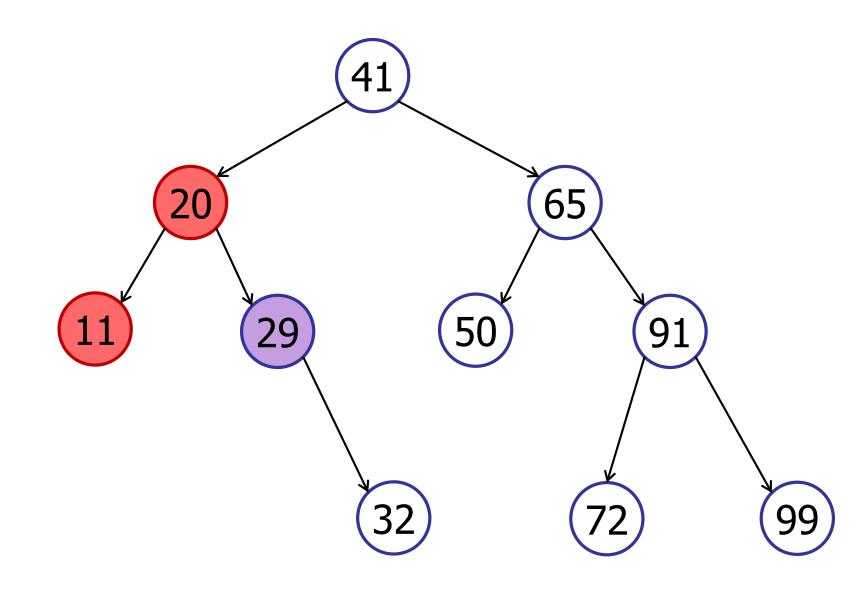


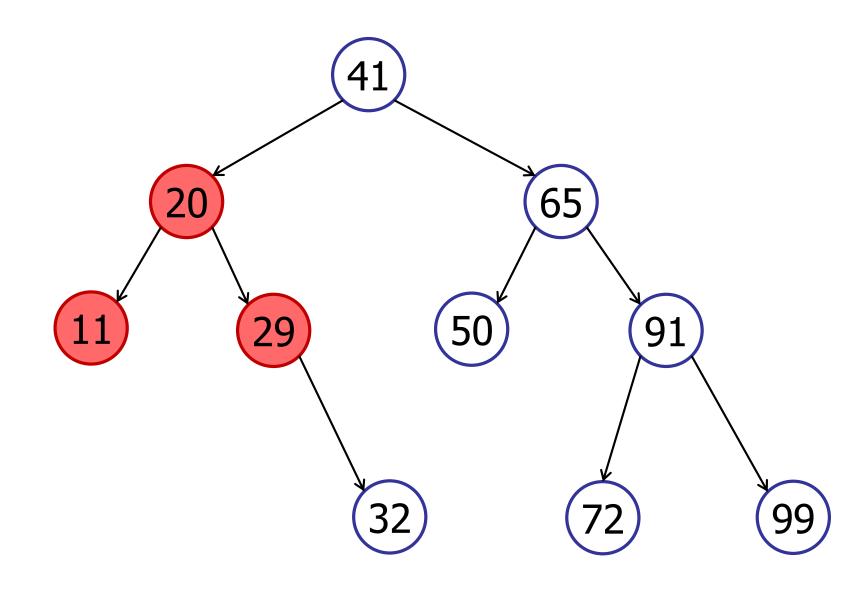


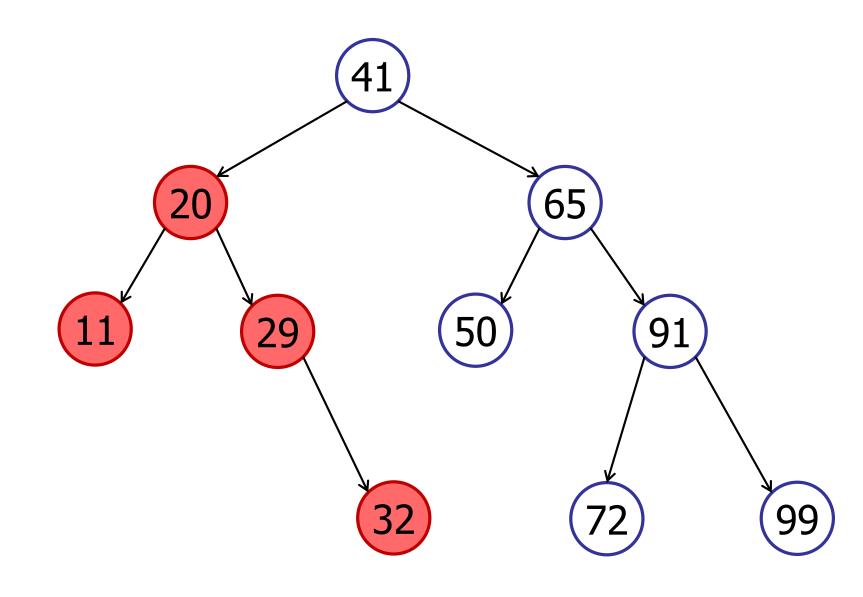


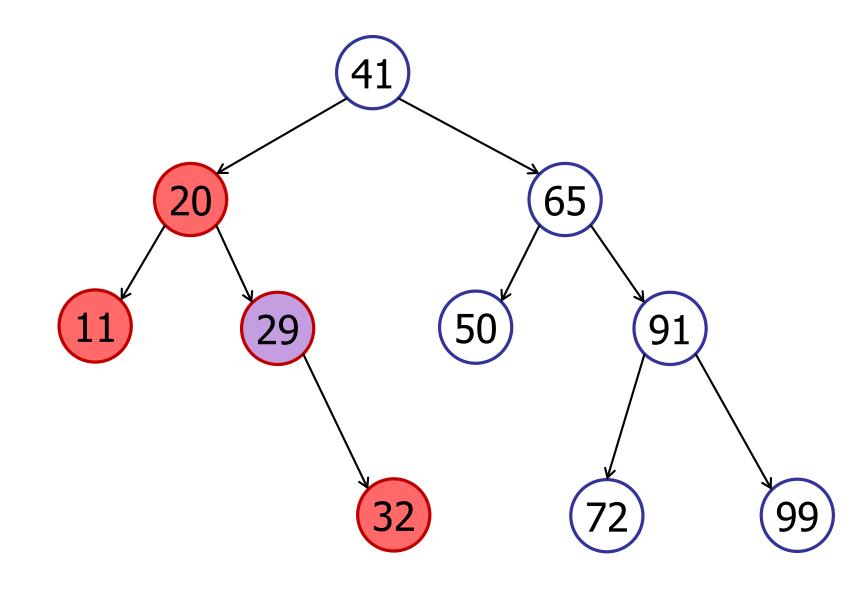


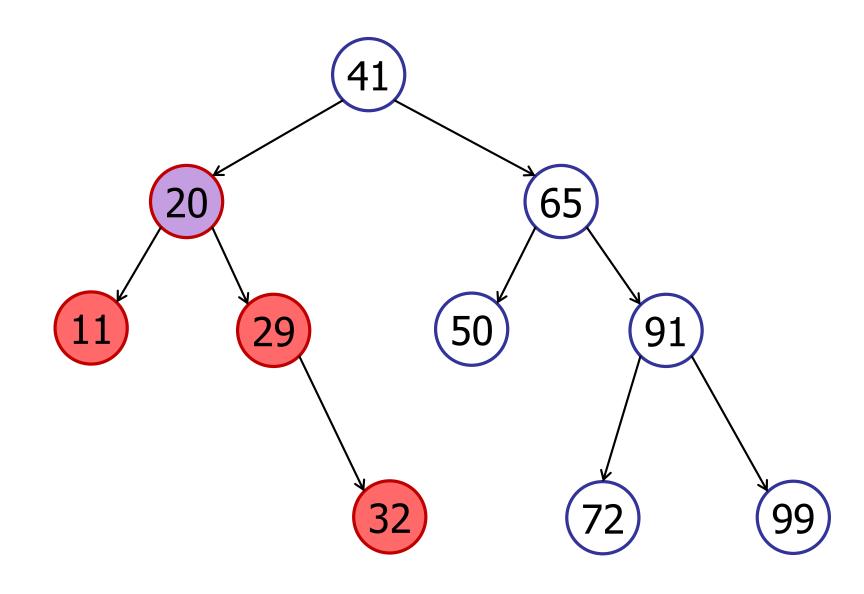


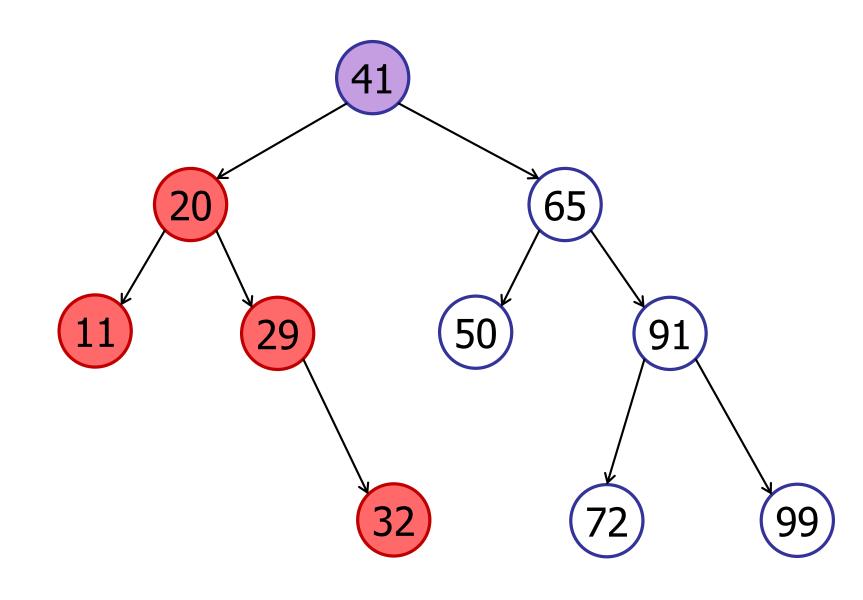


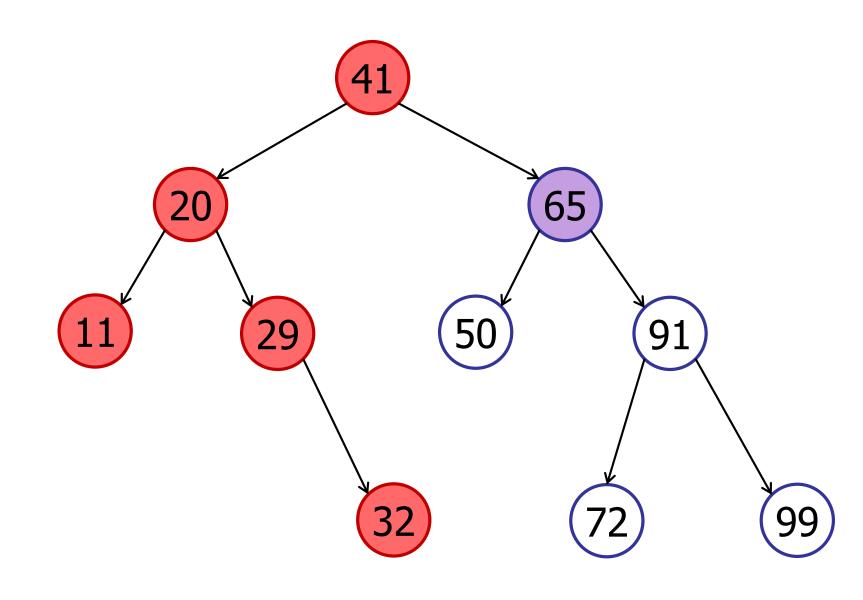


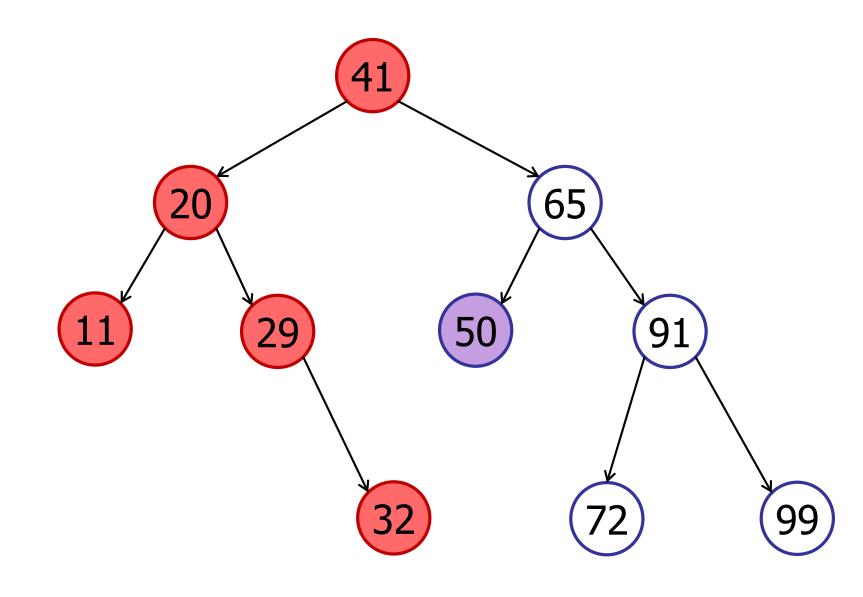


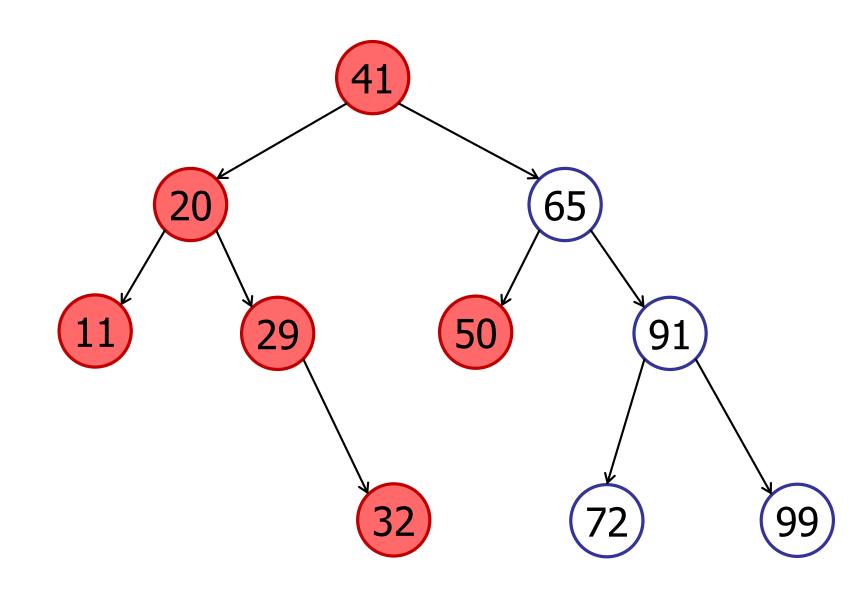


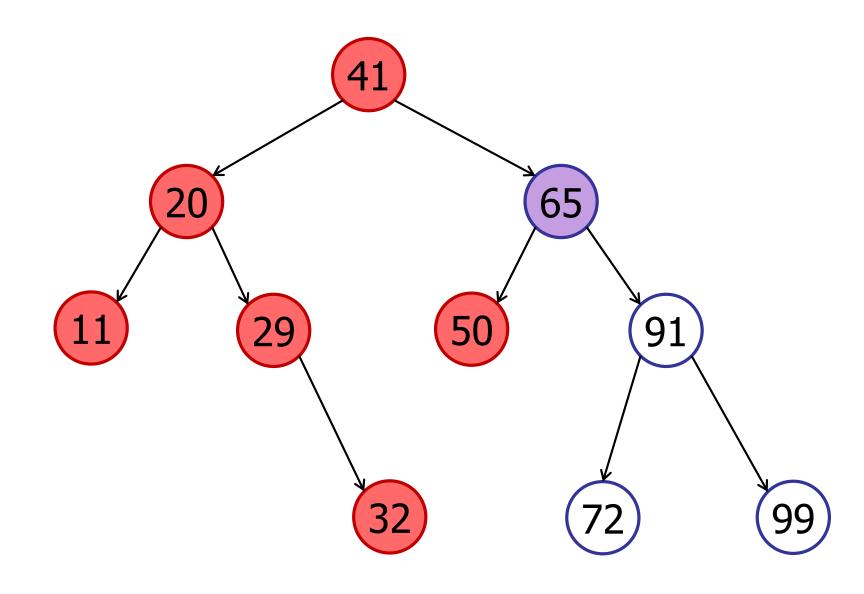


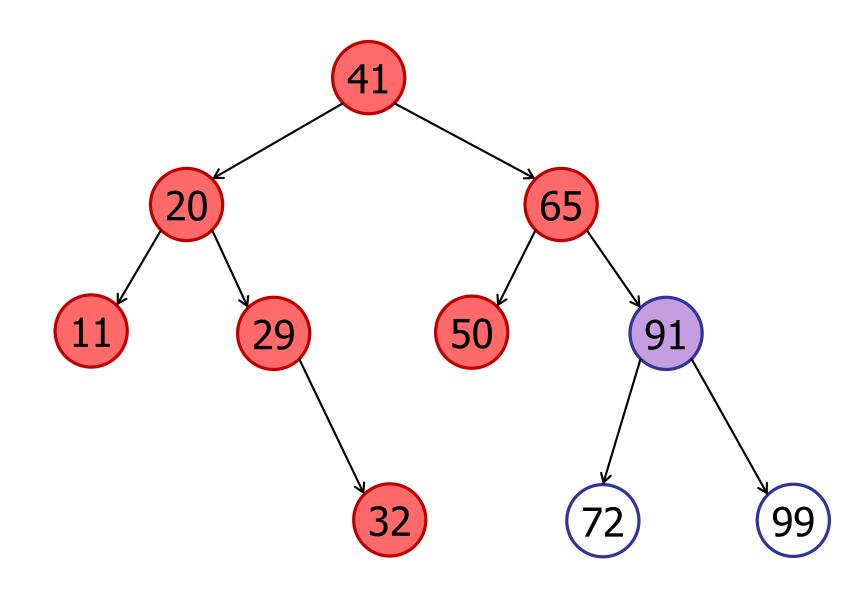


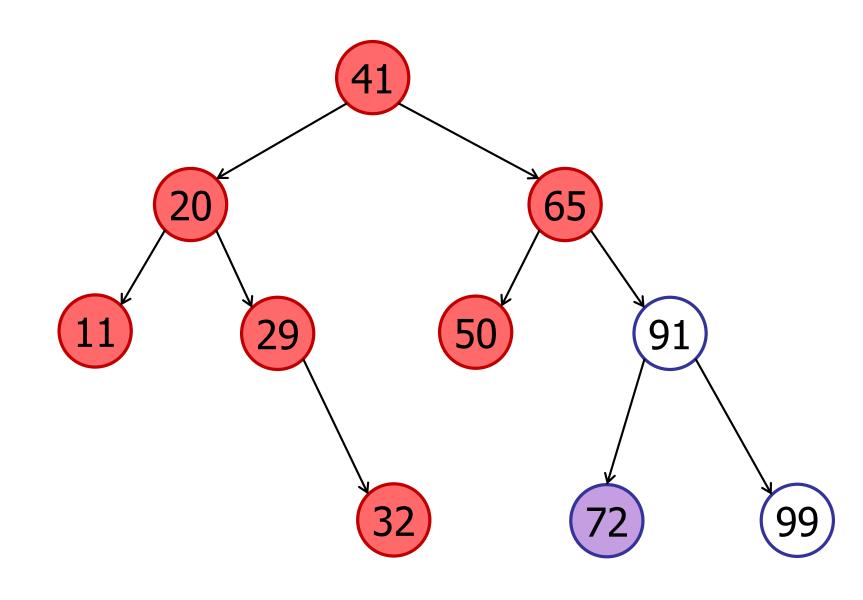


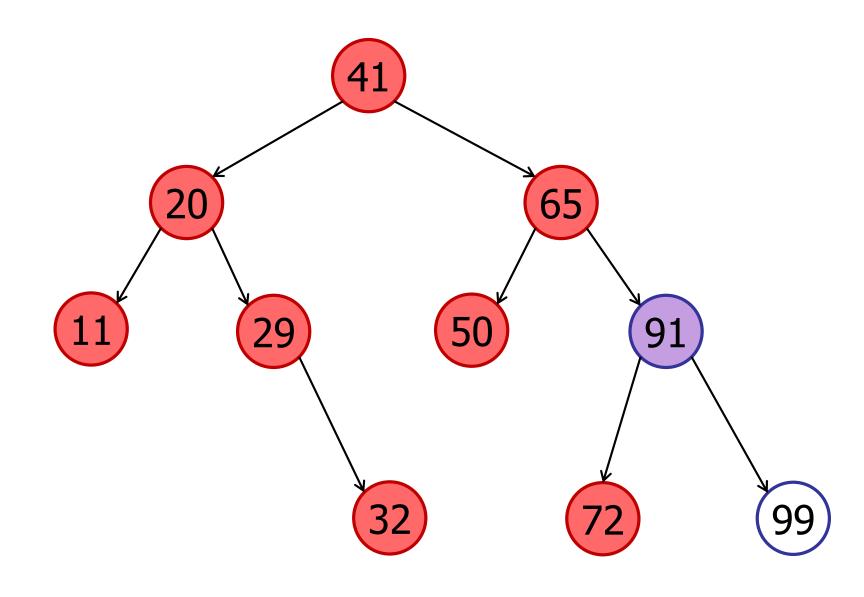


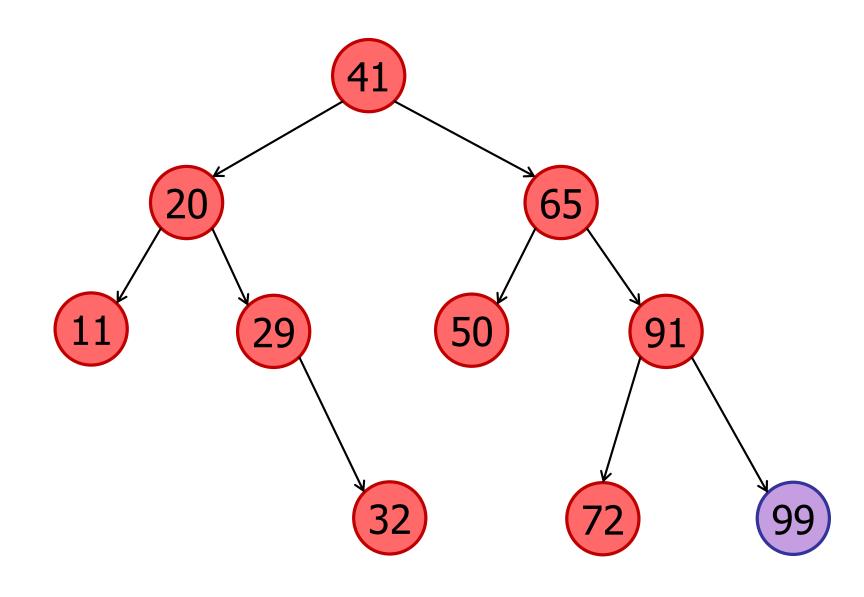


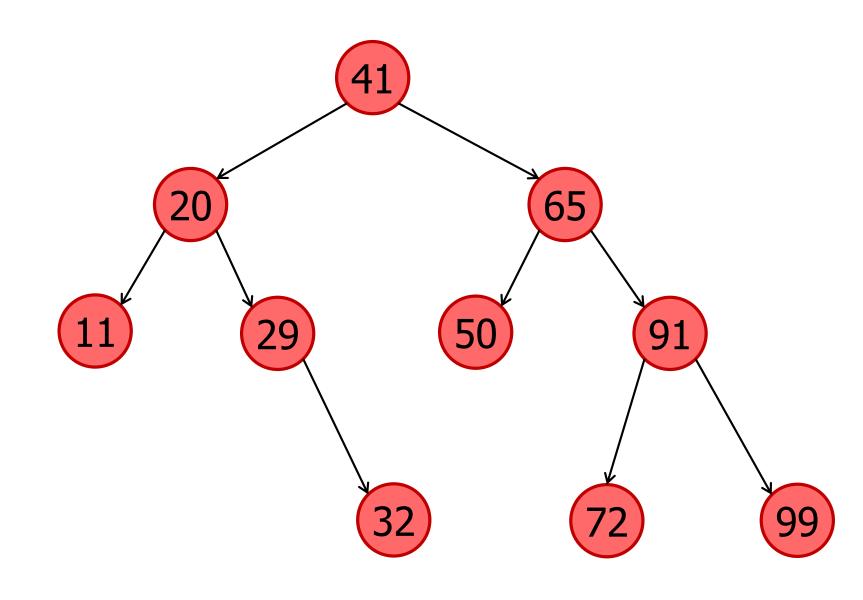












```
public void in-order-traversal() {
       // Traverse left sub-tree
       if (leftTree != null)
               leftTree.in-order-traversal();
       visit(this);
       // Traverse right sub-tree
       if (rightTree != null)
              rightTree.in-order-traversal();
```

#### How long does an in-order-traversal take?

- 1. O(1)
- 2. O(log n)
- 3. O(n)
- 4. O(n log n)
- 5.  $O(n^2)$
- 6.  $O(2^n)$

#### How long does an in-order-traversal take?

- 1. O(1)
- 2. O(log n)
- 3. O(n)
  - 4. O(n log n)
  - 5.  $O(n^2)$
  - 6.  $O(2^n)$

#### in-order-traversal(v)

```
public void in-order-traversal() {
       // Traverse left sub-tree
       if (leftTree != null)
               leftTree.in-order-traversal();
       visit(this);
       // Traverse right sub-tree
       if (rightTree != null)
              rightTree.in-order-traversal();
```

