

Jamie Sweeney - 213728s

Question 4

Use 'n' to denote λ -term for n .

$$\text{size}(\text{'n'}) = 7n + 2$$

Question 5

- T - Not linear
- F - Not linear.
- Z - Linear.
- S - Not linear.
- pair - Linear.
- fst - Not linear.
- snd - Not linear.

Question 6

Base case:

M is a variable, x .

Then $M[x := N] = N$.

Since N is linear so too is $M[x := N]$.

Inductive step:

M is an application AB .

$$(AB)[x := N] = (A[x := N])(B[x := N])$$

By induction hypothesis $A[x := N]$ and $B[x := N]$ are linear

Since $A[x := N]$ and $B[x := N]$ are linear so too is $(AB)[x := N]$

Inductive step:

M is an abstraction $\lambda y. A$

$$M[x := N] = \lambda y. (A[x := N])$$

By induction hypothesis $A[x := N]$ is linear if N and A are linear

Since $A[x := N]$ is linear so too is $M[x := N]$

Question 7

Base case:

M is a variable x .

$$\text{size}(M) = 1$$

$$\text{Then } M[x := N] = N$$

$$\text{We have } \text{size}(M[x := N]) = \text{size}(N)$$

$$\text{And } \text{size}(M) + \text{size}(N) = \text{size}(N) + 1$$

$$\text{Since } \text{size}(N) < \text{size}(N) + 1, \text{ we have } \text{size}(M[x := N]) < \text{size}(M) + \text{size}(N)$$

Inductive step:

M is an abstraction $\lambda y.A$

$$\text{size}(M) = \text{size}(A) + 1$$

$$M[x := N] = \lambda y.(A[x := N])$$

$$\text{size}(M[x := N]) = \text{size}(A[x := N]) + 1$$

$$\text{By our induction hypothesis, } \text{size}(A[x := N]) < \text{size}(A) + \text{size}(N)$$

$$\text{So, } \text{size}(A[x := N]) + 1 < \text{size}(A) + \text{size}(N) + 1$$

Substitute LHS with $\text{size}(M[x := N])$ to get,

$$\text{size}(M[x := N]) < \text{size}(A) + \text{size}(N) + 1$$

Substitute with $\text{size}(M)$ to get,

$$\text{size}(M[x := N]) < \text{size}(M) + \text{size}(N)$$

So hypothesis holds.

Inductive step:

M is an application AB

$$\text{size}(M) = \text{size}(A) + \text{size}(B)$$

$$M[x := N] = (A[x := N])(B[x := N])$$

$$\text{size}(M[x := N]) = \text{size}(A[x := N]) + \text{size}(B[x := N])$$

$$\text{By our induction hypothesis, } \text{size}(A[x := N]) < \text{size}(A) + \text{size}(N)$$

$$\text{And } \text{size}(B[x := N]) < \text{size}(B) + \text{size}(N)$$

$$\text{So, } \text{size}((A[x := N])(B[x := N])) < \text{size}(A) + \text{size}(B) + 2 \cdot \text{size}(N)$$

Substitute LHS with $\text{size}(M[x := N])$ to get,

$$\text{size}(M[x := N]) < \text{size}(A) + \text{size}(B) + 2 \cdot \text{size}(N)$$

Substitute with $\text{size}(M)$ to get,

$$\text{size}(M[x := N]) < \text{size}(M) + 2 \cdot \text{size}(N)$$

Therefore

$$\text{size}(M[x := N]) < \text{size}(M) + \text{size}(N)$$

So hypothesis holds.

Question 8

Question 9

For any β reduction we have:

$$(\lambda x.M)N \rightarrow_{\beta} M[x := N]$$

$$\text{Size(LHS)} = \text{size}((\lambda x.M)N) = 1 + \text{size}(M) + \text{size}(N)$$

$$\text{Size(RHS)} = \text{size}(M[x := N])$$

$$\text{We know that } \text{size}(M[x := N]) < \text{size}(M) + \text{size}(N)$$

So,

$$\text{Size(RHS)} < \text{size}(M) + \text{size}(N)$$

$$\begin{aligned} \text{Since } \text{size}(M) + \text{size}(N) &< 1 + \text{size}(M) + \text{size}(N) \\ \text{size(RHS)} &< \text{Size(LHS)} \end{aligned}$$

So for any linear M , if $M \rightarrow_{\beta} N$.

$$\text{size(RHS)} < \text{Size(LHS)}$$

$$\text{size}(N) < \text{Size}(M)$$

Question 10

Linear λ -calculus seems quite limited in the scope of computations that are available. From Q4 we see that only Z and pair are linear. I would expect that the terms expressible in linear λ -calculus would be those with no 'branches' in computation or conditional terms.