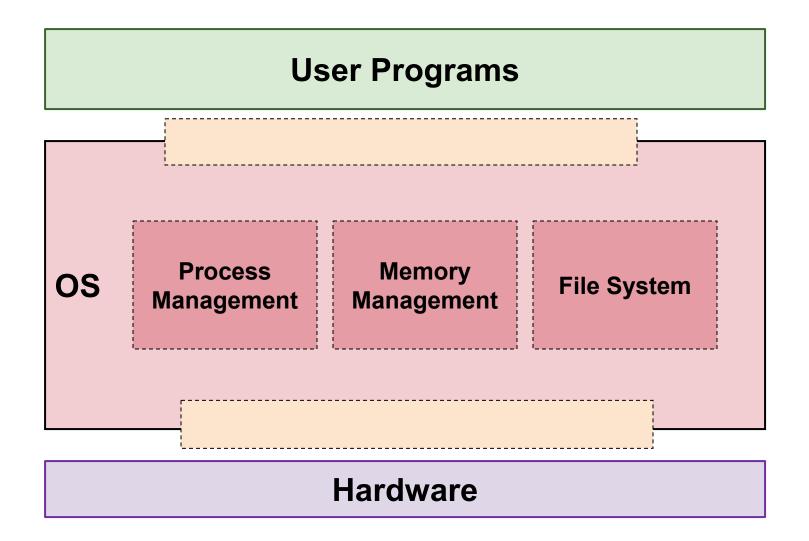
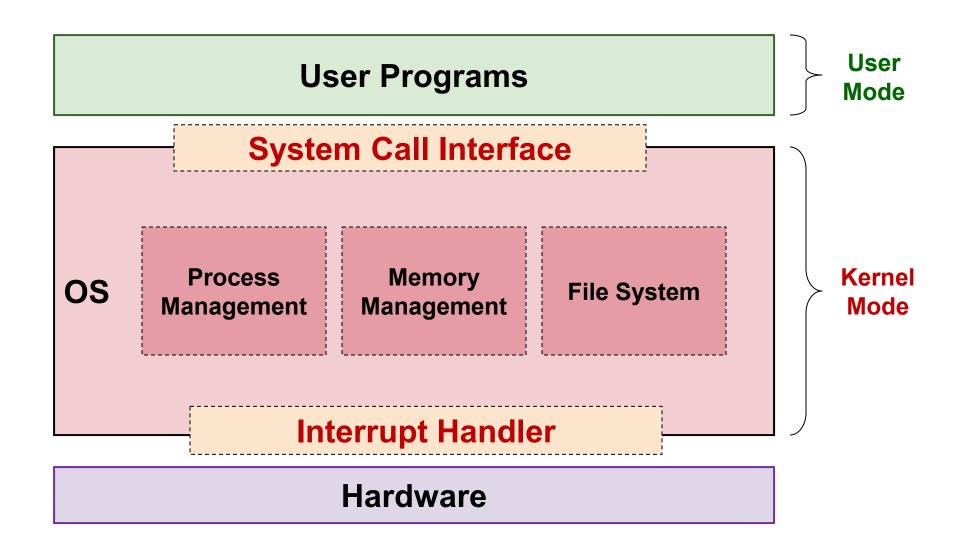
Big Bang & Revision

Lecture 12 (Live Version)

OS Structure (Monolithic)



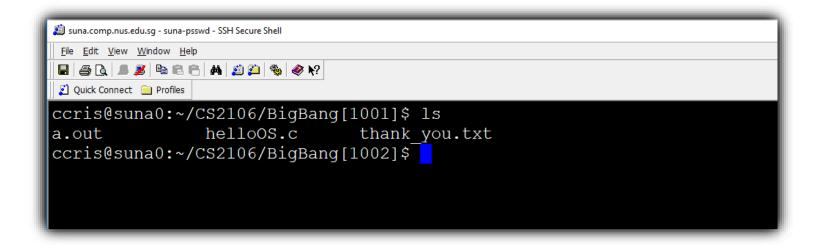
OS Structure (Monolithic)



23 + 11 + lots of hours just to learn about what happened in 1 second

THE STORY OF SHELL INTERPRETER

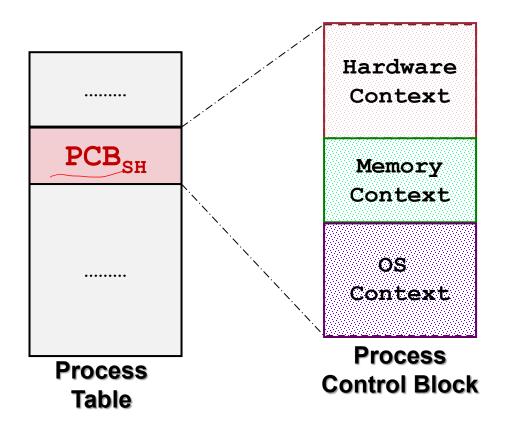
Ah.... Ah.... so simple..... or is it?