#### Running time: O(n)

visits each node at most once

#### in-order-traversal(v)

- left-subtree
- SELF
- right-subtree

#### pre-order-traversal(v)

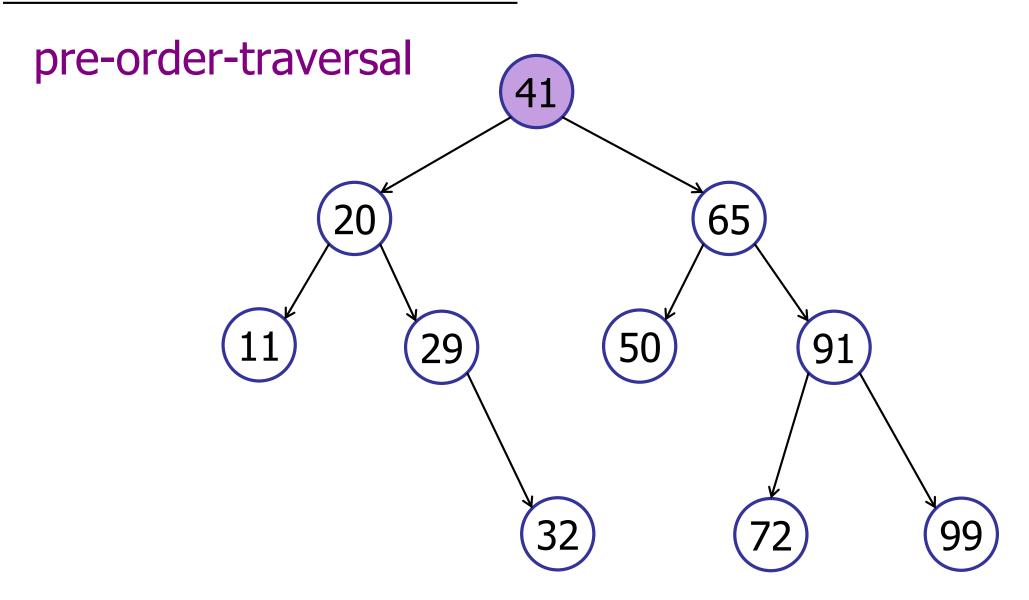
- SELF
- left-subtree
- right-subtree

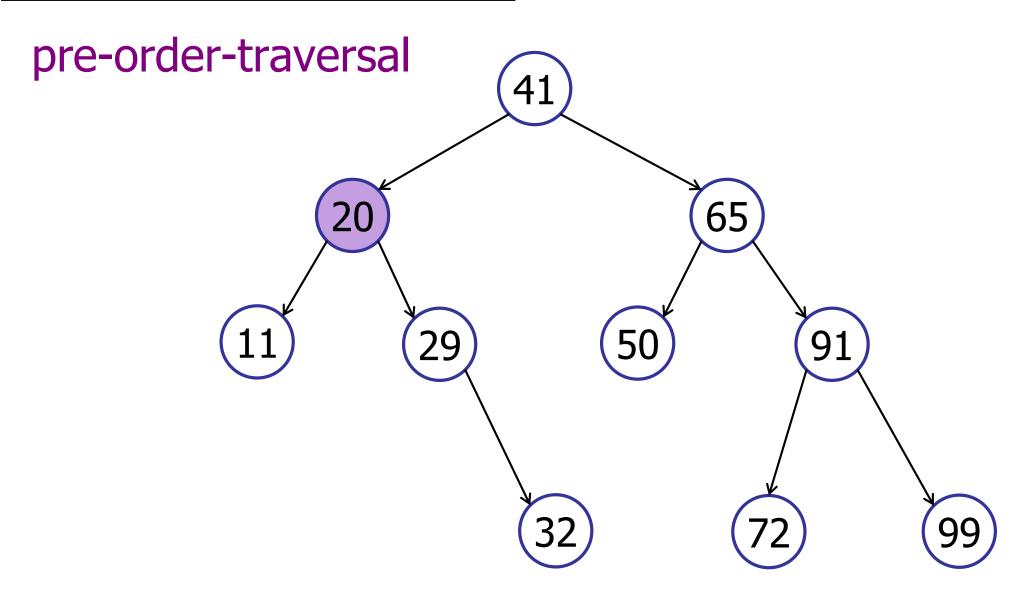
#### post-order-traversal(v)

- left-subtree
- right-subtree
- SELF

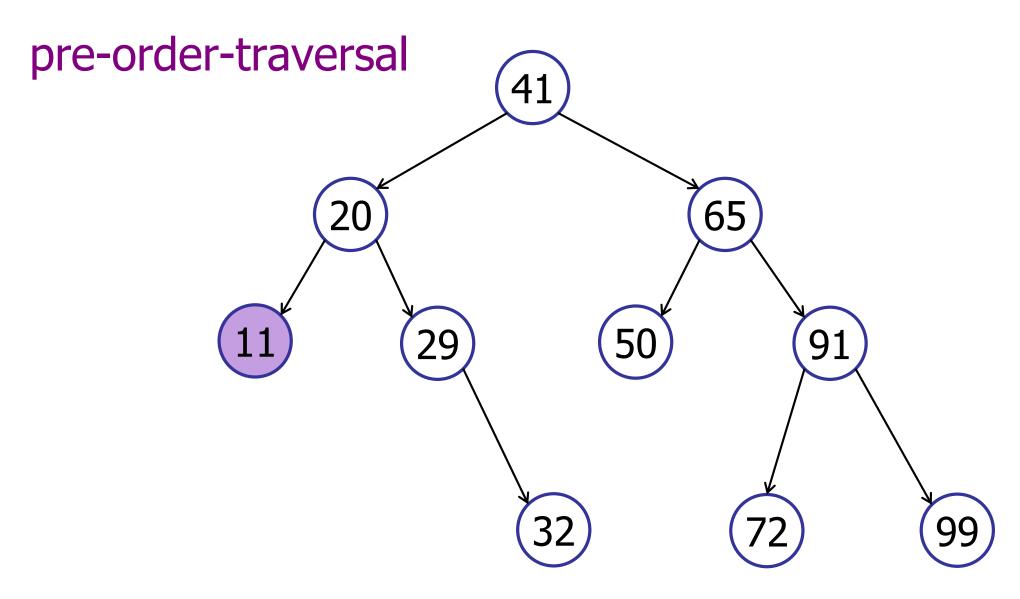
#### pre-order-traversal(v)

```
public void pre-order-traversal() {
      visit(this);
       // Traverse left sub-tree
       if (leftTree != null)
               leftTree.in-order-traversal();
       // Traverse right sub-tree
       if (rightTree != null)
              rightTree.in-order-traversal();
```

