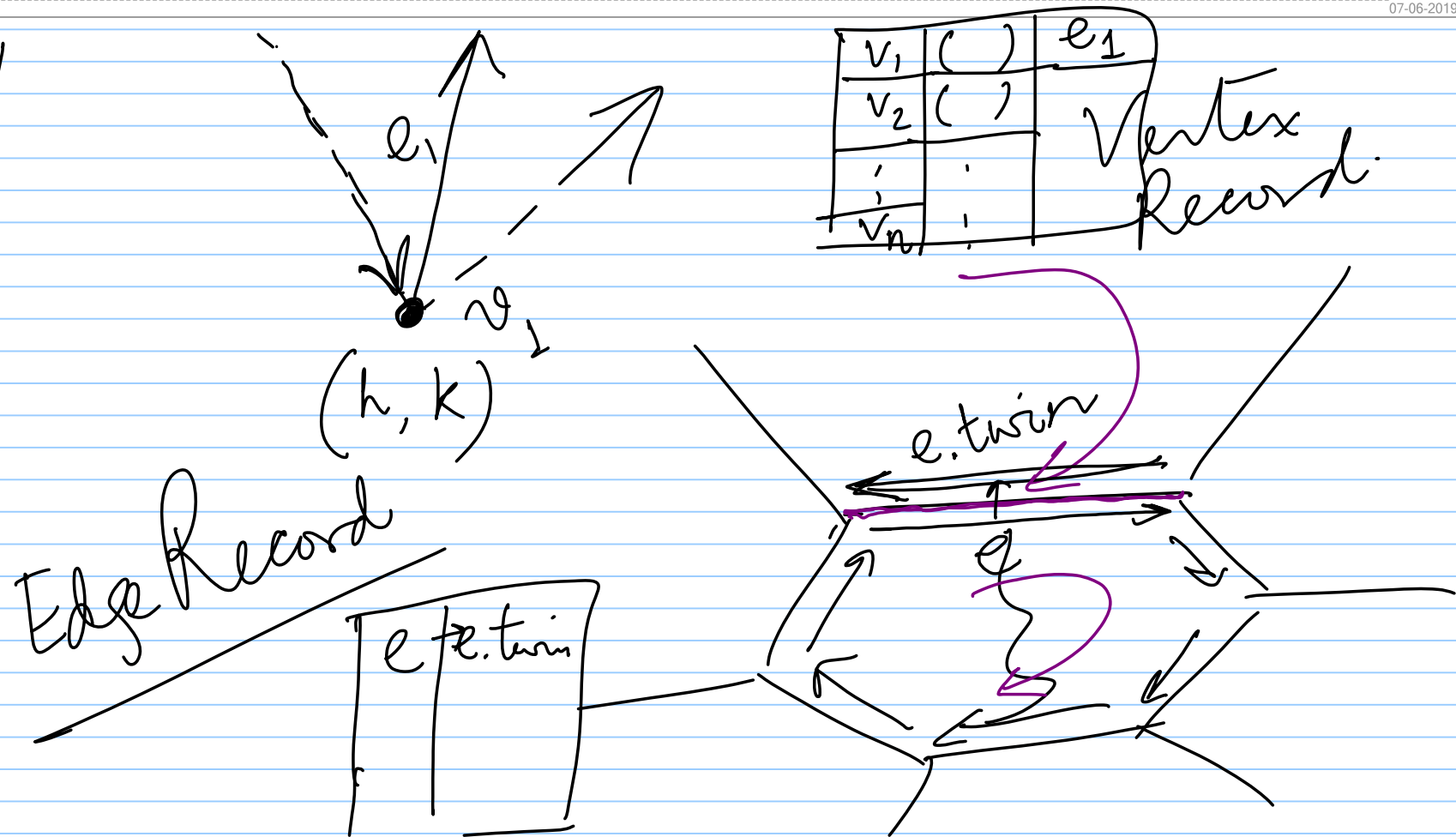


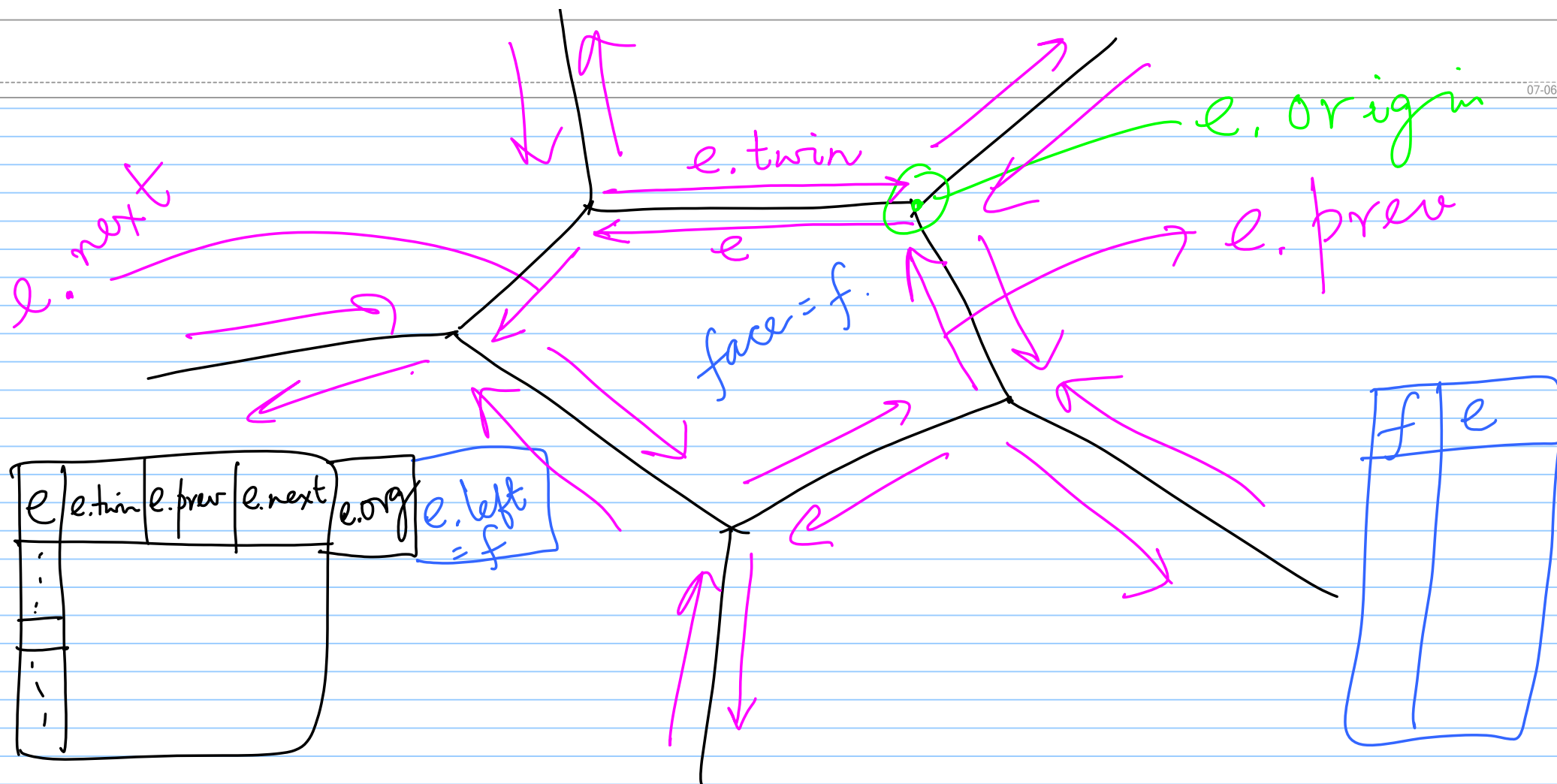
Point location