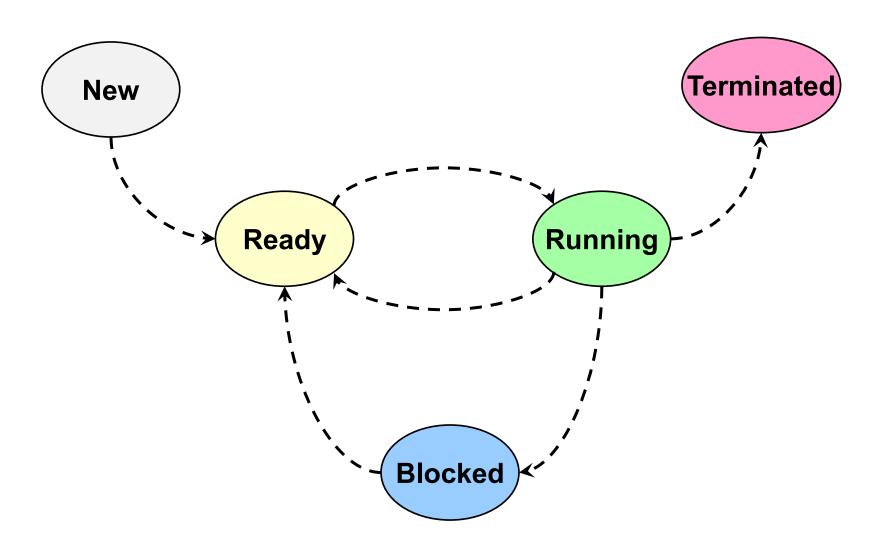
Here we go



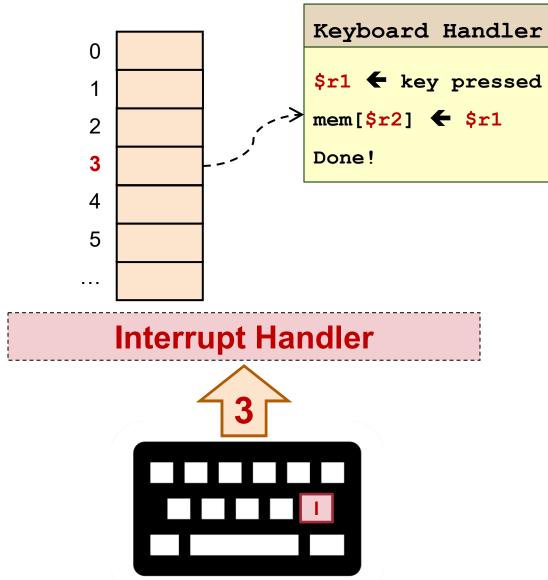
Process Table & Process Control Block

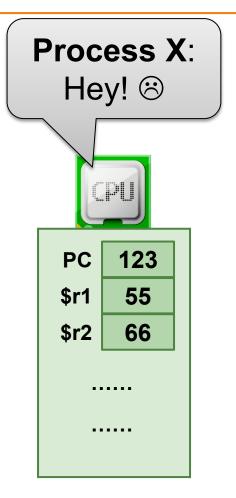


The Interpreter is in.....

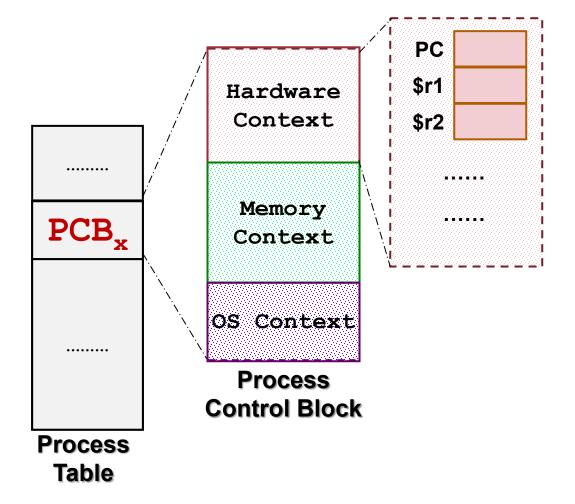


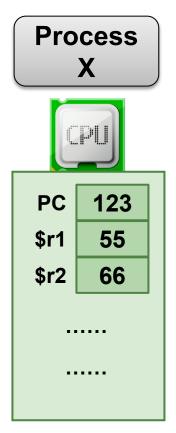
User press "1"

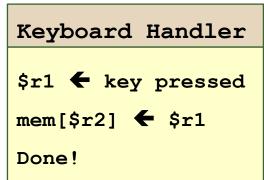




Sorry to **interrupt** you....



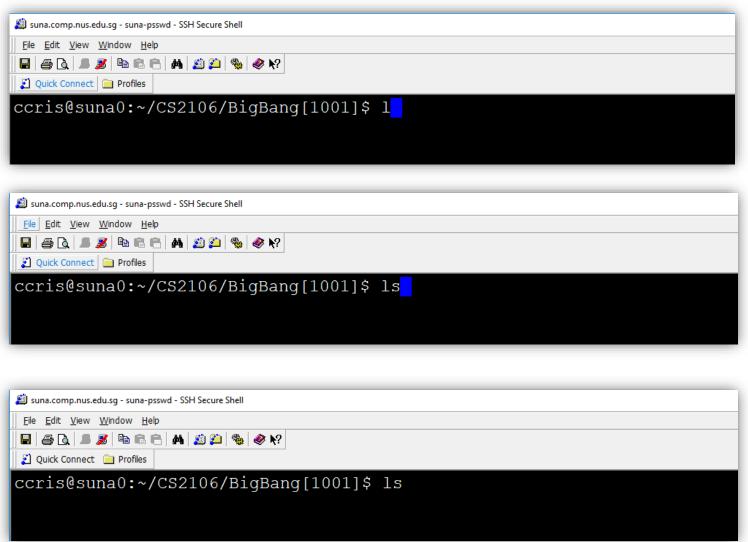




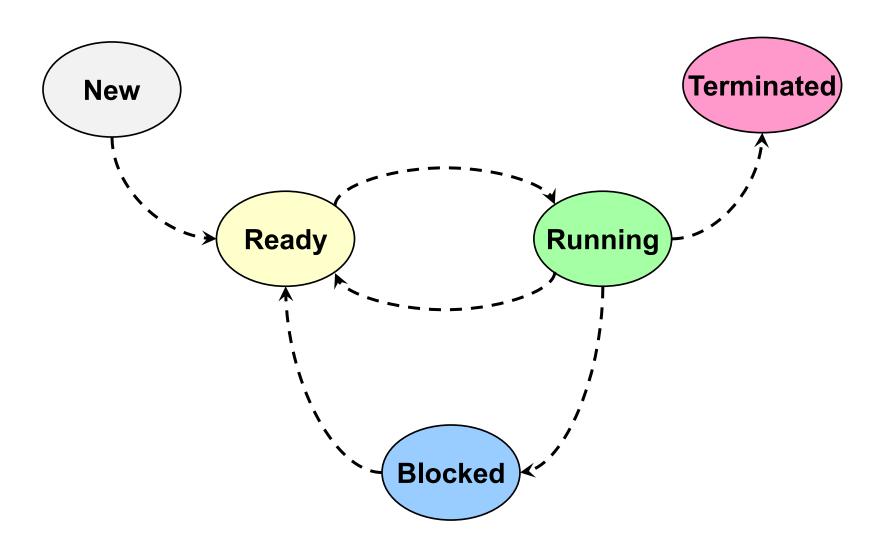
Interrupt steps

- Give the sequence of steps for handling an interrupt.
 - Interrupt occurs
 - Save registers/CPU state
 - Perform the handler routine
 - Restore registers/CPU state
 - Return from interrupt