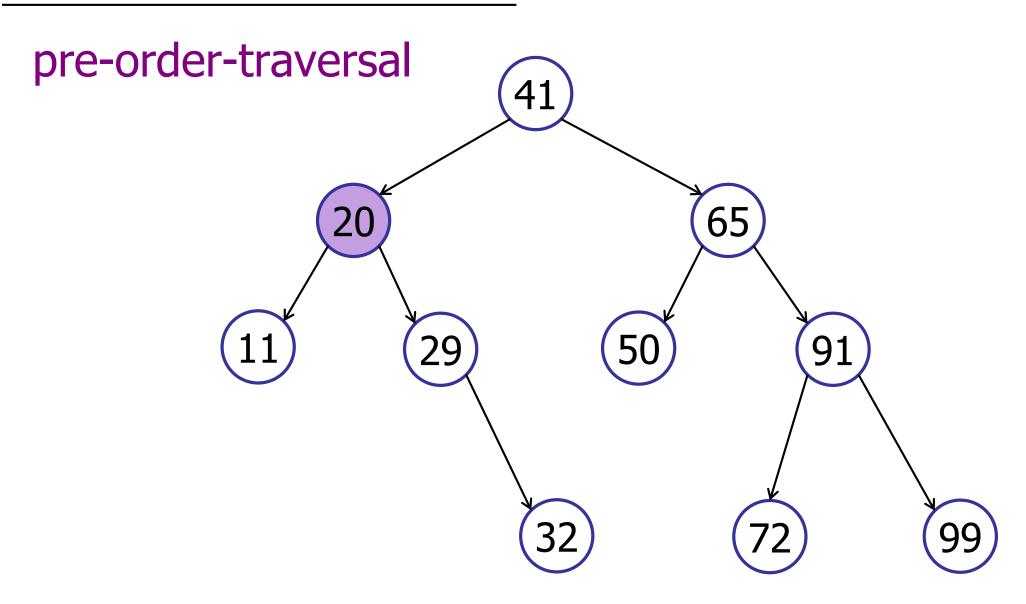




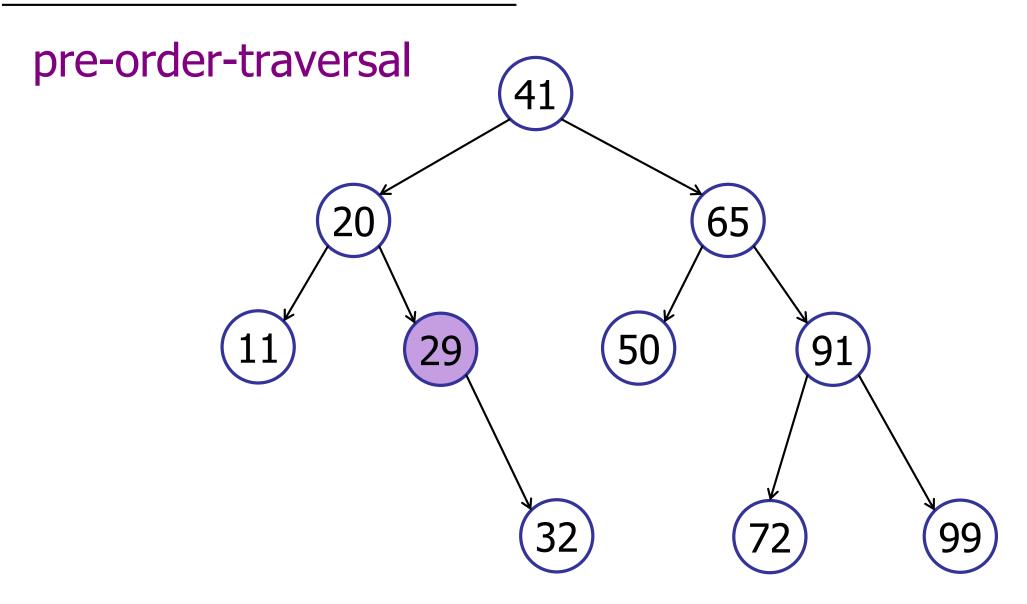
41 20



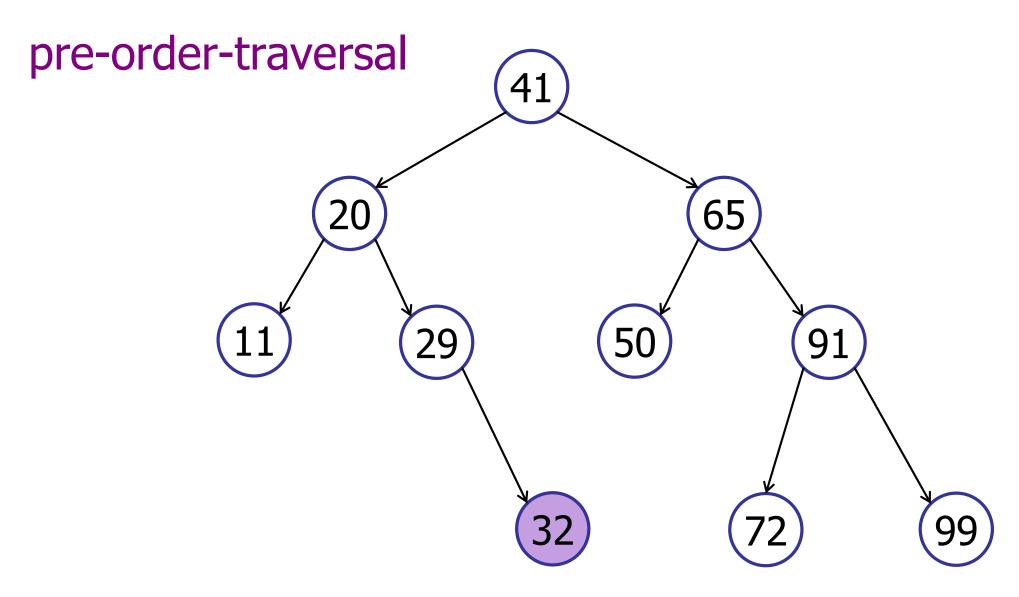
41 20 11

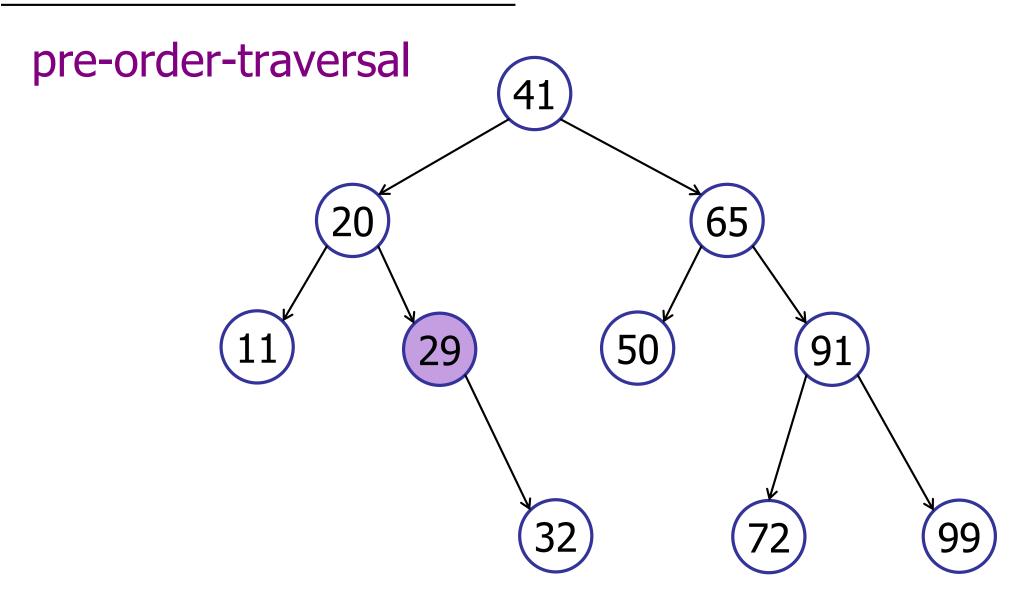


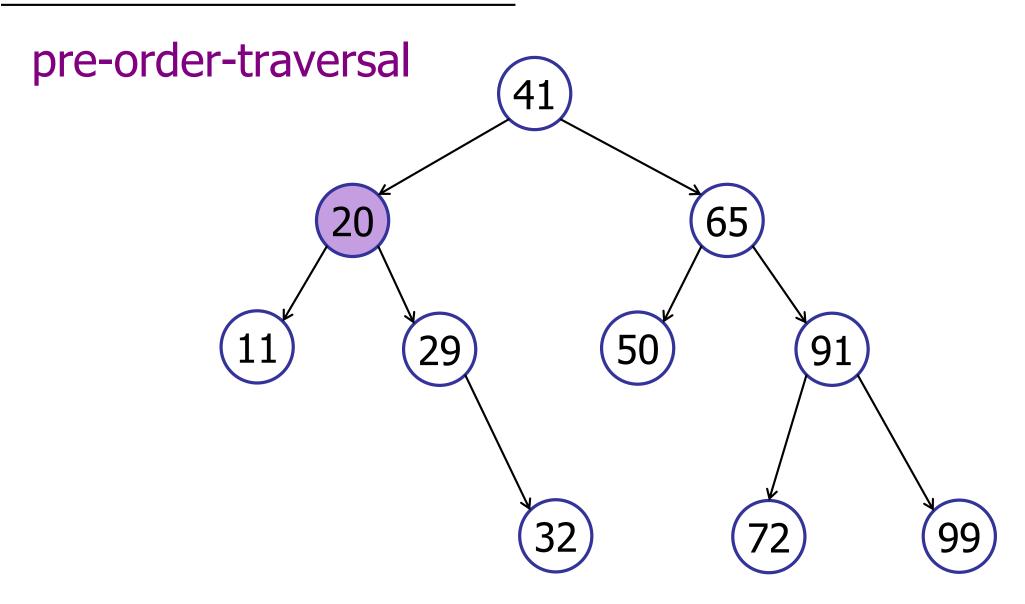
41 20 11

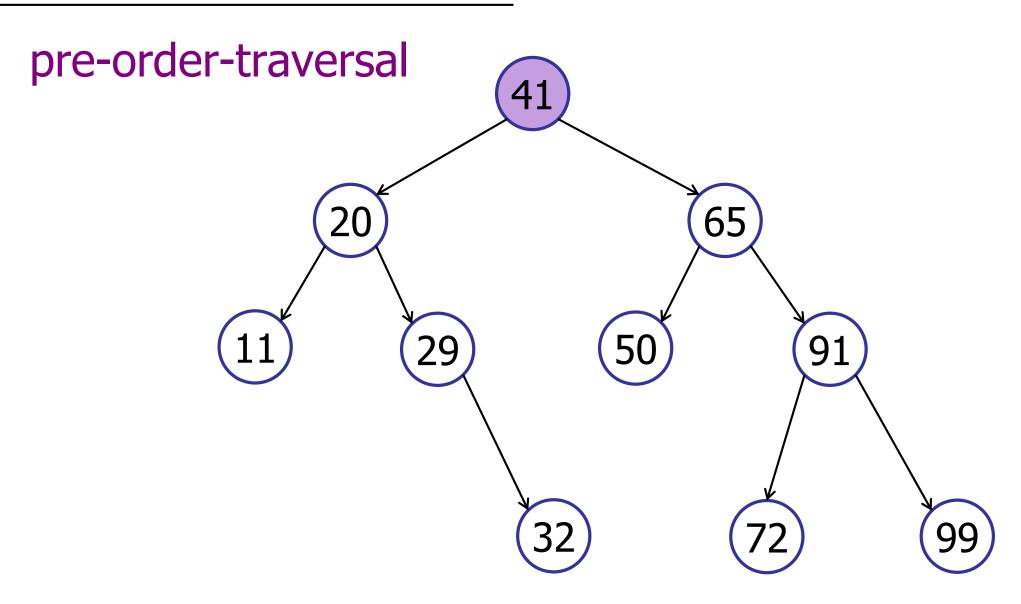


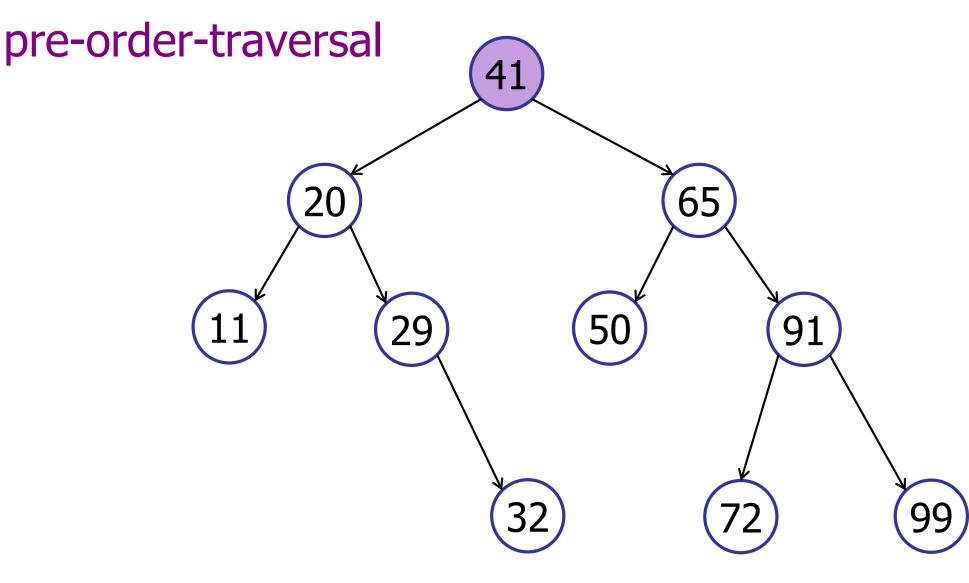
41 20 11 29







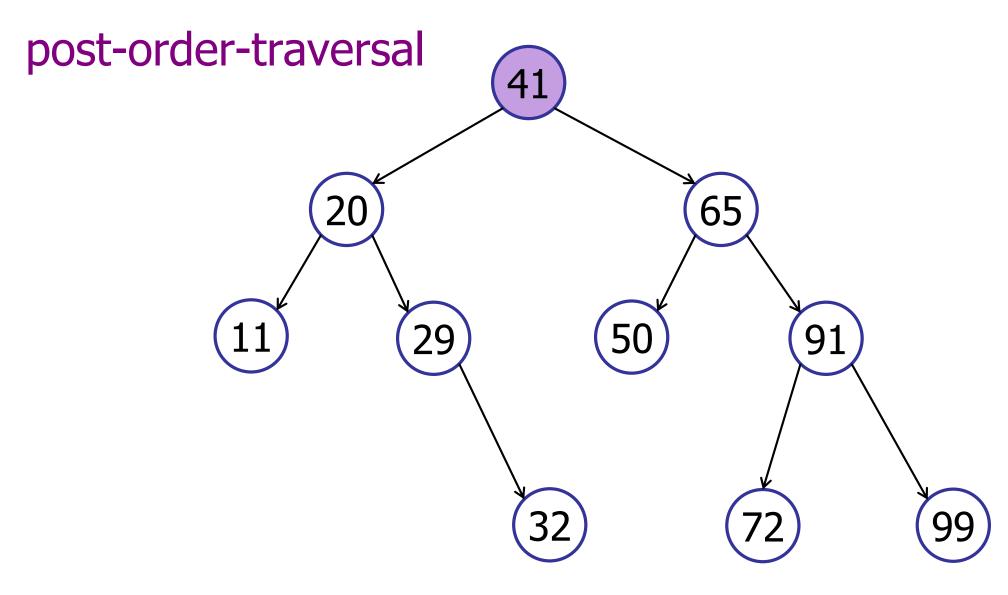




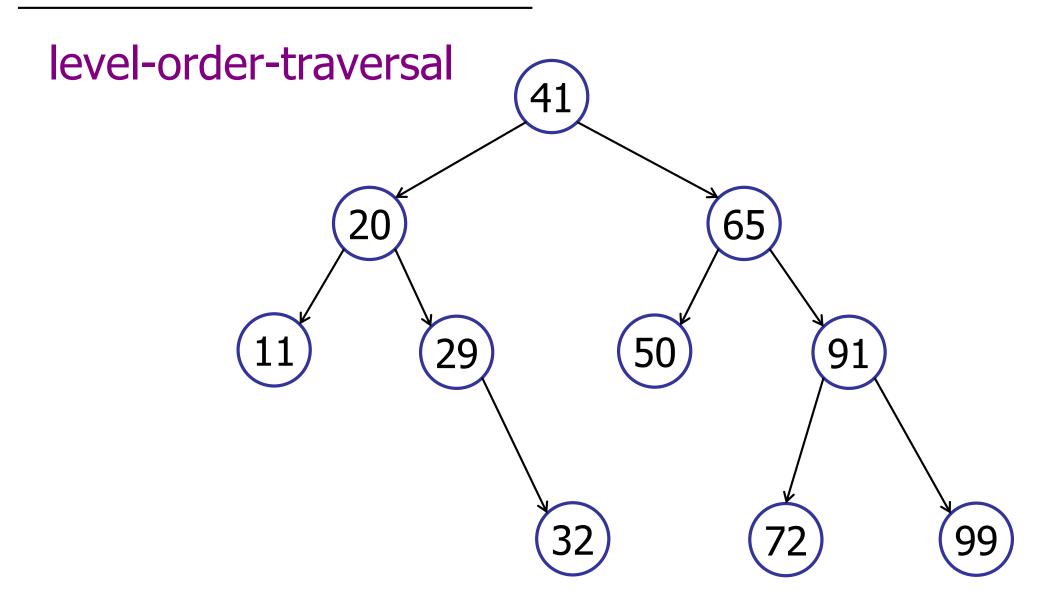
41 20 11 29 32 65 50 91 72 99

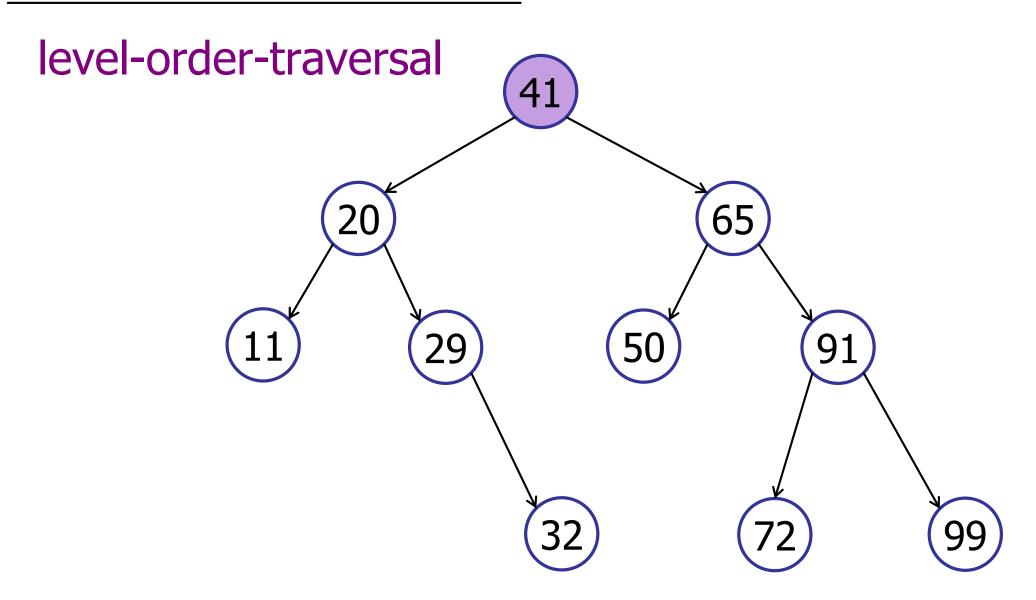
#### post-order-traversal(v)

