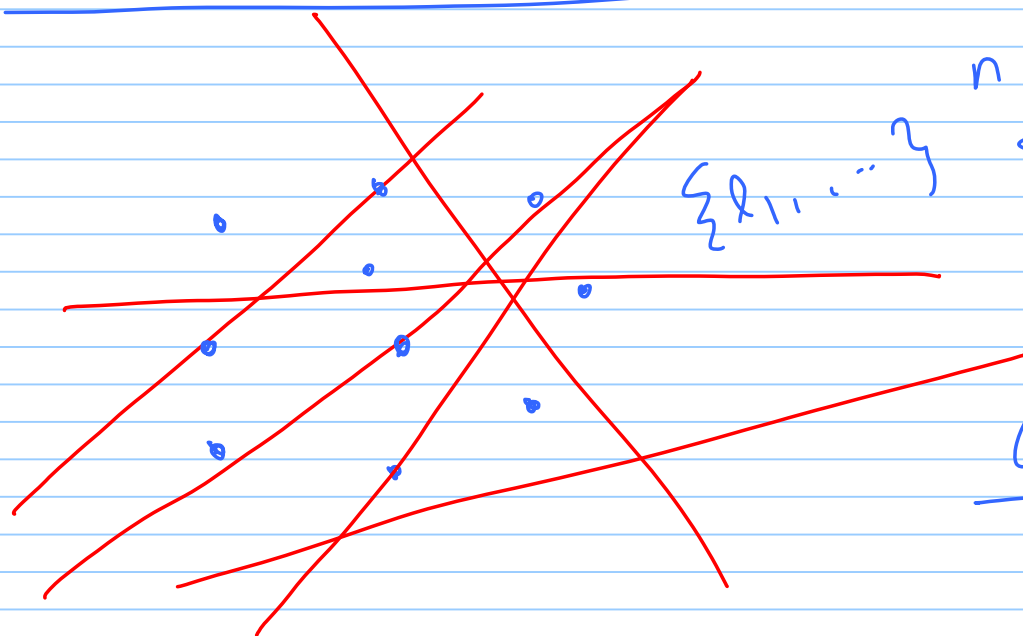


Incidence Geometry

Point line incidences



$m = |P| \leftarrow \{p_1, \dots, p_m\}$ points in \mathbb{R}^2
 $n = |L| \leftarrow$ lines in \mathbb{R}^2

$$I(P, L) = \{(p_i, l_j) \mid \begin{array}{l} p_i \in P, \\ l_j \in L, \\ p_i \in l_j \end{array}\}$$

Q Estimate $I(P, L)$?

Q2 $A \subseteq \mathbb{R}$

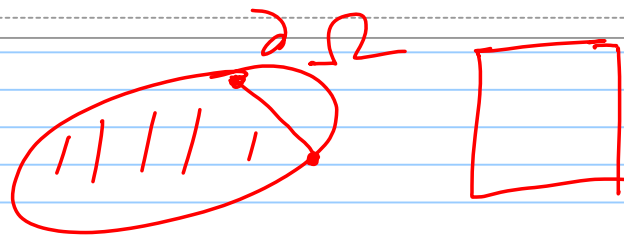
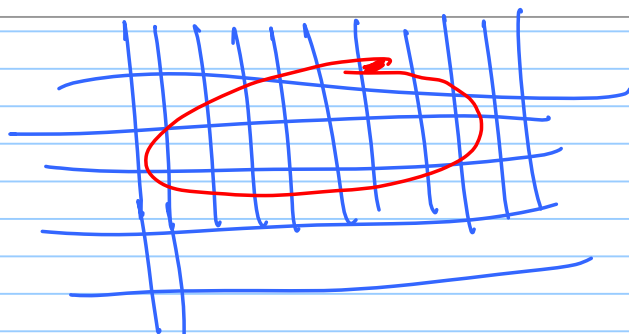
$$\boxed{A+A} = \{a+b \mid a, b \in A\} \quad \swarrow$$

$$\boxed{A \cdot A} = \{ab \mid a, b \in A\} \quad \swarrow$$

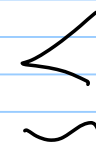
Estimate size $|A+A|$, $|A \cdot A|$ in terms of size of $|A|$.

Q3 Let P is a set of n points in \mathbb{R}^2 . Estimate no. of unit area \triangle that can be formed with vertices from P .

