

Lecture 7: Fitting to data, Normal equations, Space of linear systems

Jamie Haddock

Table of contents

1	Fitting functions to data	1
1.1	Example: worldwide temperature anomaly	1
2	Least-squares formulation	5
2.1	Least-squares formulation	5
2.2	The normal equations	6
2.3	Pseudoinverse and definiteness	7
2.4	Implementation	7
2.5	Conditioning and stability	7

1 Fitting functions to data

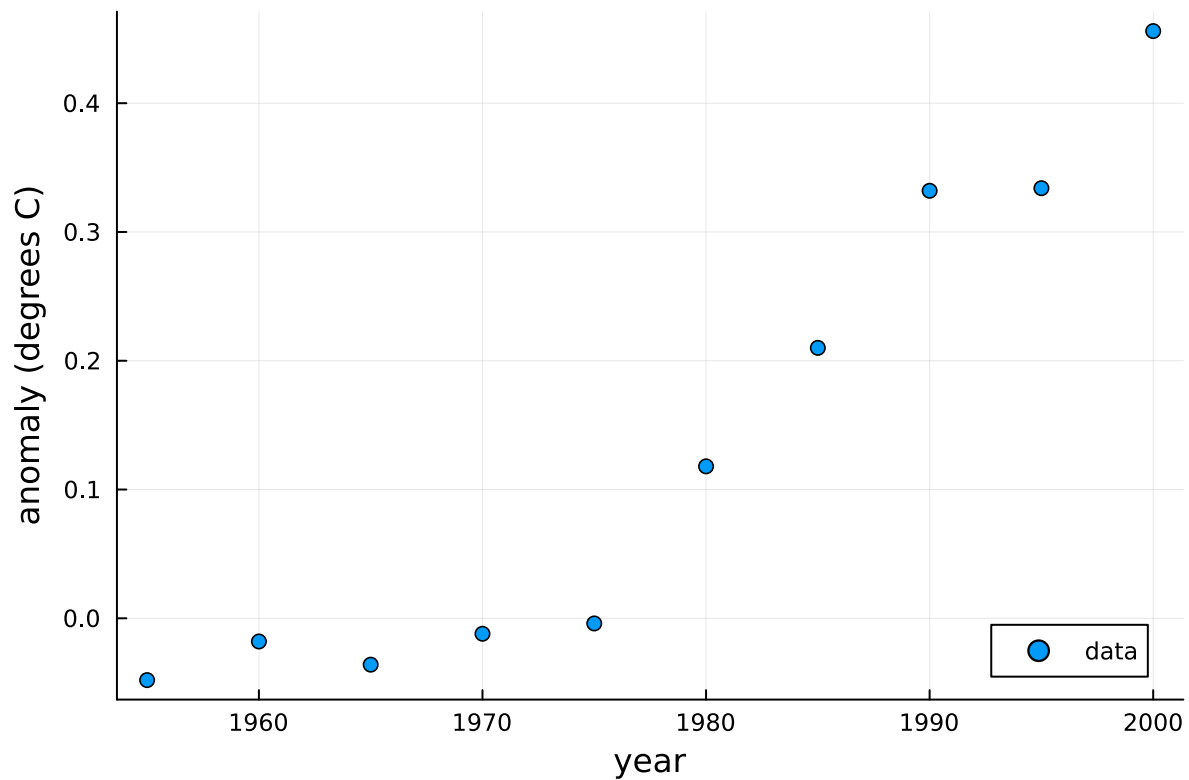
Previously, we saw how interpolating data with a polynomial can be solved by a linear system of equations. However, interpolation is often not an appropriate model for learning a functional relationship from data!

