CS2040S Data Structures and Algorithms

(e-learning edition)

Augmented Trees!
Part 2

Today

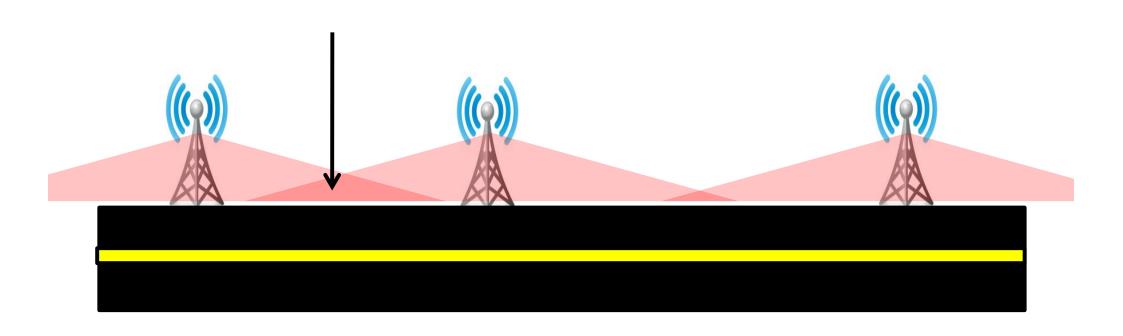
Three examples of augmenting balanced BSTs

1. Order Statistics

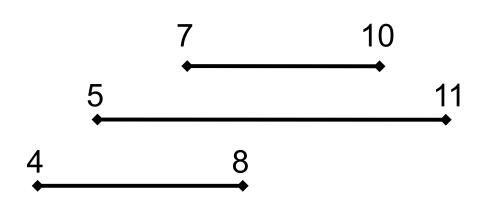
2. Intervals

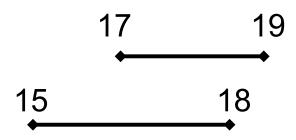
3. Orthogonal Range Searching

Find a tower that covers my location.

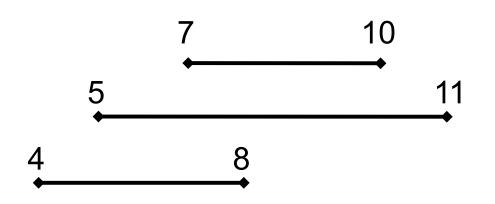


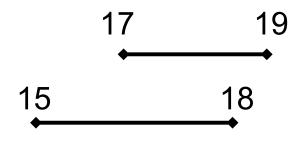
Find a tower that covers my location.





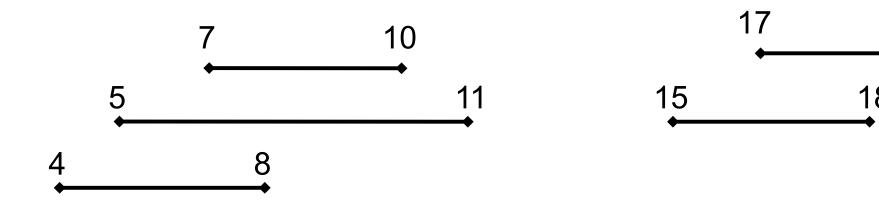
Find a tower that covers my location.





insert(begin, end) delete(begin, end)

Find a tower that covers my location.

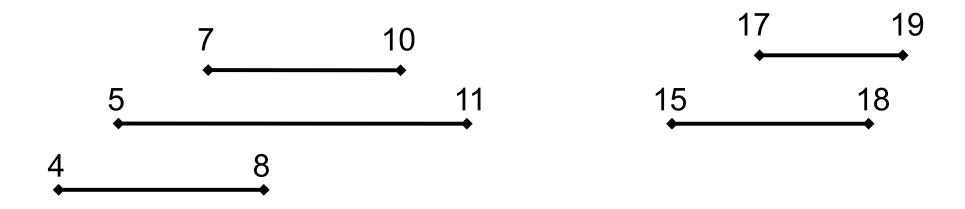


19

insert(begin, end) delete(begin, end)

query(x): find an interval that overlaps x.

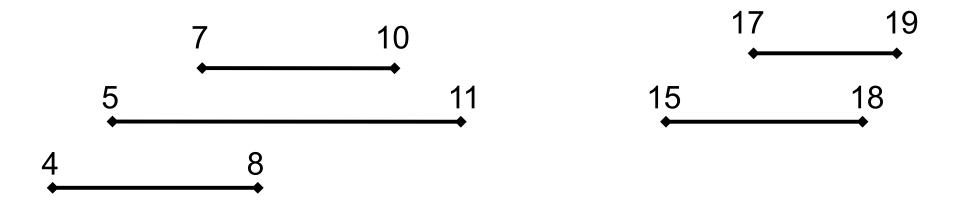
Find a tower that covers my location.



Idea 1: Keep intervals in a list.

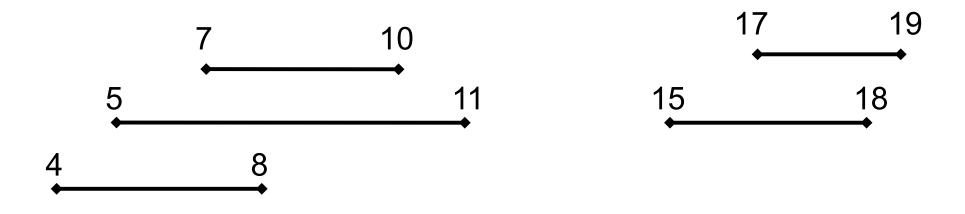
Query: scan entire list.

Find a tower that covers my location.



Idea 2: O(1) queries??

Find a tower that covers my location.

