

CS2040S

Data Structures and Algorithms

(e-learning edition)

Graphs!

(Part 4)

Searching a Graph

Goal:

- Start at some vertex **s** = start.
- Find some other vertex **f** = finish.

Or: visit **all** the nodes in the graph;

Two basic techniques:

- Breadth-First Search (BFS)
- Depth-First Search (DFS)

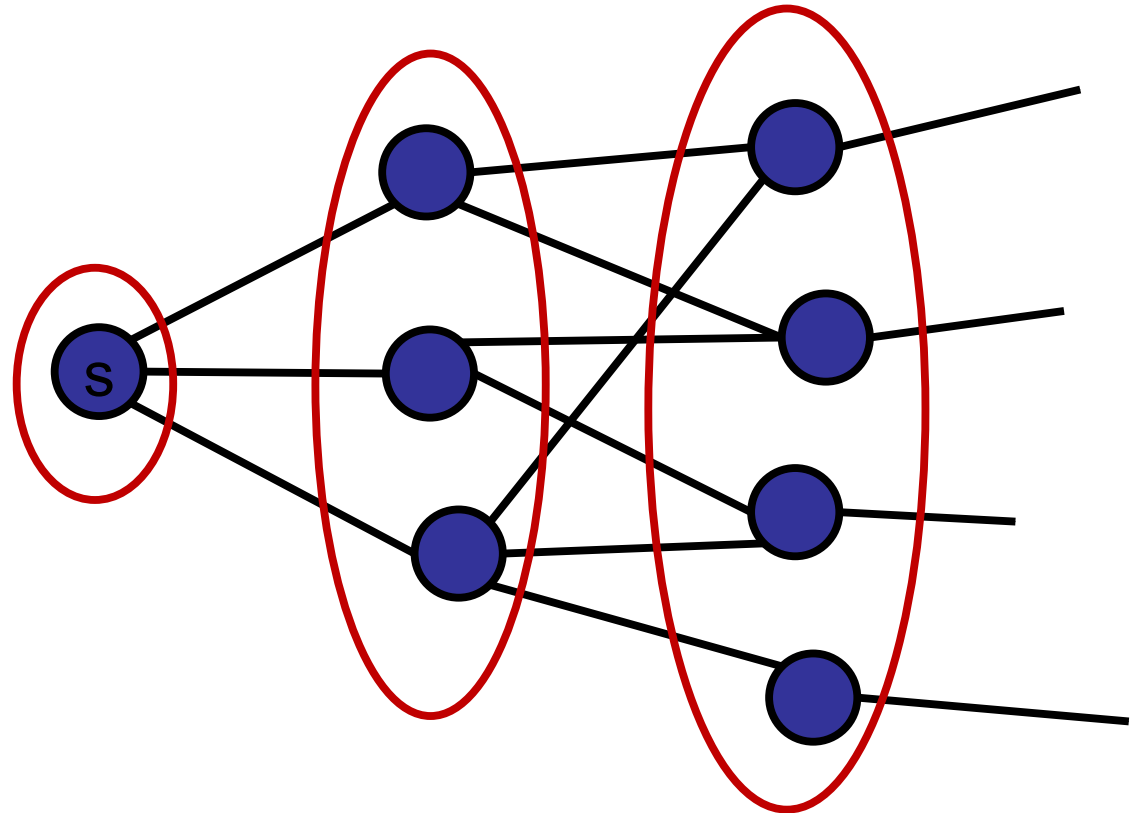
Graph representation:

- Adjacency list

Searching a graph

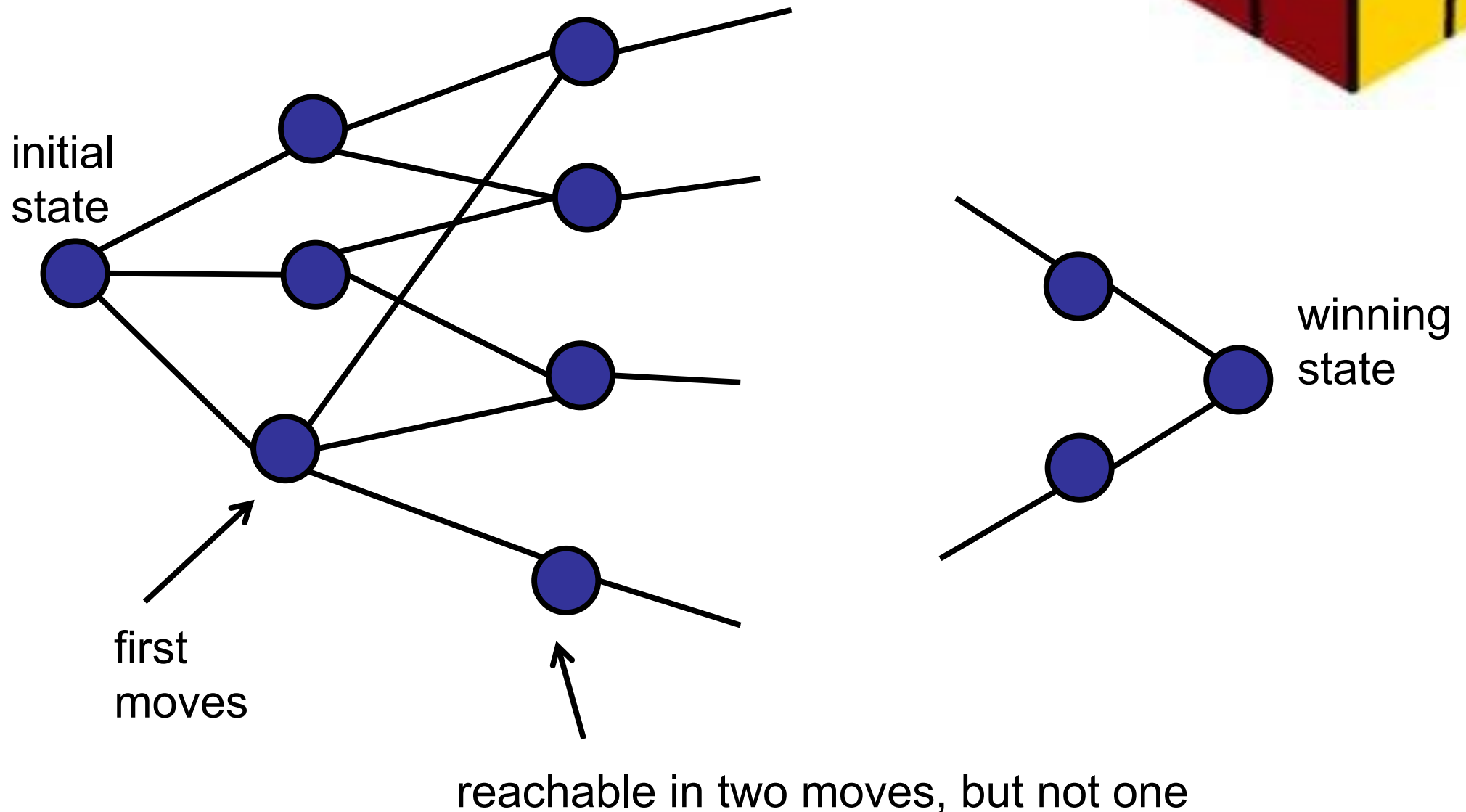
Breadth-First Search:

- Explore level by level



2 x 2 x 2 Rubik's Cube

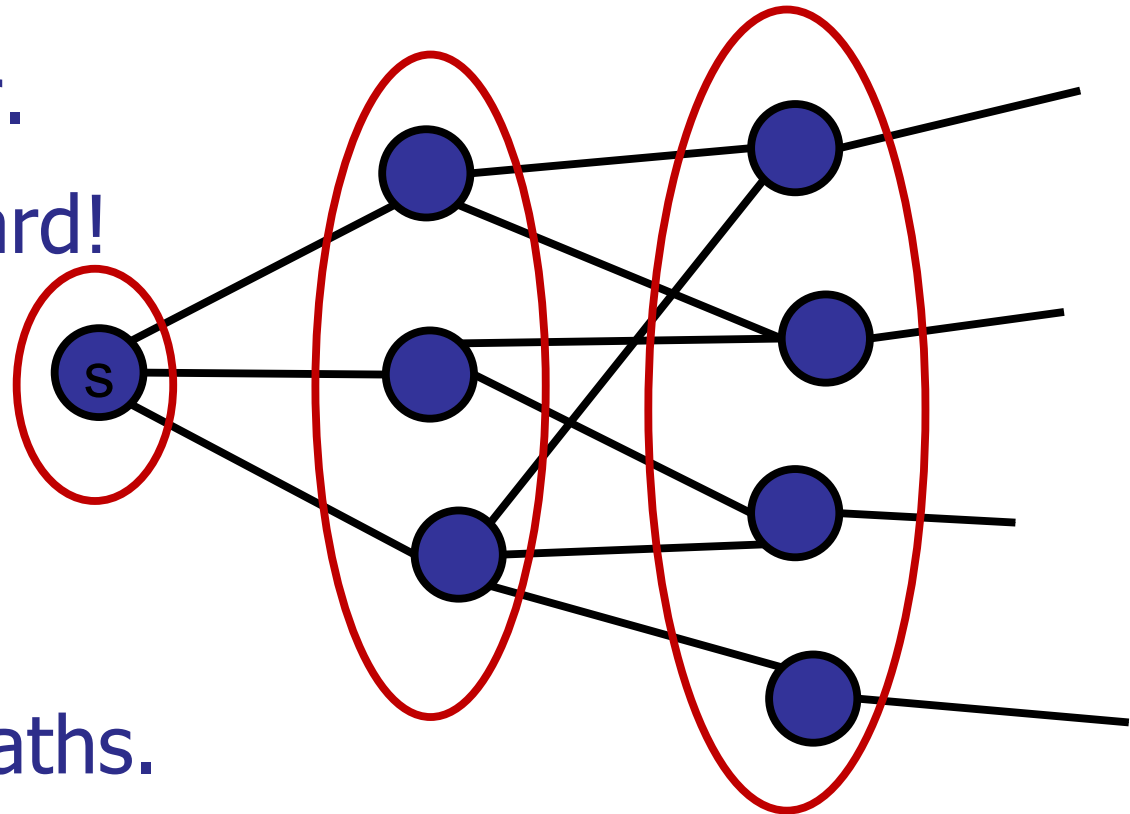
Geography of Rubik's configurations:



Searching a graph

Breadth-First Search:

- Explore level by level
- Frontier: current level
- Initially: $\{s\}$
- Advance frontier.
- Don't go backward!

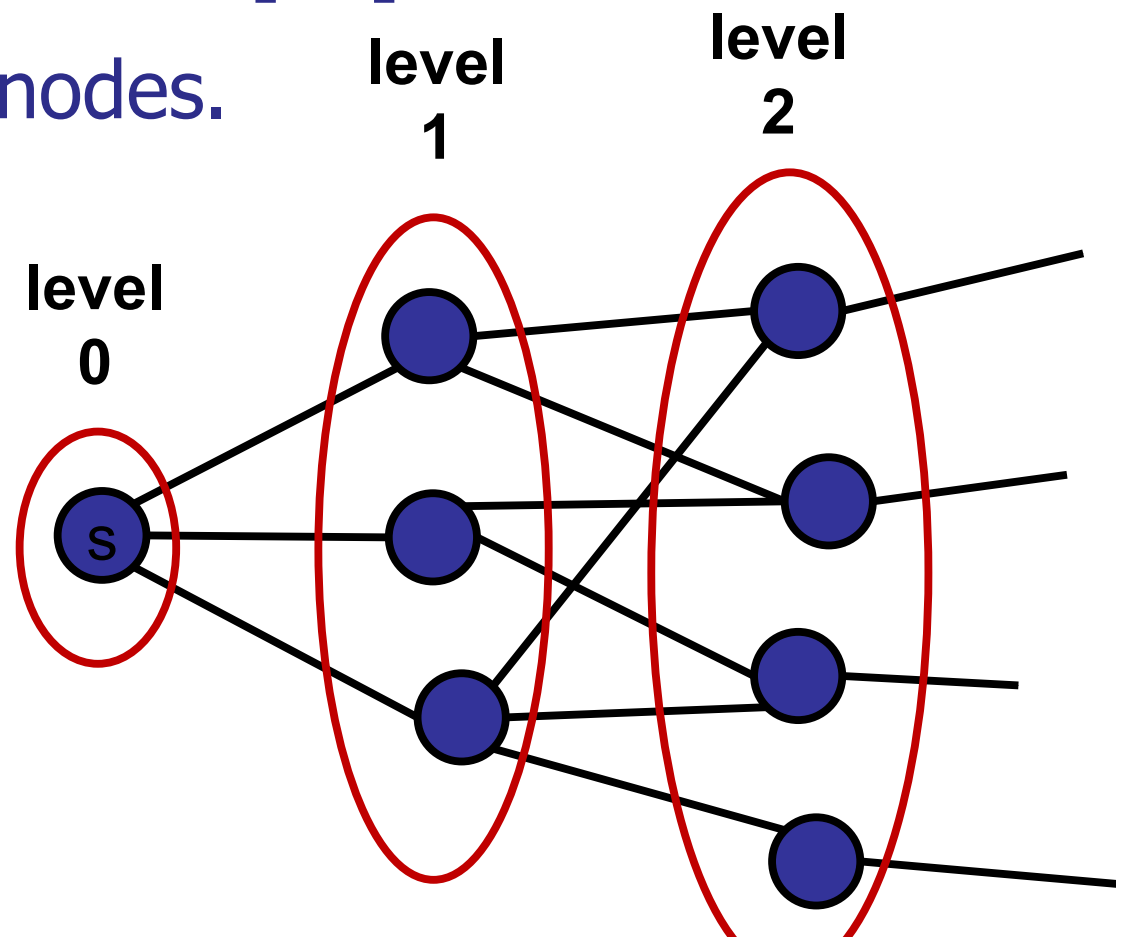


- Finds shortest paths.