1.1 Example: worldwide temperature anomaly

```
using Plots

year = 1955:5:2000
temp = [ -0.0480, -0.0180, -0.0360, -0.0120, -0.0040, 0.1180, 0.2100, 0.3320, 0.3340, 0.4560 ]

scatter(year, temp, label="data", xlabel="year", ylabel="anomaly (degrees C)", leg=:bottomright)
```

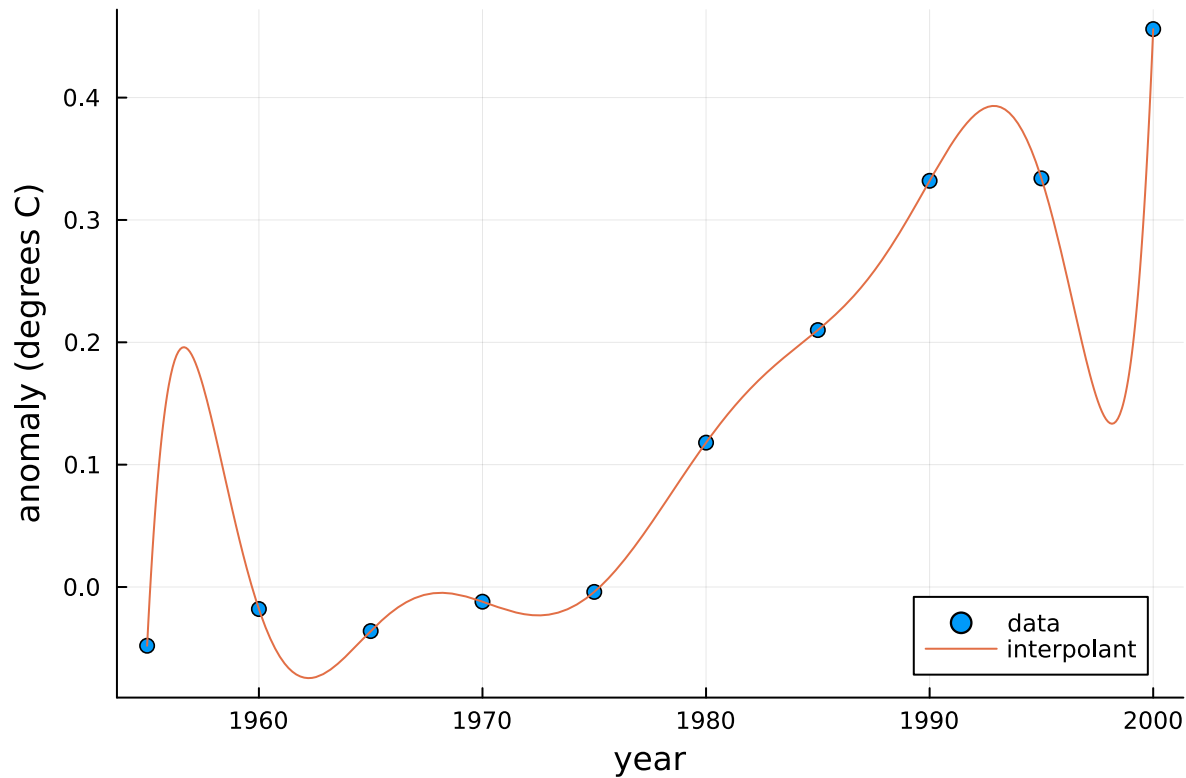


```
t = @. (year-1950)/10
n = length(t)
V = [ t[i]^j for i in 1:n, j in 0:n-1 ]
c = V\temp
```

```
10-element Vector{Float64}:
-14.114000001832462
 76.36173810552113
-165.45597224550528
 191.96056669514388
-133.27347224319684
 58.015577787494486
-15.962888891734785
 2.6948063497166928
-0.2546666667177082
 0.010311111113288083
```

```
using Polynomials
```

```
p = Polynomial(c)
f = yr -> p((yr-1950)/10)
plot!(f, 1955, 2000, label="interpolant")
```



For this application, this functional relationship is far too complex! This is known as *overfitting*.

We can get better results (in this case and many others) by relaxing the interpolant requirement – this is equivalent to lowering the degree of the fitting polynomial.