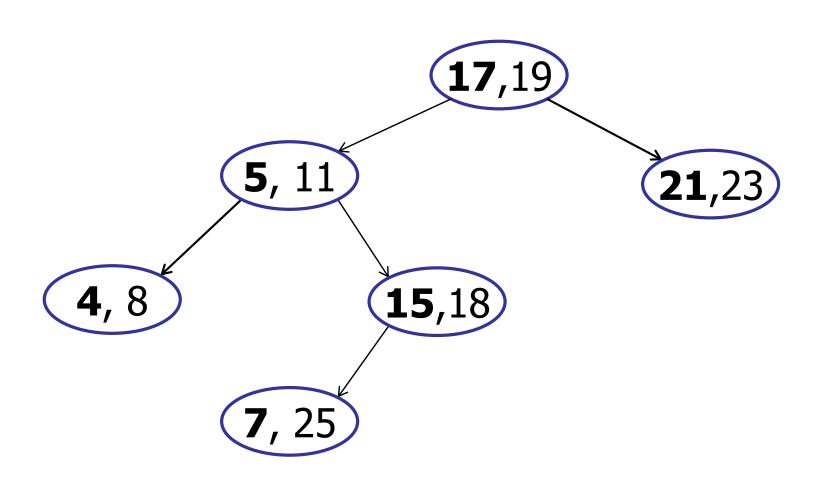


Idea 2: O(1) queries

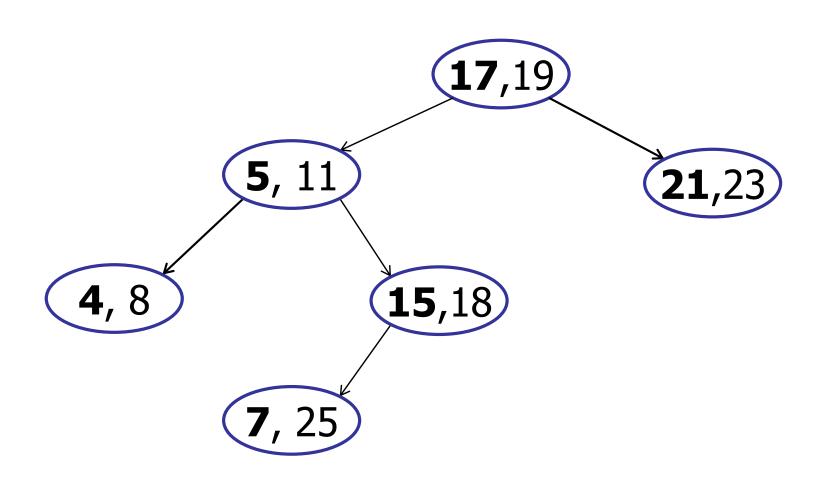
| | | | A | A | A | A | A | В | В | C | | | | D | D | D | D | Ε | |
|---|---|---|---|---|---|---|---|---|----|----|----|----|----|----|----|----|----|----|----|
| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 |

Idea 3: Interval Trees

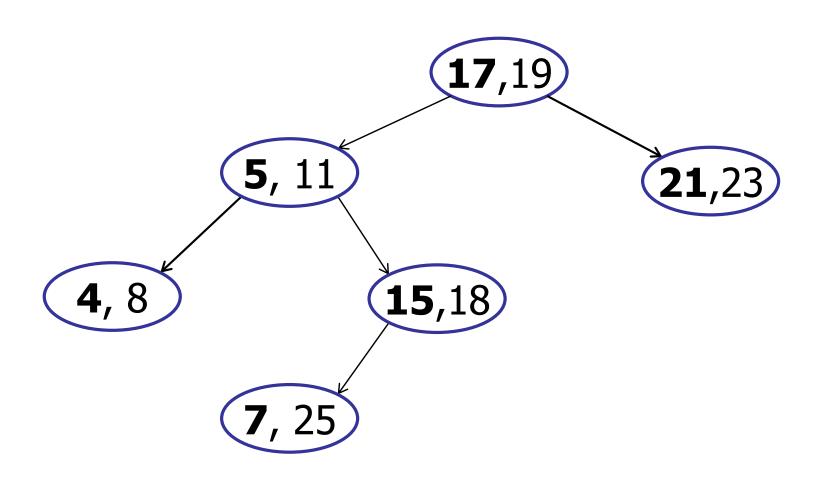
Each node is an interval



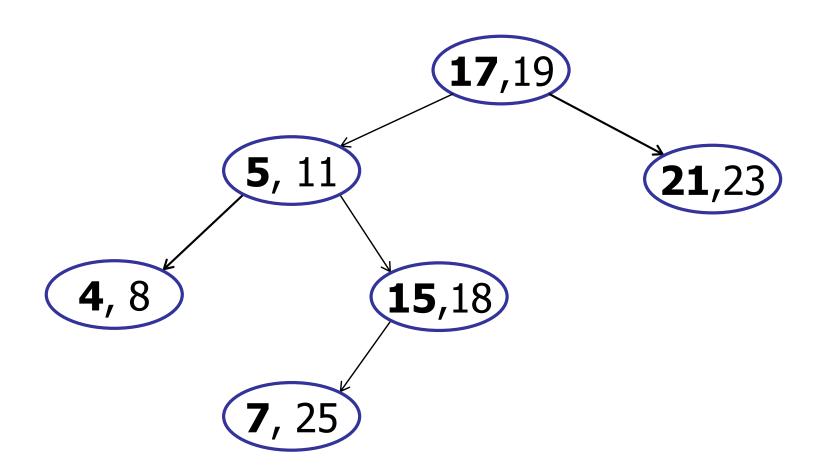
Sorted by left endpoint



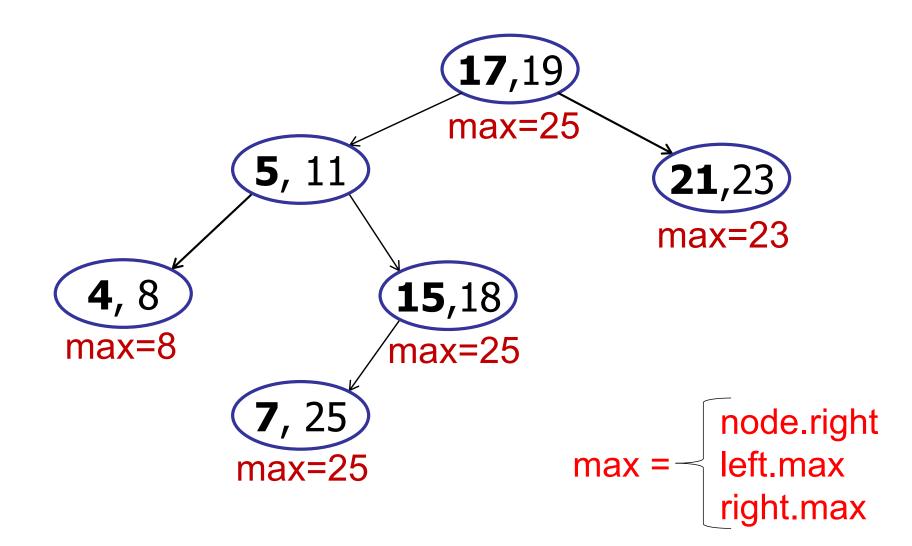
search-interval(25) = ?

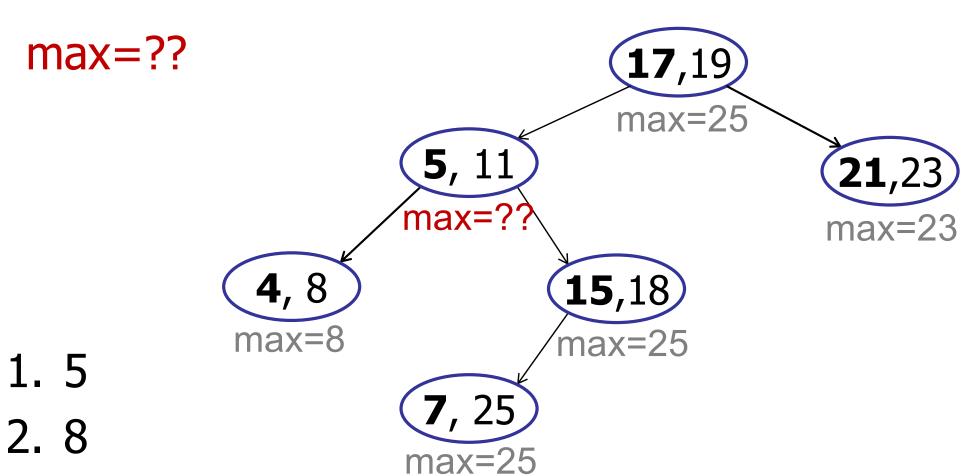


Augment: ??