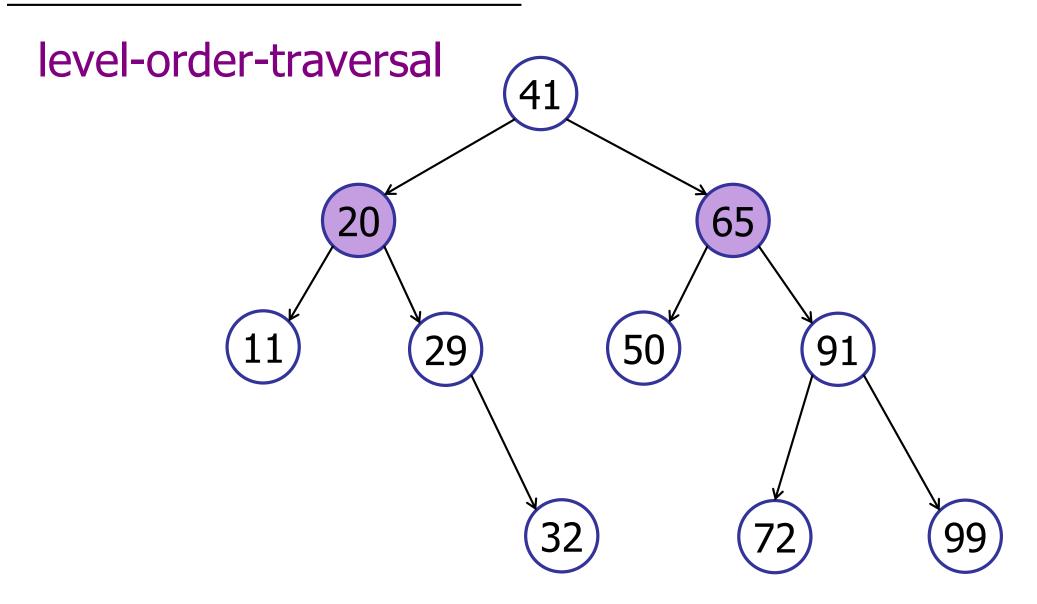
```
public void post-order-traversal() {
       // Traverse left sub-tree
       if (leftTree != null)
               leftTree.in-order-traversal();
       // Traverse right sub-tree
       if (rightTree != null)
              rightTree.in-order-traversal();
      visit(this);
```



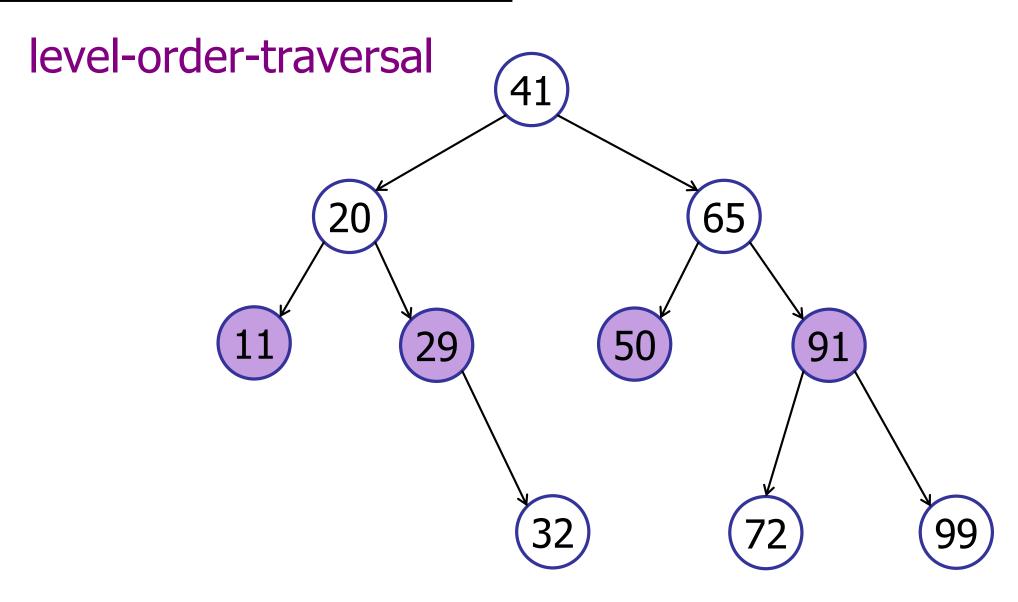
11 32 29 20 50 72 99 91 65 41



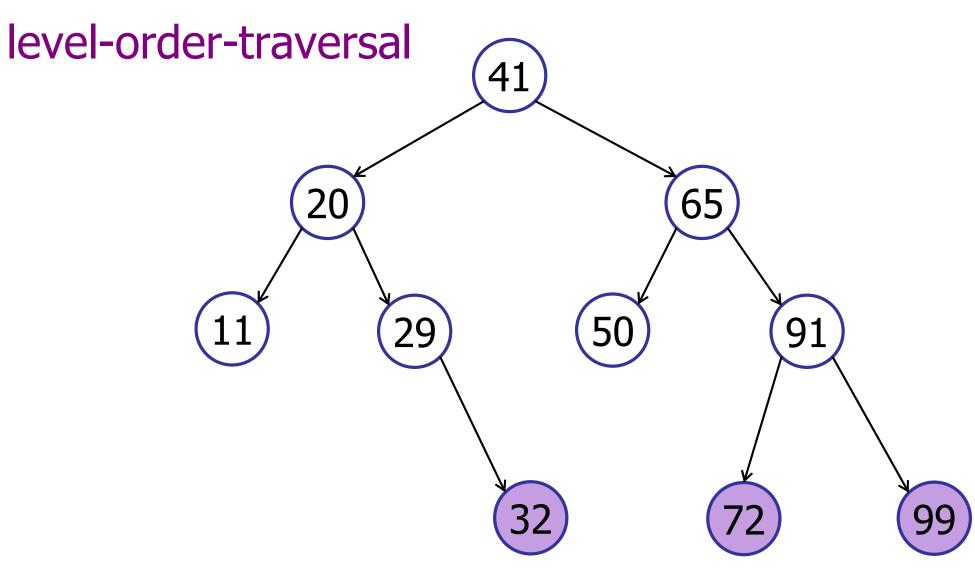




41 20 65



41 20 65 11 29 50 91



41 20 65 11 29 50 91 32 72 99

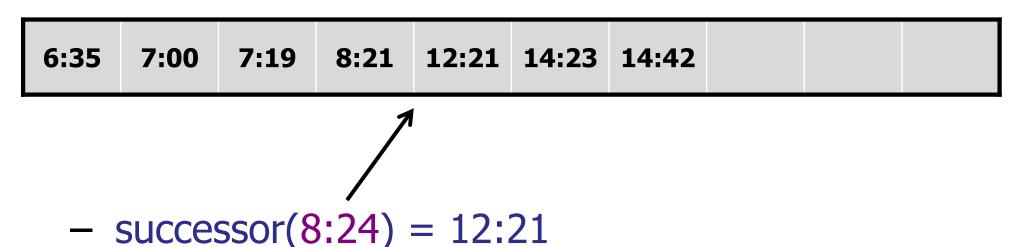
#### Several varieties:

- pre-order
- in-order
- post-order
- level-order

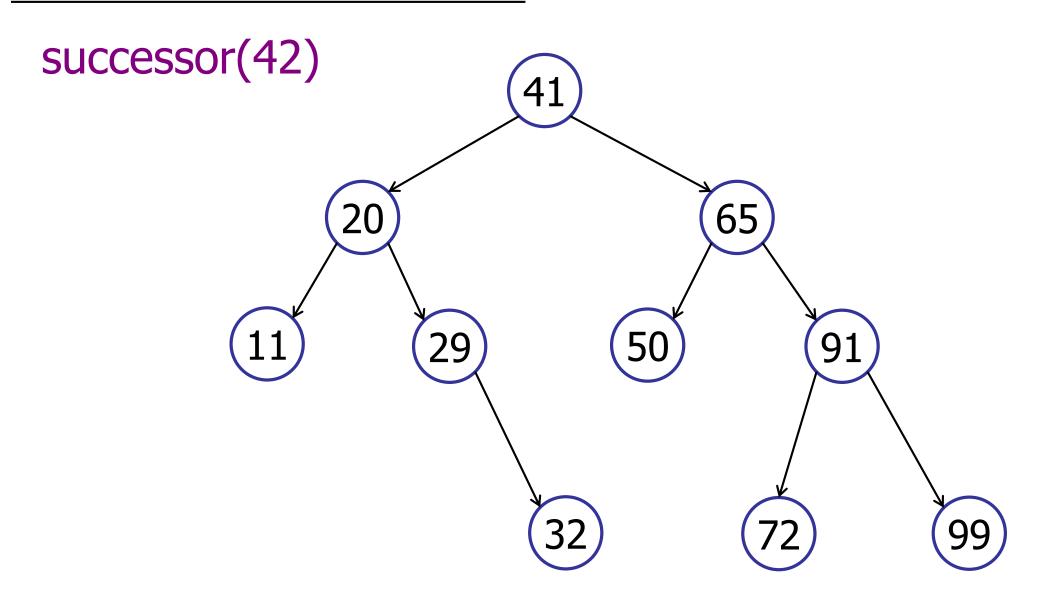
- 1. Terminology and Definitions
- 2. Basic operations:
  - height
  - searchMin, searchMax
  - search, insert
- 3. Traversals
  - in-order, pre-order, post-order
- 4. Other operations

# Airport Scheduling

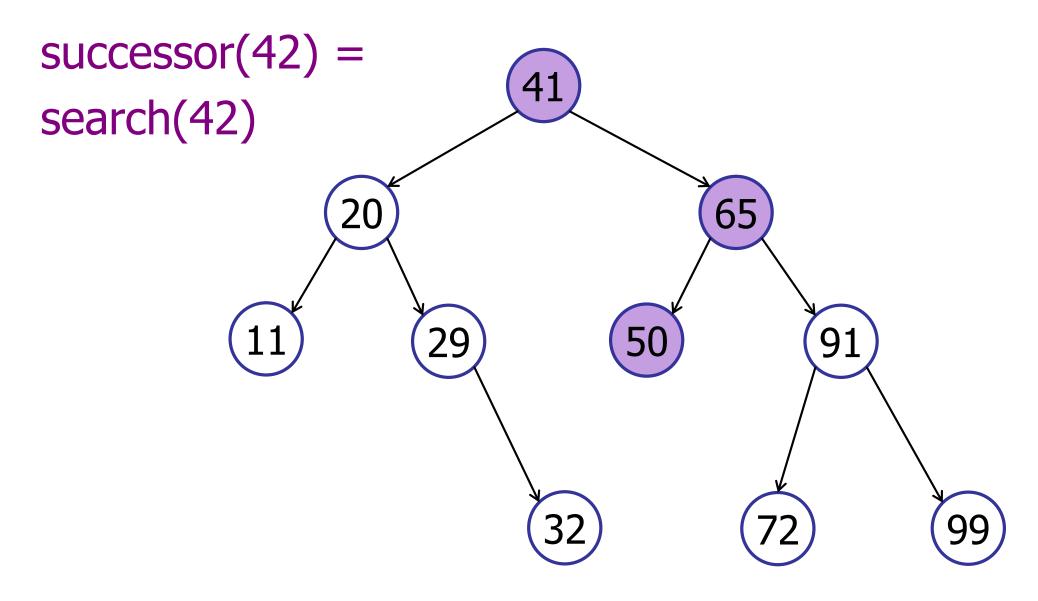
#### **Dictionary**



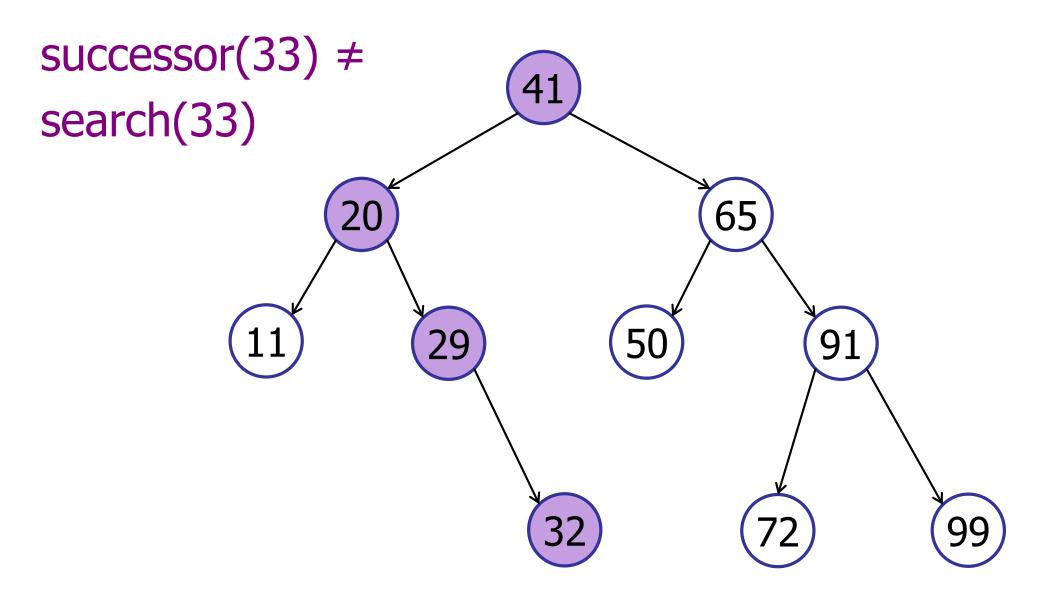
How do we implement this?