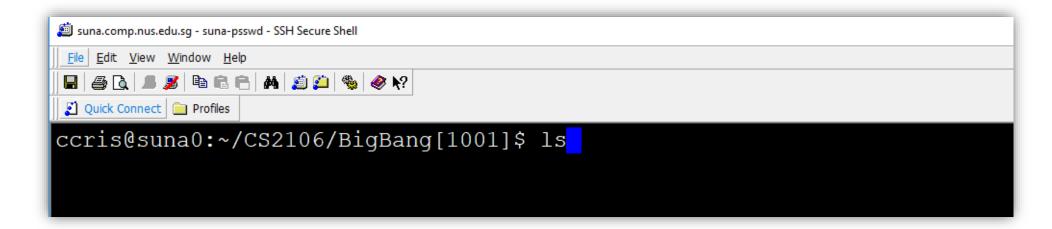
Rinse and Repeat.....



The interpreter is now...



User entered "ls", the interpreter will...



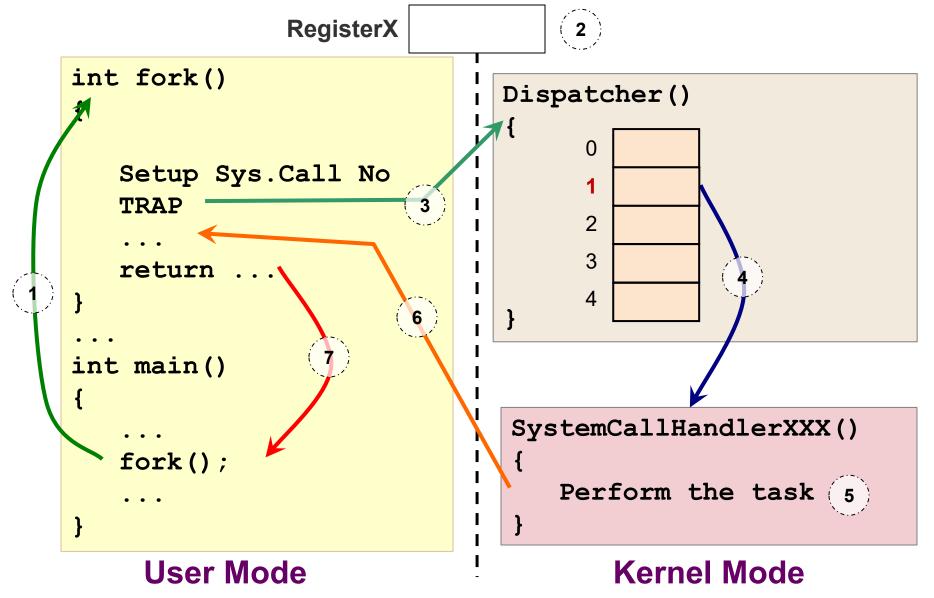
Typical Steps for Shell Interpreter

```
UserCmd ← read from keyboard
fork()
if I am the parent (i.e. the shell)
   wait (child to finish)
else
    exec ( UserCmd )
```

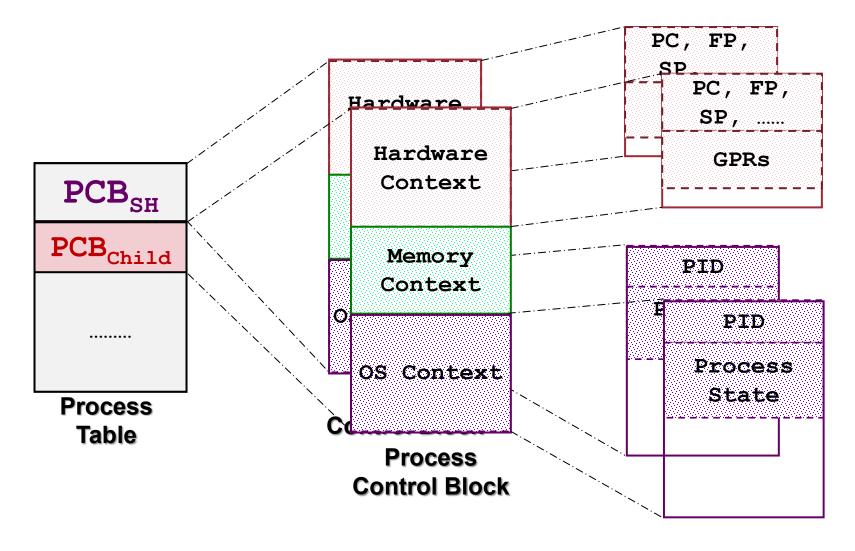
fork() involves a system call

```
RegisterX
int fork()
                            Dispatcher()
int main()
                             SystemCallHandlerXXX()
   fork();
                                   Kernel Mode
    User Mode
```

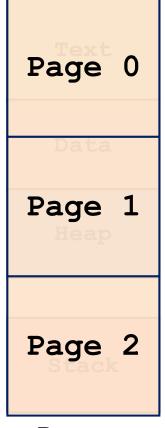
fork() involves a system call



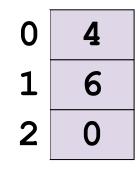
What is the **effect of fork**()?