Let (t_i, y_i) for $i = 1, \dots, m$ be the given points, and let the polynomial be given by

$$y \approx f(t) = c_1 + c_2 t + \dots + c_{n-1} t^{n-2} + c_n t^{n-1},$$

with $n < m$.

We seek an approximation such that

$$\begin{bmatrix} y_1 \\ y_2 \\ y_3 \\ \vdots \\ y_m \end{bmatrix} \approx \begin{bmatrix} f(t_1) \\ f(t_2) \\ f(t_3) \\ \vdots \\ f(t_m) \end{bmatrix} = \begin{bmatrix} 1 & t_1 & \dots & t_1^{n-1} \\ 1 & t_2 & \dots & t_2^{n-1} \\ 1 & t_3 & \dots & t_3^{n-1} \\ \vdots & \vdots & \dots & \vdots \\ 1 & t_m & \dots & t_m^{n-1} \end{bmatrix} \begin{bmatrix} c_1 \\ c_2 \\ \vdots \\ c_n \end{bmatrix} = \mathbf{V}\mathbf{c}.$$

Note that this matrix has the same structure as the Vandermonde matrix but is $m \times n$ with $m \geq n$, and the system is **overdetermined** – it has more conditions than variables.

Overdetermined systems are often **inconsistent**, like this one, and have no exact solution (although it is not impossible for such a system to be consistent). The best approximation of such a system is also given by the `\` operator in Julia.

```
V = [ t.^0 t ] # Vandermonde-ish matrix
@show size(V)
```

```

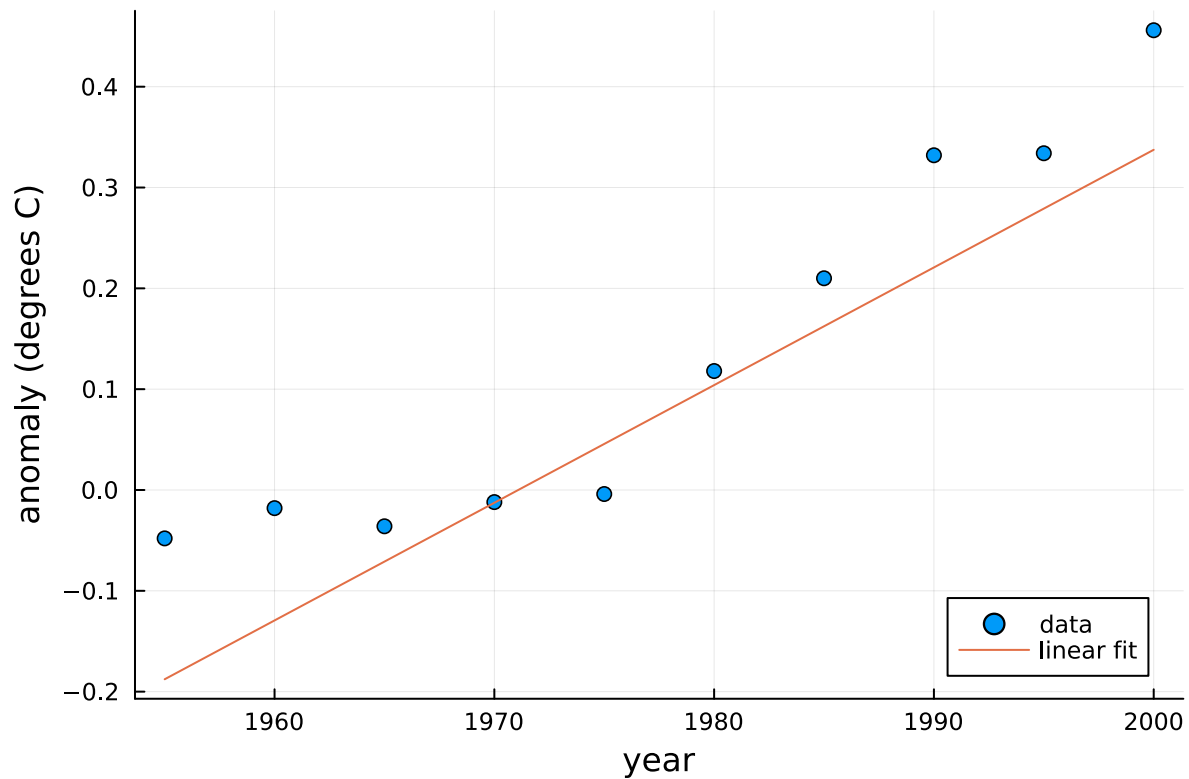
c = V\temp
p = Polynomial(c)

size(V) = (10, 2)

-0.18773333333333356 + 0.11670303030303034 · x

f = yr -> p((yr-1955)/10)
scatter(year,temp,label="data",xlabel="year",ylabel="anomaly (degrees C)",leg=:bottomright)
plot!(f,1955,2000,label="linear fit")

