

```

patrick@nightmare:~/OS_examples$ ls
'Data Interleave' 'Data Race Illustrated' pthread_eg.c test_fork.c test_fork_exec.c
patrick@nightmare:~/OS_examples$ gcc test_fork.c -o tf -lpthread
patrick@nightmare:~/OS_examples$ ./tf
Bye from the child process! PID= 0
Child 29894 Completed with status = 0
patrick@nightmare:~/OS_examples$ gcc test_fork_exec.c -o tfe -lpthread
patrick@nightmare:~/OS_examples$ ./tfe

----- BEFORE COMMAND-----
total 60
drwx----- 3 patrick patrick 4096 May 28 21:03 'Data Interleave'
drwx----- 3 patrick patrick 4096 May 28 21:08 'Data Race Illustrated'
-rw----- 1 patrick patrick 1051 May 28 19:52 pthread_eg.c
-rw----- 1 patrick patrick 619 May 27 18:50 test_fork.c
-rw----- 1 patrick patrick 576 May 27 18:48 test_fork_exec.c
-rwx----- 1 patrick patrick 16928 May 29 19:34 tf
-rwx----- 1 patrick patrick 16920 May 29 20:21 tfe

-----Child Complete-----
patrick@nightmare:~/OS_examples$ vim test_fork

```

```

#include <sys/types.h>
#include <sys/wait.h>
#include <stdio.h>
#include <unistd.h>
int main()
{
    pid_t pid;
    /* fork a child process */
    pid = fork();

    if (pid < 0) { /* error occurred */
        fprintf(stderr, "Fork Failed");
        return 1;
    }
    else if (pid == 0) { /* child process */
        printf("\n----- BEFORE COMMAND-----\n");
        execlp("/bin/ls", "ls", "-l", NULL);
        printf("\n-----AFTER COMMAND-----\n");
    }
    else { /* parent process */
        /* parent will wait for the child to complete */
        wait(NULL);
        printf("\n-----Child Complete-----\n");
    }
}

```

- With child process, have two print statements (BEFORE and AFTER)
 - AFTER never printed – because execlp is executed
 - Once execlp executed – overlays parent program, destroys whatever was in the parent address space
 - Now have new address space – load ls program into that and execute it
- Once exec operates – parent no longer exists
 - Have new program running until termination
 - Makes no sense to put code after exec statement – will never be executed

```
package interleaving;

public class STest {
    public static void main(String[] args) {
        int numberThreads = 100;
        AThread [] threads = new AThread[numberThreads];

        for (int i=0;i<numberThreads;i++)
            threads[i]= new AThread(i);

        System.out.println("Starting simulation with " + numberThreads + " threads");

        for (int i=0;i<numberThreads;i++)
            threads[i].start(); // threads don't run until you start them
        System.out.println("Parent thread completed");
    }
}

"STest.java" [dos] 17L, 501C 11,2-16
```

crucial for being able to solve many of the synchronization problems.

- With this code, parent stops before child => issues!

```
package dataRace;

public class TotalCounter {
    private int sum;

    TotalCounter() {sum=0;}

    public void increment(int incrVal) {
        sum+=incrVal;
    }

    public int getValue() {
        return sum;
    }
}

"TotalCounter.java" 15L, 199C 8,23-30
```

```

package dataRace;

public class SumThread extends Thread {
    private TotalCounter counter;
    private int toAdd;

    SumThread(TotalCounter tC, int toAdd) {
        this.counter = tC;
        this.toAdd=toAdd;
    }

    public void run() {
        for (int i=0;i<toAdd;i++) {
            counter.increment(1);
        }
    }
}

```

"SumThread.java" [noeol] 16L, 291C 6,1-8

```

package dataRace;

public class RaceTest {
    public static void main(String[] args) {
        int numberThreads =100;
        int incr = 100; //divide adding task among threads
        TotalCounter counter = new TotalCounter();
        SumThread [] threads = new SumThread[numberThreads];
        for (int i=0;i<numberThreads;i++)
            threads[i]= new SumThread(counter,incr); //create the threads

        System.out.println("Starting simulation with " + numberThreads + " threads");
        for (int i=0;i<numberThreads;i++)
            threads[i].start(); // threads don't run until you start them
    /*
        for (int i = 0; i < numberThreads; i++)
            try {
                threads[i].join();
            }
        catch (InterruptedException e) {}
    */
        System.out.println("Counter is :"+ counter.getValue());
    }
}

