

Motivation

Writing Lua applications...

- Used for fast prototyping of networked applications...
 - ...extended to control / sensor interaction
- Must be highly portable
- Running on embedded / low power hardware

Applications naturally grow up to need concurrence

- Performance: blocking on I/O is bad
- Some tasks are naturally concurrent: serializing them is a loss/loss

Concurrence

Preemptive

- Controlled by OS: Transparent for the applications
- Implemented at the hardware level
- Example: UNIX, *

Cooperative

- Controlled by apps: Explicit in the source code
- Implemented in user space
- Example: RISC OS, Windows 3.1

Concurrence

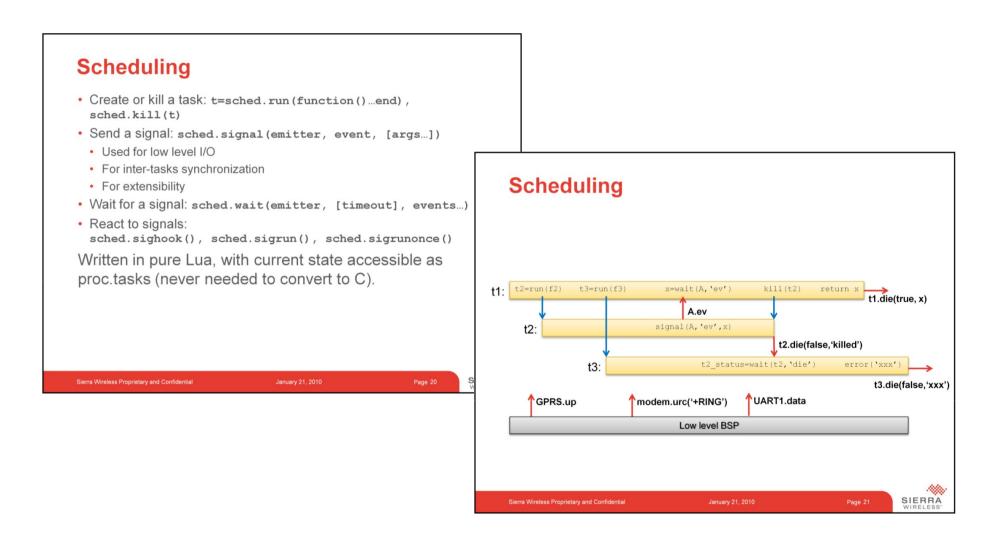
Message Passing

- "Processes"
- Examples: UNIX pipes, web services
- API model: serialization (e.g. JSON) + messaging (e.g. Sockets)
- Pros: robust, easily distributable
- Cons: serialization&messaging overhead, hard to design

Shared State

- "Threads"
- Examples: pthreads, [obligatorios del curso de redes]
- API model: access control (e.g. semaphors)
- Pros: no data copying
- Cons: high contact surface, error prone, hard to track bugs (when preemptive)

Sierra Wireless luasched



Luasched vs Lumen

- Sierra Wireless
- Enterprise grade
 M2M development
 and tooling
- Modified Lua + modified luasocket
- Tooling modules in C (shell, LTN12)

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- "API changes and other random breakages occur"
- Plain lua + luasocket
 & nixio backends
- All tooling in pure Lua+lumen sched.

Lumen scheduler

- Inspired in Sierra Wireless luasched.
- Cooperative
- Message passing and shared memory supported
 - Messaging trough signals and pipes/streams
 - Access control trough critical section marking
- Pure Lua, no external dependencies: runs on Lua 5.1, 5.2, 5.3 LuaJIT
- Implemented using coroutines
- Runs in a single VM instance (state), available as a Lua module
 local sched = require 'sched'
- MIT licence

Lumen vs Others

- Pros
- Lightweight
- Pure Lua: Hardware & OS-independent
- Simple usage model
- Well suited for I/O bound logic
- Tools available: logging, remote shell, webserver...
- Cons
- Single Lua state single core.
 (Multiple States can interchange messages, tough)
- Slow context switching (compared to preemptive), not well suited to CPU-bound logic

Lumen scheduler

```
local sched=require 'sched'
-- task emits signals
sched.run(function()
   for i = 1, 10 do
       sched.signal('an_event', i)
       sched.sleep(1)
   end
end)
-- task receives signals
sched.run(function()
   local waitd = {'an event'}
   while true do
      local ev, data = sched.wait(waitd)
       print (data)
   end
end)
sched.go()
```

Lumen tasks

Create and run a task

```
local f = function(...)
    -- blabla
end
local taskd = sched.run( f, ... )
```

Create a task, run later

```
local taskd = sched.new_task( f )
taskd:run(...) --start
```

Current task

```
local my_taskd = sched.running_task
```

Tasks control

Pause tasks

```
taskd:set_pause(true) --pause
taskd:set_pause(false) --unpause
```

Yielding and waiting

```
taskd:wait()
taskd:sleep(1.5)
```

Finishing a task == finishing the function

```
local f = function(...)
   error ... --here
   return ... --here
end --here
```

Killing a task

```
taskd:kill()
```

Lumen signals - emitting

- Tasks emit signals
- A signals is a event, plus parameters
 - The event can be of any type
 - Parameters in any number, of any type
- Tasks should emit signals for everything that could be interesting to other tasks
- Examples:

```
sched.signal( 'new client', skt )
local EVENT_DIE = {}
sched.signal(EVENT_DIE, true, 10, 10)
sched.signal(EVENT_DIE, false, 'div by 0')
```

