

CS 354.
 Assembly Worksheet -1
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x86 general-purpose registers

(most significant)	(least)		
[.....]]	eax	32 bits
	[.....]	ax	16 bits
	[.....]	ah	8 bits
	[.....]	al	8 bits
[.....]]	ebx	
	[.....]	bx	
	[.....]	bh	
	[.....]	bl	
[.....]]	ecx	
	[.....]	cx	
	[.....]	ch	
	[.....]	cl	
[.....]]	edx	
	[.....]	dx	
	[.....]	dh	
	[.....]	dl	
[.....]]	esi	
[.....]]	edi	

Referred to as %eax, %ebx, %ecx, %edx, %esi, %edi, etc.

INSTRUCTION: mov SOURCE, DESTINATION

definition: moves "SOURCE" into "DESTINATION"

commonly has trailing character that indicates size of move, e.g.,

movb - move a byte

movw - move 2 bytes.

movl - move "long" or 4 bytes (that's an L after mov, not a one)

movq - quad or 8 bytes

our focus: movl (mostly)

Initial (limited) usage

- source=number ("immediate") destination=register

e.g., mov \$10, %eax

- source=register

destination=register

e.g., mov %eax, %ebx

Later, we will add different types of operands for mov

INSTRUCTION: **addl SOURCE, DESTINATION**

definition: adds SOURCE and DESTINATION, puts result into DESTINATION
i.e., $DESTINATION = DESTINATION + SOURCE$

limited usage (for now):

- source=number ("immediate") destination=register
- source=register destination=register

INSTRUCTION: **subl SOURCE, DESTINATION**

definition: $DESTINATION = DESTINATION - SOURCE$

limited usage (for now):

- source=number ("immediate") destination=register
- source=register destination=register

INSTRUCTION: **imull SOURCE, DESTINATION**

definition: $DESTINATION = DESTINATION * SOURCE$

alternate:

imull AUX, SOURCE, DESTINATION

definition: $DESTINATION = AUX * SOURCE$

limited usage (for now):

- source=number ("immediate") destination=register
- source=register destination=register
- (aux=immediate)

INSTRUCTION: **idivl DIVISOR**

definition: contents of %edx:%eax (64 bit number) divided by DIVISOR
quotient -> %eax
remainder -> %edx

limited usage (for now):

- divisor=register

Notes: A bit weird in its usage of VERY SPECIFIC registers!

Problem #1

Write assembly to:

- move value 1 into %eax
- add 10 to it and put result into %eax

*movl \$1, %eax**movl 10(%eax), %eax**OR movl %10, %ebx**addl %ebx, %eax***Problem #2**Expression: $3 + 6 * 2$

Use one register (%eax), and 3 instructions to compute this piece-by-piece

*movl \$6, %eax**imull \$2, %eax**addl \$3, %eax***Problem #3***movl \$0, %edx**movl \$7, %eax**movl \$3, %ebx**idivl %ebx**movl %eax, %ecx**movl \$0, %edx**movl \$9, %eax**movl \$2, %ebx**idivl %ebx**movl %edx, %eax**addl %ecx, %eax**[%edx 0]: [%eax 7]**7/2 [%edx: rem 1] [%eax: quo 2]**%ecx = 2**%edx = 0**%eax = 9**%ebx = 2**9/2 [%edx: rem 1] [%eax: quo 4]**%eax = 1**%ecx = 3*

Write simple C expression that is equivalent to these instructions

$$7//2 + 9\%3 = 3$$

Many x86 instructions can refer to **memory addresses**;
these addresses take on many different forms.

ABSOLUTE/DIRECT addressing

definition: just use a number as an address

`movl 1000, %eax` *careful to remember to use 0x*
gets contents (4 bytes) of memory at address 1000, puts into %eax

NOTE: DIFFERENT than `movl $1000, %eax`
(which just moves the VALUE 1000 into %eax)

INDIRECT addressing

definition: address is in register

`movl (%eax), %ebx`
treat contents of %eax as address, get contents from that address,
put into %ebx

BASE + DISPLACEMENT addressing

definition: address in register PLUS displacement value (an offset)

`movl 8(%eax), %ebx`
address = 8 + contents of eax
get contents from that address, put into %ebx

INDEXED addressing

definition: use one register as base, other as index

`movl 4(%eax, %ecx), %ebx`
address = 4 + contents[eax] + contents[ecx]
get contents from that address, put into %ebx

SCALED INDEXED addressing (most general form)

definition: use one register as base, other as index, scale index by
constant (e.g., 1, 2, 4, 8)

`movl 4(%eax, %ecx, 8), %ebx`
address = 4 + contents[eax] + 8*contents[ecx]
get contents from that address, put into %ebx

Problem #4 (from CSAPP 3.1)

Memory

Address	Value
0x100	0xFF
0x104	0xAB
0x108	0x13
0x10C	0x11

Registers

%eax	0x100
%ecx	0x1
%edx	0x3

Value of:

%eax 0x100

0x104 0xAB

\$0x108 0x108

(%eax) 0xFF

4(%eax) addy: 0x100+4 = 0x104 & = 0xAB

9(%eax, %edx) addy: 9 + 0x100 + 0x3 = 0x10C & = 0x11

260(%ecx, %edx) 0x104 + 0x1 + 0x3 = 108 & = 0x13

0xFC(, %ecx, 4) 0xFC + 4 * 0x1 = 0x100 & = 0xFF

(%eax, %edx, 4) 0x11

New register to help with stack: esp (extended stack pointer)

Referred to as %esp

[.....]	eax	32 bits
[.....]	ax	16 bits
[.....]	ah	8 bits
[.....]	al	8 bits
[.....]	ebx	
[.....]	bx	
[.....]	bh	
[.....]	bl	
[.....]	ecx	
[.....]	cx	
[.....]	ch	
[.....]	cl	
[.....]	edx	
[.....]	dx	
[.....]	dh	
[.....]	dl	
[.....]	esi	
[.....]	edi	
[.....]	esp	32 bits
[.....]	eip	32 bits

Points to "top of stack" when program is running

Changes often (room for local variables, function call/return, etc.)

Can use normal instructions to interact with it, e.g., addl, subl

Can also use special instructions (we'll see this later)

Problem #5

Use instructions to:

- Increase size of stack by 4 bytes
- Store an integer value 10 into the top of the stack
- Retrieve that value and put it into %ecx
- Add 5 to it
- Put final value into %eax

- Decrease size of stack by 4 bytes.

Edited by: Gerald.