Searching a graph

Breadth-First Search:

- Build levels.
- Calculate level[i] from level[i-1]
- Skip already visited nodes.



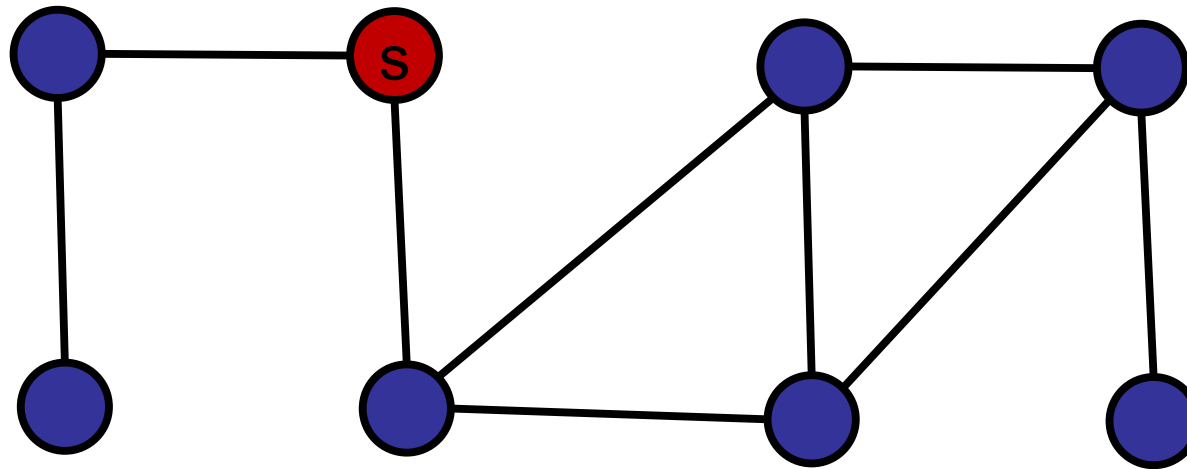
Breadth-First Search

```
BFS(Node[] nodeList, int startId) {  
    boolean[] visited = new boolean[nodeList.length];  
    Arrays.fill(visited, false);  
  
    int[] parent = new int[nodeList.length];  
    Arrays.fill(parent, -1);  
  
    Collection<Integer> frontier = new Collection<Integer>;  
    frontier.add(startId);  
  
    // Main code goes here!  
  
}
```

Breadth-First Search

```
while (!frontier.isEmpty()) {
    Collection<Integer> nextFrontier = new ... ;
    for (Integer v : frontier) {
        for (Integer w : nodeList[v].nbrList) {
            if (!visited[w]) {
                visited[w] = true;
                parent[w] = v;
                nextFrontier.add(w);
            }
        }
    }
    frontier = nextFrontier;
}
```


Breadth-First Search Example



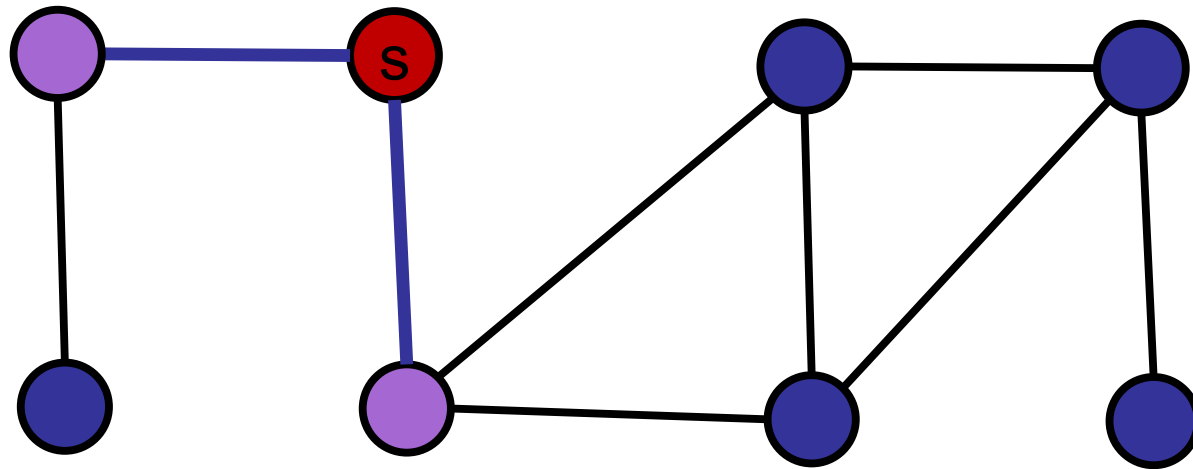
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Breadth-First Search Example



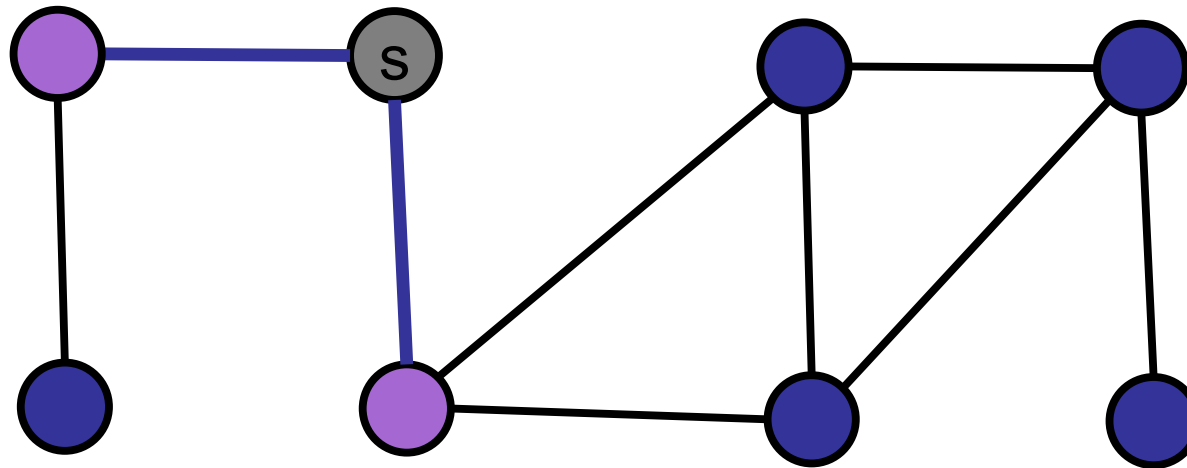
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Breadth-First Search Example



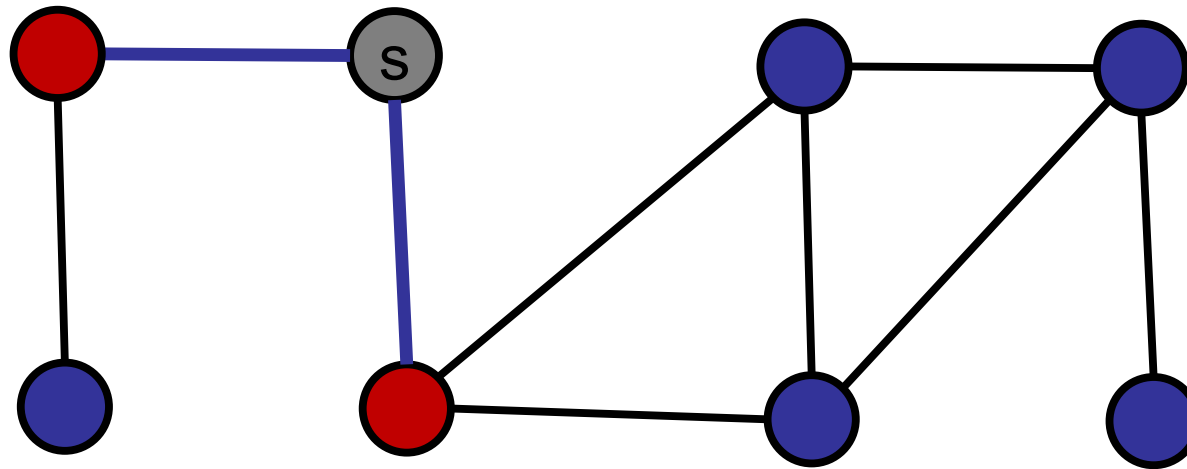
Red = active frontier

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Breadth-First Search Example



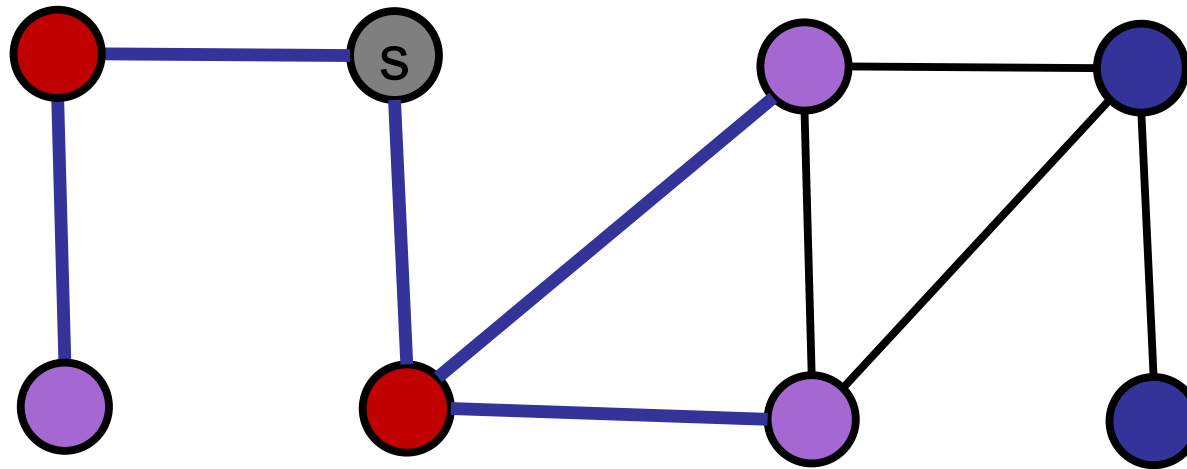
Red = active frontier

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Breadth-First Search Example



