2022 NOVEMBER

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1. Span	n Core						
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* features :
- In charge of essential I/O functionalities Task dispatching
- Task dispattching
0 11 11 1
Embedded with RDD8 (Ruilient distributed
dataset).
" -> handles partitioning data across all nodes
ALV LOAKE BOATLY
12 > It hadds them in memory pool of the cluster as a single unit.
as a single unit.
2. Spark SQL
3 works to acrese structured & semi-structured
info.
& enables powerful, interactive, analytical app.
-> Enables powerful, interactive, analytical app. across both streaming & histopical data.
4
> Features :
5- Cost based optimizes
- Mid query fault taulerent
- Mid query fault tablerant Full compatibility with existing nive data
3. Spark streaming
- Add on to we spark API which allows
e la la ble bigh a thunuahout, fautt tours
stream proceering of live data streams.

- Accese data from sources like or TEP cocket MIWIFSSMIWIFSSMIWIFSSMIWIFSSMIWIFSS DEC

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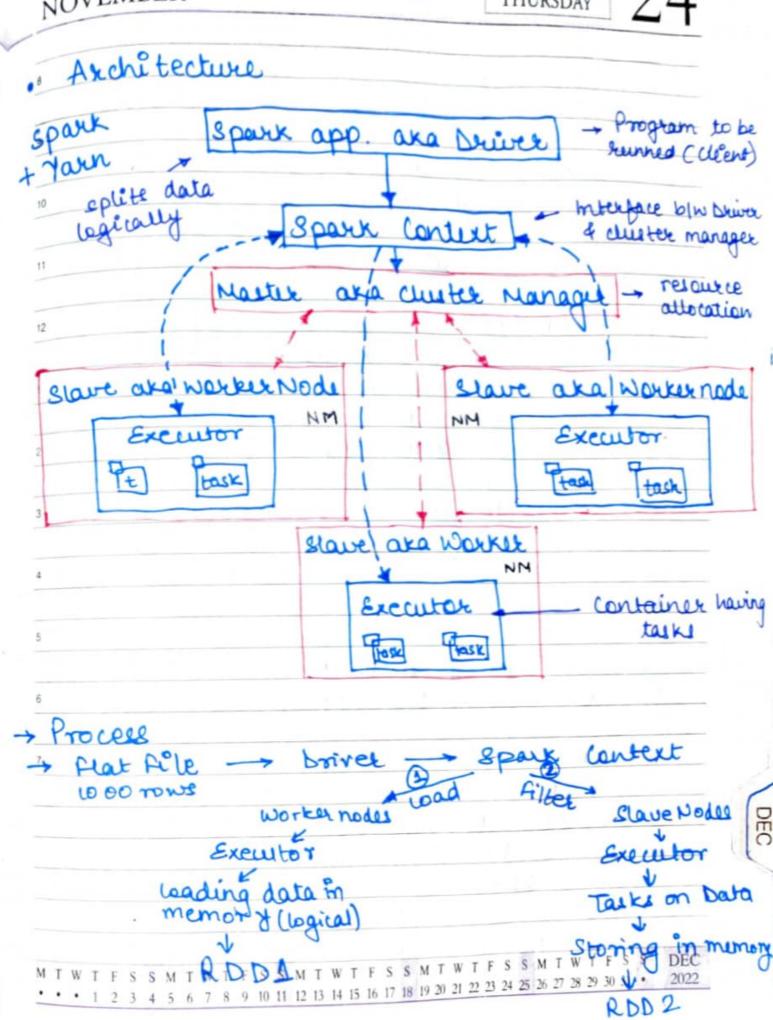
## WEDNESDAY

- 4. Spark Muib
  4. Spark Muib
  Library that performs
  Library that performs
  high-quality algorithm & high spread
  high-quality algorithm & high spread
- 5. Graph X

  3 API for graphs & graph parallel execution.

  3 It is network graph analysics engine &

  " data store.
- · 12 Features of spark
  - 1. Swift proceeding spud, made possible by reducing no. of read-write to disk.
- 2. Synamic in Nature is easily develop a parallel application
- 3: In- Memory compulation in spark
  - Processing speed is increased
- from disk every time, thus time is saved.
- 4. Cault tolerant
  - Provided through RDDs.
- 5. heal Time Stream processing
- 6. Lazy Evaluation
- All the transformation we make in RPD are
- NOV MINTES SMIWIFS SMI



1. Drivet

"> It is a process that rune the main() for all the application & creates spark context object the application & creates spark Application of a spark Application of the states of spark dustry & maintains are of the states of spark dustry of the manager in order to actually get physical resources in order to actually get physical resources.

2. Spark Context

> Coordinate the spark applications, running
as independent sets of processes on a duster.

> It acquires executors on nodes in cluster.

> Then, it sends your app. code to the exactor

> At last, the spark Context sends tack to the

executors to run.

3. Custer Manager.

4. Norker Nede I slave node. 5 role is to run the approde in dustre.

5. Executor sunde ruentes to driver.

- An executor is a process launched for an application on worker node.

- It suns tasks & keeps data in memory or disk storage across them.

- It read & write data to external sources

- Every app, contain its executor

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- A unit of work that will be sent to one executor.

DEC

## NOVEMBER

(Resilient Distributed bataset)

It is a fundamental data structure of spark.

It is an immutable collection of objects which

computes on the different node of the duster.

across servers to be computed on diff

Resilient, i.e fault tolorant with the help

of RDD lineage graph (DAG). to recompute miseing or damaged partitions due to node

failures.