Q4

 \mathbb{Z}^2 

$\partial\Omega \cap \mathbb{Z}^2$ vs $\Omega \cap \mathbb{Z}^2$
Isoperimetric problems.

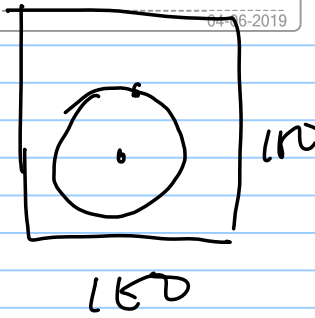


Thm 1 (Szemerédi - Trotter Thm 1983)

$$|I(P, L)| \leq c(|P||L|)^{2/3} + |P| + |L| \quad \checkmark$$

$L \leftarrow$ circles.

• Erdős Unit Distance Problem (1946)



Q How many unit distance pairs can there be in a set of " n " points in \mathbb{R}^2 ?

Trivial Estimate

Thm 2 (Erdős 1946)

$$\frac{n^{1 + \frac{1}{\log \log n}}}{n} \leq u(n) \leq n^{3/2}$$

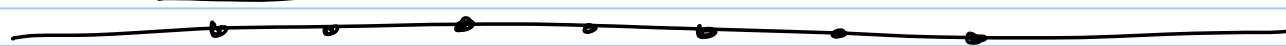
$$u(n) \sim n \times \sqrt{2 \frac{\log \log n}{\log n}}$$

Thm 3 (Spencer - Szemerédi - Trotter 84)

$$\underline{u(n)} \leq n^{4/3}$$

← ST theorem

$$\underline{O(n)}$$

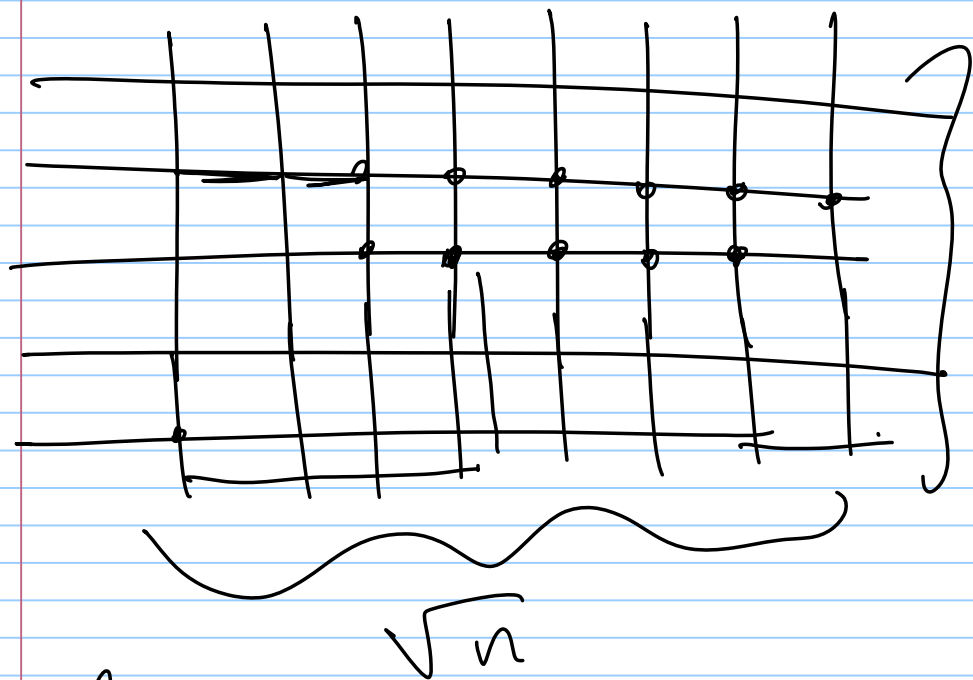
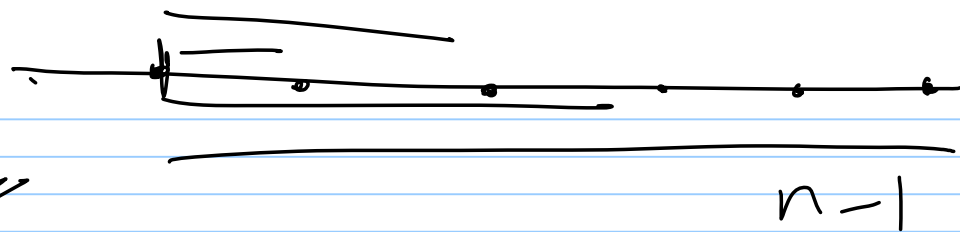


• Erdős Distinct distance problem (1946)

Q Minimum number of distinct distances that
can be determined by "n" points in \mathbb{R}^2 ?