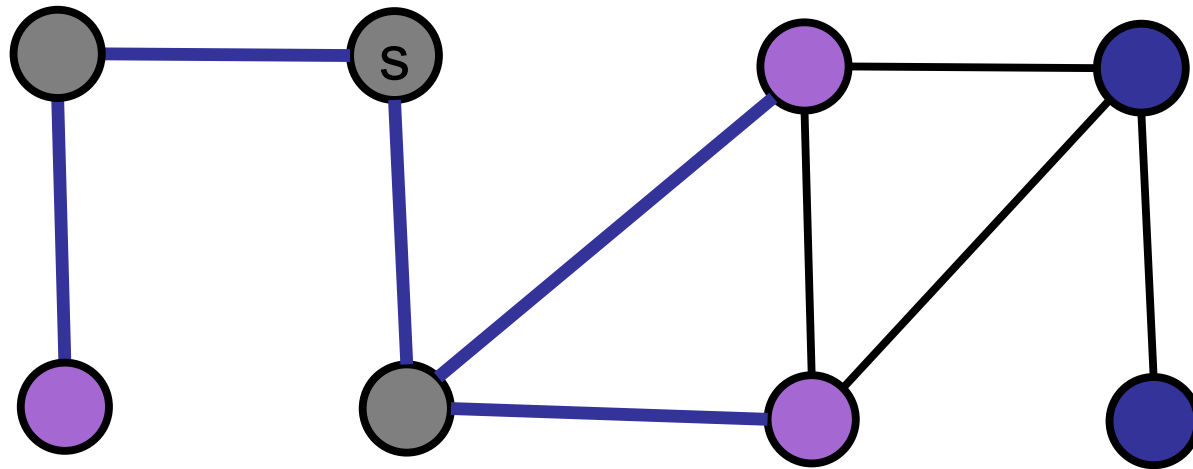
Red = active frontier

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Gray = visited

Blue = unvisited

Breadth-First Search Example



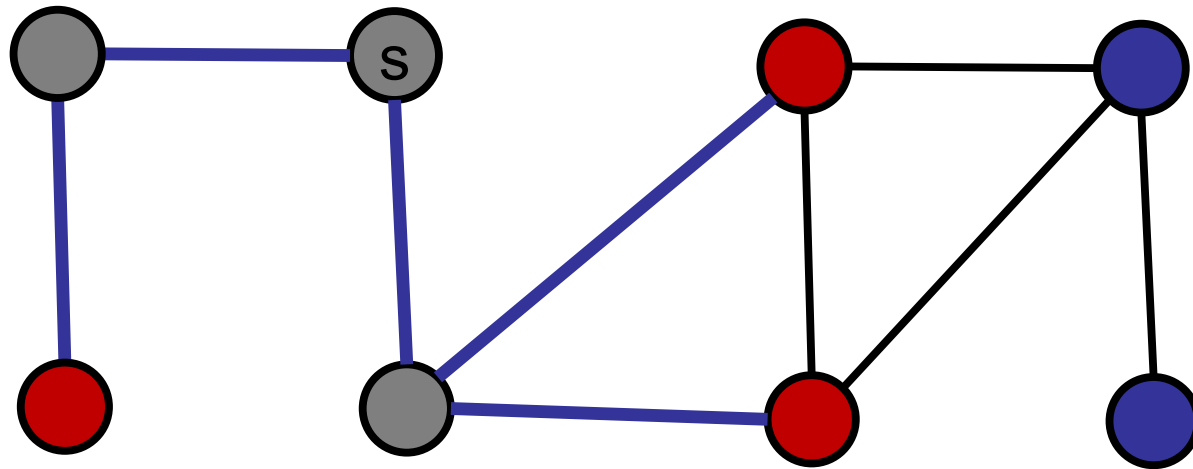
Red = active frontier

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Gray = visited

Blue = unvisited

Breadth-First Search Example



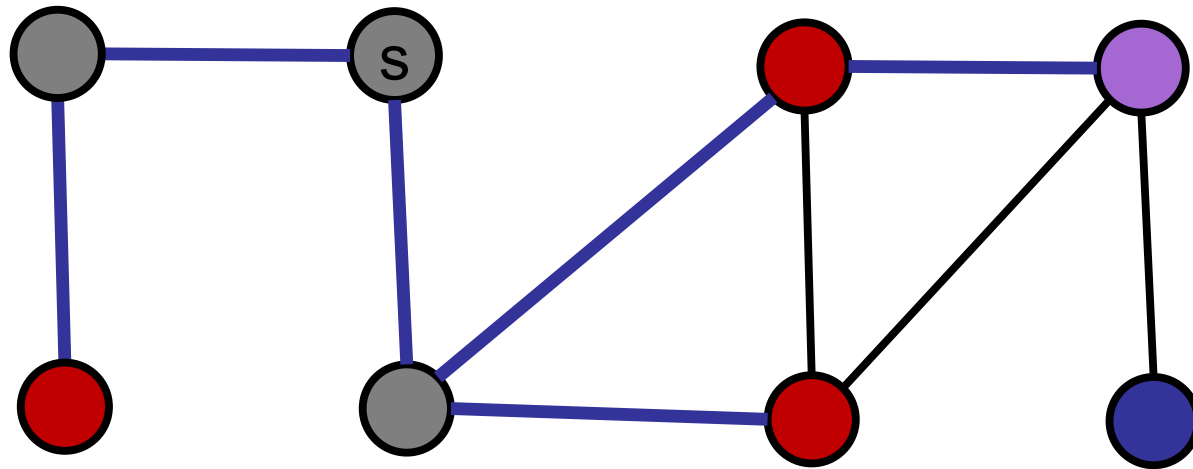
Red = active frontier

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Gray = visited

Blue = unvisited

Breadth-First Search Example



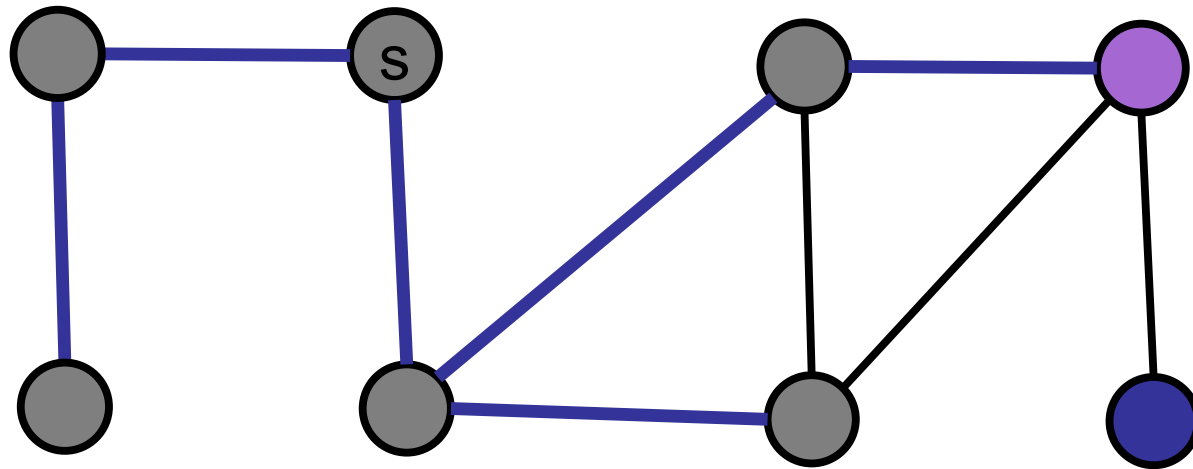
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Breadth-First Search Example



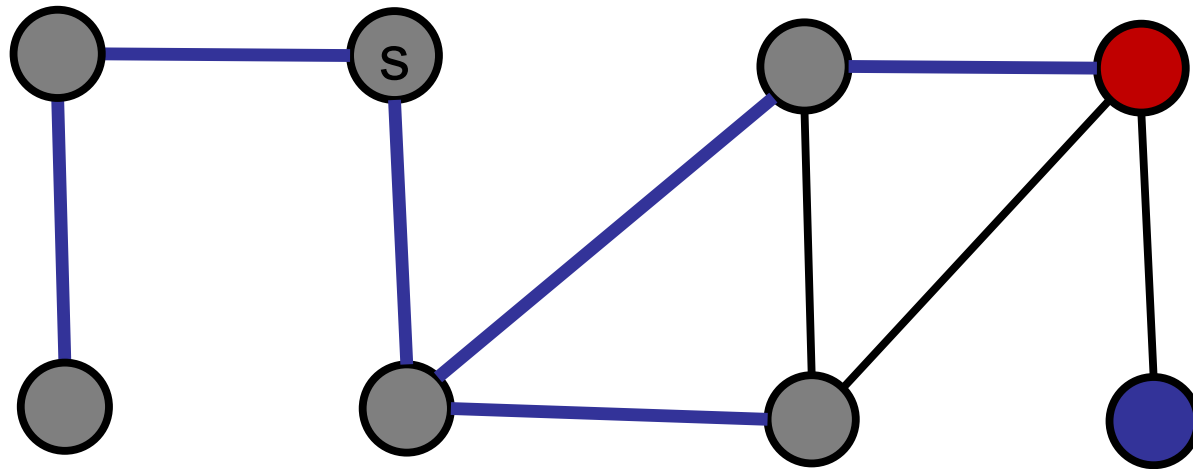
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Breadth-First Search Example



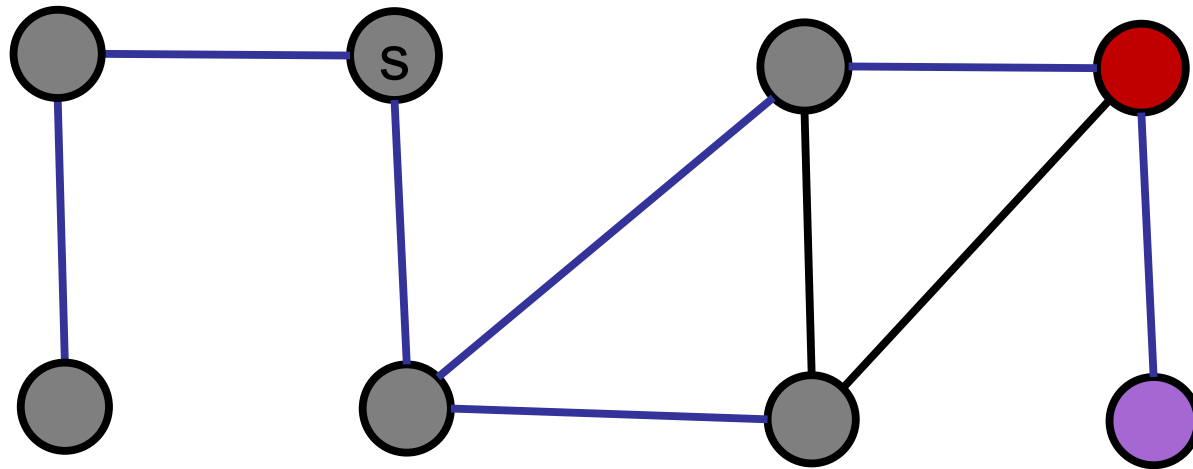
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Breadth-First Search Example



