# SOLVING SYSTEMS OF LINEAR INEQUALITIES WITH RANDOMIZED PROJECTIONS

Jamie Haddock

Graduate Group in Applied Mathematics, Department of Mathematics, University of California, Davis

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Joint work with Jesus De Loera and Deanna Needell

# **OPTIMIZATION**

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Today we'll consider a specific form of optimization problem...

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But, even this can be simplified...

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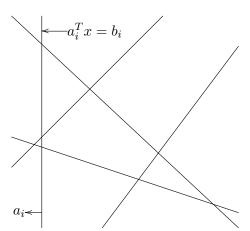
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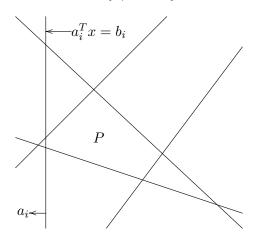
Find x such that  $Ax \leq b$  or conclude one does not exist.

It can be shown that (LP) and (LF) are equivalent.

Reminder: linear equations represent a *hyperplane* (in the proper dimension), so linear inequalities define a *halfspace*.



LF can be interpreted as seeking a point within a (possibly nonempty) polyhedron  $P = \{x | Ax \le b\}$ :



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## PROJECTION METHODS

If we want all of the linear inequalities to be satisfied (meaning we want our point to lie on the correct side of all the hyperplanes), then we need that each of the linear inequalities is satisfied.

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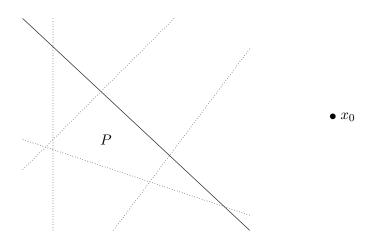
If we want all of the linear inequalities to be satisfied (meaning we want our point to lie on the correct side of all the hyperplanes), then we need that each of the linear inequalities is satisfied. Tautology.

## Projection Methods

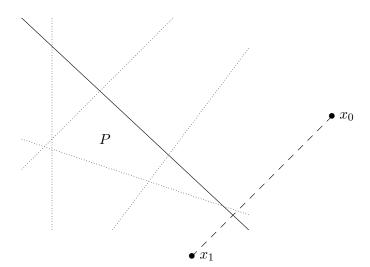
If we want all of the linear inequalities to be satisfied (meaning we want our point to lie on the correct side of all the hyperplanes), then we need that each of the linear inequalities is satisfied.

So... If we have some point that isn't satisfying one of the inequalities, we should force it to satisfy that inequality!

# PROJECTION METHODS



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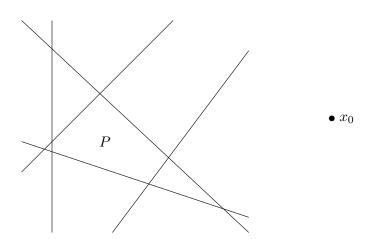


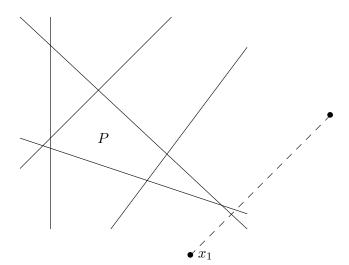
# MOTZKIN'S RELAXATION METHOD(S)

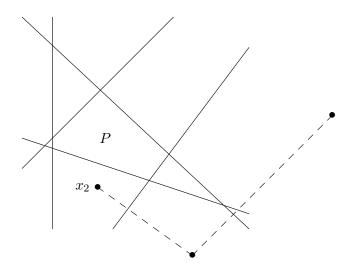
#### **МЕТНО**D

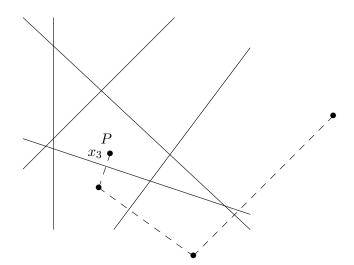
Suppose  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$  and  $P := \{x \in \mathbb{R}^n : Ax \leq b\}$  is nonempty. Fix  $0 < \lambda \leq 2$ . Given  $x_0 \in \mathbb{R}^n$ , iteratively construct approximations to P in the following way:

- 1. If  $x_k$  is feasible, stop.
- 2. Choose  $i_k \in [m]$  as  $i_k := \underset{i \in [m]}{\operatorname{argmax}} a_i^T x_{k-1} b_i$ .
- 3. Define  $x_k := x_{k-1} \lambda \frac{a_{i_k}^T x_{k-1} b_{i_k}}{\|a_{i_k}\|^2} a_{i_k}$ .
- 4. Repeat.