```



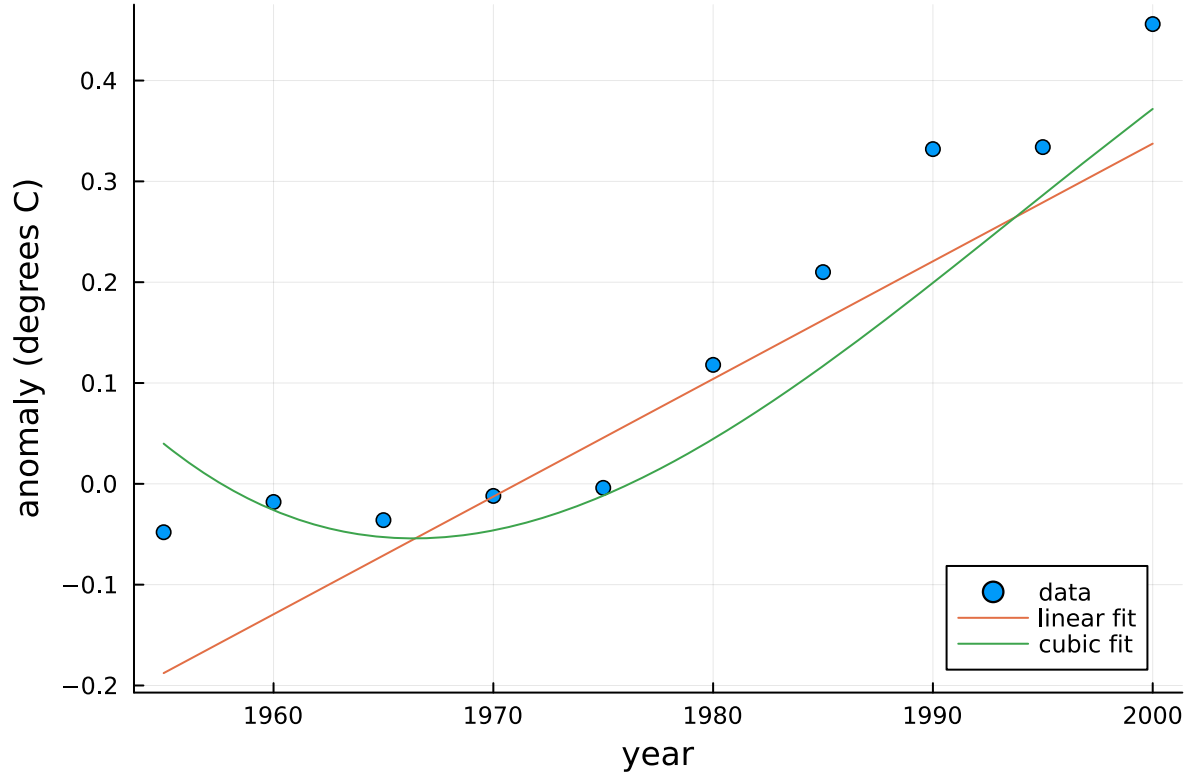
A cubic polynomial fits the data even better.

```

V = [ t[i]^j for i in 1:length(t), j in 0:3 ]
@show size(V)
p = Polynomial( V\temp )
plot!(f,1955,2000,label="cubic fit")