Augment: maximum endpoint in subtree





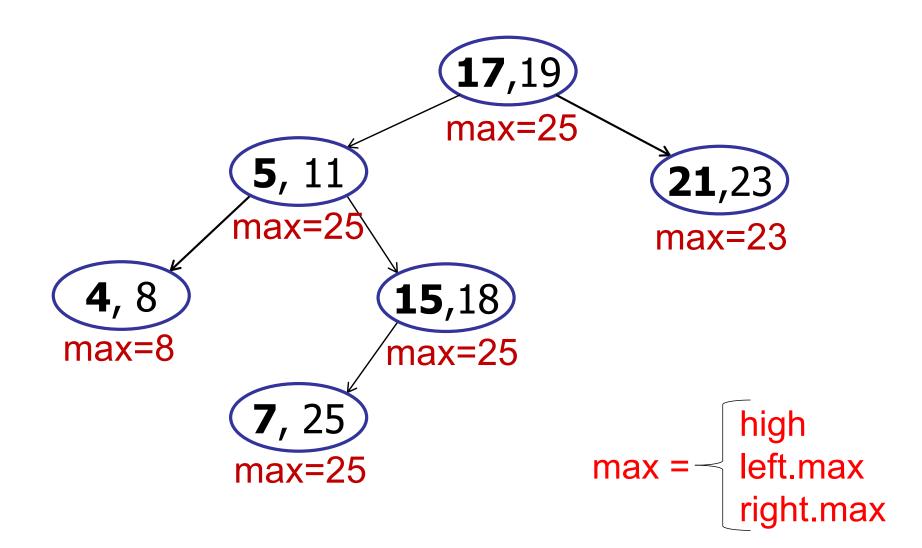
3. 11

4. 18

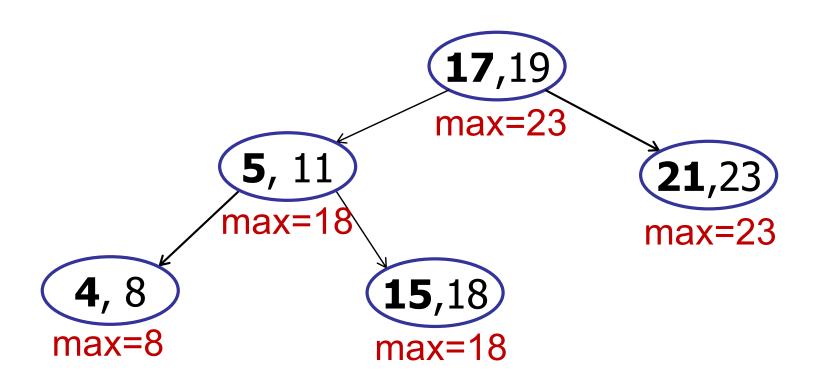
✓5. 25

6. 19

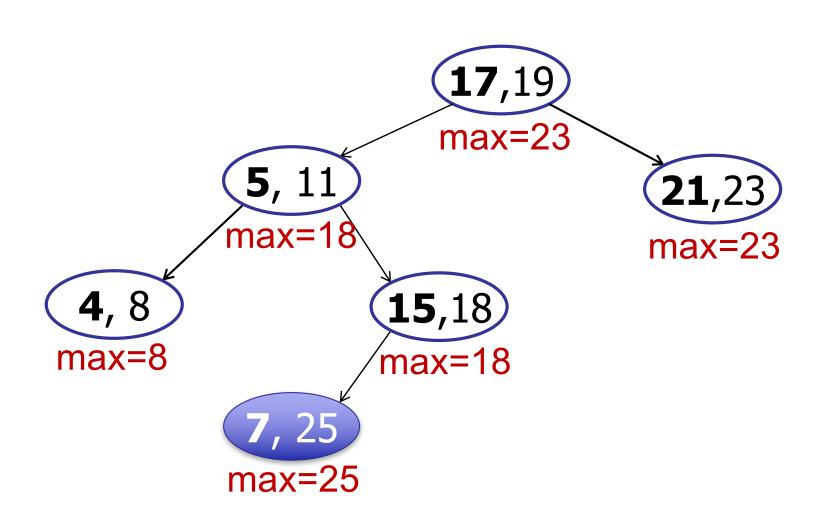
Augment: maximum endpoint in subtree



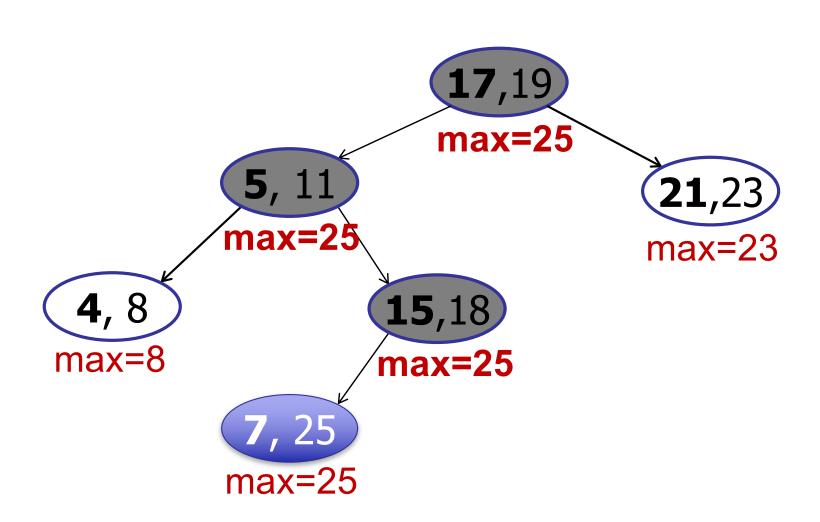
Insertion: example – insert(7, 25)