```

"RaceTest.java" 24L, 774C 14,23-44

- These three programs show how data races occur

```
import java.util.concurrent.Semaphore;

public class SumThread extends Thread {
    private int id;
    private TotalCounter counter;
    private int toAdd;
    private Random snooze_time;

    SumThread( int n, TotalCounter tC, int toAdd) {
        this.id = n;
        this.counter = tC;
        this.toAdd=toAdd;
        this.snooze_time = new Random();
    }

    public void run() {
        for (int i=0; i<toAdd; i++) {
            try { //sleep a random time to force interleavings
                sleep(snooze_time.nextInt(100));
                counter.increment(1);
            } catch (InterruptedException e) {
                e.printStackTrace();
            } // not actually doing anything 'coz lazy
        }
    }
}
```

```
package dataRace;
import java.util.concurrent.Semaphore; //for semaphores

public class TotalCounter {
    private int sum;
    private Semaphore mutex;

    TotalCounter() {
        sum=0;
        mutex=new Semaphore(1);
    }

    public void increment(int incrVal)
        throws InterruptedException {
        mutex.acquire();
        sum+=incrVal;
        mutex.release();
    }

    public int getValue() {
        return sum;
    }
}

*TotalCounter.java* 23L, 384C
```

- Semaphores protect shared data – guarantees that only one thread can increment at any specific time

```

package rendezvous;
import java.util.Random;
import java.util.concurrent.Semaphore;

public class RThread extends Thread {
    private Semaphore signal1;
    private Semaphore signal2;
    private Random snooze_time;
    private String name;

    RThread(Semaphore s1, Semaphore s2, String nm) {
        this.signal1 = s1;
        this.signal2 = s2;
        this.snooze_time = new Random();
        this.name = nm;
    }

    public void run() {
        try { //sleep a random time
            sleep(snooze_time.nextInt(100));
            System.out.println("Hello " + name);
            signal1.release();
            signal2.acquire();
            System.out.println("Goodbye " + name);
        }
        catch (InterruptedException ex) {
            "RThread.java" [noel] 28L, 689C
        }
    }
}

```

```

package rendezvous;
import java.util.concurrent.Semaphore; //for semaphores

public class STest {
    public static void main(String[] args) {
        Semaphore signal1 = new Semaphore(0); //start closed
        Semaphore signal2 = new Semaphore(0); //start closed

        //note signal order changed
        RThread threadA = new RThread(signal1, signal2, "Anne");
        RThread threadB = new RThread(signal2, signal1, "Bob");

        System.out.println("Starting simulation");
        threadA.start(); // start thread
        threadB.start(); // start thread
        System.out.println("Parent thread completed");
    }
}

```

"STest.java" 18L, 571C

- Rendezvous with semaphores
- Need to create two semaphores that're closed – that first instruction needs control when executed
- Need to swap S1 and S2 when passed in as parameter for one of the threads


```

package BarrierS;

public class BarrierTest { // class to run barrier simulation
    public static void main(String[] args) {
        int n = 4; //size of barrier
        int t = 6; // number of threads

        Barrier sharedBarrier = new Barrier(n);
        BThread[] threads = new BThread[t];
        for (int i=0; i<t; i++)
            threads[i]=new BThread(i, sharedBarrier); // call constructor

        System.out.println("Starting simulation with " + t + " threads, barrier size " + n);
        for (int i=0; i<t; i++)
            threads[i].start(); // start thread
        System.out.println("Parent thread completed");
    }
}

```

"BarrierTest.java" 18L, 564C

- Turnstile – wait, signal
- Works because initialise barrier to $-(n-1)$ – when -ve, nothing can progress, inhibit anything from passing the barrier before the semaphore becomes positive
 - Spins – but we can release
- Trick – start with -ve semaphore value, so every new thread coming is pushes up that value
 - Until we can pass through busy waiting loop
- BarrierTest – have to wait until a number of threads arrive at the barrier before others pass through
 - Once barrier is opened, all threads can pass through (even those that weren't explicitly waiting at the barrier)

```

package dishWashS;
import java.util.Random;

public class Washer extends Thread {
    private int dirty_dishes;
    private WetDishRack shared_rack;
    private Random washing_time;
    private int sleep;

    Washer(int n, WetDishRack the_rack, int sleep) {
        this.dirty_dishes = n;
        this.shared_rack = the_rack;
        this.washing_time = new Random();
        this.sleep=sleep;
    }

    public void run() {
        try {
            int counter=0;
            while (dirty_dishes>0) {
                counter++;
                System.out.println("---Washer is washing dish #" + counter);
                sleep(washing_time.nextInt(sleep)); //wait a bit
                shared_rack.addDish(counter);
                System.out.println("---Washer added dish #" + counter + " to rack.");
                dirty_dishes--;
            }
        }
    }
}