```

```
size(V) = (10, 4)
```



2 Least-squares formulation

This problem here is to fit $y_i \approx f(t_i)$ where $f(t) = c_1 + c_2 t^1 + \dots + c_n t^{n-1}$. This is a special case of the more general case with generic basis functions $f(t) = c_1 f_1(t) + c_2 f_2(t) + \dots + c_n f_n(t)$.

2.1 Least-squares formulation

In either case, the fitting problem solved here is

$$\min R(c_1, \dots, c_n) = \sum_{i=1}^m [y_i - f(t_i)]^2 =: \mathbf{r}^\top \mathbf{r} = \|\mathbf{r}\|^2,$$

where

$$\mathbf{r} = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_{m-1} \\ y_m \end{bmatrix} - \begin{bmatrix} f_1(t_1) & f_2(t_1) & \dots & f_n(t_1) \\ f_1(t_2) & f_2(t_2) & \dots & f_n(t_2) \\ \vdots & \vdots & \dots & \vdots \\ f_1(t_{m-1}) & f_2(t_{m-1}) & \dots & f_n(t_{m-1}) \\ f_1(t_m) & f_2(t_m) & \dots & f_n(t_m) \end{bmatrix} \begin{bmatrix} c_1 \\ c_2 \\ \vdots \\ c_n \end{bmatrix}.$$

Definition: Linear least-squares problem

Given $\mathbf{A} \in \mathbb{R}^{m \times n}$ and $\mathbf{b} \in \mathbb{R}^m$, with $m > n$, find

$$\operatorname{argmin}_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{b} - \mathbf{Ax}\|_2^2.$$

Here, *argmin* is short for *argument minimizing* – this asks one to find the variable, $\mathbf{x} \in \mathbb{R}^n$, that minimizes the objective function.

Note that if we find a solution to the linear system, then $\mathbf{r} = \mathbf{0}$.

2.2 The normal equations

Theorem:

If \mathbf{x} satisfies $\mathbf{A}^\top(\mathbf{Ax} - \mathbf{b}) = \mathbf{0}$ then \mathbf{x} solves the linear least-squares problem – \mathbf{x} minimizes $\|\mathbf{b} - \mathbf{Ax}\|_2$.

Proof:

We'll show that any other vector $\mathbf{x}' = \mathbf{x} + \mathbf{y}$ has objective function value at least as large as the objective function at \mathbf{x} . Note that

$$\begin{aligned}\|\mathbf{Ax}' - \mathbf{b}\|_2^2 &= [(\mathbf{Ax} - \mathbf{b}) + (\mathbf{Ay})]^\top [(\mathbf{Ax} - \mathbf{b}) + (\mathbf{Ay})] \\ &= \|\mathbf{Ax} - \mathbf{b}\|_2^2 + 2\mathbf{y}^\top \mathbf{A}^\top (\mathbf{Ax} - \mathbf{b}) + \|\mathbf{Ay}\|_2^2 \\ &= \|\mathbf{Ax} - \mathbf{b}\|_2^2 + \|\mathbf{Ay}\|_2^2 \\ &\geq \|\mathbf{Ax} - \mathbf{b}\|_2^2.\end{aligned}$$