- Distributed, data residing on multiple nodes - bataset, represents records of data you work

· Features

- In-memory computation

- stores everything in RAM

- Lazy Evaluations

- spark computes transformation when an action requires a result for driver program.

fault Tollrance.

- Because of lineage

- Immutability

- cannot be changed or updated

- Partitioning

> fundamental unit of paralleliem in spark Ros

on existing partitions PELLI 2 13 15 CP 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 .

\* Users can state which RDDs they will reuse & choose a storage strategy for them.

· limitations · No inbuilt optimization engine. · No inbuilt optimization engine.
- Handling states
- Performance limitation  - Callection & Java suialization which as  - Collection & Java suialization which as
expensive where all
RAM then it uses disk for remaining
Ban Robert
· 2 Ways to Create RDDs in Spark.
1. Parallelized collection (parallelizing)  by taking an existing collection in the properties it to spark context's parallect  method.
vsed in initial stage, as it creates RDD quirkly
2. External bataoit (Referencing a dataset).
of line.
+ loading CSV, JSON, textfile.
3. Creating RDD from existing RDD transformation is the way to viewle an

## . RDD Persistence & Caching

- Persistance is an optimization technique in which saves the reput of RDD evaluation.
- so that we can use it it required.
- it reduces the computation areashead.
- > when we use the cache () method we can store
- 12 all the RDD in-memory
- we can persist the RBD in memory & use it efficiently across parallel operations.
- The diff. blw cache () & persist () is that using cache () the default storage level is NEMORY\_ONLY while using persist () we can use various storage levels.
  - > When we persist RDD each node stores any partition of it that it computes in memory & makes it restable for future use. This process speed up the computation.
  - is kept in the memory on the rode. The cache memory of the spark is fault tolerant so whenever any partition of RDD is lost, it can be reconstred by transformation operation that originally created it.

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TUESDAY well of persist () (spill over induk) ( sevialized Java doje memory only sex 4. Memory and disk sek 5. Diek buly Spark automatically drops out the old data partition in LRU ( least recently us) fashion we can remove manually using RDD unposition 2 RDD Operations 1. Transformation new RDD from existing Applying transformation entire parent RDDs of the final RDD dependency graph. Are lazy in nature i.e they get executed call an action. Two types of transformation: 1. Narrow transformation -> All elements that are seg, to compute the xecords in single partition live in the single partition of parent RDD. , filter (), flat Mop (), map

2. Wide transformation All elements that req. to compute the records in the single partition may live in many partitions of parent RDO.

Sea. groupby Key (), reduceby Key (), Join () intersection (), distinct (), cartesian (), Repartition (), coalise () 22. Action - works with actual data when action is perform - boes not form new RDD. - The values of action who stored to soprest drivers or to the external storage system - Action is one of the way of sending from Executor to the driver. > count() - no. of elements is returned collect () - returns the entire RDDs content to driver program take (1) - returns n no. of elements from RDD. top () - extract top eternents from RDD. count By Value () - returns, occurrence of each element L'heduce () - operation like addition, takes two L'fold () - Takes zero value as input. Telement as diff. - reduce theous an exception for empty collection, but fold is defined for empty collection

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aggregate(), foreach()

WEDNESDAY

\* RDD Lineage

RDD, that new RDD auso carries a pointer

to the parent RDD.

- Au the dependencies b/w RDDs those are logged in a graph, redher than actual data, called lineage graph.

It is a graph of all the parent RDDs of

an RDB.

\* bataframe

- Data organised into named columns.

3 similar to table in RDBMS.

-> can say that it is a relational table with good optimization technique.

43 Processes large amount of structured data

-> contains schema (illustration of the structure of data).

-> Immutability, in-memory, resilient, distributed

computing capability.

-> sources data from structured data file, tables in rive, external dbs or existing RODS.

- Available in scala, Java, Python & R.

> bata frame over RDD

1 Provides memory management

- data is stored in of hear memory in MTWTFS SMTWTFS SMTWTFS SMTWTFS SMTWTFS S 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 • • •

a optimized Execution plan
another optimized where execution alon is
2 optimized Execution plan - query optimizer, where execution plan is - vieated for the execution of a query.
10
· Limitation of RDD
Limitation of RDD  1 Does not have any built-in optimization
2 No provision to handle structured data.
·
· features of DataFrame
7. distributed collection of data organized in
2 named column, equivalent to table in RDBM
- Seals with structured Esconstructured data
· formats. eg. Auro, csv, elastic search, cassandre
- beals with storage eysterne-HDFS, HWE, MYSGIEL
& catalust eupoports optimization, 4 phases:
1. Analyze logical plan to solve references
2. Logical plan optimization. 3. Prysical planning
3. Physical planning
· 4. code generation to compute part of query to
java byterade.
, 0
· Limitation
- Does not provide provision for compile
a structure in order to manipulate.
a structure in order to manipulate.

03

FRIDAY

Lataret

- Stronly typed & is map to relational scheme.

- Represents structured queries with encoders.

- Provides both type eafrey & object-oriented programming interface.

- Optimized Querry - (catalyst query Optimizes).

- aplimized guerry - (catalyst guerry Optimizer).

- Analysis at compile Time. - check syntax & analyse
at compile time.

- Persistent storage - serializable & gueryable.
- Faster computation ( than RDD).

2 tess memory consumption - structure of data in dataset is known.