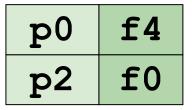
Memory Space of a Process



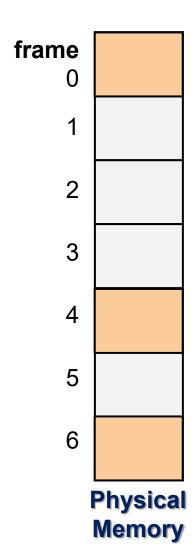
Process Memory Space



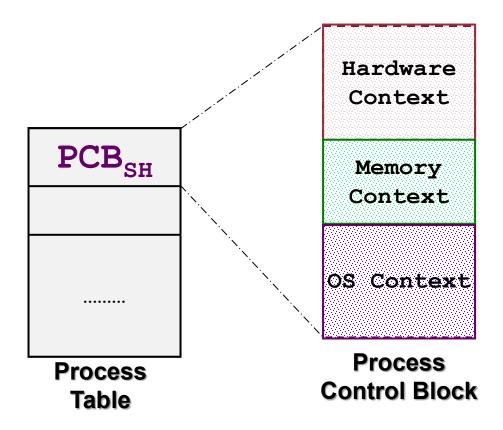
Page Table



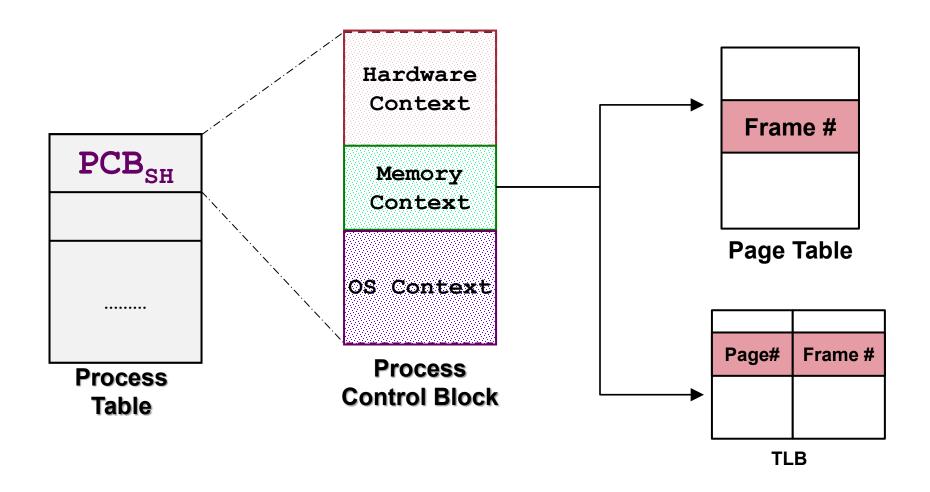
TLB

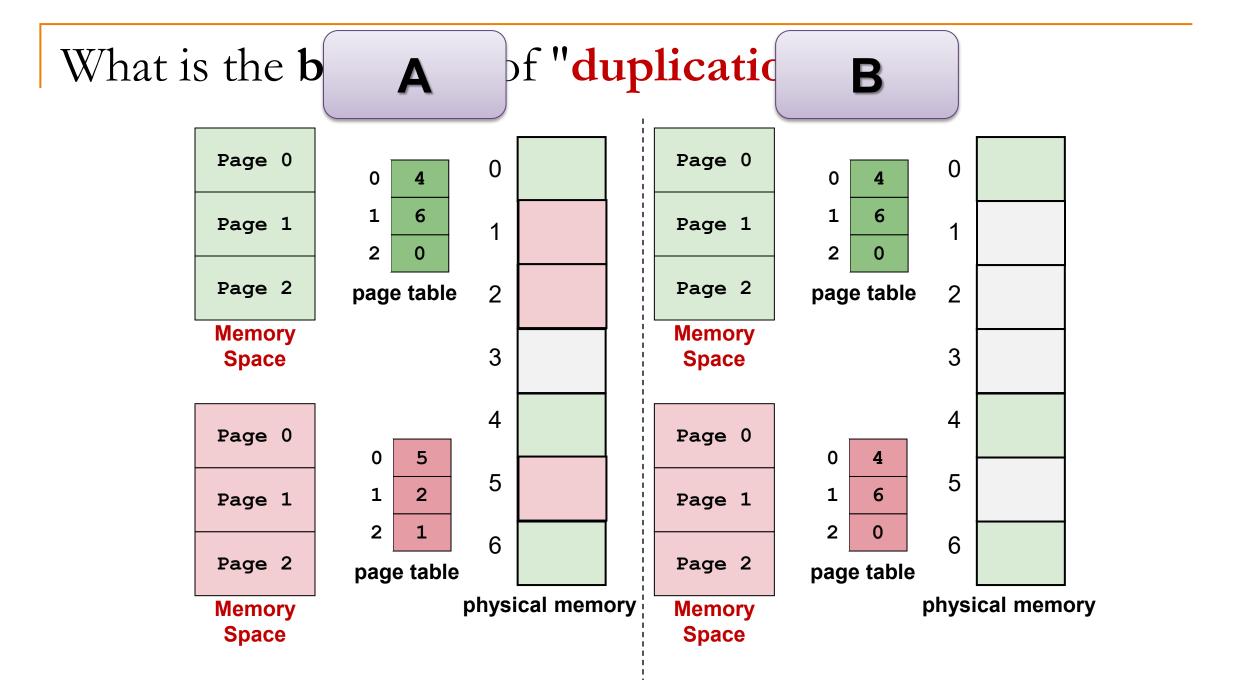


Memory Context = ?



Memory Context = ?