Definition: Normal equations

Given $\mathbf{A} \in \mathbb{R}^{m \times n}$ and $\mathbf{b} \in \mathbb{R}^m$, the **normal equations** for the linear least-squares problem $\arg\min \|\mathbf{b} - \mathbf{Ax}\|$ are $\mathbf{A}^\top(\mathbf{Ax} - \mathbf{b}) = \mathbf{0}$, or equivalently,

$$\mathbf{A}^\top \mathbf{Ax} = \mathbf{A}^\top \mathbf{b}.$$

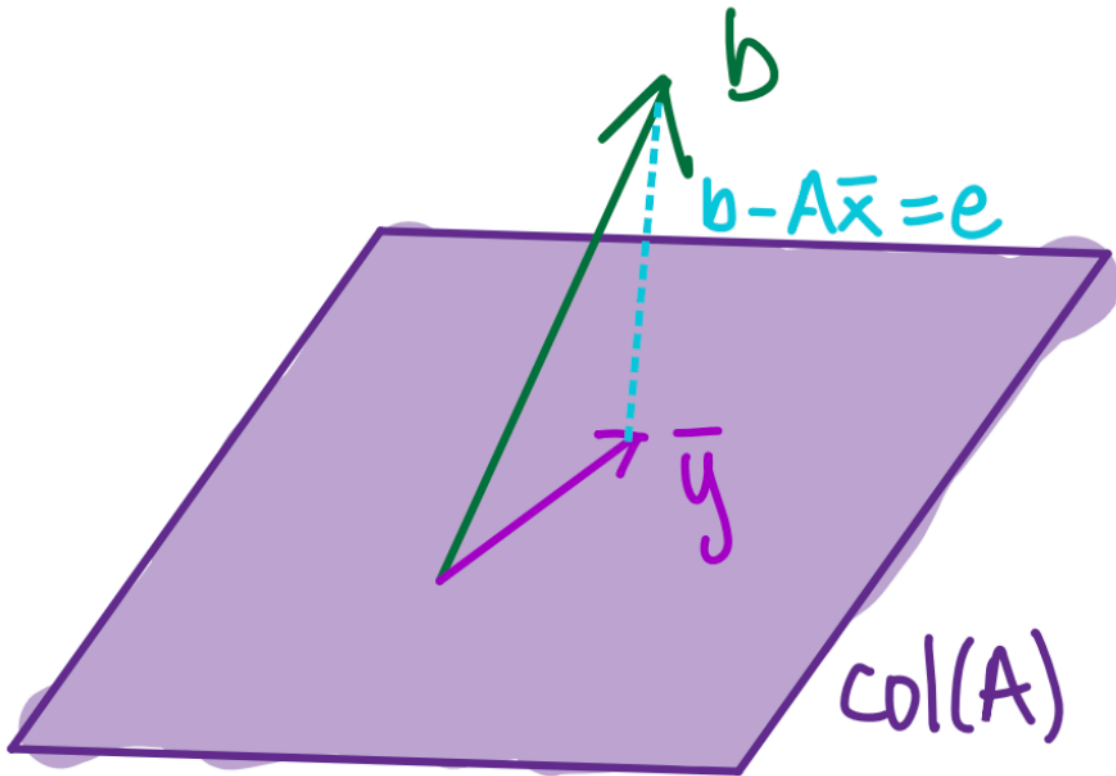


Figure 1: Least-squares geometry

2.3 Pseudoinverse and definiteness

The normal equations show us that we can solve the least-squares problem by solving this system of linear equations.

Definition: Pseudoinverse

If $\mathbf{A} \in \mathbb{R}^{m \times n}$ with $m > n$, its **pseudoinverse** is the $n \times m$ matrix

$$\mathbf{A}^\dagger = (\mathbf{A}^\top \mathbf{A})^{-1} \mathbf{A}^\top.$$

The overdetermined least-squares problem has solution $\mathbf{x} = \mathbf{A}^\dagger \mathbf{b}$. In Julia, the backslash operator `\` is mathematically equivalent to left multiplication by the inverse matrix in the square case and by the pseudoinverse in the overdetermined rectangular case.

Theorem:

For any real $m \times n$ matrix \mathbf{A} with $m \geq n$, the following are true:

1. $\mathbf{A}^\top \mathbf{A}$ is symmetric.
2. $\mathbf{A}^\top \mathbf{A}$ is singular if and only if the columns of \mathbf{A} are linearly dependent; that is, if the rank of \mathbf{A} is less than n .
3. if $\mathbf{A}^\top \mathbf{A}$ is nonsingular, then it is positive definite.