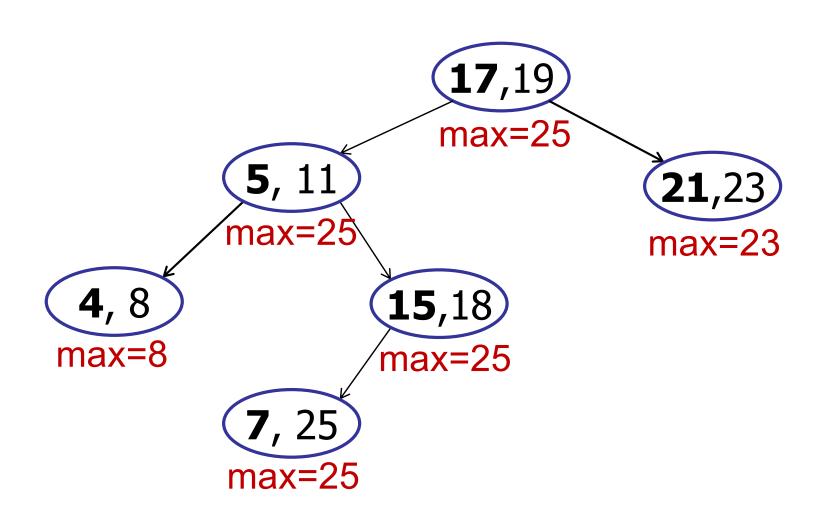
Insertion: example – insert(7, 25)



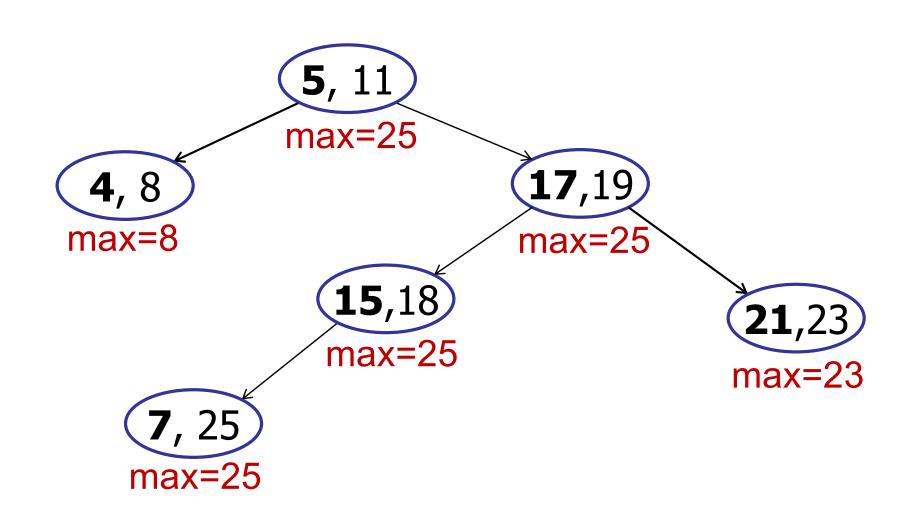
Insertion: example – insert(7, 25)



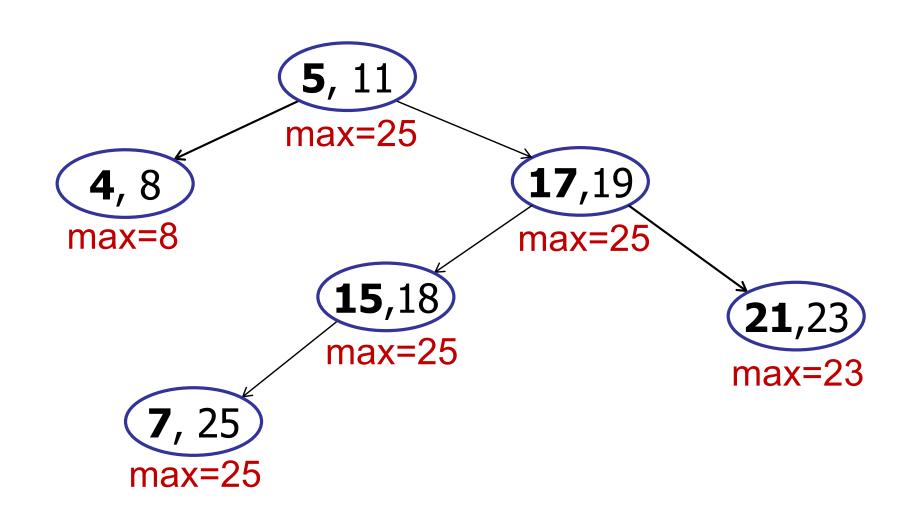
Insertion: out-of-balance



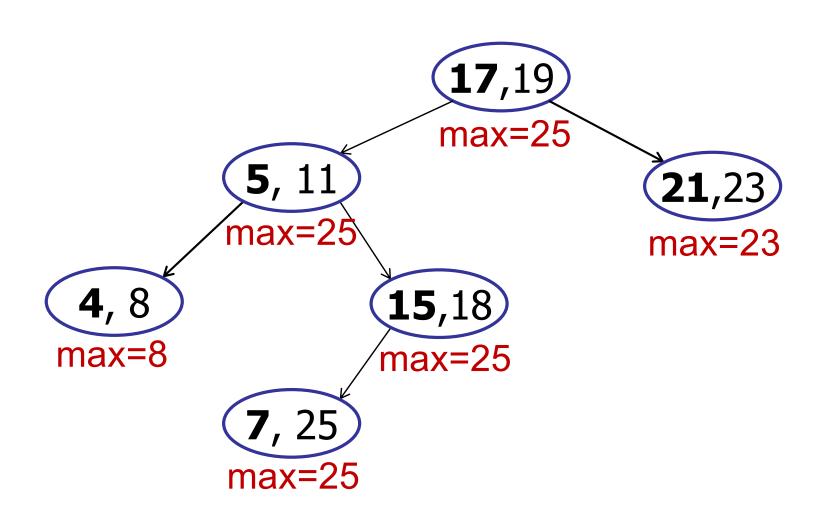
Insertion: right-rotate (17, 19)



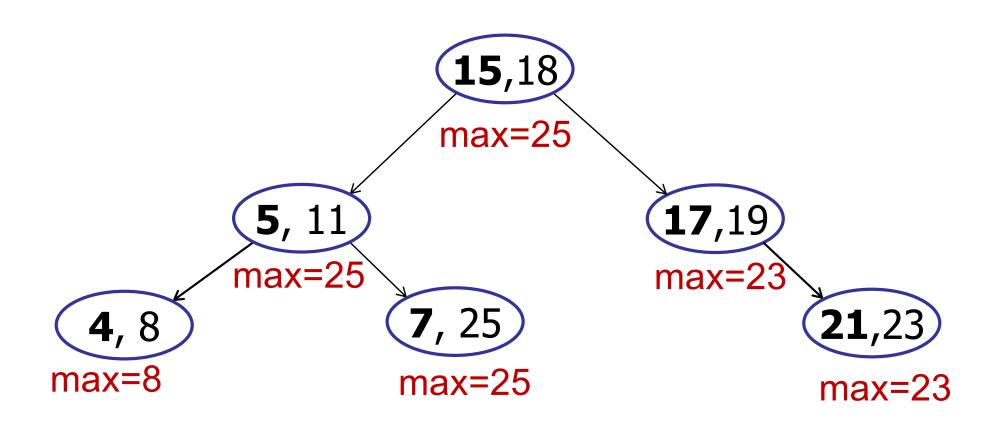
Insertion: right-rotate (17, 19), OOPS!