Key 42 is not in the tree



Key 42 is not in the tree



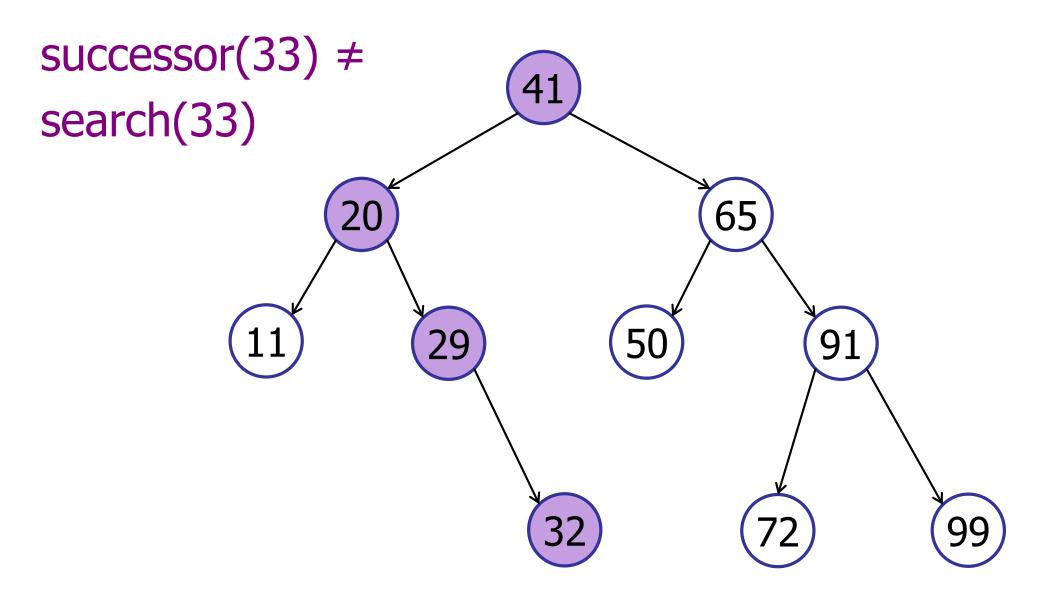
Key 33 is not in the tree

Basic strategy: successor(key)

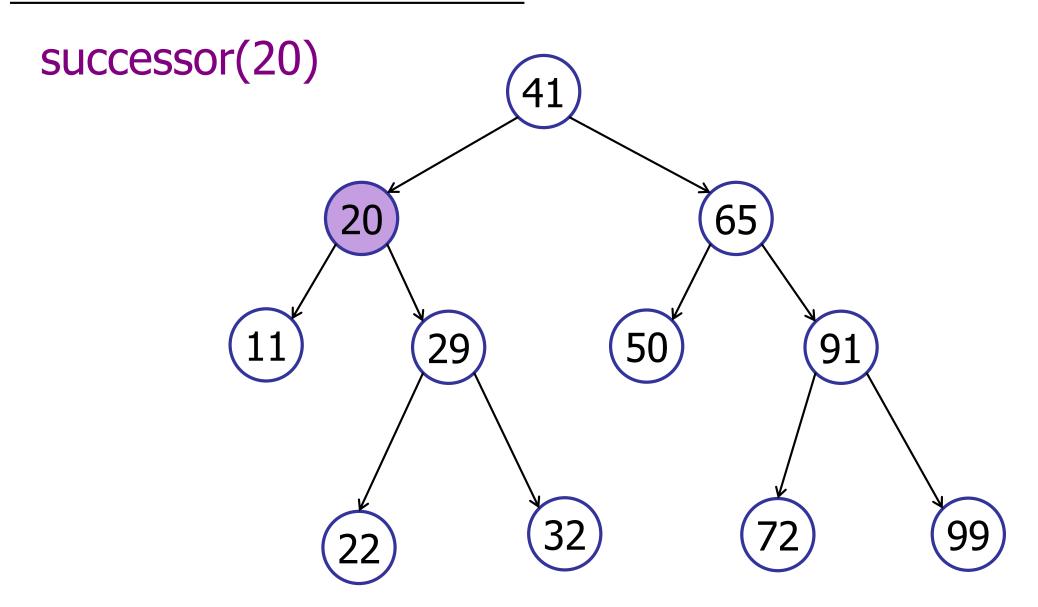
1. Search for key in the tree.

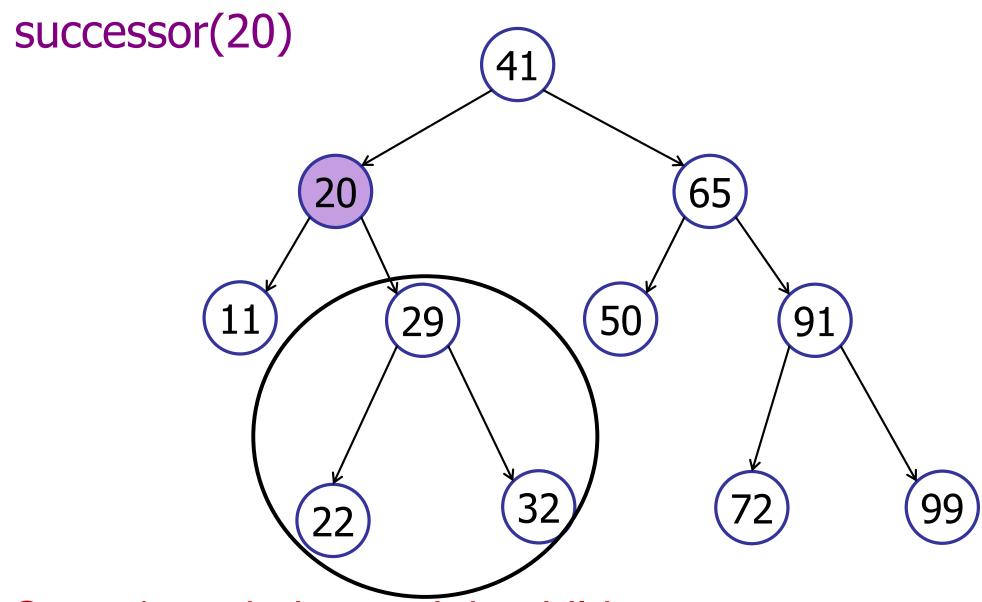
2. If (result > key), then return result.

3. If (result <= key), then search for successor of result.

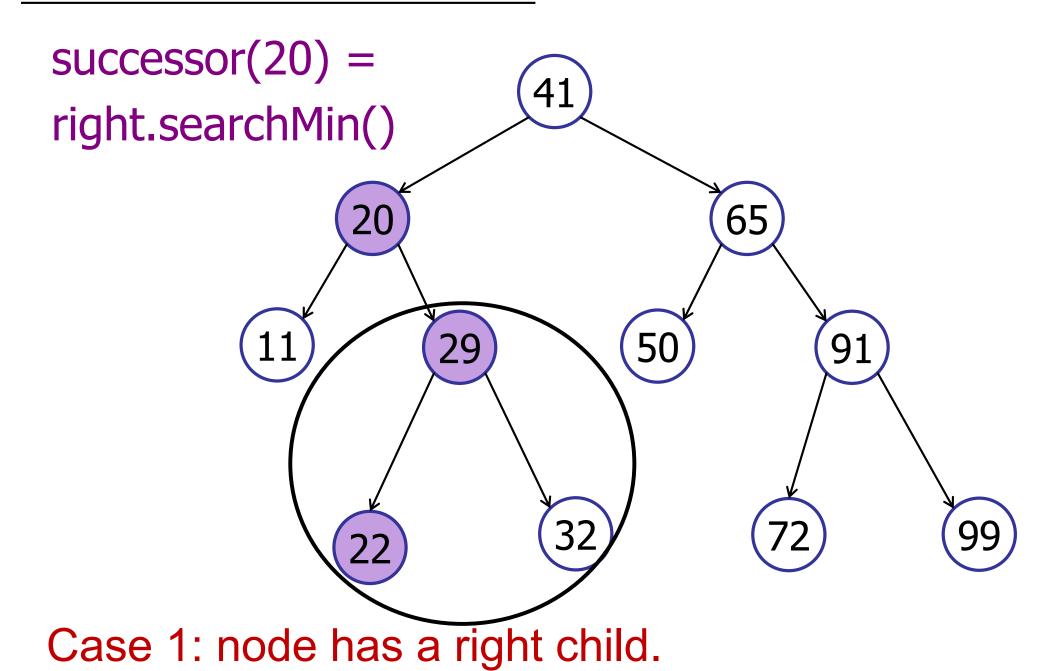


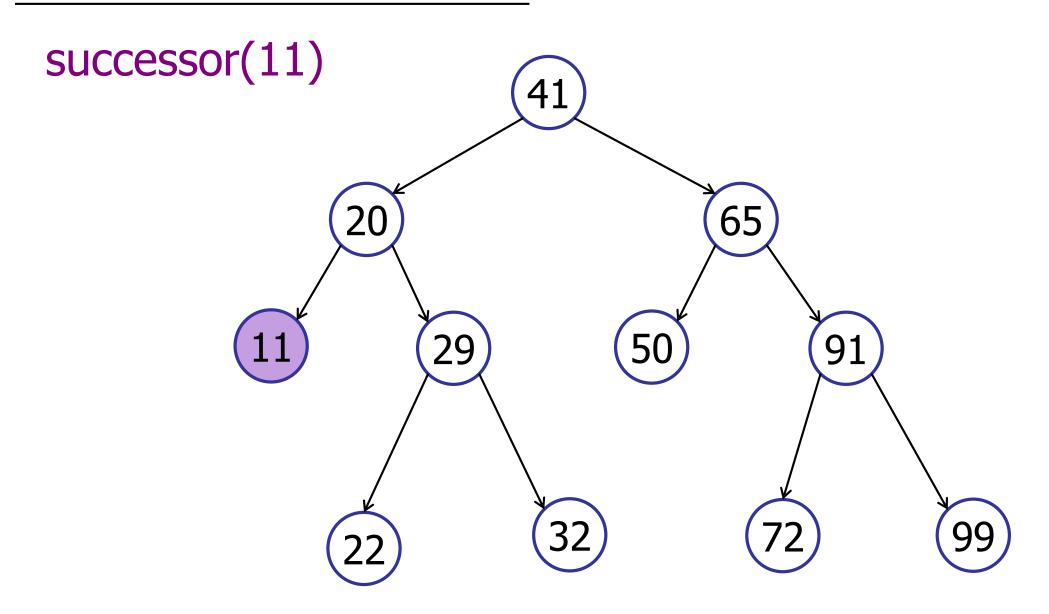
Key 33 is not in the tree



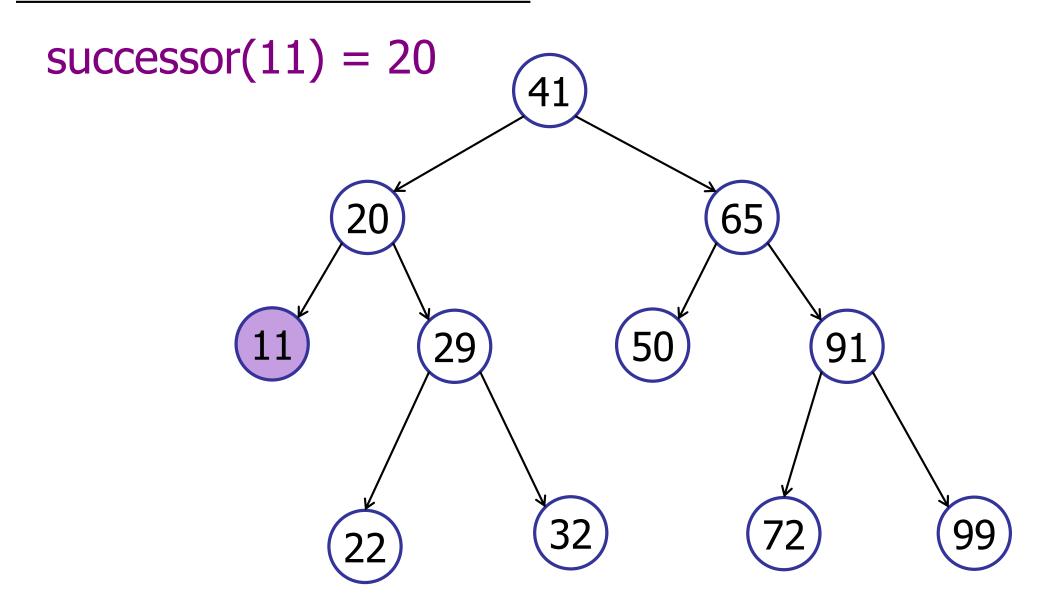


Case 1: node has a right child.