2.4 Implementation

The algorithm for solving least-squares by the normal equations is:

1. Compute $\mathbf{N} = \mathbf{A}^\top \mathbf{A}$.
2. Compute $\mathbf{z} = \mathbf{A}^\top \mathbf{b}$.
3. Solve the $n \times n$ linear system $\mathbf{N}\mathbf{x} = \mathbf{z}$ for \mathbf{x} .

```
"""
    lsnormal(A,b)
Solve a linear least-squares problem by the normal equations.
"""
function lsnormal(A,b)
    N = A'*A; z = A'*b;
    R = cholesky(N).U
    w = forwardsub(R',z)
    x = backsub(R,w)
    return x
end
```

lsnormal

Theorem:

Solution of linear least squares by the normal equations takes $\sim (mn^2 + \frac{1}{3}n^3)$ flops.

2.5 Conditioning and stability

Julia does *not* solve the linear least-squares problem through the normal equations in the algorithm used by `\`. Using the normal equations is unstable.

Definition: Matrix condition number (rectangular case)

If \mathbf{A} is $m \times n$ with $m > n$, then its condition number is defined to be

$$\kappa(\mathbf{A}) = \|\mathbf{A}\|_2 \|\mathbf{A}^\dagger\|_2.$$

If the rank of \mathbf{A} is less than n , then $\kappa(\mathbf{A}) = \infty$.

The normal equations are a square system, so we know from the square case that perturbations to the data \mathbf{A} and \mathbf{b} can be amplified by a factor of $\kappa(\mathbf{A}^\top \mathbf{A})$.

Theorem: Condition number in the normal equations

If \mathbf{A} is $m \times n$ with $m > n$, then

$$\kappa(\mathbf{A}^\top \mathbf{A}) = \kappa(\mathbf{A})^2.$$

We'll be able to prove this when we see some techniques later in the semester. The takeaway is that solving the normal equations doubles the instability of solving the least-squares problem – we shouldn't do this!

```
using LinearAlgebra
```

```
t = range(0,3,length=400)
f = [ x->sin(x)^2, x->cos((1+1e-7)*x)^2, x->1. ]
A = [ f(t) for t in t, f in f ]
    = cond(A)
```

```
1.8253225426741675e7
```

```
x = [1., 2, 1]
b = A*x;
```

```
x_BS = A\b
@show observed_error = norm(x_BS - x)/norm(x);
@show error_bound = *eps();
```

```
observed_error = norm(x_BS - x) / norm(x) = 1.0163949045357309e-10
error_bound = * eps() = 4.053030228488391e-9
```

Given the condition number of this matrix, we expect that solving the linear system, we will lost at most 7 digits of accuracy – this agrees with what we see!

However, if we solve the normal equations, we have a much larger condition number and may not be left with more than two accurate digits.

```
N = A'*A
x_NE = N\(A'*b)
@show observed_err = norm(x_NE - x)/norm(x);
@show acc_digits = -log10(observed_err);
```

```
observed_err = norm(x_NE - x) / norm(x) = 0.021745909192780664
acc_digits = -(log10(observed_err)) = 1.6626224298403076
```

Exercise: Venn diagram of linear systems

Draw a “venn diagram” of the space of all linear systems and mark the sets of consistent and inconsistent systems, the sets of systems with a unique solution or infinitely many solutions, and the sets of overdetermined, square, and underdetermined systems.

Answer:

	$m > n$ over determined	$m = n$ square	$m < n$ under determined	
Systems $Ax = b$	$\text{rk}(A) = n$ unique solution	$\text{rk}(A) = n$		
	$\text{rk}(A) < n$	$\text{rk}(A) < n$ infinitely many solutions	$\text{rk}(A) < m$	consistent
	$\text{rk}(A) \leq n$	$\text{rk}(A) < n$	$\text{rk}(A) < m$	inconsistent