Duplicating Memory Space the HARD WAY



Page 1

Page 2

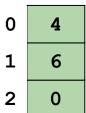
Memory Space

Page 0

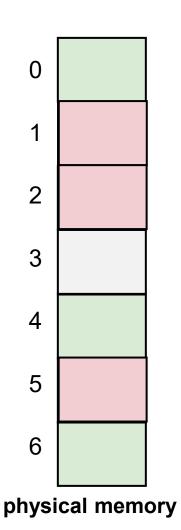
Page 1

Page 2

Memory Space



page table



Copy on Write

Page 0

Page 1

Page 2

Memory Space

Page 0

Page 1

Page 2

Memory Space 0

1

6

2 0

page table

0

4

6

2

page table

0

1

2

3

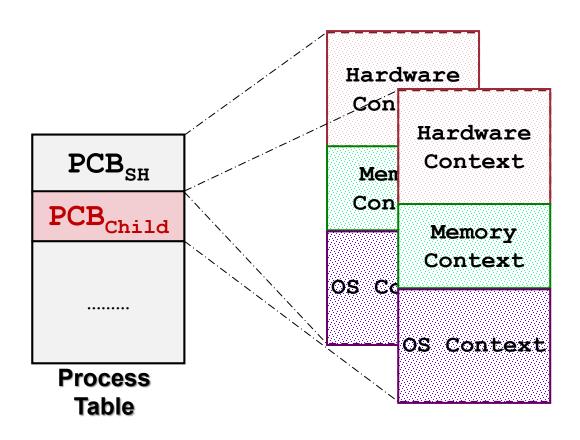
4

5

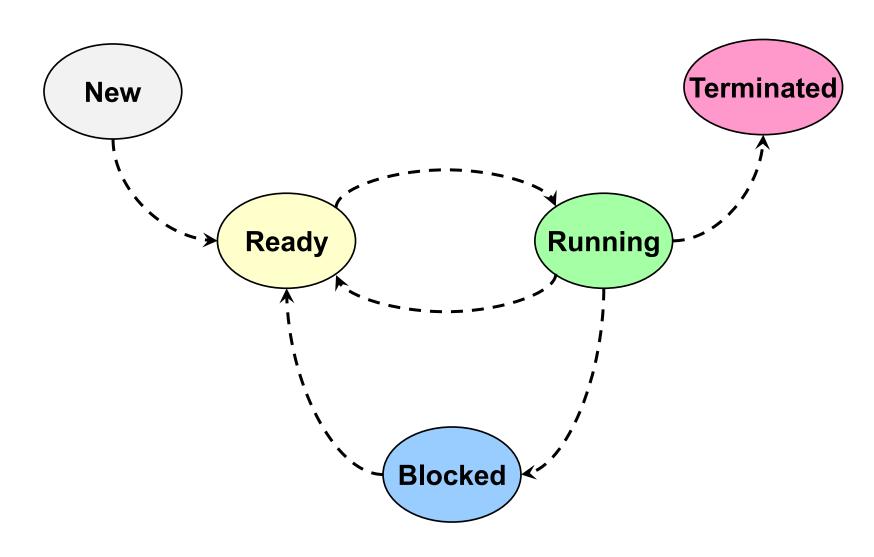
6

physical memory

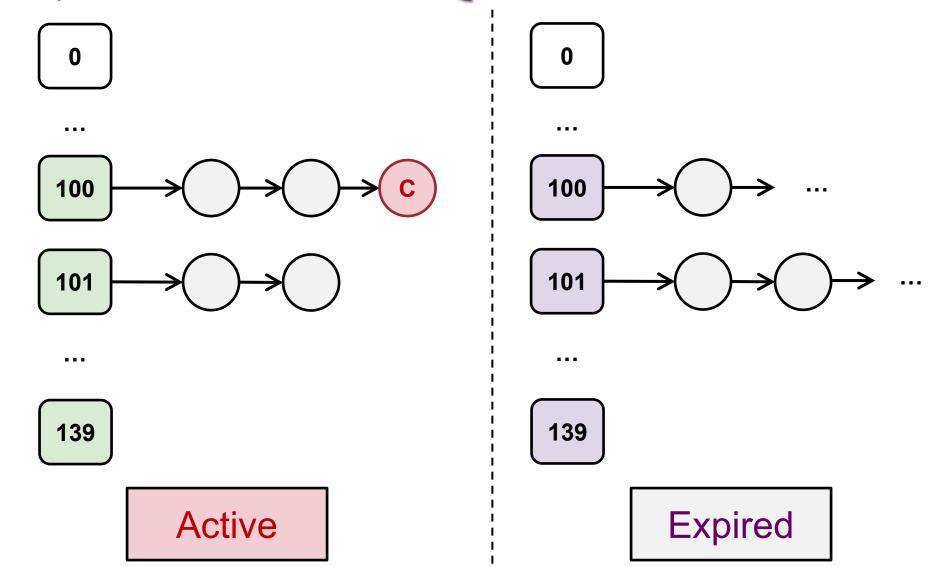
So, Effect of **fork**()



Child, welcome to the world!



Child, welcome to the Queue!



fork() finishes and returns

```
RegisterX
int fork()
   Setup Sys.Call No
   TRAP
   return ...
int main()
   fork();
```

User Mode

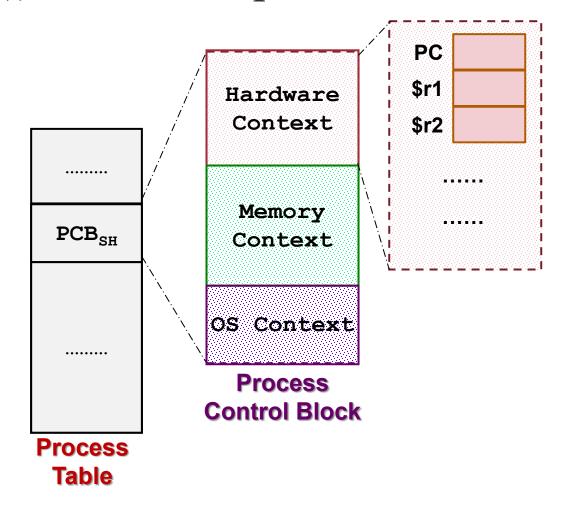
```
SystemCallHandlerXXX()
{
    Perform the task
}
```

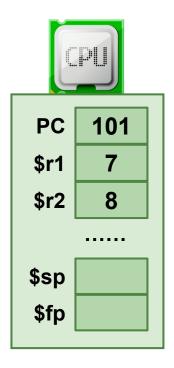
Kernel Mode

Typical steps for Shell Interpreter

```
UserCmd ← Read from keyboard
fork()
if I am the parent (i.e. the shell)
   wait ( child to finish )
else
    exec( UserCmd )
```

wait(): The interpreter will....