$d(n)$ ← denotes this number.

Q $d(n) \leq n-1$



$$T = \sum d(p_i, p_j)^2 \mid p_i, p_j \in S$$

Thm (Landau-Ramseyjan)

$$\Theta\left(\frac{n}{\sqrt{\log n}}\right)$$

Q

Thm (Erdős 1946) $d(n) = \underline{\underline{O\left(\frac{n}{\sqrt{\log n}}\right)}}$

Conjectured $d(n) \geq C \frac{n}{\sqrt{\log n}}$

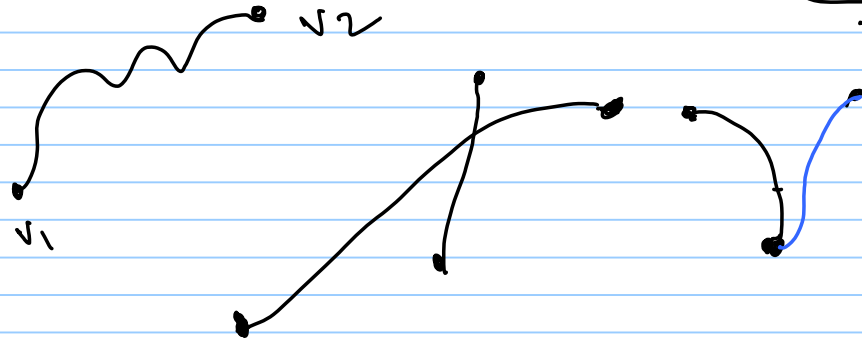
$$f(n) = \Theta(t(n))$$
$$c_1 t(n) \leq f(n) \leq c_2 t(n)$$

Thm (Guth-Katz 2010) $d(n) = \Omega\left(\frac{n}{\log n}\right)$

Proved S T Theorem

Planar graphs

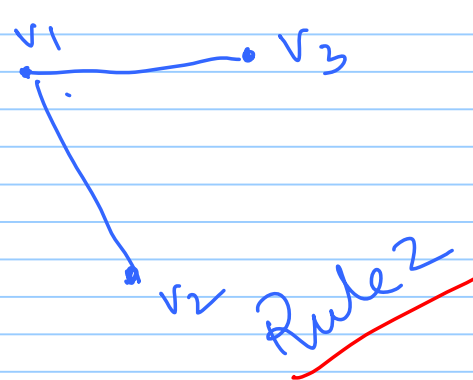
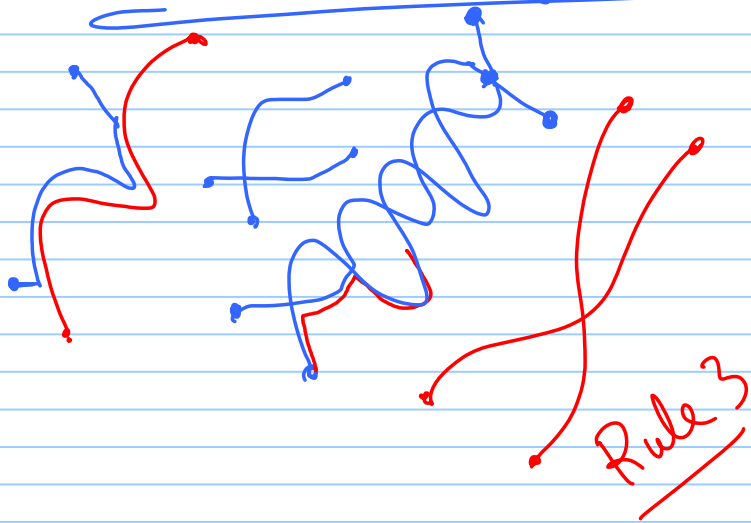
$G := (V, E)$
 $\searrow \{v_1, v_2\}$