in a planar subdivision

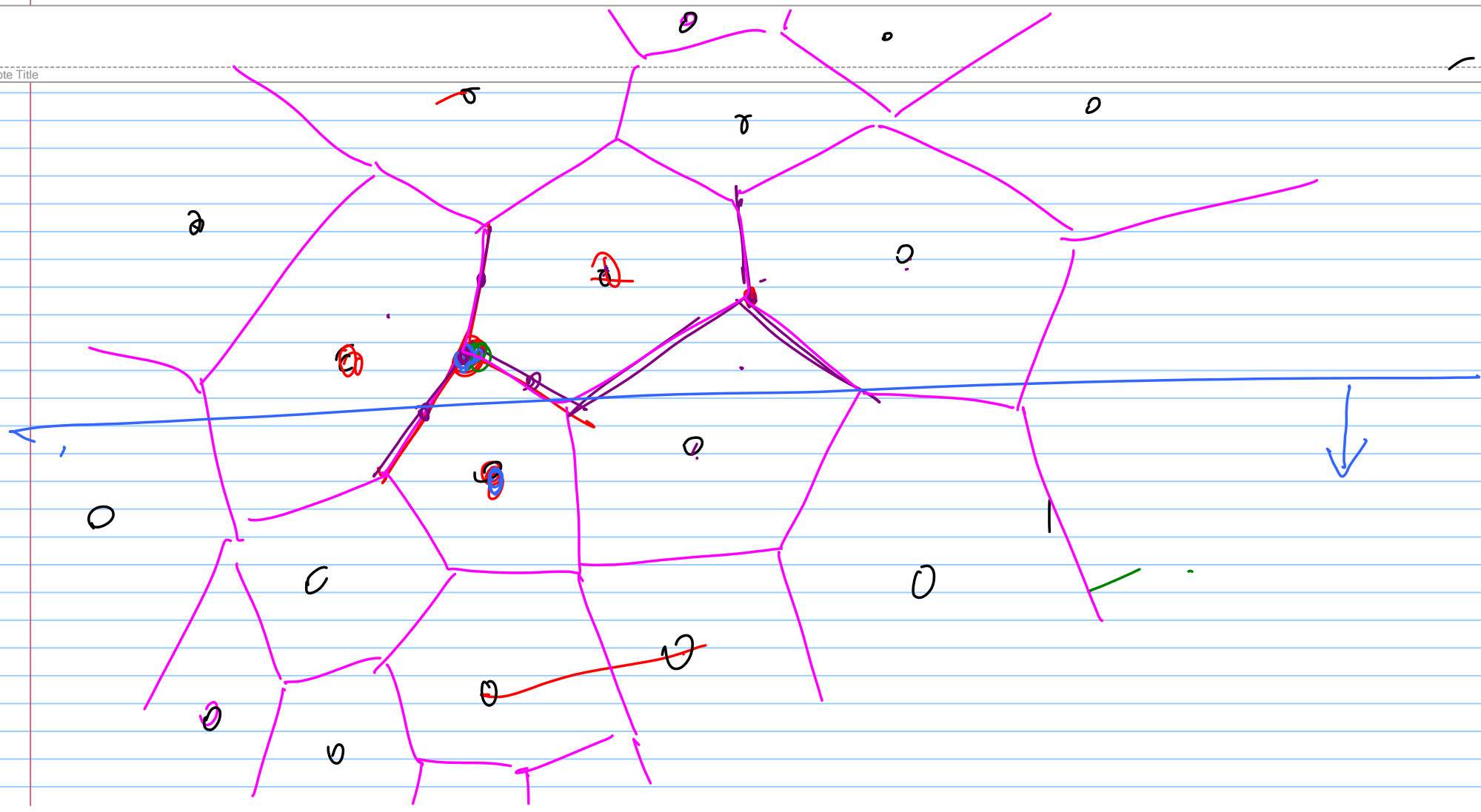
PSLG: Planar straight line graph

DCEL: Doubly Connected  
Edge List

- Keep records for
- 1) Vertices
  - 2) Edge
  - 3) Faces.