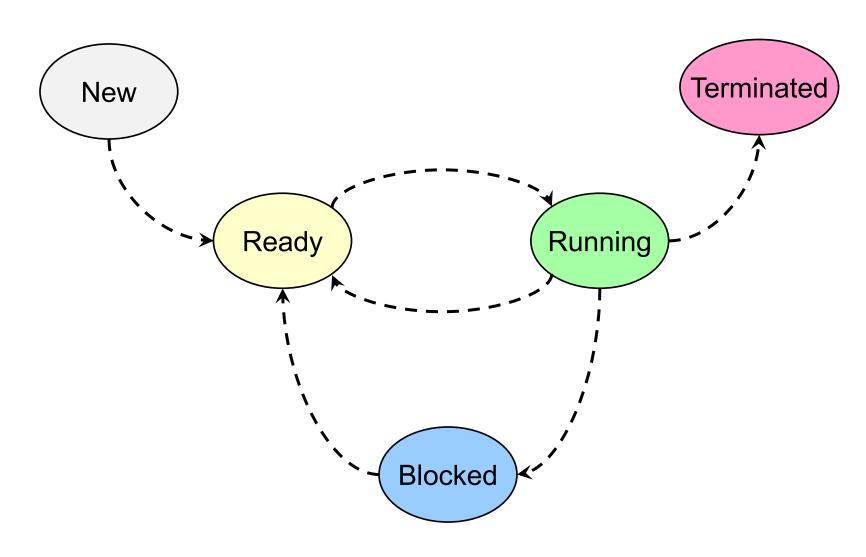




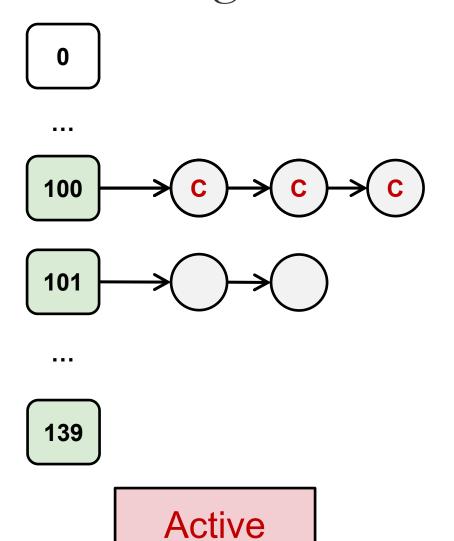
p0	f4
p2	f0

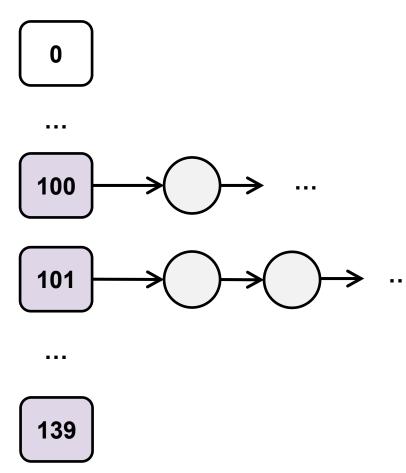
TLB

The interpreter is now...



Hmm... who gets to run?

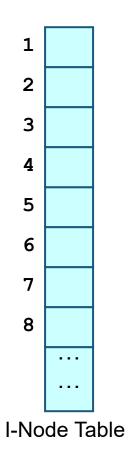




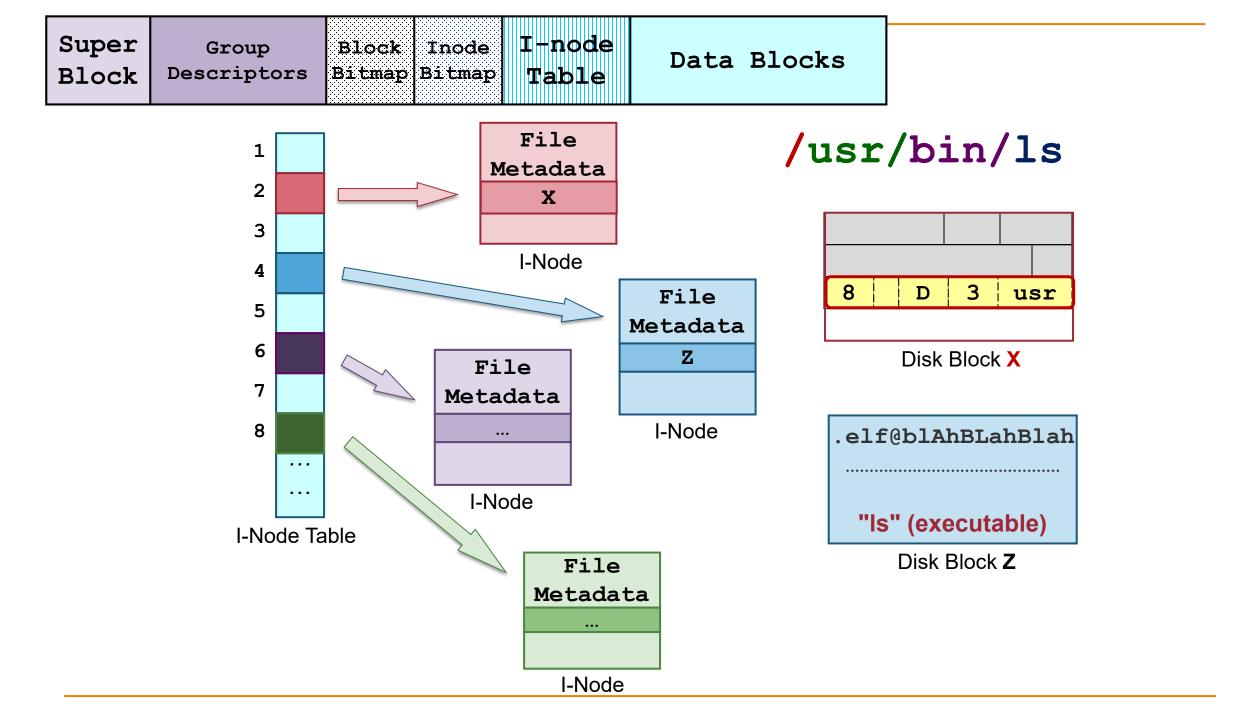
Expired

Typical steps for Shell Interpreter

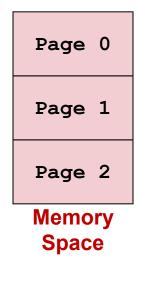
```
UserCmd ← Read from keyboard
fork()
if I am the parent (i.e. the shell)
   wait (child to finish)
else
   exec( UserCmd )
```

