CS566HW7-2024

October 23, 2024

0.1 Lab 7. Introduction to algorithms

This is seventh homework for CS 566.

0.2 Task 1. Solve the problem "Insert Into BST" from https://leetcode.com/problems/insert-into-a-binary-search-tree/description/ using Python3.

Use the box below, to paste the working code. The format of the code should be identical to LeetCode platform. (4 points)

```
[1]: from collections import deque
     from typing import Optional, List
     # Definition for a binary tree node.
     class TreeNode:
         def __init__(self, val=0, left=None, right=None):
             self.val = val
             self.left = left
             self.right = right
     class Solution:
         def insertIntoBST(self, root: Optional[TreeNode], val: int) ->__
      →Optional[TreeNode]:
             if not root:
                 return TreeNode(val)
             if val>root.val:
                 root.right = self.insertIntoBST(root.right,val)
             elif val<root.val:</pre>
                 root.left = self.insertIntoBST(root.left,val)
             return root
```

0.2.1 Do not modify the testing code below. If you get message "Mistake in test case #", it means that you algorithm is incorrect.

```
[2]: # do not modify testing code
def preorder_step(node, result):
    if not node:
        return
```

```
result.append(node.val)
    preorder_step(node.left, result)
    preorder_step(node.right, result)
def preorderTraversal(root: Optional[TreeNode]) -> List[int]:
    ls = []
    preorder_step(root, ls)
    return 1s
#test case 1
root = TreeNode(4, TreeNode(2, TreeNode(1), TreeNode(3)), TreeNode(7))
expected = [TreeNode(4, TreeNode(2, TreeNode(1), TreeNode(3)), TreeNode(7, __
 →TreeNode(5))),
            TreeNode(5, TreeNode(2, TreeNode(1), TreeNode(3)), TreeNode(7, __
⊸TreeNode(4)))]
actual = Solution().insertIntoBST(root, val)
pr_actual = preorderTraversal(actual)
pr_expected = [preorderTraversal(tree) for tree in expected]
assert pr_actual in pr_expected, "Mistake in test case 1"
print("OK")
```

OK

0.2.2 Write analysis of the Memory Complexity and Time Complexity using Aymptotic Notation O. (1 point)

Memory Analysis: O(h), where h is height of tree

Time Analysis: O(h), where h is height of tree

Use the box below, to paste the working code. The format of the code should be identical to LeetCode platform. (4 points)

```
[3]: from typing import Optional

class TreeNode:
    def __init__(self, val=0, left=None, right=None):
        self.val = val
        self.left = left
        self.right = right

class Solution:
```

```
def deleteNode(self, root: TreeNode, key: int) -> TreeNode:
       # base
       if not root:
           return root
       # search node to delete
       if key>root.val:
           root.right = self.deleteNode(root.right,key)
       elif key<root.val:</pre>
           root.left = self.deleteNode(root.left,key)
       # found node to delete
       else:
           # for O child or 1 right child, return right child
           if root.left is None:
               return root.right
           # for 1 left child, return left child
           elif root.right is None:
               return root.left
           # two children, find min in right subtree
           else:
               # go to leftmost node in rightsubtree
               curr = root.right
               while curr.left:
                   curr = curr.left
               # swap val with delete node val
               root.val = curr.val
               # update rightsubtree recursively, where we delete the swapped_
⇔value
               root.right = self.deleteNode(root.right,root.val)
      return root
```

```
[4]: # do not modify testing code
def preorder_step(node, result):
    if not node:
        return

    result.append(node.val)
    preorder_step(node.left, result)
    preorder_step(node.right, result)

def preorderTraversal(root: Optional[TreeNode]) -> List[int]:
    ls = []
    preorder_step(root, ls)
    return ls
```

OK

0.3.1 Write analysis of the Memory Complexity and Time Complexity using Asymptotic Notation O. (1 point)

Memory Analysis: O(h), where h is height of tree

Time Analysis: O(h), where h is height of tree

0.3.2 Task 3. Theoretical Question.

As we have seen on lecture Binary Search Trees can be created differently for the same list of n keys. Some tress can be close to optimal (with the height close to log(n)), and some tress can be very suboptimal (with the height close to n/2), but still all these trees are valid Binary Search Trees for given list of n keys.

Derive mathematically how many valid Binary Search Trees can be created for the list of n distinct keys. Hint: you will come up with a very useful formula which is used in combinatorics. Write it on a list of paper and embed it here as a screenshot or write with markdown inside of colab notebook.

Good Luck!

Using n as number of distinct keys, we can use Catalan numbers to determine the number of distinct BSTs, which is given by:

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$