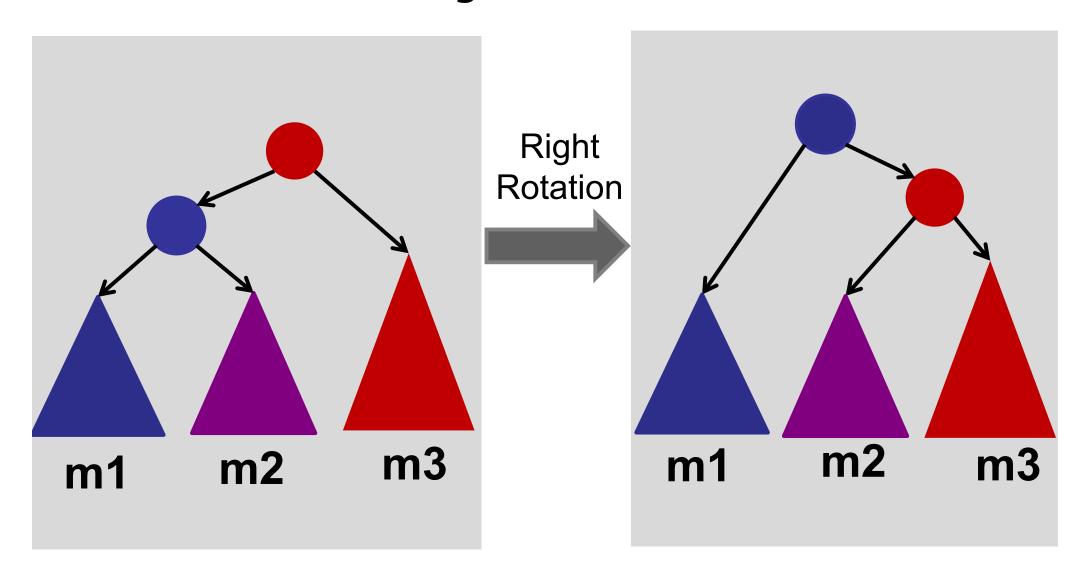
Insertion: out-of-balance



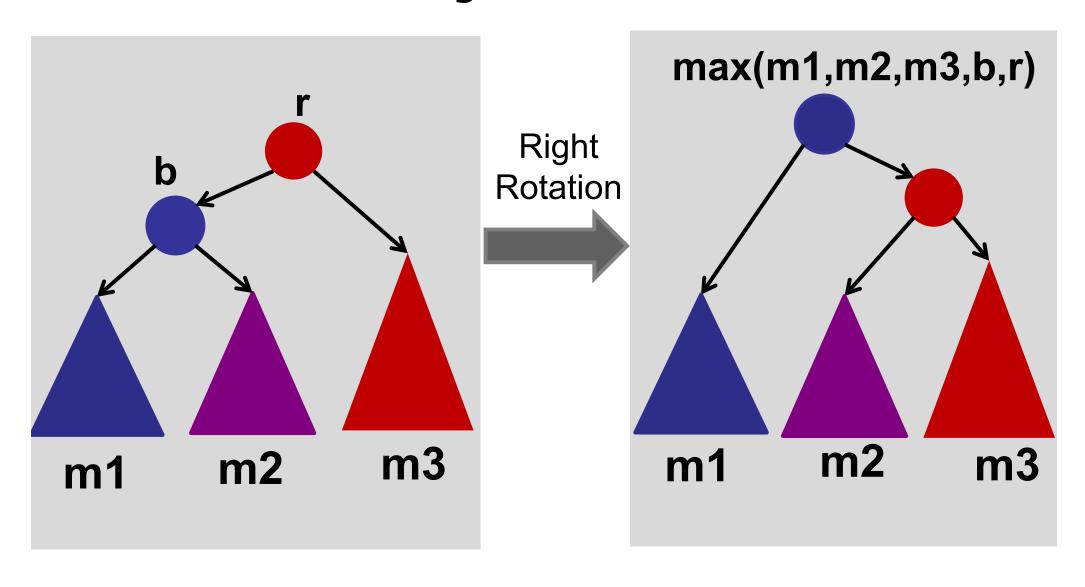
Insertion: left-rotate, right-rotate



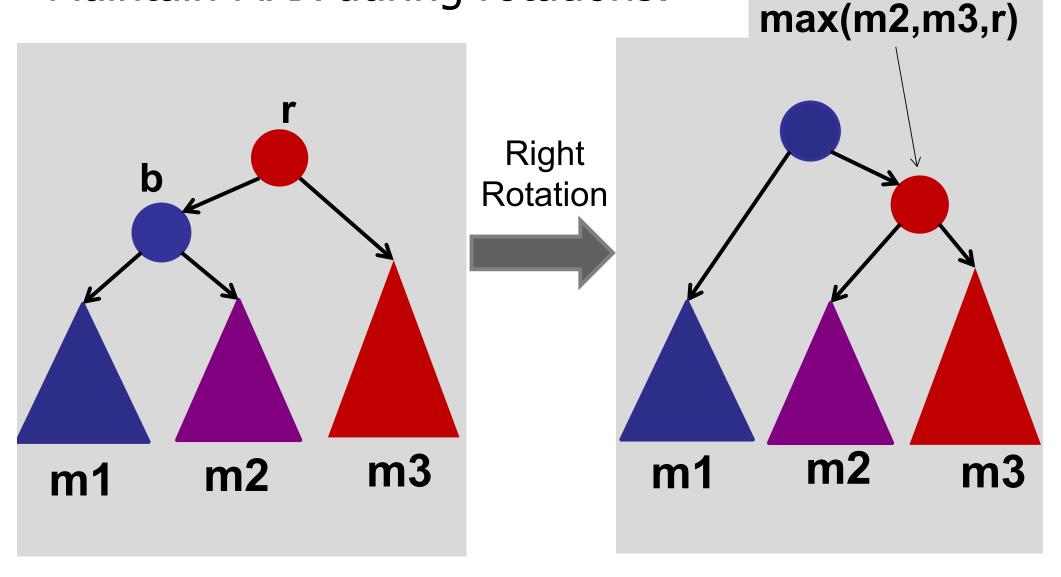
Maintain MAX during rotations:

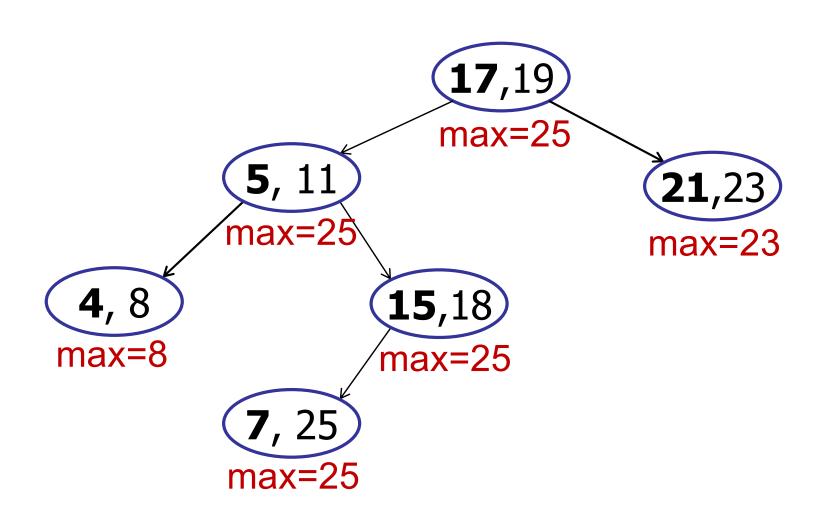


Maintain MAX during rotations:

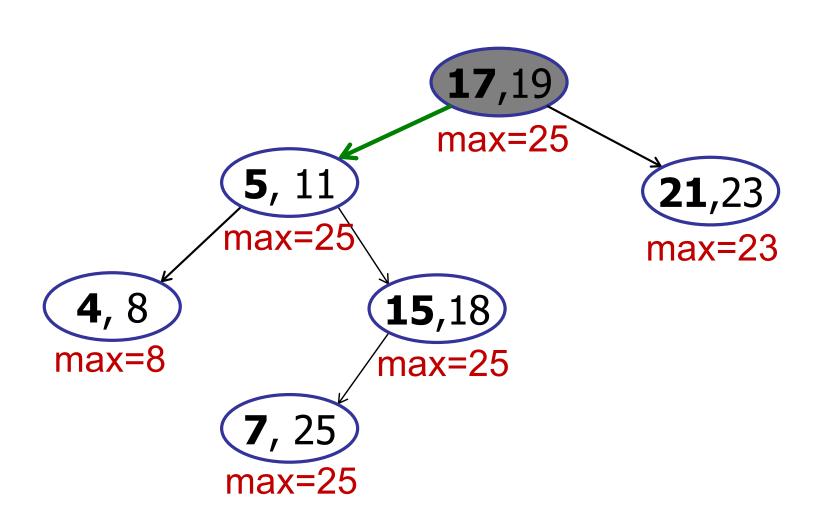


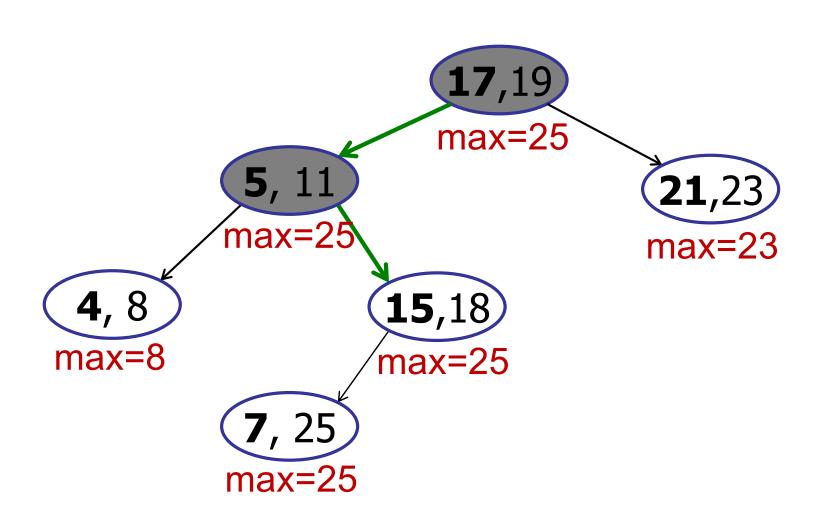
Maintain MAX during rotations:

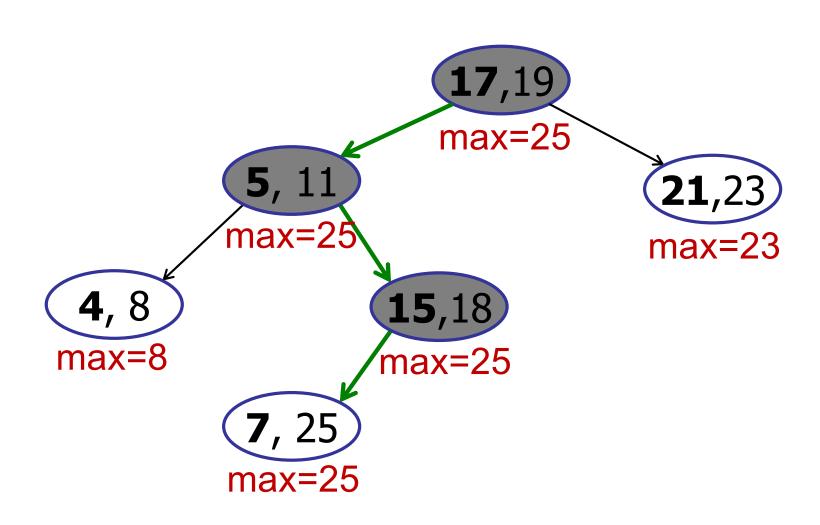


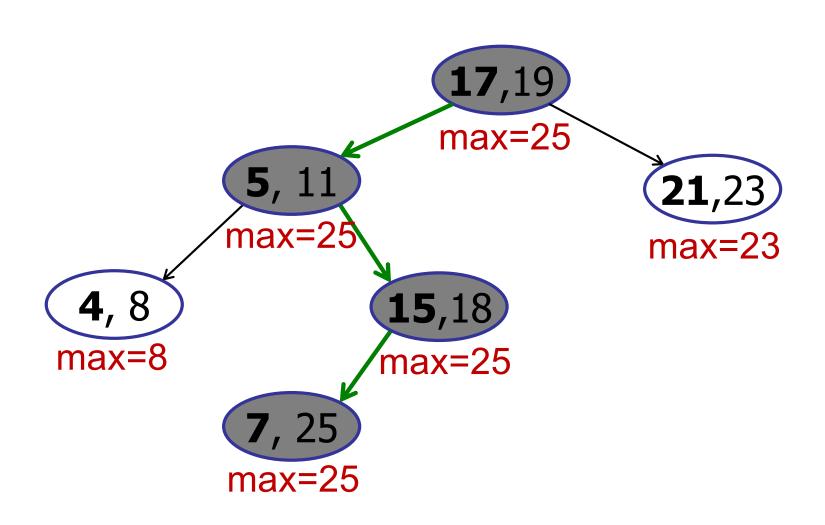


```
interval-search(x): find interval containing x
interval-search(x)
    c = root;
    while (c!= null and x is not in c.interval) do
          if (c.left == null) then
                 c = c.right;
          else if (x > c.left.max) then
                c = c.right;
          else c = c.left;
    return c.interval;
```