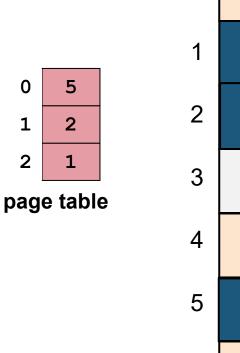


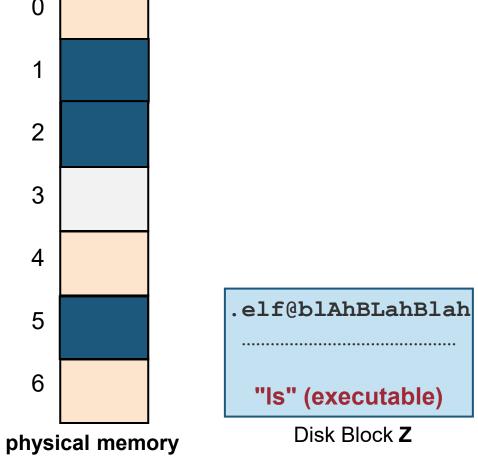
/usr/bin/ls



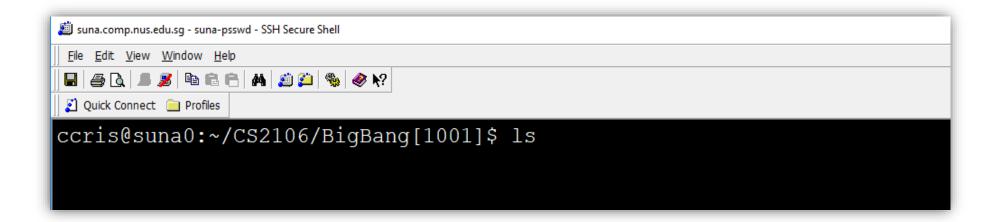
Memory Content Replaced



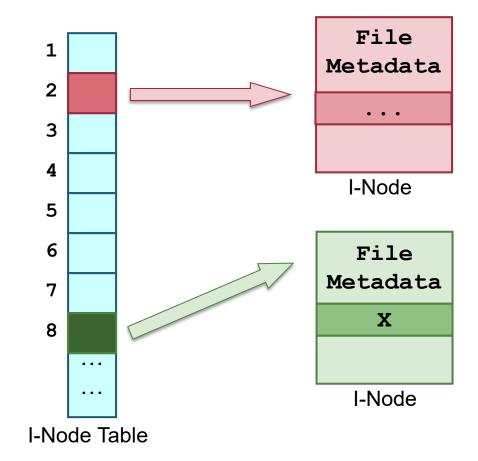




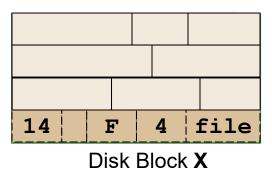
Child is now "ls", what next?