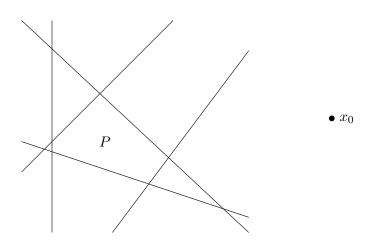


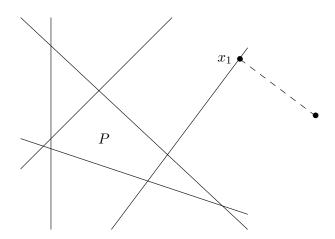
## RANDOMIZED KACZMARZ METHOD

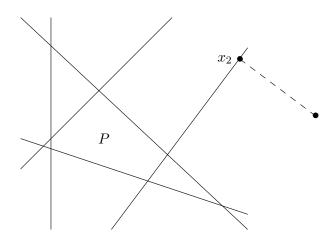
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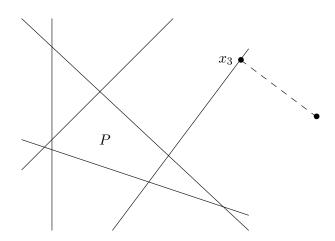
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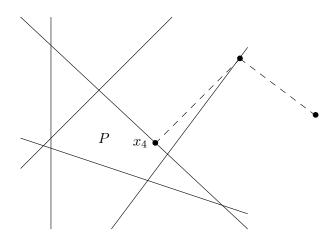
- 1. If  $x_k$  is feasible, stop.
- 2. Choose  $i_k \in [m]$  with probability  $\frac{||a_{i_k}||^2}{||A||_F^2}$ .
- 3. Define  $x_k := x_{k-1} \frac{(a_{i_k}^T x_{k-1} b_{i_k})^+}{\|a_{i_k}\|^2} a_{i_k}$ .
- 4. Repeat.











# **MOTIVATION**

#### Motzkin's Method

Pro: convergence produces monotone decreasing distance

sequence

Con: computationally expensive for large systems

#### **Kaczmarz Method**

Pro: computationally inexpensive, able to analyze the expected

convergence rate

Con: slow convergence near the polyhedral solution set

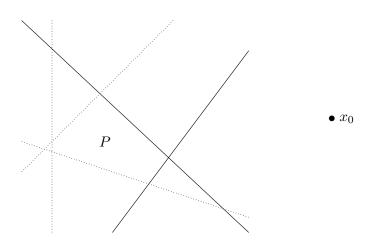
# A Hybrid Method

# METHOD (SKMM)

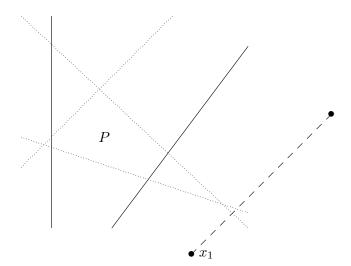
Suppose  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$  and  $P := \{x \in \mathbb{R}^n : Ax \leq b\}$  is nonempty. Fix  $0 < \lambda \leq 2$ . Given  $x_0 \in \mathbb{R}^n$ , iteratively construct approximations to P in the following way:

- 1. If  $x_k$  is feasible, stop.
- 2. Choose  $\tau_k \subset [m]$  to be a sample of size  $\beta$  constraints chosen uniformly at random from among the rows of A.
- 3. From among these  $\beta$  rows, choose  $i_k := \underset{i \in \tau_k}{\operatorname{argmax}} a_i^T x_{k-1} b_i$ .
- 4. Define  $x_k := x_{k-1} \lambda \frac{(a_{i_k}^T x_{k-1} b_{i_k})^+}{\|a_{i_k}\|^2} a_{i_k}$ .
- 5. Repeat.

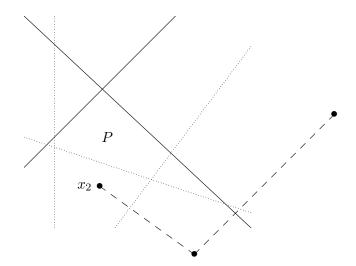
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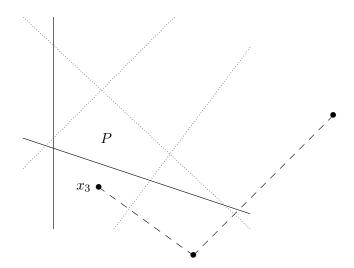
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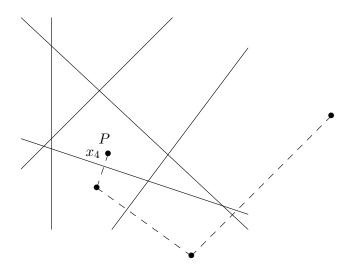
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Note that both previous methods are captured by the class of SKMM methods:

- 1. The Kaczmarz method is SKMM where the sample size,  $\beta=1$  and the relaxation parameter,  $\lambda=1$ .
- 2. Motzkin's Relaxation methods are SKMM where the sample size,  $\beta=m$ .