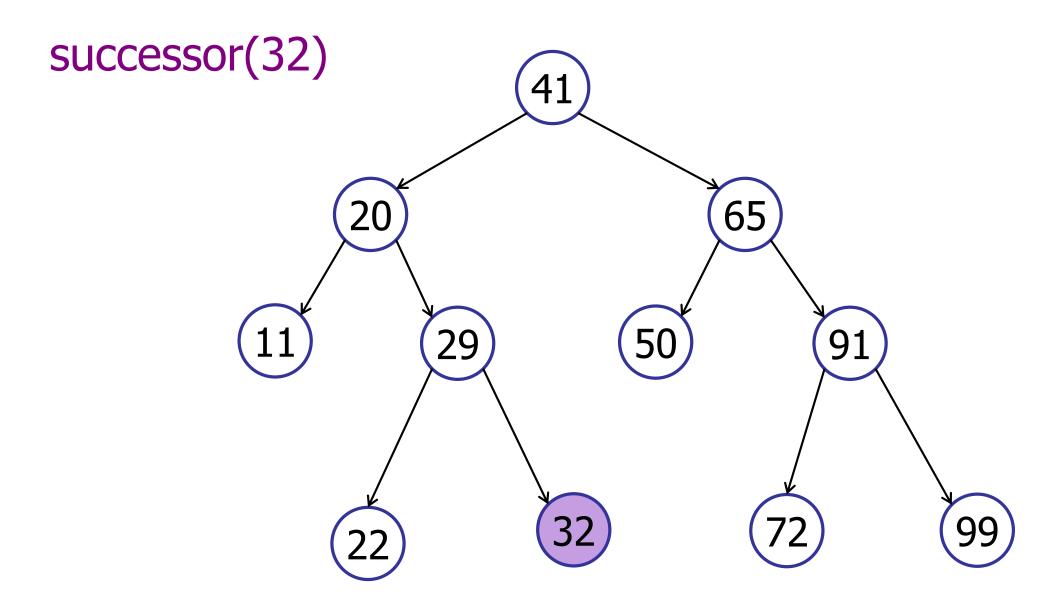




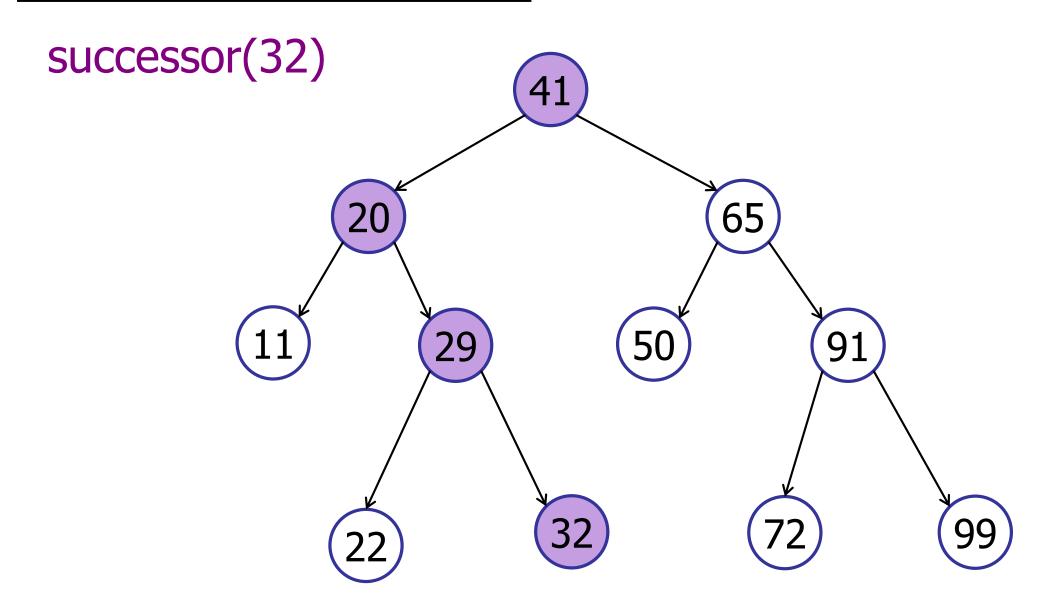
Case 2: node has no right child.



Case 2: node has no right child.



Case 2: node has no right child.



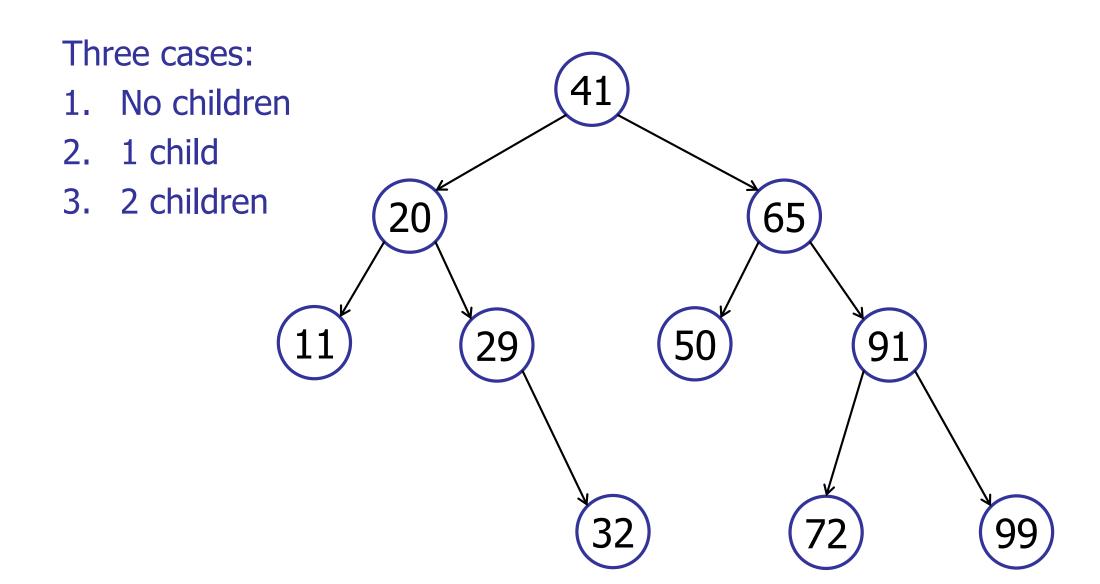
Case 2: node has no right child.

#### Find the next TreeNode:

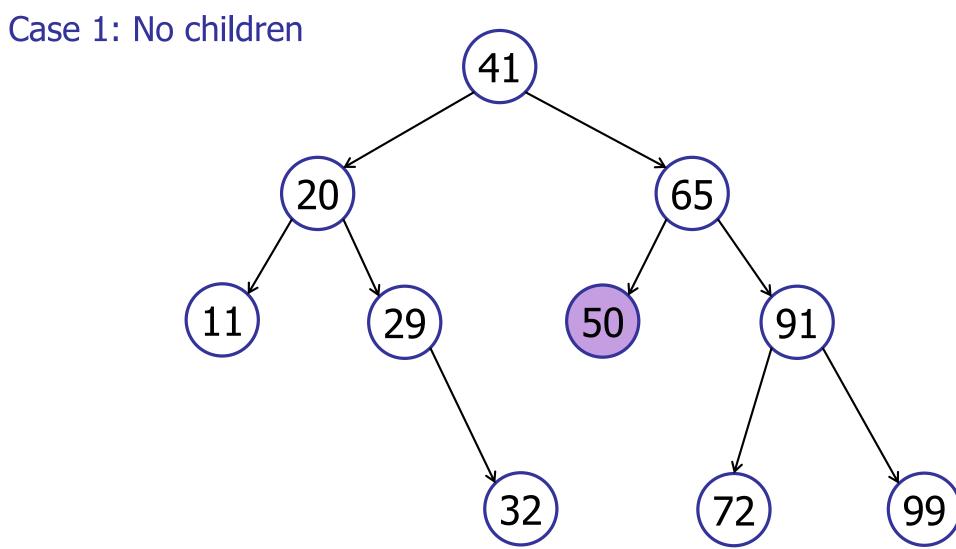
```
public TreeNode successor() {
       if (rightTree != null)
             return rightTree.searchMin();
      TreeNode parent = parentTree;
      TreeNode child = this;
      while ((parent != null) && (child = parent.rightTree))
             child = parent;
             parent = child.parentTree;
       return parent;
```

- 1. Terminology and Definitions
- 2. Basic operations:
  - height
  - searchMin, searchMax
  - search, insert
- 3. Traversals
  - in-order, pre-order, post-order
- 4. Other operations

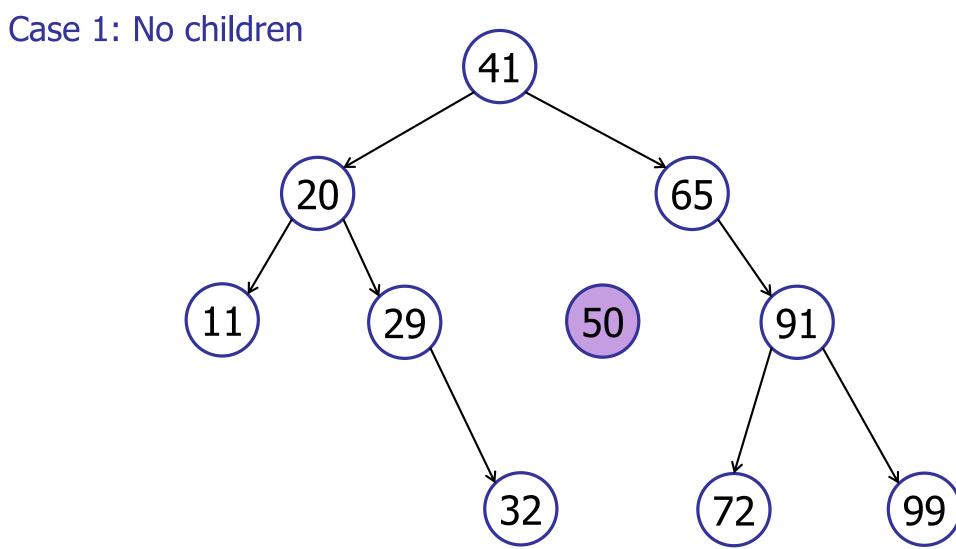
#### delete(v)



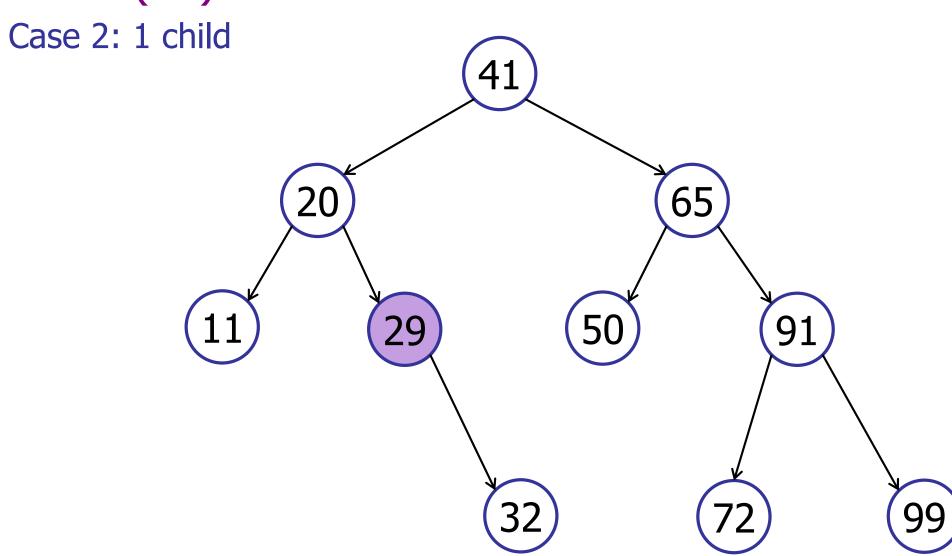
#### delete(50)