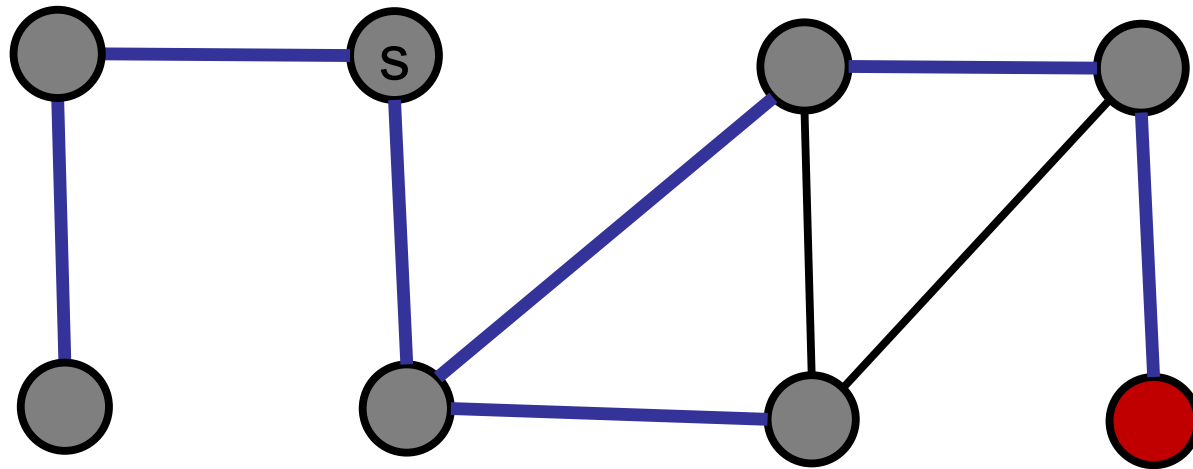
Red = active frontier

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Gray = visited

Blue = unvisited

Breadth-First Search Example



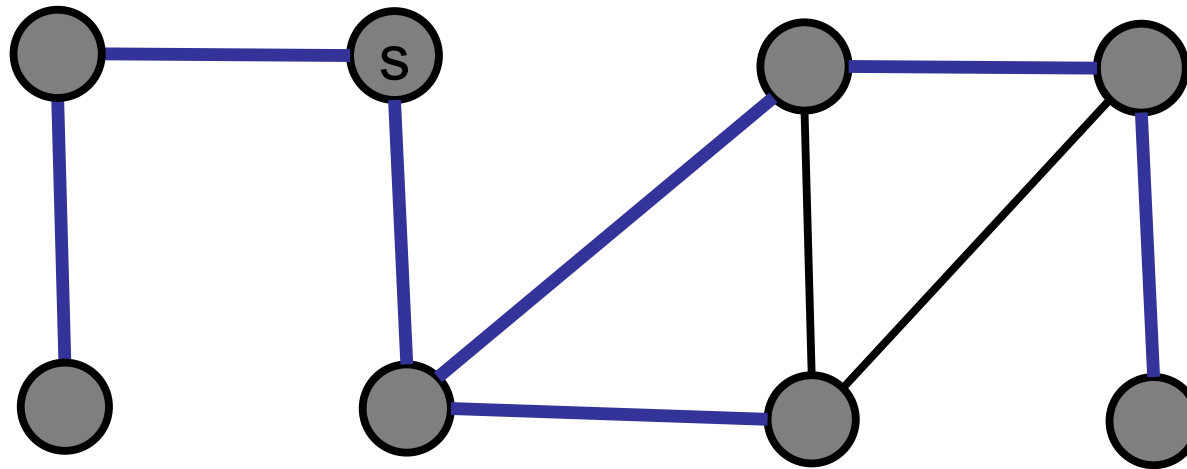
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Breadth-First Search Example



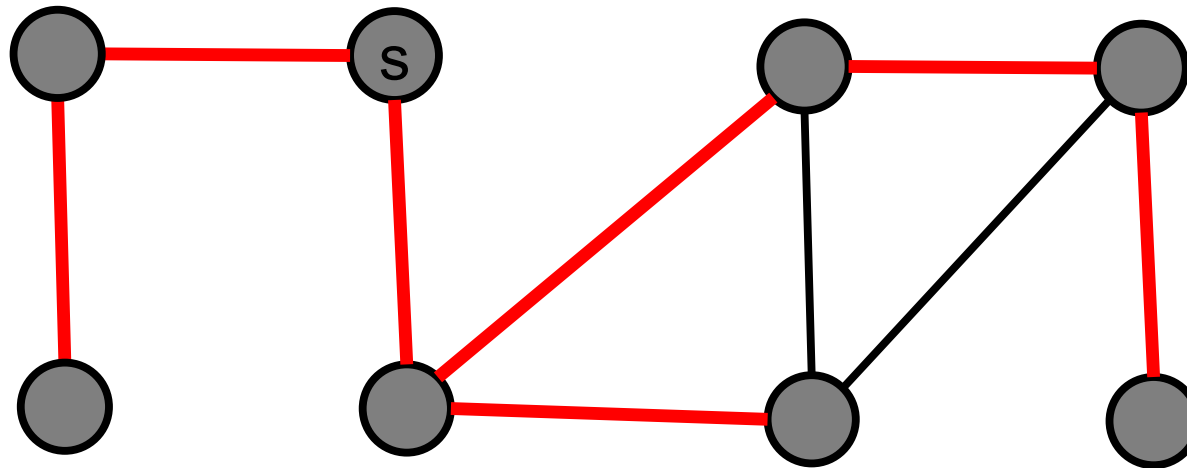
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Breadth-First Search Example




Red = active frontier

Purple = next

Gray = visited

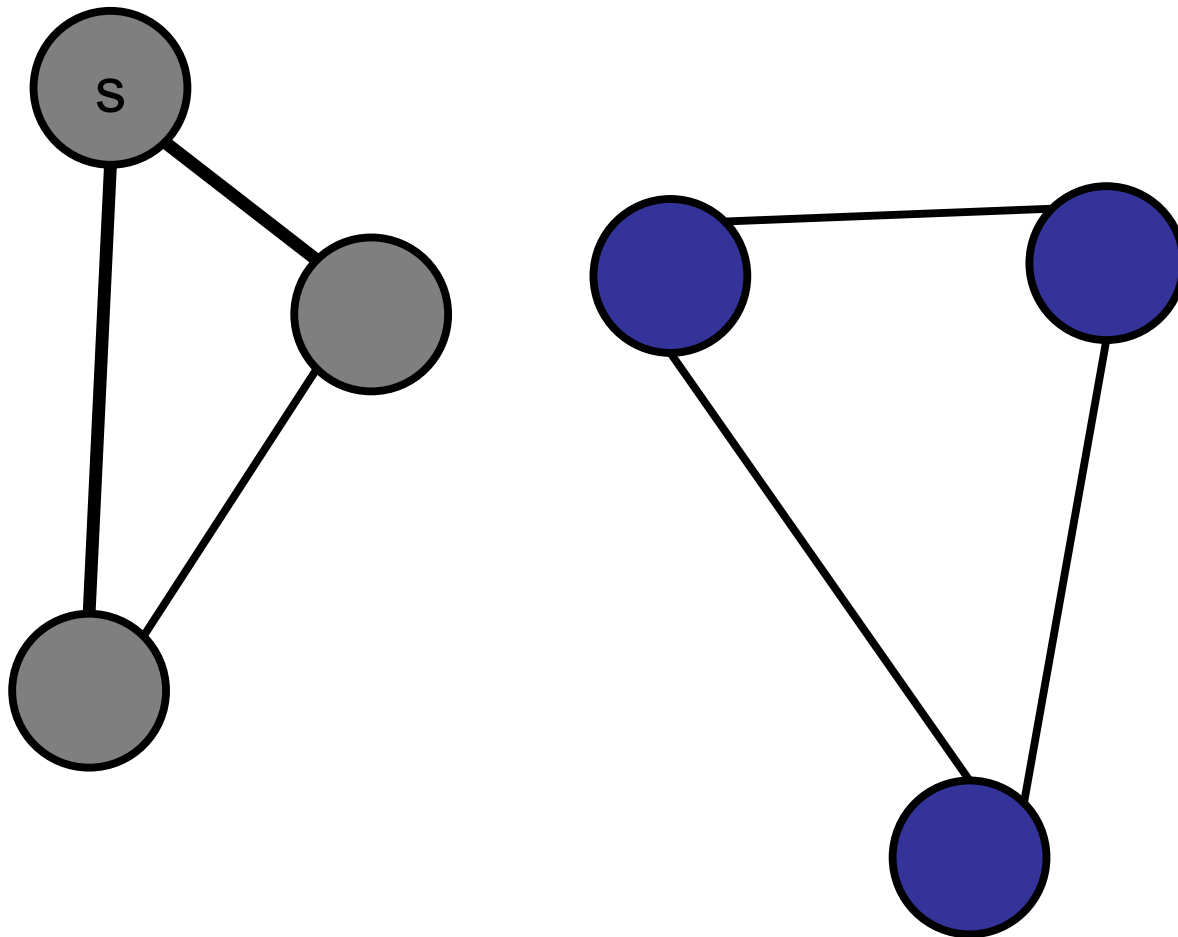
Blue = unvisited

When does BFS fail to visit every node?

1. In a clique.
2. In a cycle.
-  3. In a graph with two components.
4. In a sparse graph.
5. In a dense graph.
6. Never.

BFS on Disconnected Graph

Example:



Breadth-First Search


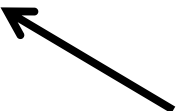
```
BFS(Node[] nodeList) {  
    boolean[] visited = new boolean[nodeList.length];  
    Arrays.fill(visited, false);  
  
    int[] parent = new int[nodeList.length];  
    Arrays.fill(parent, -1);  
  
    for (int start = 0; start < nodeList.length; start++) {  
        if (!visited[start]){  
            Bag<Integer> frontier = new Bag<Integer>;  
            frontier.add(startId);  
  
            // Main code goes here!  
        }  
    }  
}
```

The running time of BFS is:

1. $O(V)$
2. $O(E)$
- ✓ 3. $O(V+E)$
4. $O(VE)$
5. (V^2)
6. I have no idea.

Breadth-First Search

Analysis:

- Vertex v = “start” once.  $O(V)$
- Vertex v added to nextFrontier (and frontier) once.
 - After visited, never re-added.
- Each $v.nbrlist$ is enumerated once.
 - When v is removed from frontier.  $O(E)$

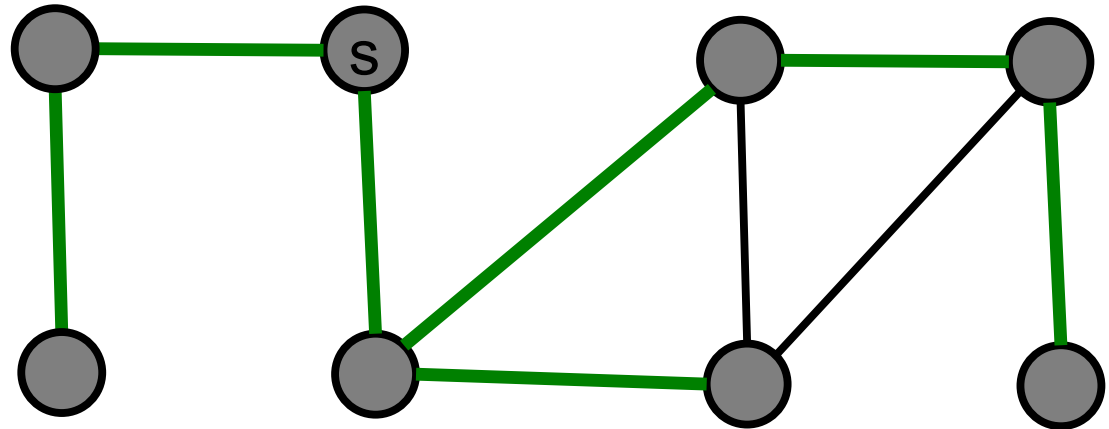
Breadth-First Search

```
while (!frontier.isEmpty()) {  
    Collection<Integer> next = new Collection<Integer>;  
    for (Integer v : frontier) {  
        for (Integer w : nodeList[v].nbrList) {  
            if (!visited[w]) {  
                visited[w] = true;  
                parent[w] = v;  
                next.add(w);  
            }  
        }  
    }  
    frontier = next;  
}
```

Breadth-First Search

Shortest paths:

- Parent pointers store shortest path.



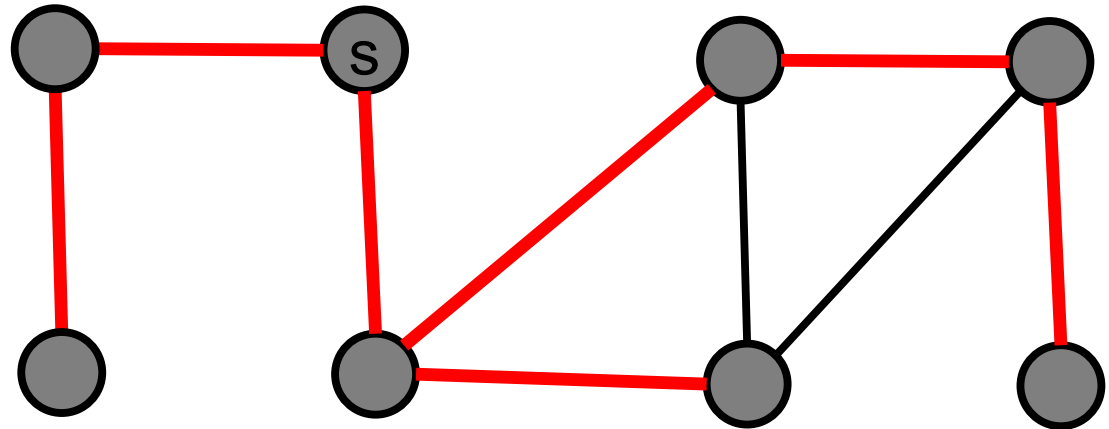
Which is true? (More than one may apply.)

1. Shortest path graph is a cycle.
- ✓ 2. Shortest path graph is a tree.
3. Shortest path graph has low-degree.
4. Shortest path graph has low diameter.
5. None of the above.