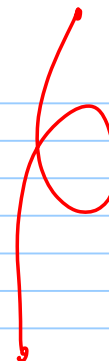
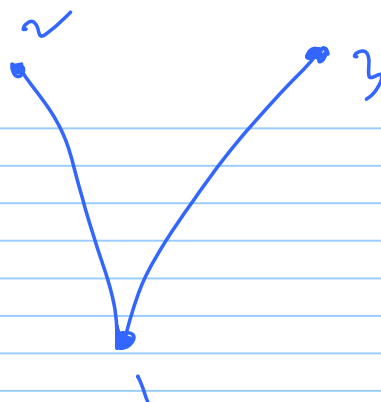
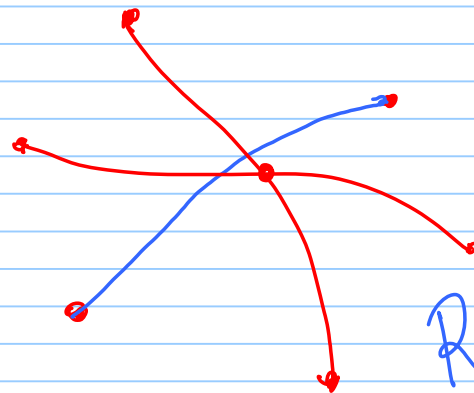
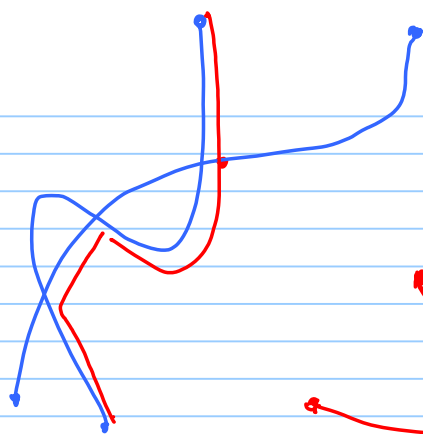
Lemma 1 (Bounds on the # of edges in planar graph) If G is planar then $|E| \leq 3|V| - 6$.

Crossing no. of Graph $cr(G)$

$n \log^4 n$



$\{v_1, v_2\}$
 $\{v_1, v_3\}$



Rule 4

$$\textcircled{\text{Gr}(4) - 1}$$

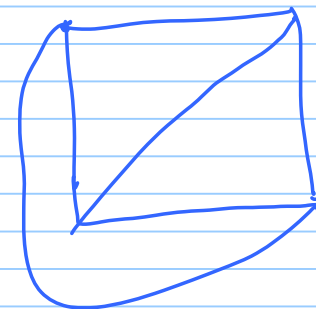
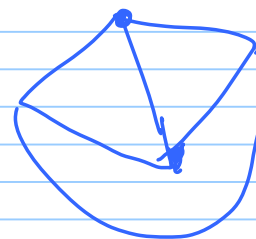
Lemma 2 (Loose crossing bound) $-6 + 6 \checkmark$ $|V| \geq 3$

$$Cr(G) \geq |E| - 3|V| + 6 \checkmark$$

$$= (3V - 6)$$

K_5 \cap $Cr(K_5) = 1$

≥ 1



Proof $G := (V, E)$ $Cr(G)$

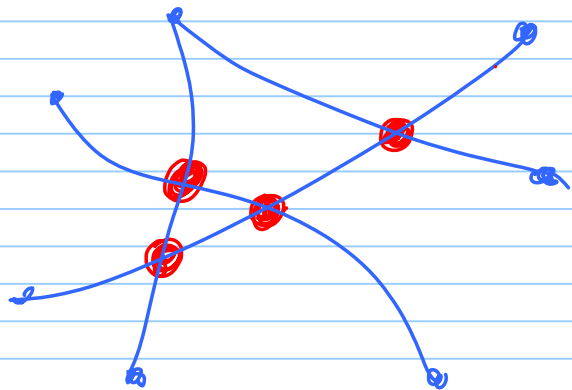
Let $E' \subseteq E$ be a maximal set of edges that does not cross in the drawing.

$$|E| - |E \setminus E'| \leq 3|V| - 6$$
$$G := (V, E') \quad |E'| \leq 3|V| - 6$$