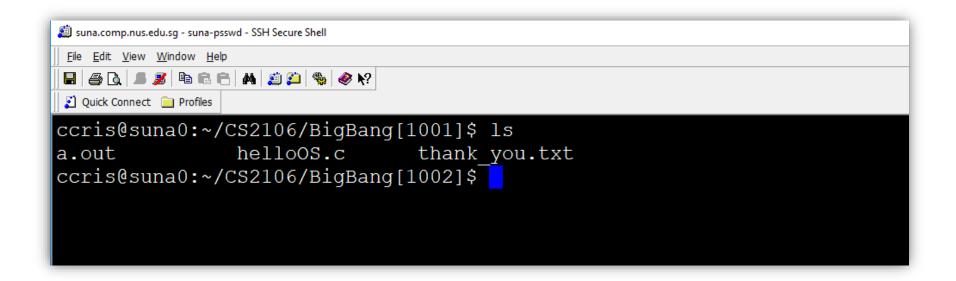
Listing /.../BigBang



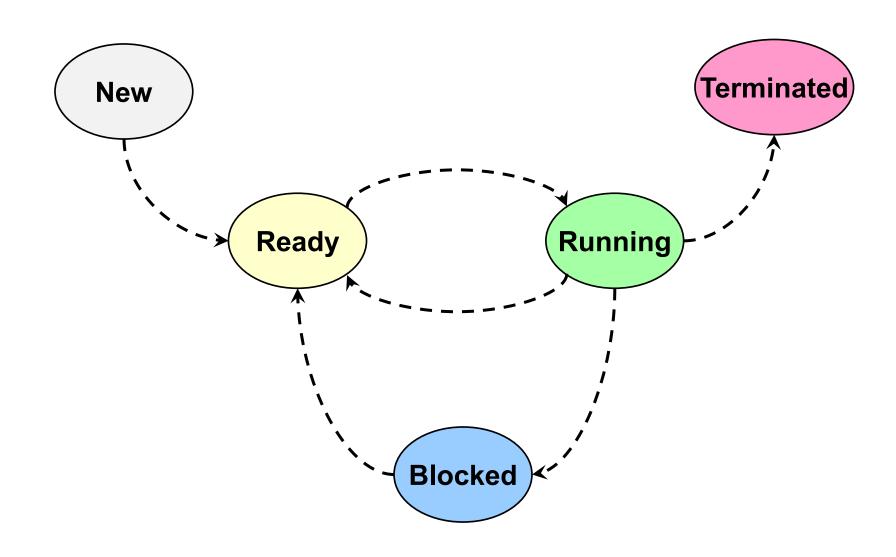
Listing /.../BigBang



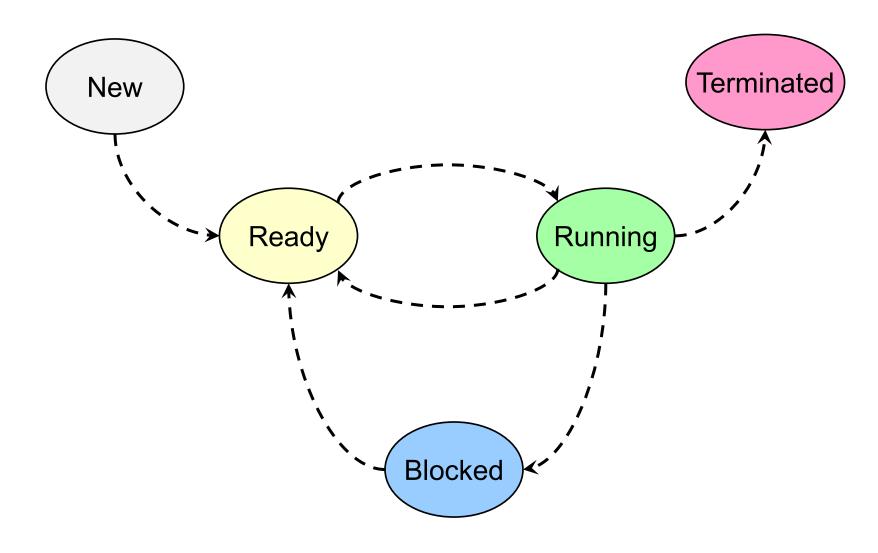
"ls" prints the directory content



Child exits



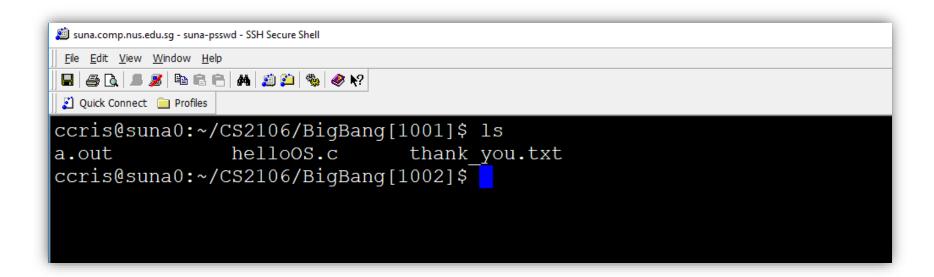
The interpreter hears about it....



Typical steps for Shell Interpreter

```
UserCmd ← Read from keyboard
fork()
if I am the parent (i.e. the shell)
   wait (child to finish)
else
   exec( UserCmd )
```

Woohoo!!



WE SHOULD SHARE!

How to **Share Memory?**

Page 0

Page 1

Page 2

Memory Space

Page 0

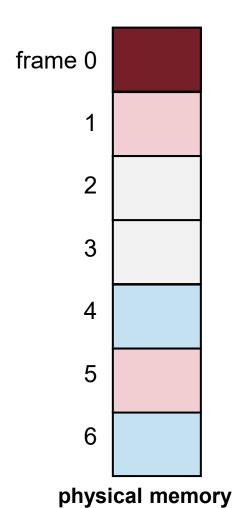
Page 1

Page 2

Memory Space

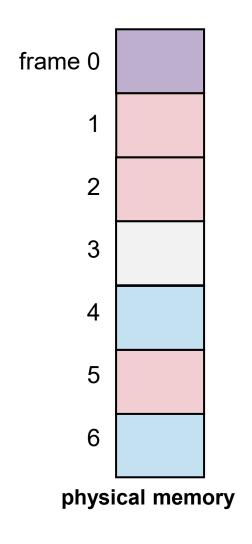
page table

page table



$$(X = X + 1; sA = sA + 1;)$$

$$sA = sA + 1;$$



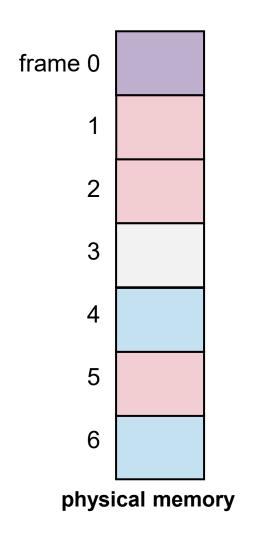
$$X = X + 1;$$

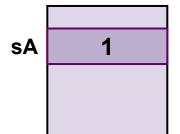
 $sA = sA + 1;$

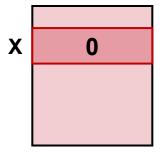
$$X = X + 1;$$

 $sA = sA + 1;$

Semaphore to the rescue!







$$X = X + 1;$$

 $sA = sA + 1;$

$$X = X + 1;$$

 $sA = sA + 1;$

Concurrency

Race conditions

Critical Section

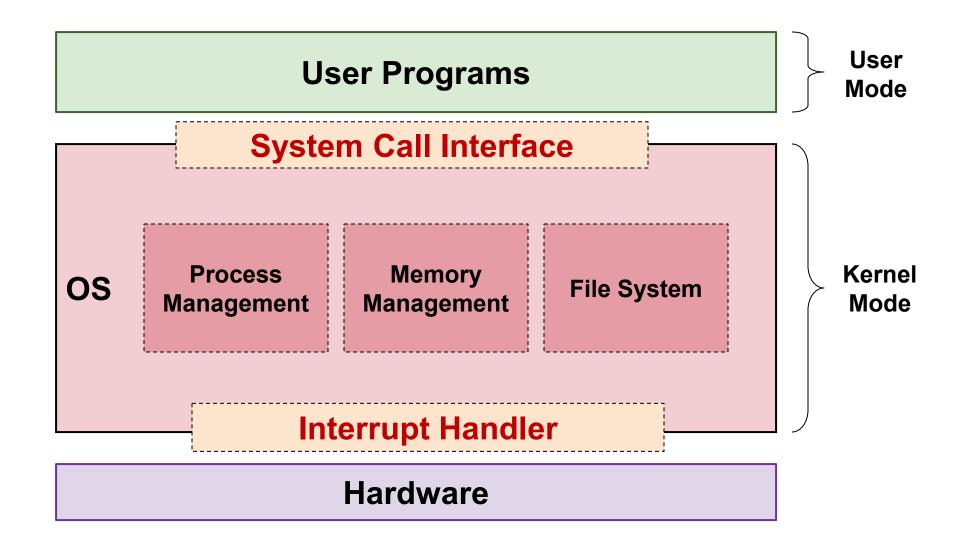
Semaphore

Classical synchronization problems

Phew..... "quick summary" now

WHAT HAVE WE LEARNED?

Operating System



WHAT ELSE HAVE WE LEARNED?

Side Benefits....

Design of complex system

Abstraction and Interface

Resource Management

Performance Trade Off (time vs space)

What's next?

System Security

Parallel Computing/ Concurrent Programming

Computer Architecture

Compilers

OH... THE EXAM ©

The plan...

- Like the midterm
 - F2F in MPSH1 an MPSH2
 - Open book with printed materials
- Backup in place:
 - LumiNUS quiz
 - Zoom proctoring
 - Record your screen
 - Refer to PDF materials
- Email us early to book a consultation slot

Important to know

- Rough percentage of coverage
 - □ Lecture 1 to Lecture $5 = \sim 25\%$
 - □ Lecture 6 to Lecture $11 = \sim 75\%$
- MCQ questions
- Short questions
 - Write short answers

Open book

It's Over!

Goodbye! Say Hi if you see us in school!