Breadth-First Search

Shortest paths:

- Parent pointers store shortest path.
- Shortest path is a tree.
- (Possibly high degree; possibly high diameter.)



What if there are
two components?

Searching a Graph

Goal:

- Start at some vertex **s** = start.
- Find some other vertex **f** = finish.

Or: visit **all** the nodes in the graph;

Two basic techniques:

- Breadth-First Search (BFS)
- Depth-First Search (DFS)

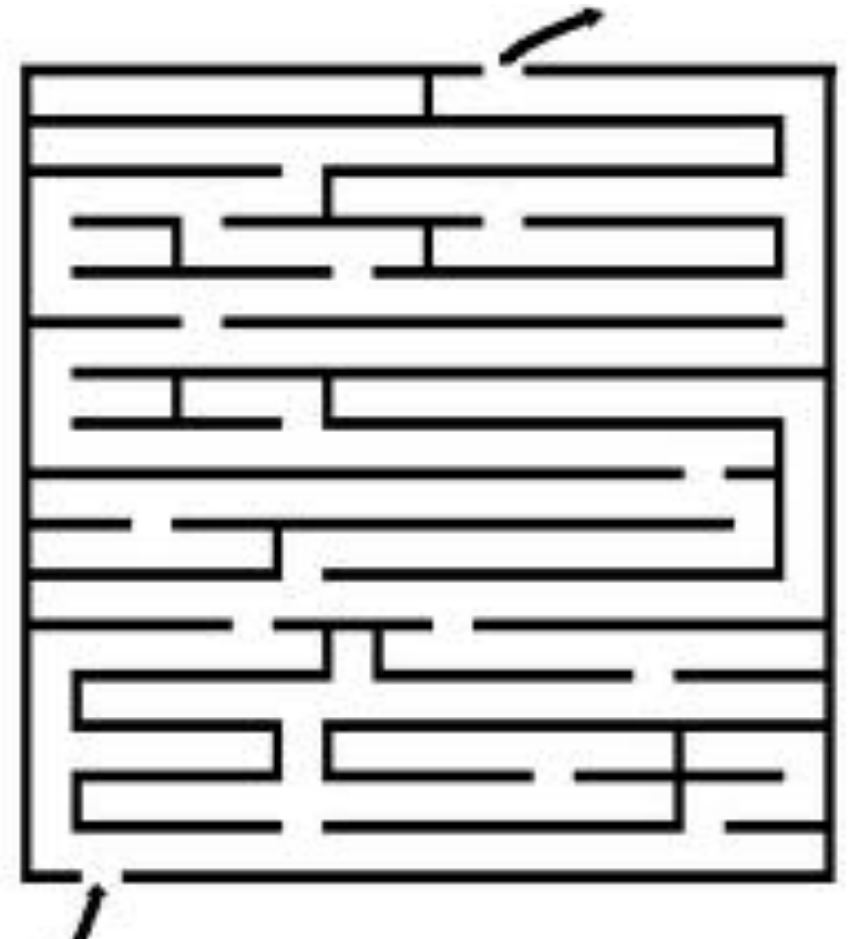
Graph representation:

- Adjacency list

Depth-First Search

Exploring a maze:

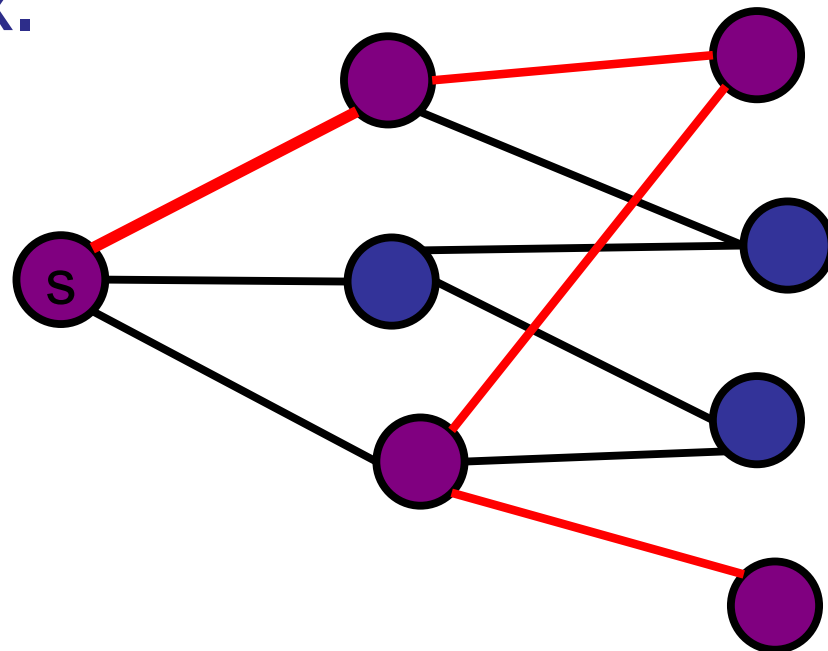
- Follow path until stuck.
- Backtrack along breadcrumbs until reach unexplored neighbor.
- Recursively explore.



Searching a graph

Depth-First Search:

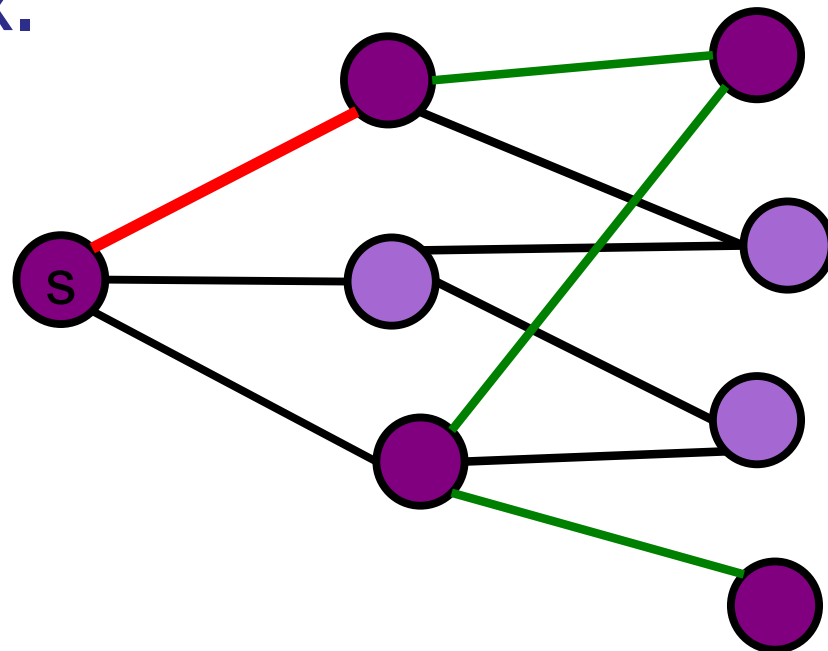
- Follow path until you get stuck
- Backtrack until you find a new edge
- Recursively explore it
- Don't repeat a vertex.



Searching a graph

Depth-First Search:

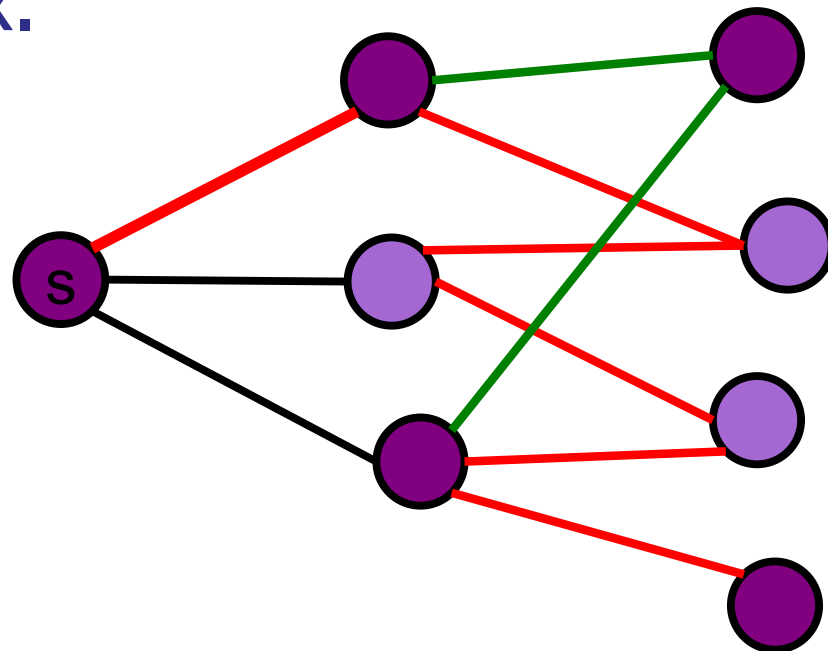
- Follow path until you get stuck
- Backtrack until you find a new edge
- Recursively explore it
- Don't repeat a vertex.



Searching a graph

Depth-First Search:

- Follow path until you get stuck
- Backtrack until you find a new edge
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- Don't repeat a vertex.



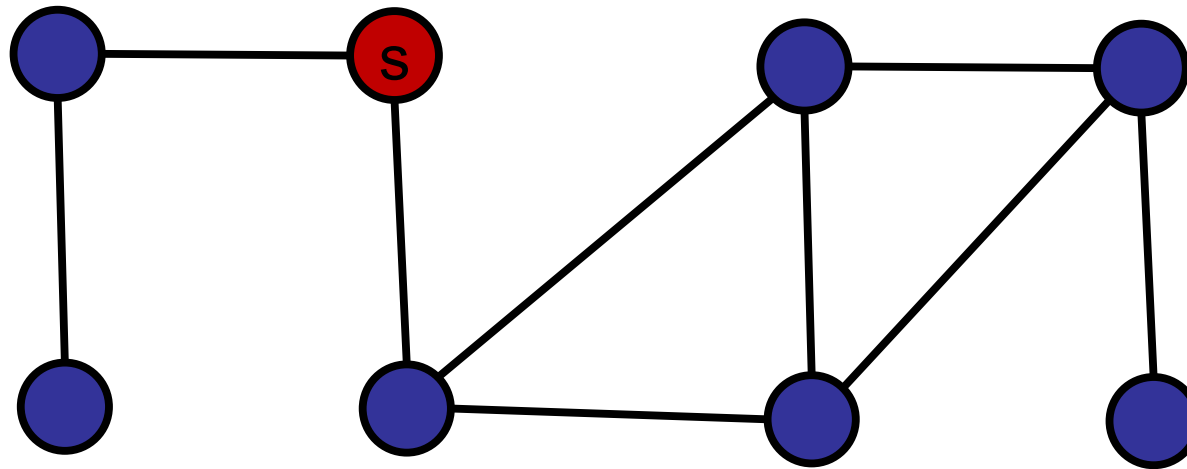
Depth-First Search

```
DFS-visit(Node[] nodeList, boolean[] visited, int startId){  
    for (Integer v : nodeList[startId].nbrList) {  
        if (!visited[v]){  
            visited[v] = true;  
            DFS-visit(nodeList, visited, v);  
        }  
    }  
}
```

Depth-First Search

```
DFS(Node[] nodeList) {  
    boolean[] visited = new boolean[nodeList.length];  
    Arrays.fill(visited, false);  
  
    for (start = 0; start < nodeList.length; start++) {  
        if (!visited[start]) {  
            visited[start] = true;  
            DFS-visit(nodeList, visited, start);  
        }  
    }  
}
```