Looking at edges in $E \setminus E'$ $Cr(G) \geq \underline{\underline{|E \setminus E'|}}$

$$Cr(G) \geq |E \setminus E'| \geq |E| - 3|V| + 6$$

□



$$(2^{n-1}) \times (3^{n-1})$$

$$G \rightarrow H \begin{cases} |V| + Gr(u) \\ |E| + 2 Gr(u) \end{cases}$$

$$|E| \geq 4|V| \quad \square$$

$$Gr(u) \geq |E| - 3|V| + 6$$

$$\Omega\left(\frac{|E|^3}{|V|^2}\right)$$

Thm (Crossing Lemma) $G := (V, E)$ If $|E| \geq 4|V|$, then

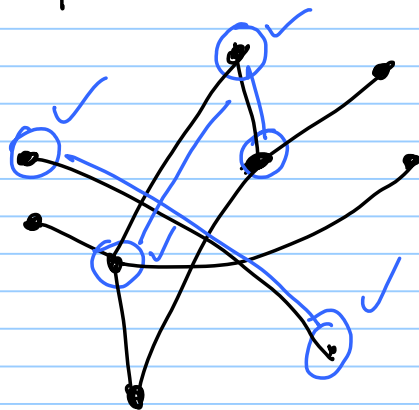
$$Cr(G) = \Omega\left(\frac{|E|^3}{|V|^2}\right) \quad \left[\begin{array}{l} |V| - |E| + |F| = 2 \\ 2|E| \geq 3|F| \end{array} \right]$$

Proof $G = (V, E)$, drawn with exactly $Cr(G)$ many crossings. Let $0 \leq p \leq 1$,

$V_p = \{ \text{set of vertices selected from } V, \text{ where each vertex is picked with probability } p \}$ @ $G_p := (V_p, E_p)$

$$\mathbb{E}[|V_p|] = np \quad | \quad \mathbb{E}[|E_p|] = |E| p^2$$

$$G_p := (V_p, E_p)$$



$$\mathbb{E}[C'(G_p)] = Cr(u) p^4$$

$$Cr(G_p) \geq |E_p| - 3|V_p|$$

$$C'(G_p) \geq |E_p| - 3|V_p|$$

$$\mathbb{E}[C'(G_p) - |E_p| - 3|V_p|] \geq 0$$

$$\mathbb{E}[C'(G_p)] - \mathbb{E}[|E_p|] - 3\mathbb{E}[|V_p|] \geq 0$$

$$\mathbb{E}[c'(G_p)] \geq \mathbb{E}[|E_p|] - 3 \mathbb{E}[|V_p|]$$

$$\boxed{Cr(G) \geq \left(\frac{|E|}{p^2} - \frac{3|V|}{p^3} \right)} \quad \leftarrow \quad p = \frac{4|V|}{|E|}$$

$$\rightarrow \geq \Omega(|E|^3 / |V|^2) \quad \square$$

Thm (ST Theorem) $I(P, L) \leq \left(|P||L| \right)^{2/3} + |P| + |L|$

$$(mn)^{2/3} + m + n$$

Proof $L = \{l_1, \dots, l_n\} / I(P, L)$
 $P = \{p_1, \dots, p_m\} \stackrel{\text{def}}{=} \{(p, l) \mid \begin{matrix} p \in P, l \in L \\ p \in l \end{matrix}\}$

$L' =$ those lines that contains at least 2 points

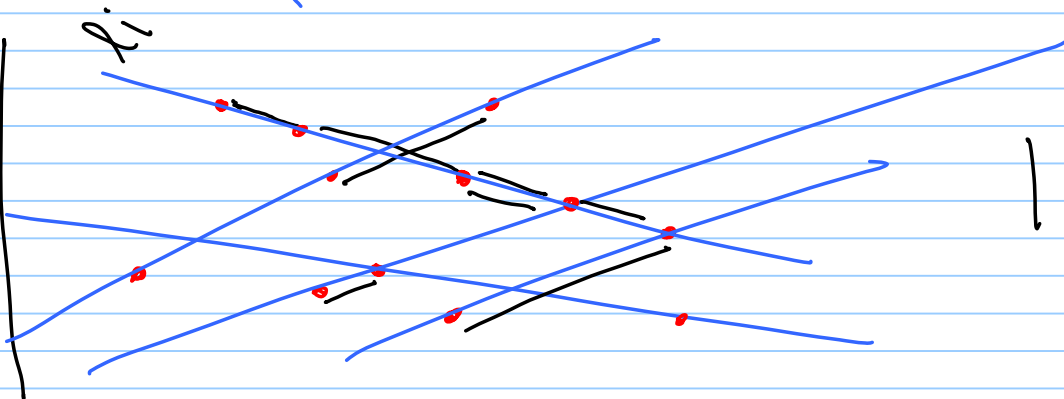
$$I(P, L) \leq \underline{I(P, L')} + n$$

