

DDH & attacks from genus theory

$$E/\mathbb{F}_q \quad \text{End}_{\mathbb{F}_q}(E) = \mathbb{O} \quad \text{order in } \mathbb{Q}(\sqrt{\ell^2 - 4g})$$

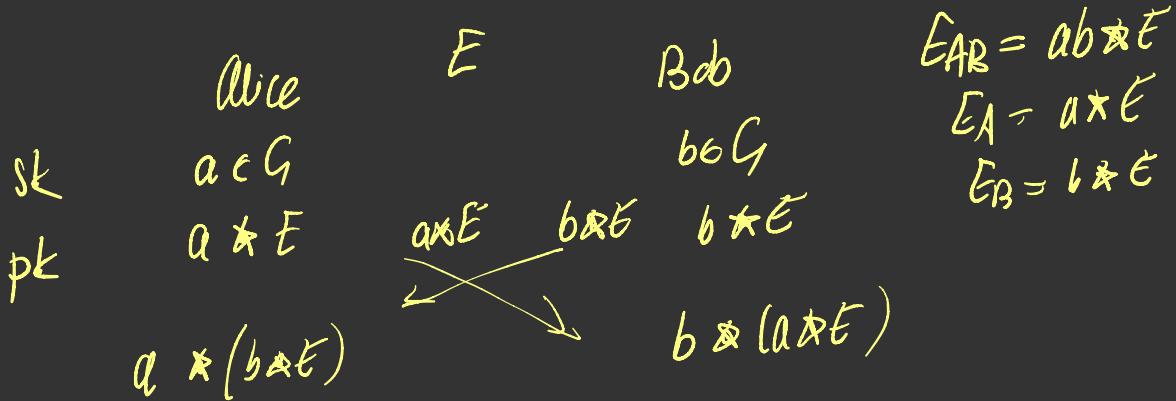
$$\# E(\mathbb{F}_q) = q + 1 - t$$

$$\ell^2 - 4g \quad \text{disc of } x^2 - tx + q \quad \leftarrow \\ G \in \mathbb{Z}[\pi] \quad \pi \text{ Frobenius, root of}$$

$$\text{Ell}_q(O, t) = \left\{ E' / \mathbb{F}_q : \begin{array}{l} \text{End}(E') = \mathbb{O} \\ \text{Fr}(E') = t \\ \# E'(\mathbb{F}_q) = q + 1 - t \end{array} \right\}_{\mathbb{F}_q = \mathbb{Z}}$$

$G = \mathcal{O}/\mathcal{O}$ acts on the set $X = \text{Ell}_q(G, t)$

Group G acting on a set X , free and transitive



Secure if Alice and Bob are the only ones
who know E_{AB}

↳ CDH assumption

Protocol $E, E_A, E_B \rightsquigarrow$ we can derive
information about E_{AB}

$$y^r = x^3 + Ax^2 + x$$

$j(E_{AB})$ as shared key

DDH: There's nothing predictable about E_{AB}
if we only see E, E_A, E_B .

CSV: There exist characters χ s.t.

$$\chi(E_{AB}) = \chi(E_A) \cdot \chi(E_B)$$

χ character on \mathbb{Q}/\mathbb{O}

χ quadratic character

$$\S 2 \quad \frac{\text{Pairings}}{E \times E} \rightarrow \mathbb{F}_2$$

bilinear map Weil pairing $E[N] \times E[N] \rightarrow \mu_N$

$$e_N(P, Q) = \sum_{n=1}^N e_n \in \mathbb{F}_2$$

m odd prime

$$E \rightarrow E' = [a] \not\cong E$$

$$[a] \in \mathcal{O}(0)$$

for every $a \in [a]$, we have an isogeny

$$\varphi_a : E \rightarrow E' \quad \deg \varphi_a = \text{norm } a$$

If we can say anything about $\deg \varphi_a \pmod{m}$

we get information about every $a \in [a]$

$$(\deg \varphi_a \pmod{m})_{a \in [a]}$$

E, E' isogenous curves $\text{val}(\#E(\mathbb{F}_q)) = 1$

$p \in E(\mathbb{F}_q)[m]$ then degree of my isogeny
 $p' \in E'(\mathbb{F}_q)[m]$ $\varphi: E \rightarrow E'$

is, up to squares mod m

$$\deg \varphi = \begin{bmatrix} \log_{T_m(p, 0)} & \log |T_m(p', p)| \\ \log T_m(p, 0) & \end{bmatrix}$$

$$\begin{array}{ccc} E & \xrightarrow{a} & E_A \\ b \downarrow & & \downarrow b \\ E_B & \xrightarrow{a} & E_{AB} \stackrel{?}{=} E_C \end{array}$$

$G = d(O)$ O order in im. quad field

\mathbb{Q}
 $X = \text{Ell}$ ← elliptic curves / \mathbb{F}_q

$\text{End}_{\mathbb{F}_q}(E) \cong O$

trace $t = \#E(\mathbb{F}_q) - q - 1$ is fixed

$G \curvearrowright X$ is free and transitive

DH key exchange

E ← starting pt

Alice

Bob

sk

a

b

pk

$a \star E$

$b \star E$

$a \star (b \star E)$

$b \star (a \star E)$

1) Secure if Alice and Bob the only ones
who hold $E_{AB} = ab \star E$ } CDH

$DDH \simeq E, E_A, E_B$ then E_{AB} "looks random"