Depth-First Search Example



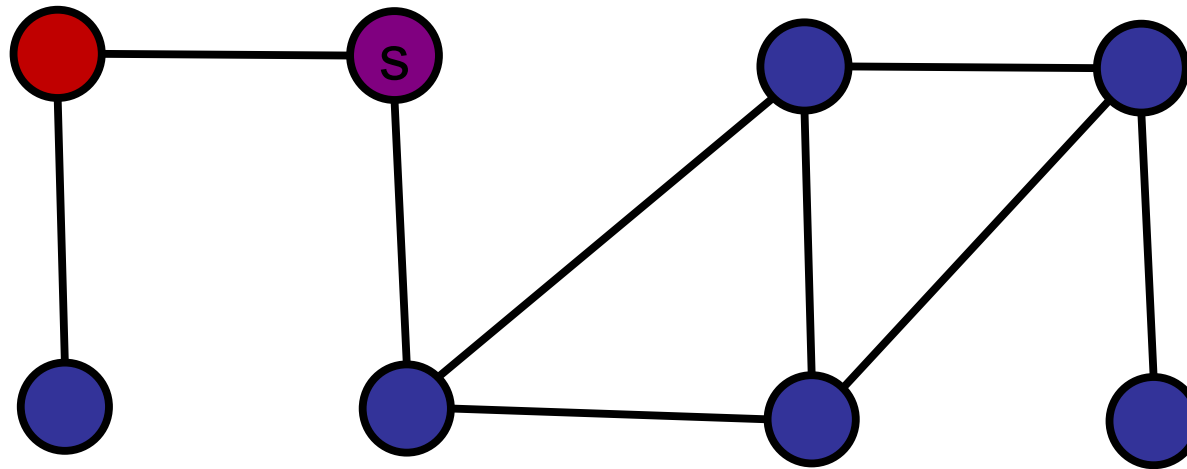
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Depth-First Search Example



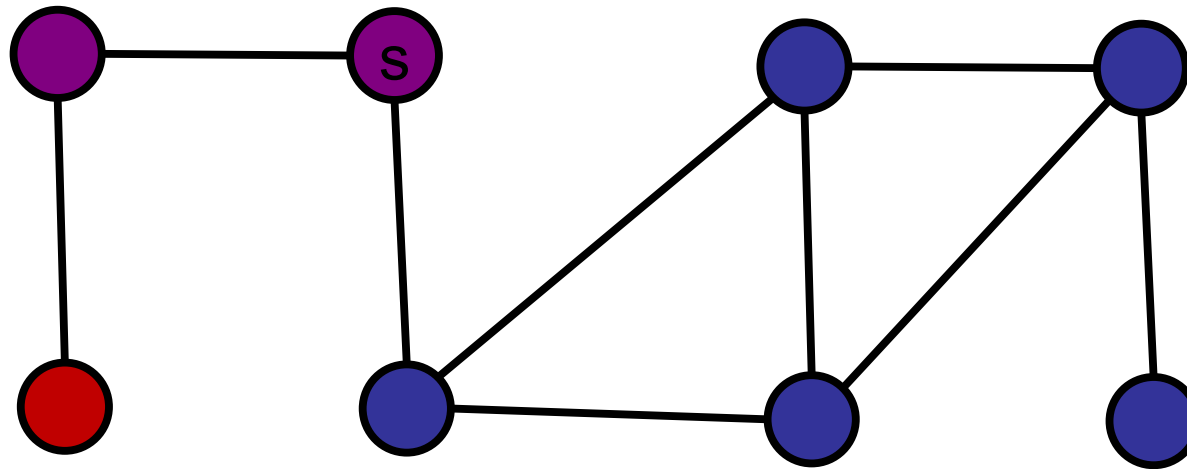
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Depth-First Search Example



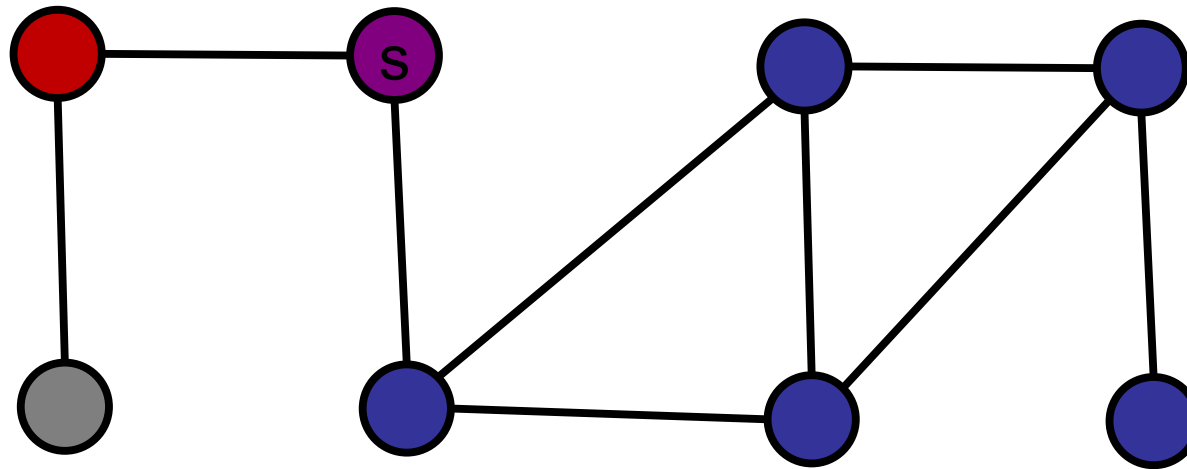
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Depth-First Search Example



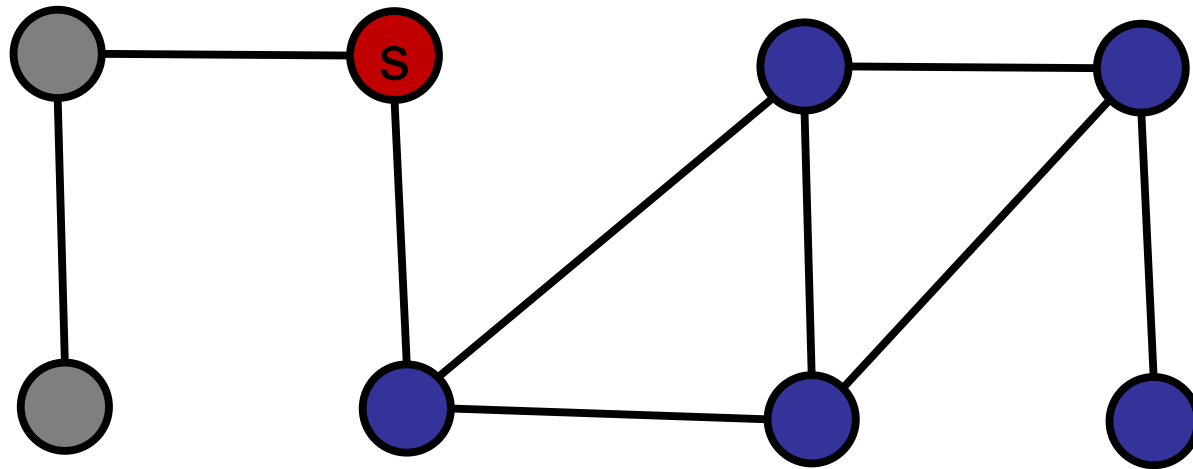
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Depth-First Search Example



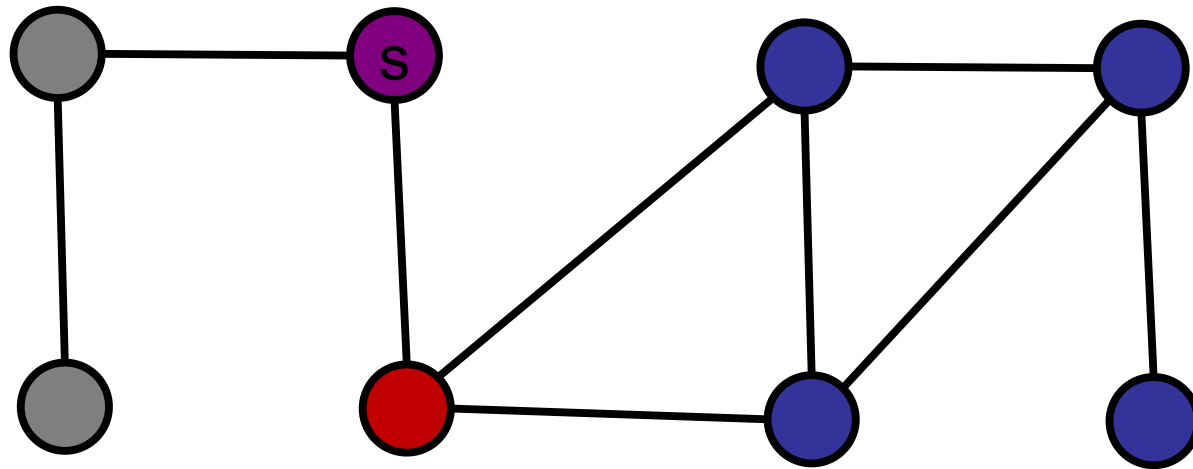
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Depth-First Search Example



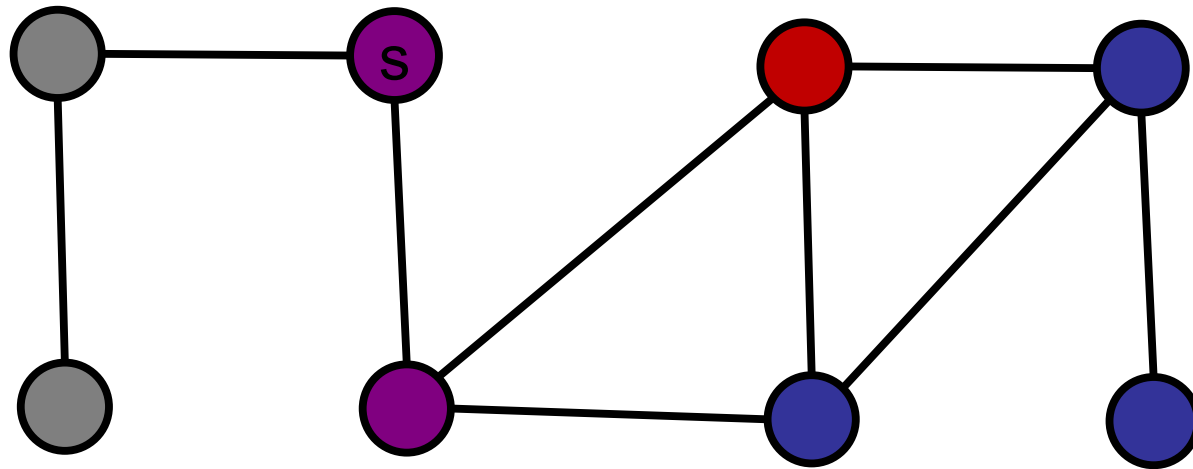
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Purple = next

Gray = visited

Blue = unvisited

Depth-First Search Example



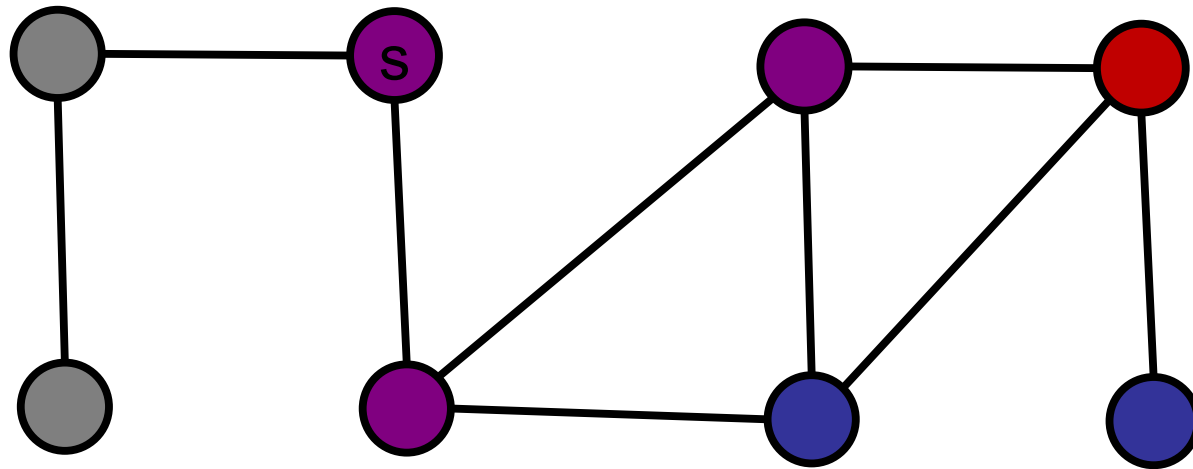
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Depth-First Search Example