$$l_i \leftarrow m_i$$

$$\sum m_i = I(P, L)$$

$$\begin{aligned} |E| &= \sum_{i=1}^n (m_i - 1) \\ &= \underline{I(P, L) - n} \end{aligned}$$

$$\begin{aligned} |E| &\leq 4m \\ I(P, L) &< 4m + n \end{aligned}$$



$$(2) |E| \geq 4m$$

$$n^2 \geq Cr(u) \geq \frac{|E|^3}{|V|^2} \stackrel{f \leq t^g}{=} \frac{(I(P,L) - n)^3}{m^2} \stackrel{f \leq g + \frac{t}{3}}{=}$$

$$\frac{(I(P,L) - n)^3}{m^2} \leq n^2$$

$$I(P,L) \leq m + n$$

$$\Rightarrow \boxed{I(P,L) \leq (nm)^{2/3} + n}$$



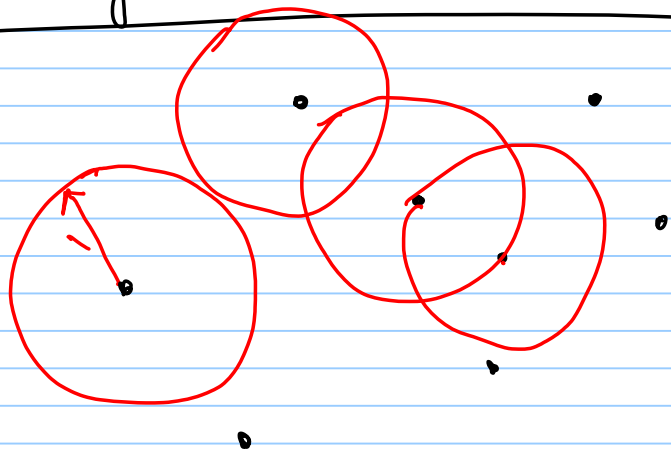
Thm (Point Circle Incidences) $L \leftarrow$ set of circles in plane