Lumen signals - receiving

- Tasks can wait for signals. Waiting for a signal is the main method for synchronization.
- Wait descriptors define a list of event(s) (can be '*')
- The wait call returns the signal's parameters
- Examples:

```
local waitd = {
    EVENT_DIE,
    EVENT_FINISH,
    timeout = 10 --optional
}
local event, p1, p2, p3 = sched.wait(waitd)
```

Lumen signals - receiving

 There are some useful shortcuts for handling signals in a event-like manner

```
sched.sigrun(waitd, f)
sched.new_sigrun_task(waitd, f)
sched.sigrunonce(waitd, f)
sched.new sigrunonce task(waitd, f)
```

They are implemented in plain wait calls

```
sched.new_sigrun_task = function ( waitd, f )
    return M.new_task (function()
        while true do
        f(M.wait(waitd))
    end
    end)
end
```

Lumen signals - buffering

 Signals are like UDP: if nobody is waiting for them, they get lost.

```
while true do
    sched.wait(waitd1)
    --do something
    sched.wait(waitd2) --events for waitd1 can get missed
    --do something
end
```

Sometimes this is the right thing. When not:

```
local waitd = sched.new_waitd({
    '*',
    buff_mode='drop_last', --or 'drop_first'
})
```

• Still, this is not enough sometimes (faster emitter than receiver).

Signal alternatives: pipes & streams

Signals never block the emitter. Not useful on a producer – consumer scheme.

Pipes

```
local pipes=require 'pipes'
local p=pipes.new(10)

--one task
p:write(a, b)

--other task
local ok,a,b=p:read()
```

Streams

```
local stream=require 'stream'
local s=stream.new(1000)

--one task
s:write('####')

--other task
local text=s:read()
```

- Implemented using plain signals
- Pipes and streams allow timeouts

Resource sharing and access

 A catalog allows to publish objects under a well known name, and wait for them to appear.

```
local catalogs = require 'catalog'
local tasks = catalogs.get_catalog('tasks')

--in one task
tasks:register('consumer', taskd)

--in another task
local consumer = tasks:waitfor('consumer', 60)
```

Some ideas for catalogs are 'events', 'pipes', 'devices'

Mutexes

Needed only seldom, because scheduling is non-preemptive: only needed if task yields control explicitly inside critical section.

```
local mutex = require 'mutex'
local n = 0
                                         Other method:
local critical = function ()
                                             local function critical()
   print( 'before', n )
                                               mx:acquire()
   n = n + 1
                                               -- do stuff
                                              mx:release()
   sched.wait()
                                             end
   n = n - 1
   print( 'after', n )
end
local mx = mutex.new()
local synched = mx:synchronized(critical)
```

I/O: selector module

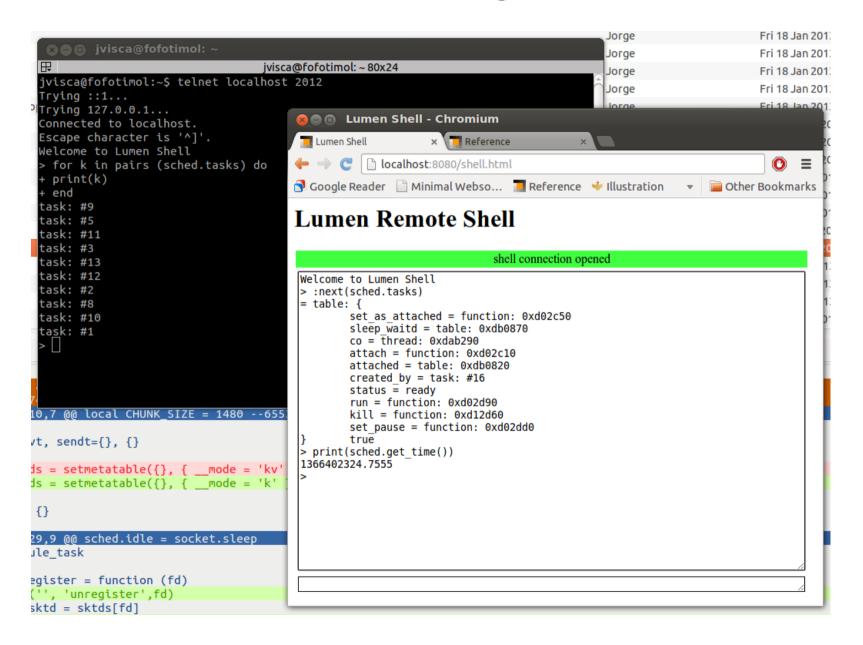
- Basic POSIX support
 - Sockets (TCP&UDP)
 - Files
 - Piping-in output from programs
- Data can be obtained in several ways, depending on its nature
 - Signals on data arrival
 - Callback to a provided function on data arrival
 - Stream
- Sending can be synchronous or asynchronous.
- Nixio and luasocket as backends
 local selector = require 'tasks/selector'
 selector.init({service='nixio'})

I/O: selector module

 Send a UDP Packet local udpsend = selector.new_udp("127.0.0.1", 8888) udpsend:send('data!') Start a TCP Server, a handler read lines local tcp_server = selector.new_tcp_server("127.0.0.1", 8888, 'line',, function(sktd, data, err) print ('arrived:', sktd, data) return true end) Receive UDP trough signals local udprecv = selector.new_udp(nil, nil, "127.0.0.1", 8888) sched.sigrun({udprecv.events.data, timeout=60}, function(_, _, ...) print("arrived:", ...) end Read from file trough a stream local fs = stream.new(buffer_size) local fd, err = selector.new_fd ('/dev/tty0', {"rdonly"}, nil, fs)

while true do print(fs:read()) end

Tooling



References

• Lumen sources:

https://github.com/xopxe/Lumen

Lumen API:

http://xopxe.github.com/Lumen/

End