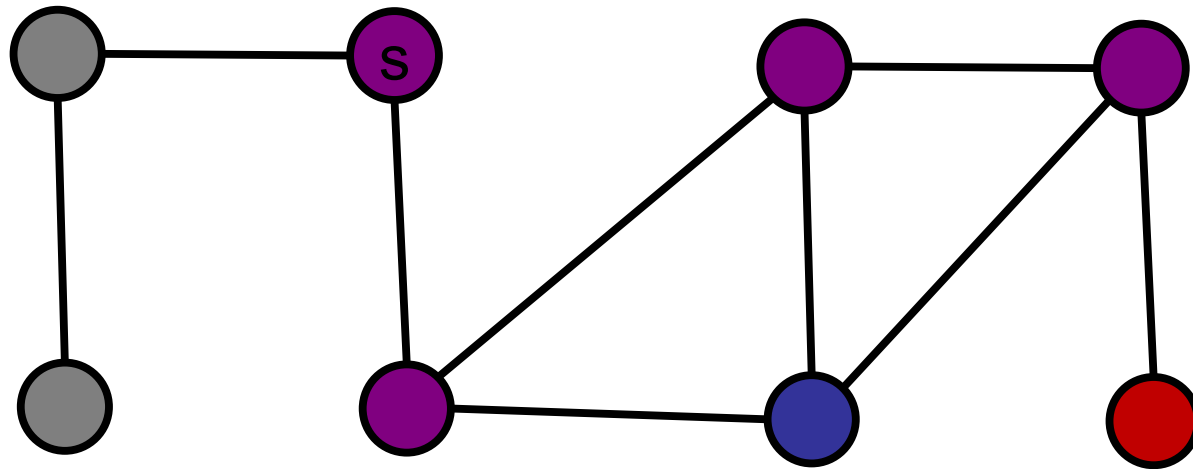
Red = active frontier

Purple = next

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Blue = unvisited

Depth-First Search Example



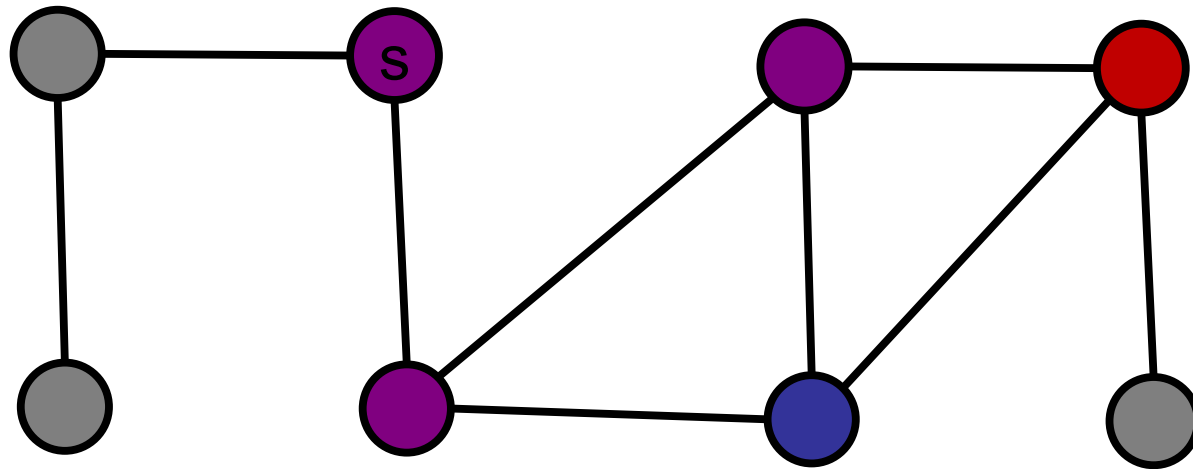
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Depth-First Search Example



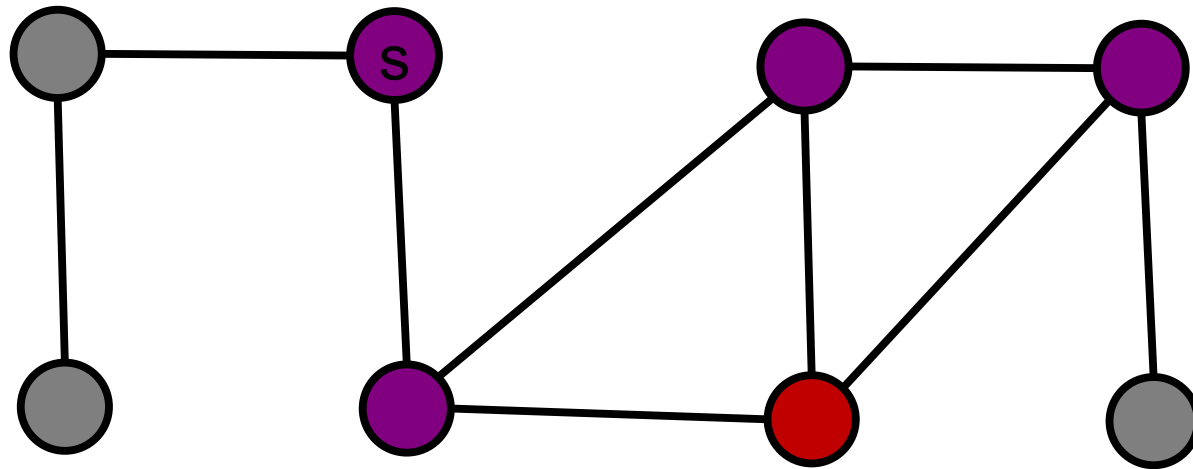
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Depth-First Search Example



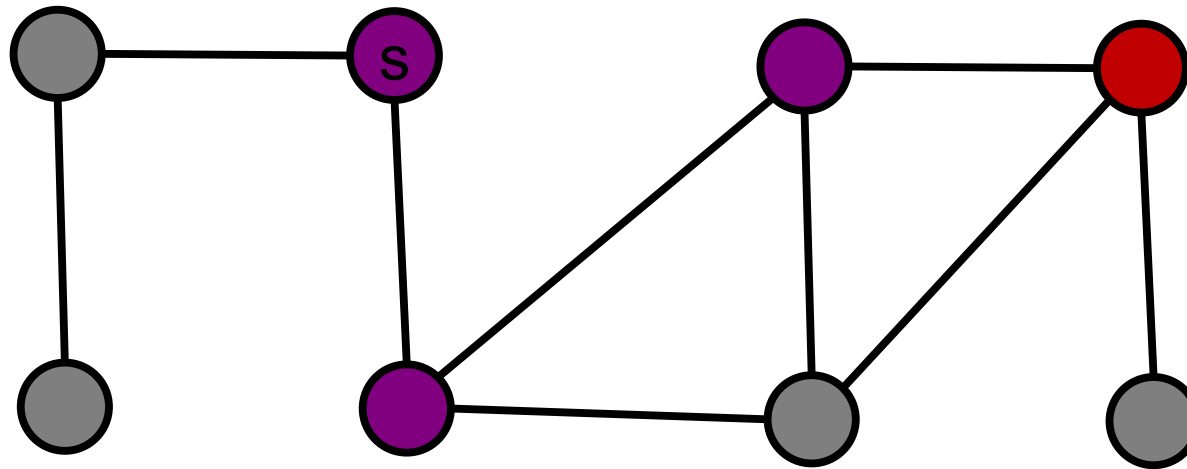
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Depth-First Search Example



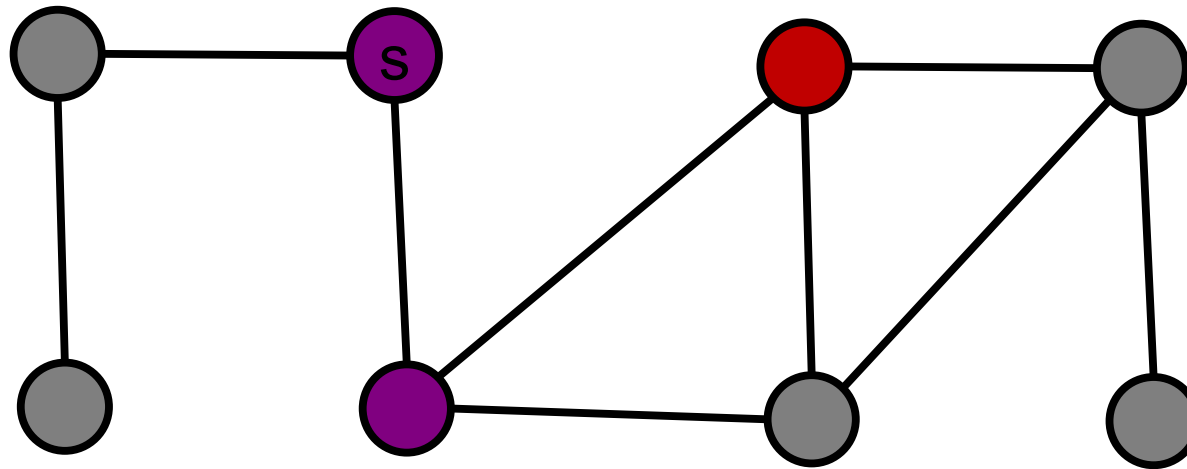
Red = active frontier

Purple = next

Gray = visited

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Depth-First Search Example



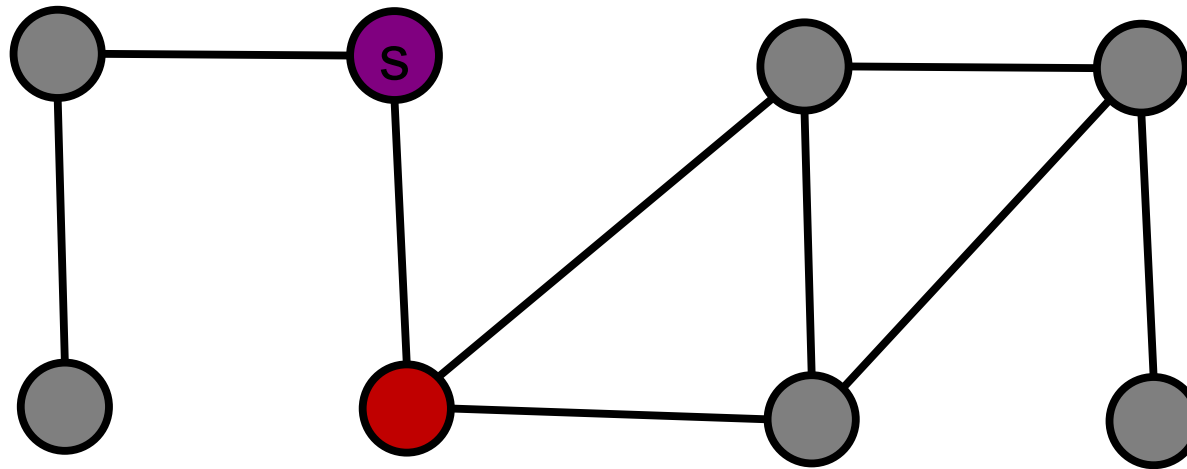
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Gray = visited

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Depth-First Search Example



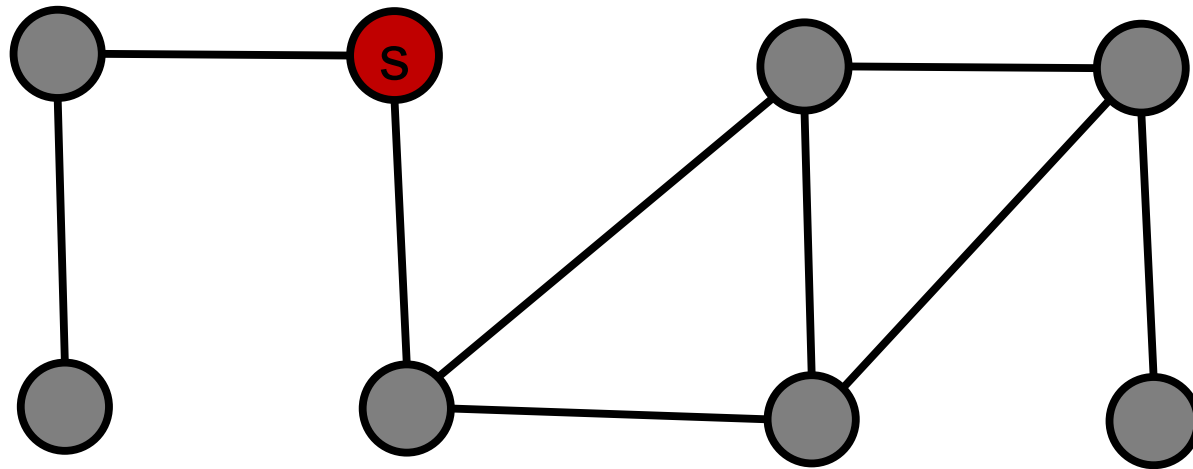
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Depth-First Search Example