$$\underline{I(P, L)} \leq (mn)^{2/3} + m + n$$

no. of unit

$P \leftarrow$ set of m points in the plane

Revisiting Erdős ~~distinct~~ distance problem



$P \leftarrow n$ circles

$L \leftarrow n$ points

$$\underline{I(P, L)} \leq n^{4/3}$$

$$v(n) = \frac{I(P, L)}{2} \leq n^{4/3}$$

No. of unit triangle problem

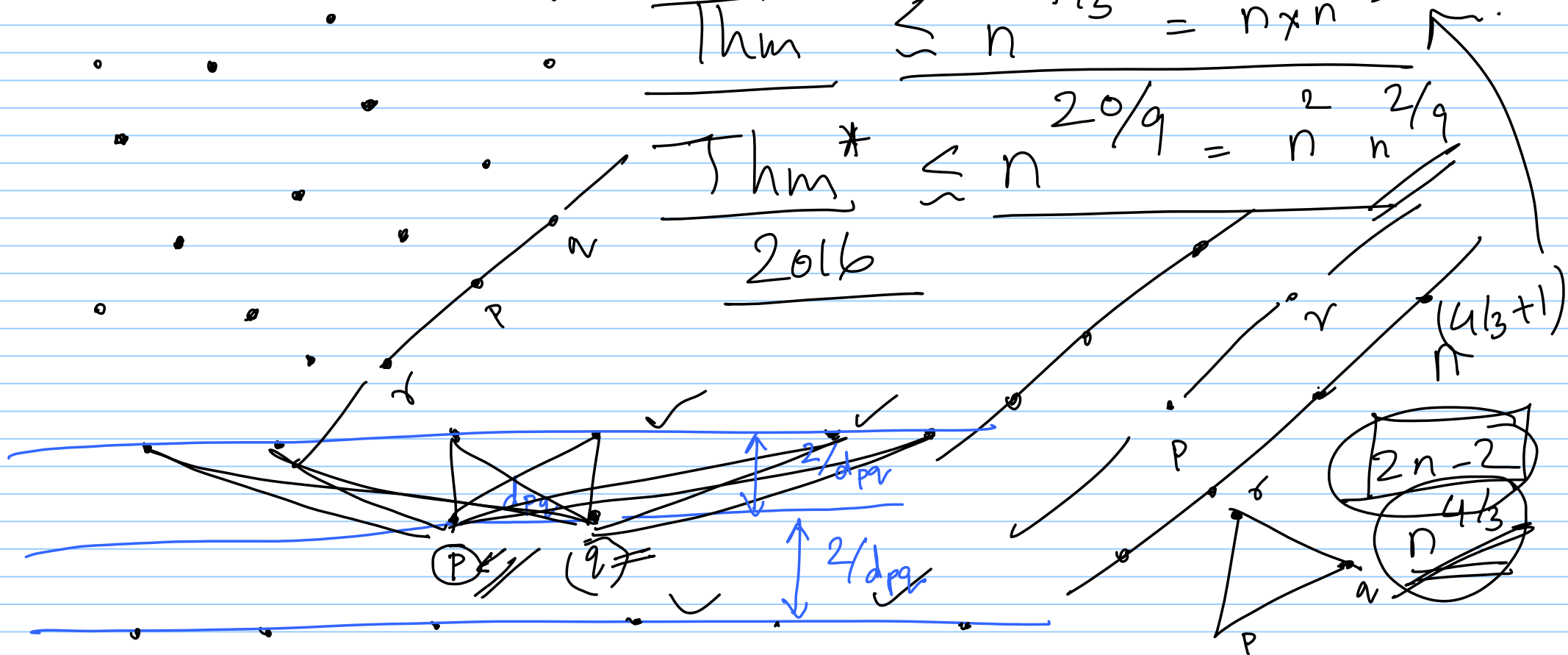
Thm

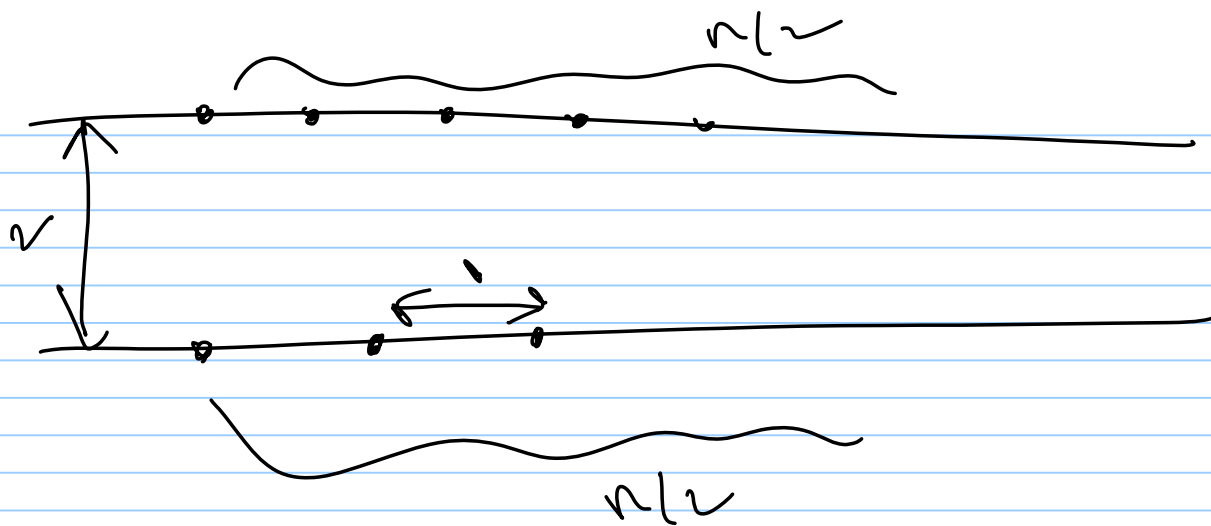
$$\lesssim n^{7/3} = n^2 n^{1/3}$$

Thm*

$$\lesssim n^{20/9} = n^2 n^{2/9}$$

2016





n^2
 $n^2 \log n$
 Erdős-Purdy
 $20/9$
 n