```

"Washer.java" 32L, 851C

```

dishWash$ vim Dryer.java -- 107x26
import java.util.Random;

public class Dryer extends Thread {
    private int dirty_dishes;
    private WetDishRack shared_rack;
    private Random drying_time;
    private int sleep;

    Dryer(int n, WetDishRack the_rack, int sleep) {
        this.dirty_dishes = n;
        this.shared_rack = the_rack;
        this.drying_time = new Random();
        this.sleep=sleep;
    }

    public void run() {
        try {
            int counter=0;
            while (dirty_dishes>0) {
                counter =shared_rack.removeDish();
                System.out.println("---Dryer removed dish #" +counter+" from rack");
                sleep(drying_time.nextInt(sleep)); //wait a bit
                System.out.println("---Dryer is done with dish #" +counter);
                dirty_dishes--;
            }
        }
    }
}

```

```

dishWash$ vim WetDishRack.java -- 107x26
package dishWash;
import java.util.concurrent.Semaphore;

public class WetDishRack {
    private Semaphore mutex; // can I access rack?
    private Semaphore empty; // does the rack have empty spaces?
    private Semaphore full; // does the rack contain dishes?
    private int[] rack;
    private int in, out;
    private int RACK_SIZE;

    WetDishRack(int rackSize) {
        this.mutex = new Semaphore(1); //
        this.empty = new Semaphore(rackSize); //
        this.full = new Semaphore(0); //
        this.rack = new int[rackSize]; //
        this.in = 0;
        this.out =0;
        this.RACK_SIZE=rackSize;
    }

    public void addDish(int dish_id) throws InterruptedException {
        empty.acquire(); //wait until empty space
        mutex.acquire(); //acquire rack
        rack[in]=dish_id; // add dish to rack
        full.release(); //
        in++;
    }

    public int removeDish() {
        full.acquire(); //
        mutex.acquire(); //
        int dish_id = rack[out];
        rack[out] = -1;
        empty.release(); //
        out++;
        return dish_id;
    }
}

```

- Dishwashing Sim
 - Producer is washer – produce dishes that go onto common dish rack (rack is the buffer)
 - Consumer is dryer – picks a dish from the rack to dry it
 - Here, have a limit to amount of dishes want to use
- Add randomness by having threads wait
- In run() – as washer, try and push dishes onto rack
 - addDish – ensures that don't try to insert wet dish onto rack with no space
- Dryer has similar behaviour
 - But removeDish instead of addDish
- Dishrack
 - Control access to rack
- No dryer removes the dish before its been added to the rack


```

package diningPhilosophers;
import java.util.concurrent.Semaphore; //for semaphores

public class Meal {

    public static void main(String[] args) {
        String[] names = {"Plato", "Aristotle", "Cicero", "Confucius", "Eratosthenes"};
        Semaphore[] chopsticks = new Semaphore[5];
        Philosopher[] philosopher = new Philosopher[5];

        for (int i = 0; i < chopsticks.length; i++)
            chopsticks[i] = new Semaphore(1);

        for (int i = 0; i < philosopher.length; i++) {
            philosopher[i] = new Philosopher(chopsticks[(i)%philosopher.length],
                                             chopsticks[(i+1)%philosopher.length], names[i]);
        }
        System.out.println("The meal begins!");
        System.out.println("=====");
        for (int i = 0; i < philosopher.length; i++)
            philosopher[i].start();
    }
}

"Meal.java" 24L, 886C

```

```

package diningPhilosophers;
import java.util.Random;
import java.util.concurrent.Semaphore; //for semaphores

public class Philosopher extends Thread {

    private Semaphore left_chopstick;
    private Semaphore right_chopstick;
    private String my_name;
    private Random snooze_time;

    Philosopher(Semaphore c_low, Semaphore c_high, String P_name) {
        this.left_chopstick = c_low;
        this.right_chopstick = c_high;
        this.my_name = P_name;
        this.snooze_time = new Random();
    }

    public void run() {
        while (true) {
            try {
                System.out.println(my_name + ": zzzzzzz " );
                sleep(snooze_time.nextInt(10)); //sleep a random time
                System.out.println(my_name + ": whatzat? " );
                eat();
            }
        }
    }
}

"Philosopher.java" [noel] 43L, 1341C

```

- Meal is main class
- Have semaphores, init to value 1 – regular mutex locks
 - Create philosophers with correct semaphore representing their l and r chopstick
- Philosopher class – have semaphores representing l and r chopsticks
 - Use sleep to have thread be suspended and another to be scheduled in
 - Larger chance of triggering sync problem if it exists
 - Eat() – routine that acquires and releases semaphores
 - Init wait for left chopstick – acquire blocks until avail
 - Then try to acquire right
 - Then eat
 - Then release left and right chopstick
- Classically deadlock occurs when they all try to acquire one chopstick at a time e.g. all pick up their left chopstick before their right
 -