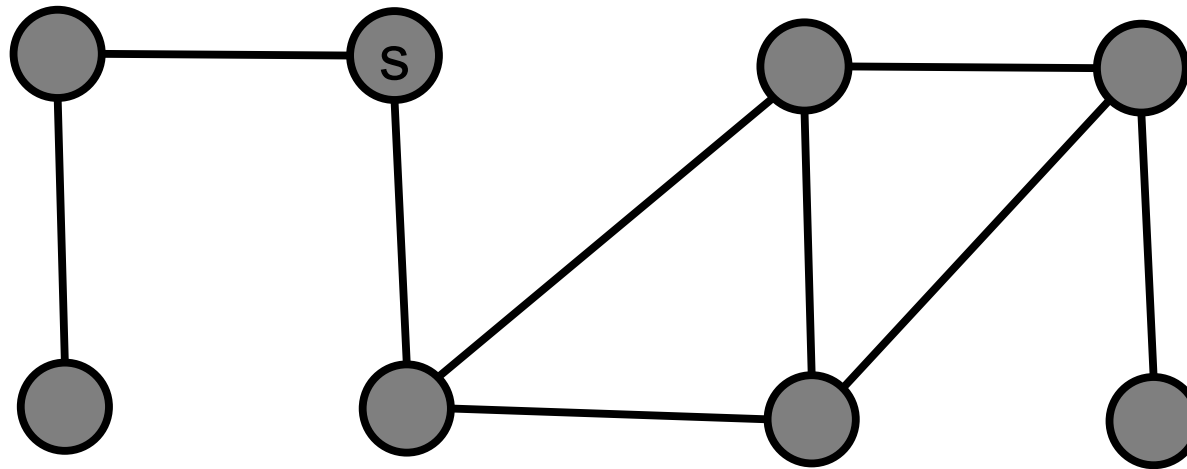
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

Depth-First Search Example



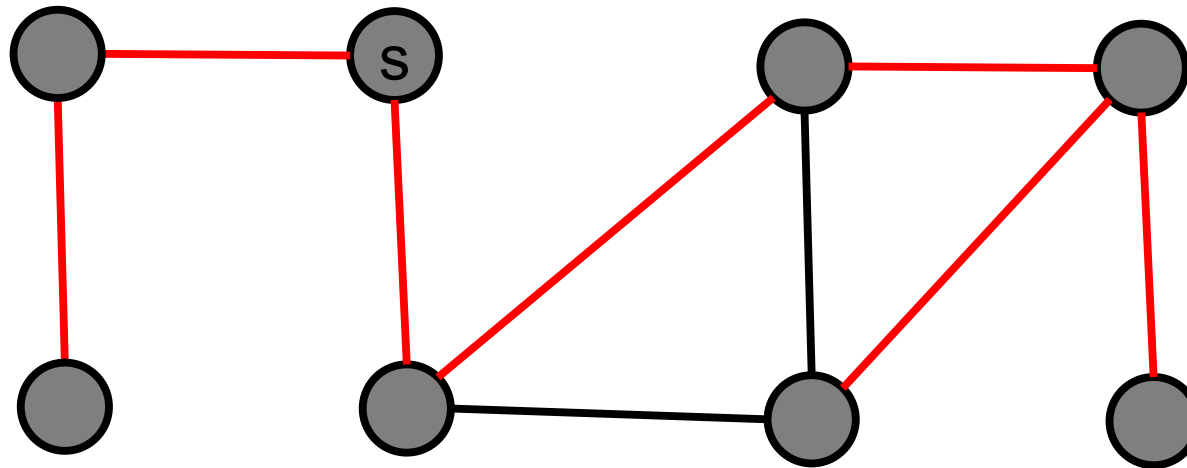
Red = active frontier

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Gray = visited

Blue = unvisited

Depth-First Search Example



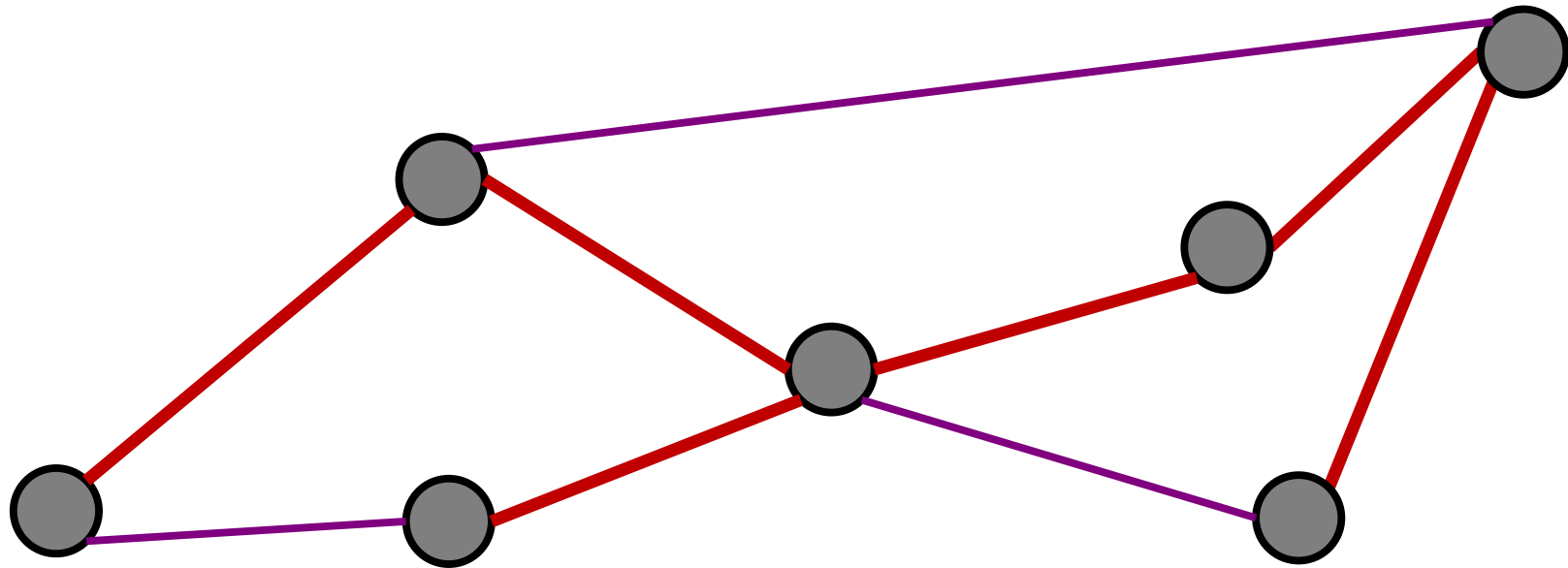
Red = active frontier

Purple = next

Gray = visited

Blue = unvisited

DFS parent edges



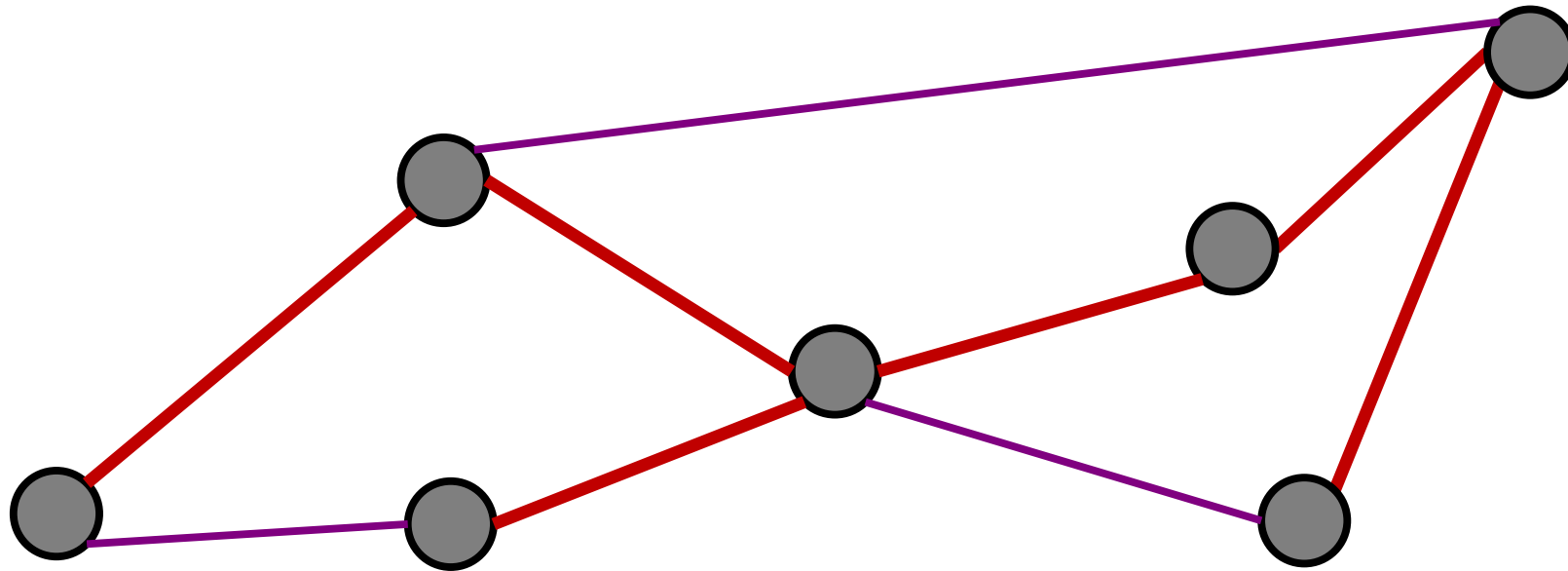
Red = Parent Edges

Purple = Non-parent edges

Which is true? (More than one may apply.)

1. DFS parent graph is a cycle.
- ✓ 2. DFS parent graph is a tree.
3. DFS parent graph has low-degree.
4. DFS parent graph has low diameter.
5. None of the above.

DFS parent edges = tree



Red = Parent Edges

Purple = Non-parent edges

Note: not shortest paths!

The running time of DFS is:

1. $O(V)$
2. $O(E)$
- ✓ 3. $O(V+E)$
4. $O(VE)$
5. (V^2)
6. I have no idea.

Depth-First Search

Analysis:

- DFS-visit called only once per node.
 - After visited, never call DFS-visit again.
- In DFS-visit, each neighbor is enumerated.

$O(V)$



$O(E)$



If the graph is stored as an adjacency matrix, what is the running time of DFS?

1. $O(V)$
2. $O(E)$
3. $(V+E)$
4. $O(VE)$
- ✓ 5. $O(V^2)$
6. $O(E^2)$

Depth-First Search

Analysis:

- DFS-visit called only once per node.
 - After visited, never call DFS-visit again.
- In DFS-visit, each neighbor is enumerated.

$O(V)$



$O(V)$



per
node

To implement an iterative version of DFS:

1. Use a queue.
- ✓ 2. Use a stack.
3. Use a bag.
4. Use a set.
5. Don't.

Graph Search

BFS and DFS are the same algorithm:

- BFS: use a queue
 - Every time you visit a node, add all unvisited neighbors to the queue.
- DFS: use a stack
 - Every time you visit a node, add all unvisited neighbors to the stack.

Graph Search

Breadth-first search:

Same algorithm, implemented with a queue:

Add start-node to queue.

Repeat until queue is empty:

- Remove node v from the front of the queue.
- Visit v .
- Explore all outgoing edges of v .
- Add all unvisited neighbors of v to the queue.

Graph Search

Depth-first search:

Same algorithm, implemented with a stack:

Add start-node to stack.

Repeat until stack is empty:

- Pop node v from the front of the stack.
- Visit v .
- Explore all outgoing edges of v .
- Push all unvisited neighbors of v on the front of the stack.

Review: Searching Graphs

BFS and DFS are the same algorithm:

- BFS: use a queue
 - Every time you visit a node, add all unvisited neighbors to the queue.
- DFS: use a stack
 - Every time you visit a node, add all unvisited neighbors to the stack.

Roadmap

Today: Graph Basics

- What is a graph?
- Modeling problems as graphs.
- Graph representations (list vs. matrix)
- Searching graphs (DFS / BFS)