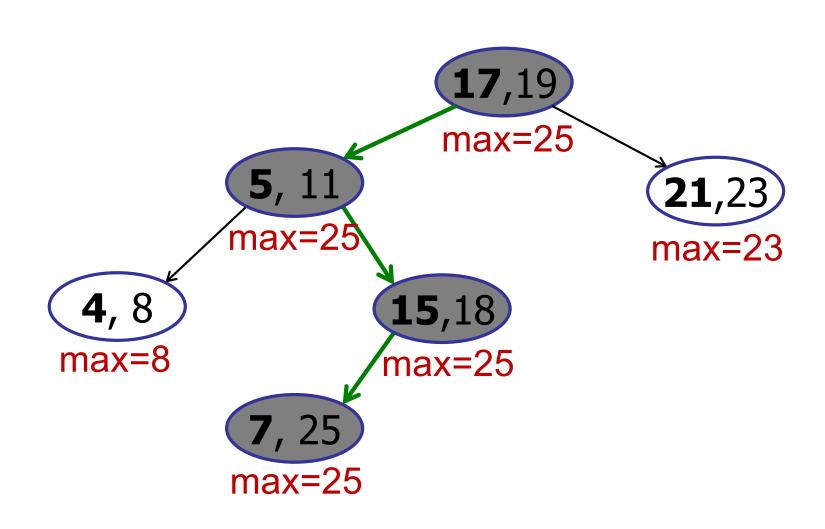




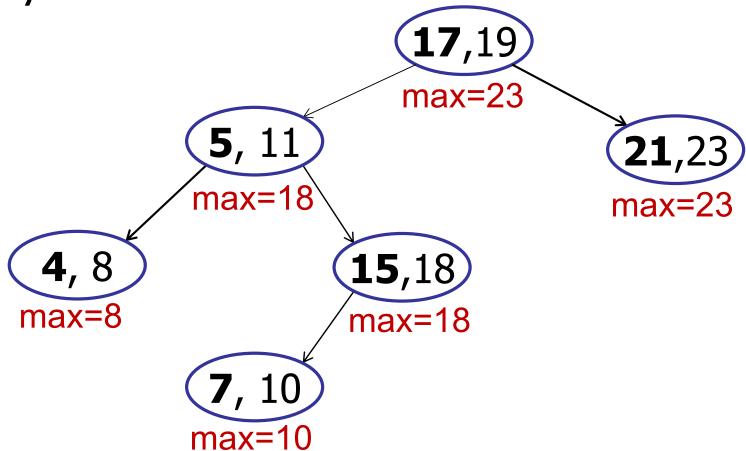


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                c = c.right;
          else c = c.left;
    return c.interval;
```

Will any search find (21, 23)?

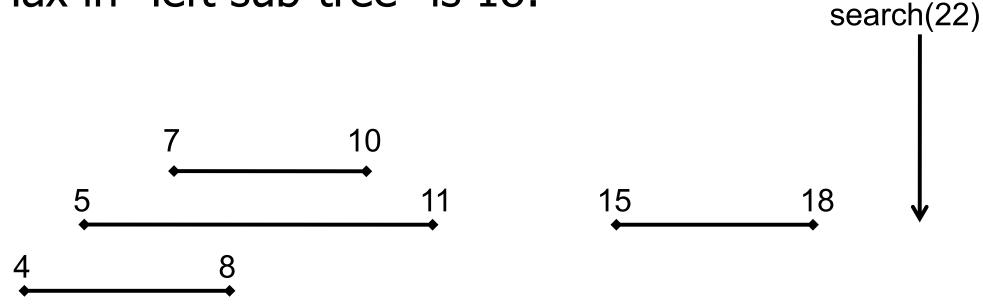


Why does it work?

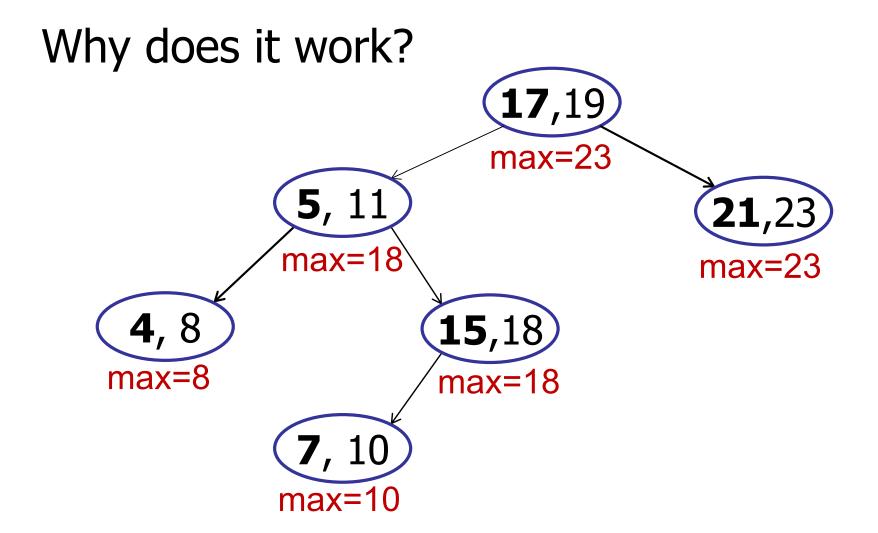


Claim: If search goes right, then no overlap in left subtree.

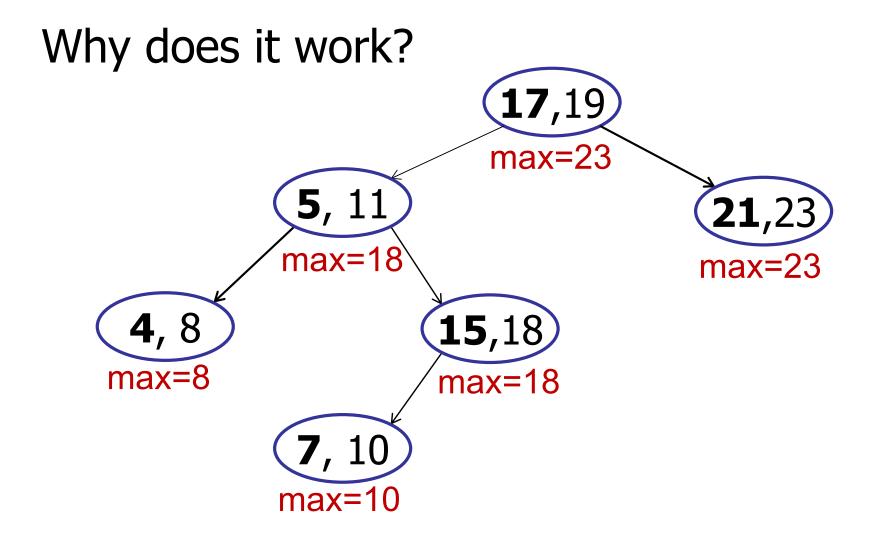
Max in "left sub-tree" is 18:



Safe to go right: 22 is not in the left sub-tree.

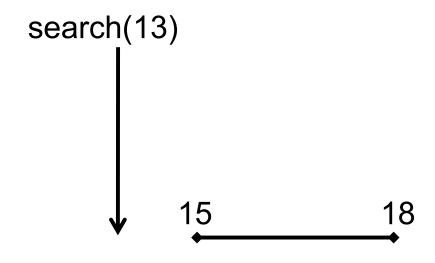


Claim: If search goes left and there is no overlap in the left subtree...



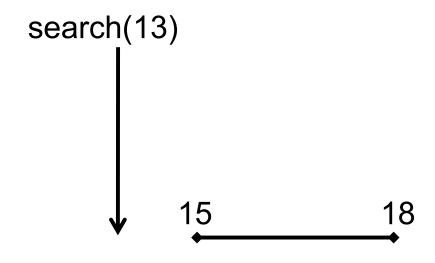
Claim: If search goes left, then safe to go left.

Max in "left sub-tree" is 18:



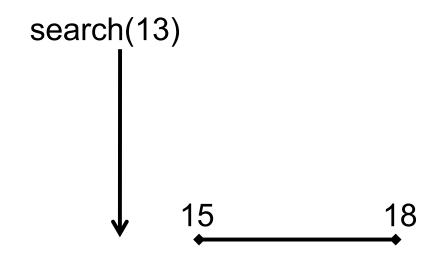
Go left: search(13) < 18

Max in "left sub-tree" is 18:



Go left: search(13) < 15 < 18

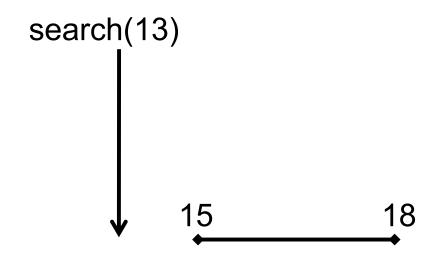
Max in "left sub-tree" is 18:



Go left: search(13) < 15 < 18

Tree sorted by left endpoint.

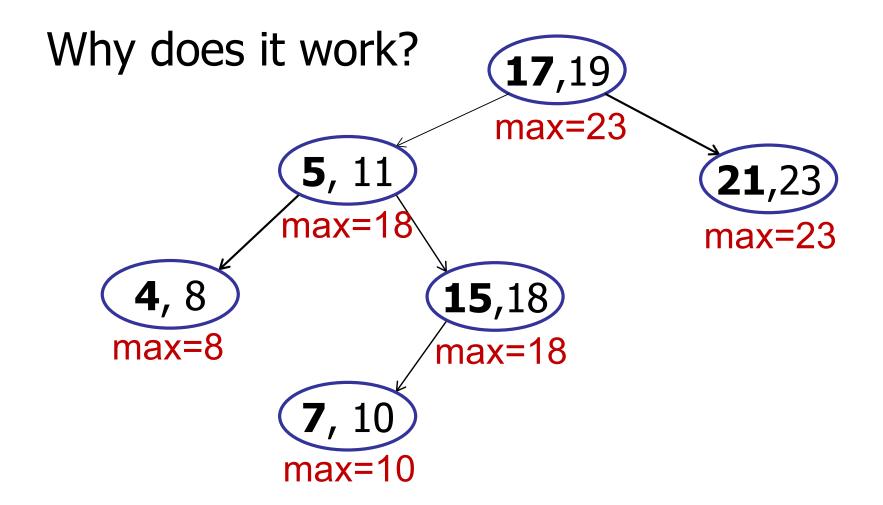
Max in "left sub-tree" is 18:



Go left: search(13) < 15 < 18

Tree sorted by left endpoint.

13 < every interval in right subtree



Claim: If search goes left and no overlap, then key < every interval in right sub-tree.

If search goes right: then no overlap in left subtree.