\curvearrowleft
public keys

you can't use E, E_A, E_B to predict anything
about E_{AB}

How many Supersingular curves are there

$$\approx \frac{p}{12} = \# \text{supersingular } j\text{-invariants in } \mathbb{F}_{p^2}$$

$$\sqrt{p} = \# \text{supersingular curves over } \mathbb{F}_p, \text{ up to iso}$$

$$\approx \# \text{cl}(\mathbb{Z}[\sqrt{p}])$$

$$h(-\Delta) \approx \sqrt{|\Delta|}$$

Wouter Weeke 3

| $p \bmod 8$ | |
|-------------|------------------|
| $1 \bmod 4$ | $h(-4p)$ |
| $3 \bmod 8$ | $h(-4p) + h(-p)$ |
| $5 \bmod 8$ | $2 \cdot h(-p)$ |

Ex 1.1. There are p different A^1 's
only \sqrt{p} of them are supersingular curves by $\frac{\sqrt{p}}{p}$

§ 2 Pairings m odd prime, E/\mathbb{F}_q

$$T_m: E(\mathbb{F}_q)[m] \times E(\mathbb{F}_{q^{\frac{m}{\gcd(m, q-1)}}}) \rightarrow \mu_m$$

$$(P, Q) \mapsto T_m(P, Q)$$

$\mu_m \subseteq \mathbb{F}_q$
 $\Leftrightarrow m | q-1$

$$\begin{aligned} \varphi: E &\rightarrow E' \\ P, Q &\mapsto (\varphi(P), \varphi(Q)) \end{aligned}$$

$$T_m((\varphi(P), \varphi(Q))) = T_m(P, Q)^{\deg \varphi}$$

$$\xrightarrow{\quad \underbrace{\quad} \quad \deg \varphi \pmod{m} \quad}$$

"hard problem": $E \rightarrow E'$ under some secret isogeny φ

find pair $P, P' = \varphi(P)$

Get around: isogenies

$\left\{ \begin{array}{l} \text{rational map: points in } E(\mathbb{F}_q) \text{ map to } E'(\mathbb{F}_q) \\ \text{group hom: } E[m] \longrightarrow E'[m] \end{array} \right.$

1) if $E(\mathbb{F}_q)[m] = \langle P \rangle$ $\Rightarrow \varphi(P) = Q \cdot P'$
 $E'(\mathbb{F}_q)[m] = \langle P' \rangle$
to unknown

E E'

(conditions apply)

computing these Tate pairings determines

$\deg \varphi \pmod{m}$, up to squares
 \pmod{m}

for any isogeny $\varphi: E \rightarrow E'$

DDH If all those conditions hold, and we can compute
 $\deg \text{ of } \varphi \text{ mod } m$ up to squares

$$E \xrightarrow{\varphi} E_A$$

$$\downarrow \quad \quad \quad \downarrow$$

$$E_B \xrightarrow{\alpha} E_{AB} \xrightarrow{\beta} E_C \quad E_C \text{ random}$$

$\deg \text{ of } \varphi \text{ mod } (E_B \rightarrow E_{AB}) \dots$

$$\chi(E, E_A) = \left(\frac{\deg \varphi}{m} \right) \quad \leftarrow \begin{array}{l} \text{Legendre} \\ \text{symbol} \end{array}$$

$$= \begin{cases} 1 & \text{if } \deg \varphi \underset{\text{mod } m}{\text{always}} \text{ square} \\ -1 & \text{if } \deg \varphi \underset{\text{mod } m}{\text{always}} \text{ non-square} \end{cases}$$

always: for any isogeny $\varphi: E \rightarrow E'$

Conditions: $\mu_m \subseteq \mathbb{F}_q$, $E(\mathbb{F}_q)[m^\infty] = \langle P \rangle$ P of order m

$$\begin{array}{ccc} \text{S1DH} & E & \xrightarrow{2^n} E_A \\ & 3^m \downarrow & \\ & E_B & E_C \stackrel{?}{=} E_{AB} \end{array}$$

E_{AB} is 3^m -isogenous to E_A

2^n -isogenous to E_B

A) Decide whether E_C is 3^m -isog to E_A

B)

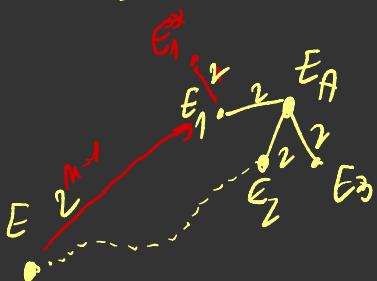
2^n -isog to E_B

CS1DH
 E, E'
 always
 $E' = aE$
 for some a

I) Galbraith - Vercauteren Comp problems in Gray...

II) Jao - Urvashi? Sol

DDH to CDH reduction Isogeny finding



$$\begin{aligned} m &= \# \text{steps in S1DH} \\ &= \text{poly size} \\ &= O(\log p) \end{aligned}$$

Input to the problems

$$P_1, \#_{p^2}, E_0/\#_{p^2}, \dots$$

$$2^n \sqrt{p}$$

$$m \approx \frac{1}{2} \log p$$

Widen it down using $\log_2 p$ bits

Polynomial time = poly in $\log_2 p$