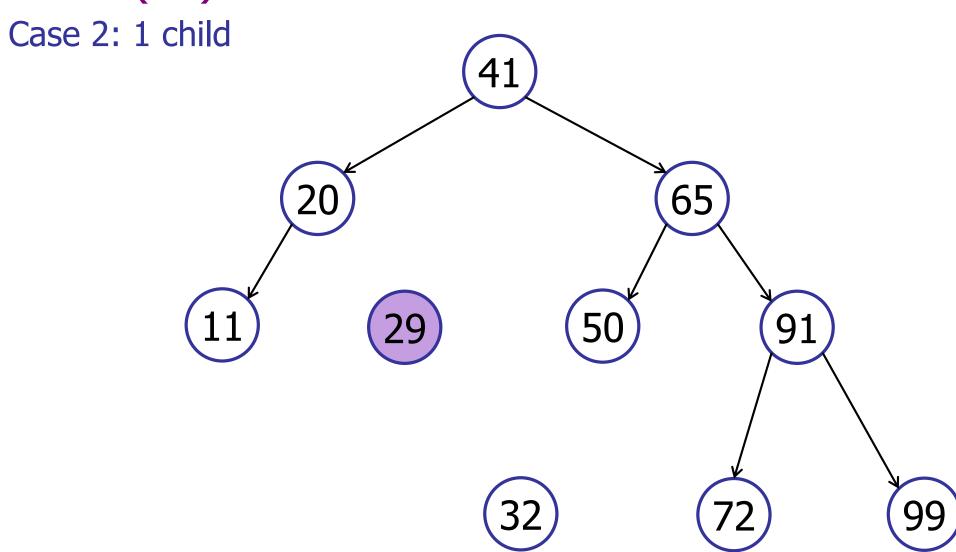
#### delete(50)



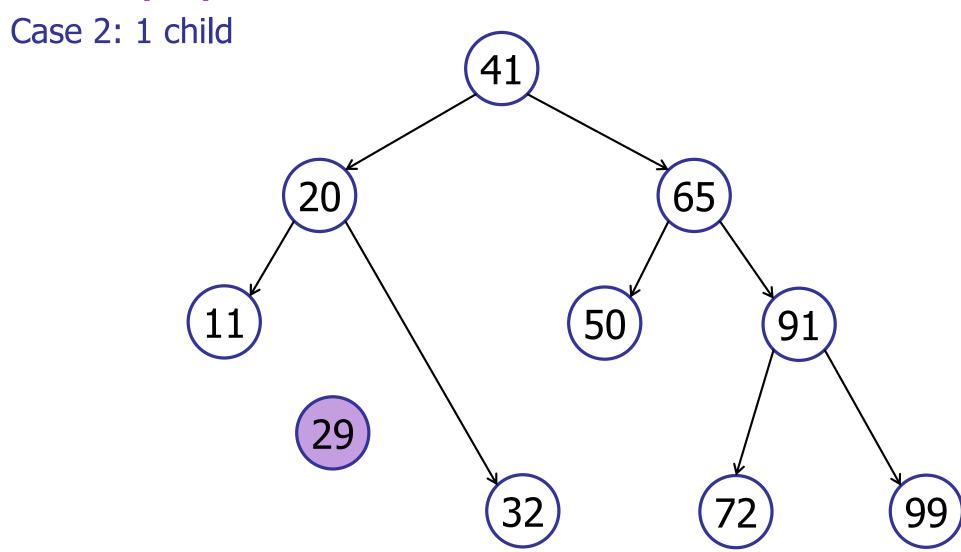
#### delete(29)

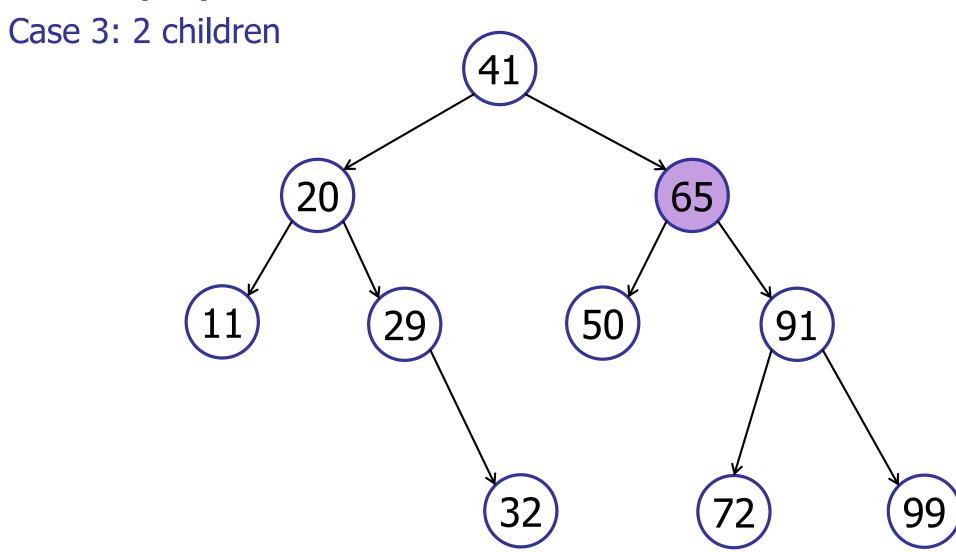


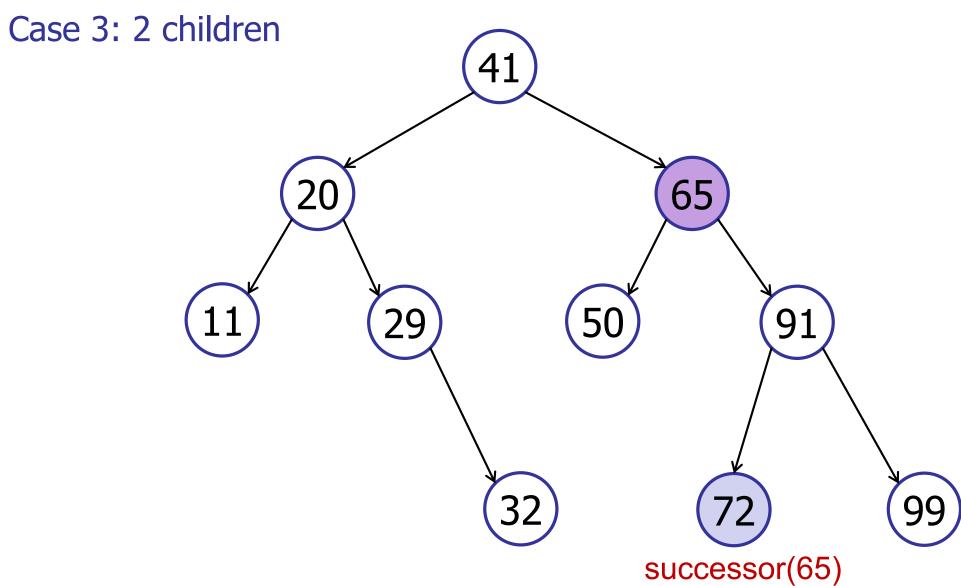
#### delete(29)



#### delete(29)







#### delete(65)

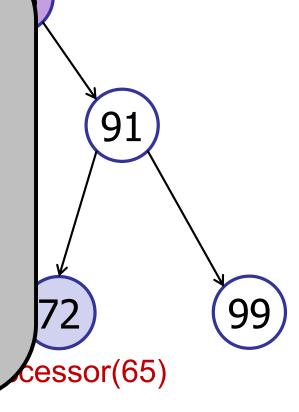
Case 3: 2 children

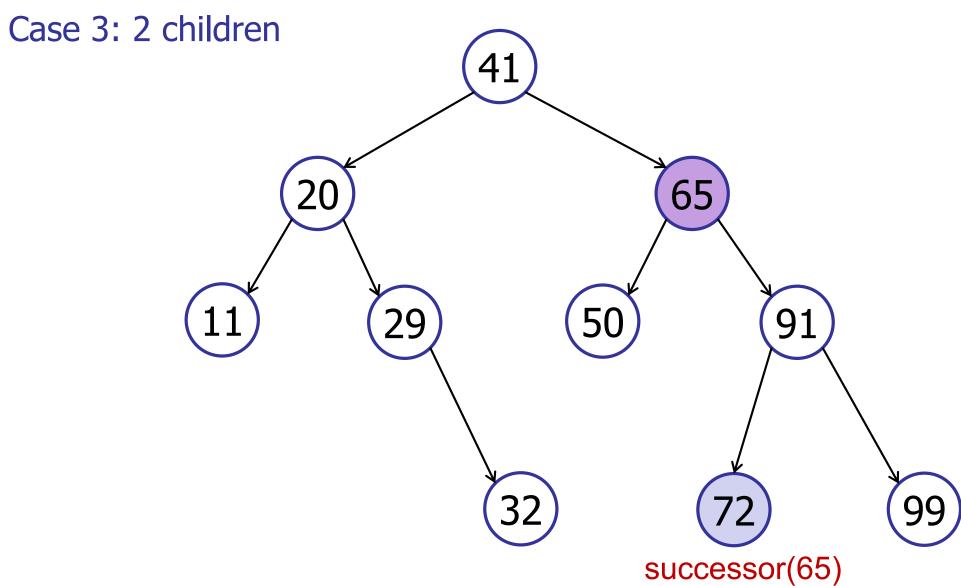
41

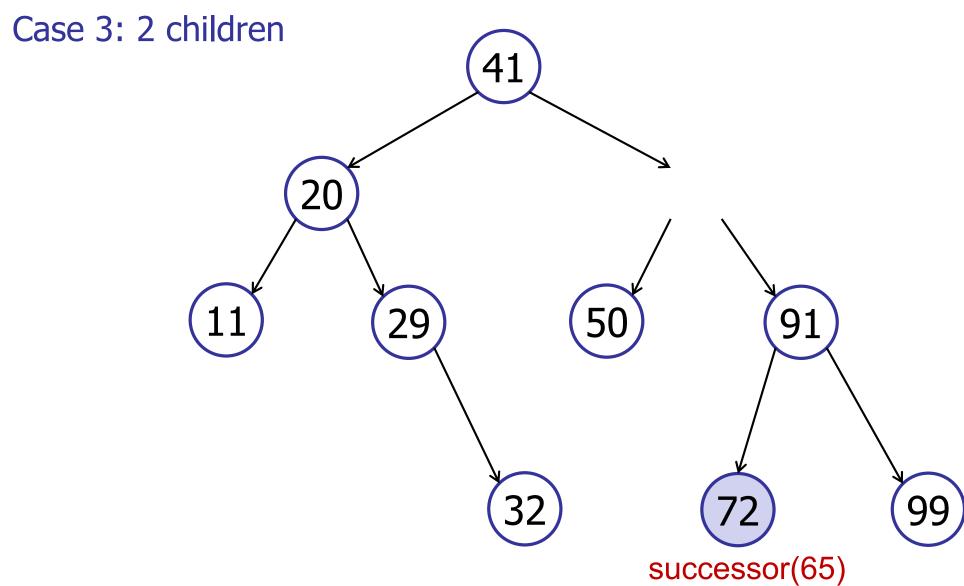
Claim: successor of deleted node has at most 1 child!

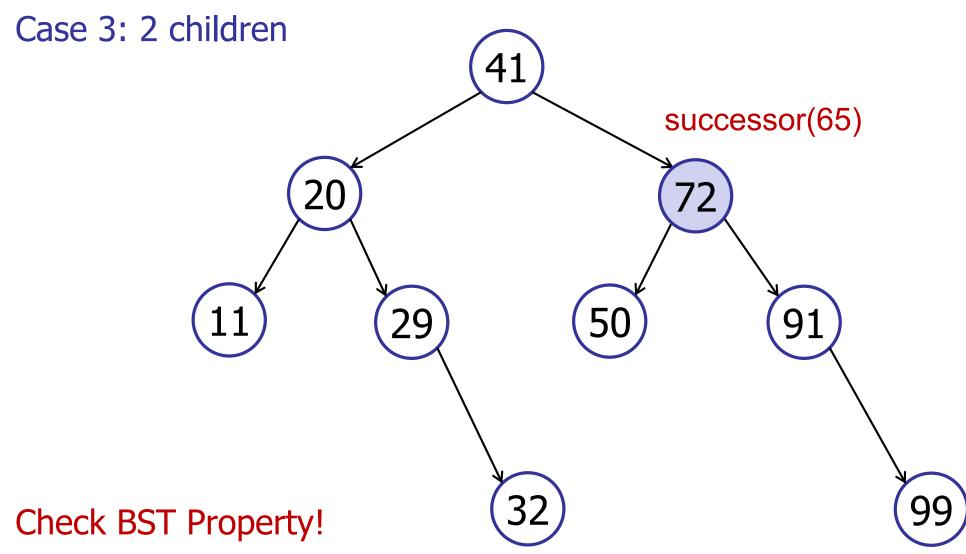
#### Proof:

- DeletedNode has two children.
- DeletedNode has a right child.
- successor() = right.findMin()
- · min element has no left child.









#### delete(v)

#### Running time: O(h)

#### Three cases:

- 1. No children:
  - remove v
- 2. 1 child:
  - remove v
  - connect child(v) to parent(v)
- 3. 2 children
  - x = successor(v)
  - delete(x)
  - remove v
  - connect x to left(v), right(v), parent(v)

#### **Modifying Operations**

- insert: O(h)
- delete: O(h)

#### **Query Operations:**

- search: O(h)
- predecessor, successor: O(h)
- findMax, findMin: O(h)
- in-order-traversal: O(n)

# Plan of the Day

#### Trees

- Terminology
- Traversals
- Operations

#### **Balanced Trees**

- Height-balanced binary search trees
- AVL trees
- Rotations