→ Either search finds key in right subtree or it is not in the tree.

If search goes left: if there is no overlap in left subtree, then there is no overlap in right subtree either.

→ Either search finds key in left subtree or it is not in the tree.

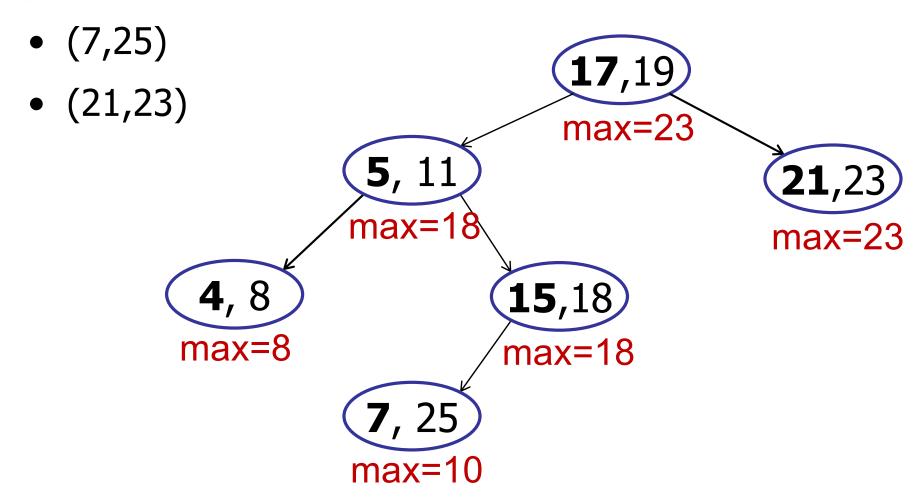
Conclusion: search finds an overlapping interval, if it exists.

The running time of interval-search is:

- 1. O(1)
- 2. O(log n)
- 3. O(n)
- 4. O(n log n)
- 5. $O(n^2)$
- 6. Can't say.

Extension: List all intervals that overlap with point?

E.g.: search(22) returns:



Extension: List all intervals that overlap with point?

All-Overlaps Algorithm:

Repeat until no more intervals:

- -Search for interval.
- -Add to list.
- Delete interval.

Repeat for all intervals on list:

Add interval back to tree.

The running time of All-Overlaps, if there are k overlapping intervals?

- 1. O(1)
- 2. O(k)
- 3. O(k log n)
- 4. O(k + log n)
- 5. O(kn)
- 6. O(kn log n)

Extension: List all intervals that overlap with point?

All-Overlaps Algorithm: O(k log n)

Repeat until no more intervals:

- -Search for interval.
- -Add to list.
- Delete interval.

Repeat for all intervals on list:

Add interval back to tree.

Best known solution: O(k + log n)

Today

Three examples of augmenting BSTs

1. Order Statistics

2. Intervals

3. Orthogonal Range Searching