### AN IMPORTANT REMINDER



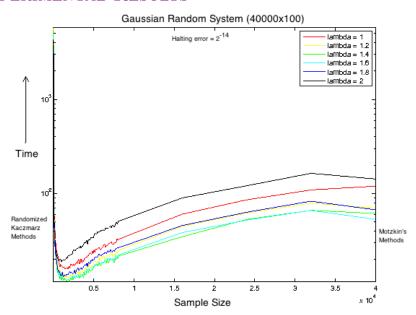
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#### AN IMPORTANT REMINDER.



These methods may not actually stop with a solution... However, we can ensure that our iterate points get arbitrarily close to the solution set, P!

### EXPERIMENTAL RESULTS



## MOTZKIN'S METHOD CONVERGENCE RATE

# THEOREM (AGMON)

For a normalized system,  $||a_i|| = 1$  for all i = 1, ..., m, if the feasible region,  $P := \{x | Ax \leq b\}$ , is nonempty then the relaxation methods converges linearly:

$$d(x_k, P)^2 \le \left(1 - \frac{2\lambda - \lambda^2}{mL_2^2}\right)^k d(x_0, P)^2.$$

# RANDOM KACZMARZ METHOD CONVERGENCE RATE

THEOREM (LEWIS, LEVENTHAL)

If the feasible region,  $P := \{x | Ax \leq b\}$ , is nonempty then the Randomized Kaczmarz method with relaxation parameter  $\lambda$  converges linearly in expectation:

$$\mathbb{E}[d(x_k, P)^2] \le \left(1 - \frac{2\lambda - \lambda^2}{\|A\|_F^2 L_2^2}\right)^k d(x_0, P)^2.$$

## SKM METHOD CONVERGENCE RATE

THEOREM (DE LOERA, H., NEEDELL)

If the feasible region (for normalized A) is nonempty, then the SKM methods with samples of size  $\beta$  converges at least linearly in expectation: In each iteration,

$$\mathbb{E}[d(x_k, P)^2] \le \left(1 - \frac{2\lambda - \lambda^2}{S_{k-1}L_2^2}\right) d(x_{k-1}, P)^2$$

where  $S_{k-1} = \max\{m - s_{k-1}, m - \beta + 1\}$  and  $s_{k-1}$  is the number of constraints satisfied by  $x_{k-1}$ . Clearly then,

$$\mathbb{E}[d(x_k, P)^2] \le \left(1 - \frac{2\lambda - \lambda^2}{mL_2^2}\right)^k d(x_0, P)^2.$$

### IMPROVED RATE

THEOREM (DE LOERA, H., NEEDELL)

If the feasible region,  $P = \{x | Ax \leq b\}$  is generic and nonempty (for normalized A), then an SKM method with samples of size  $\beta \leq m-n$  is guaranteed an increased convergence rate after some K:

$$\mathbb{E}[d(x_k, P)^2] \le \left(1 - \frac{2\lambda - \lambda^2}{mL_2^2}\right)^K \left(1 - \frac{2\lambda - \lambda^2}{(m - \beta + 1)L_2^2}\right)^{k - K} d(x_0, P)^2.$$

## FINITENESS OF MOTZKIN'S METHOD

### THEOREM (TELGEN)

Either the relaxation method\* detects feasibility of the system,  $Ax \leq b$  (with A normalized), within  $k = \left\lceil \frac{2^{4L}}{n\lambda(2-\lambda)} \right\rceil$  iterations or the system is infeasible.

<sup>\*</sup>with  $x_0 = 0$ 

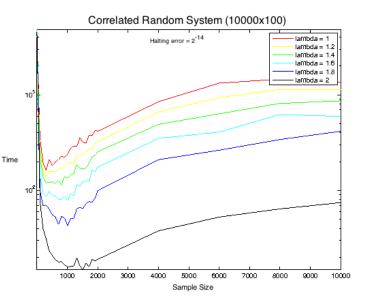
## EXPECTED FINITENESS OF SKM METHODS

THEOREM (DE LOERA, H., NEEDELL)

If the system,  $Ax \leq b$  is feasible, then with high probability the Sampling Kaczmarz-Motzkin method\* with relaxation parameter  $0 < \lambda < 2$  will detect feasibility within a given number of steps.

<sup>\*</sup>with  $x_0 = 0$ 

# CONCLUSIONS



# FUTURE WORK:

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1. Provide theoretical guidance for selection of the optimal sample size,  $\beta$ , and optimal overshooting parameter,  $\lambda$  for a given (class of) system(s).

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- 2. Describe the K after which the convergence rate is guaranteed to be improved.

### ACKNOWLEDGEMENTS

Thanks to you for attending!

Are there any questions?

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