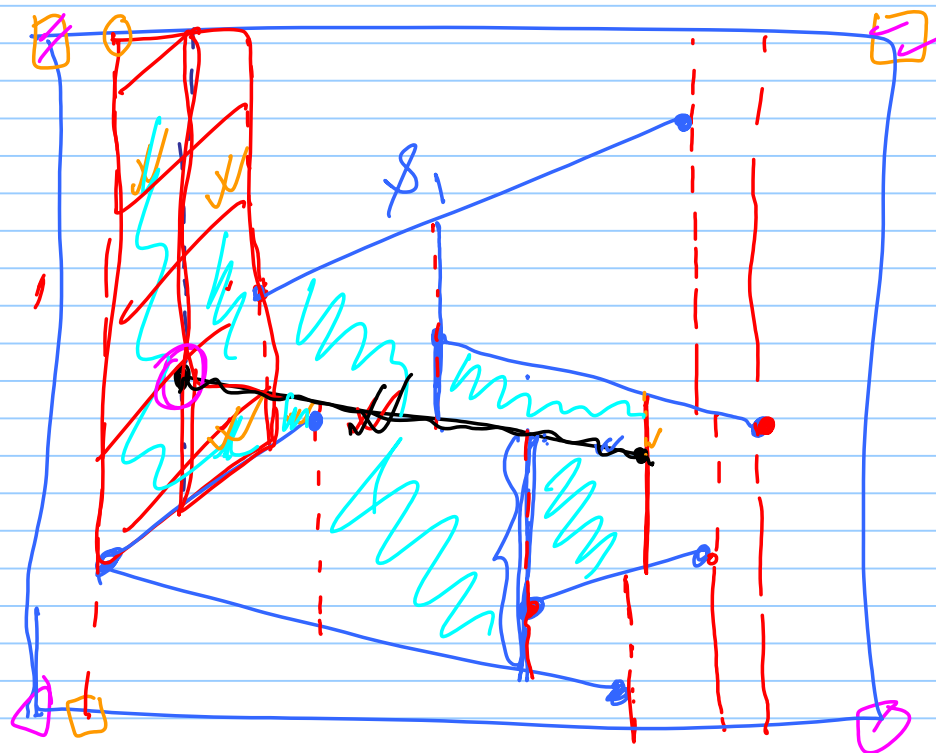
# Trapezoidal decomposition

$$6n+4$$

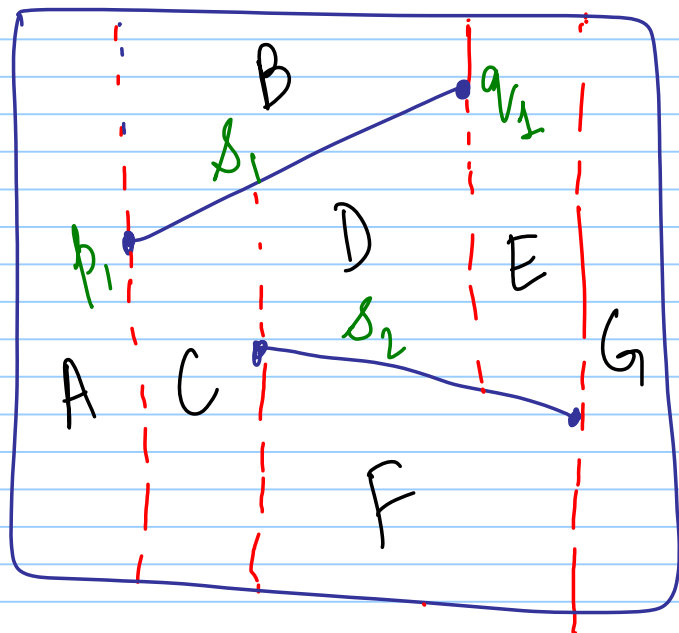
Lemma: # vertices  $\leq$

# trapezoids  $\leq$  ?

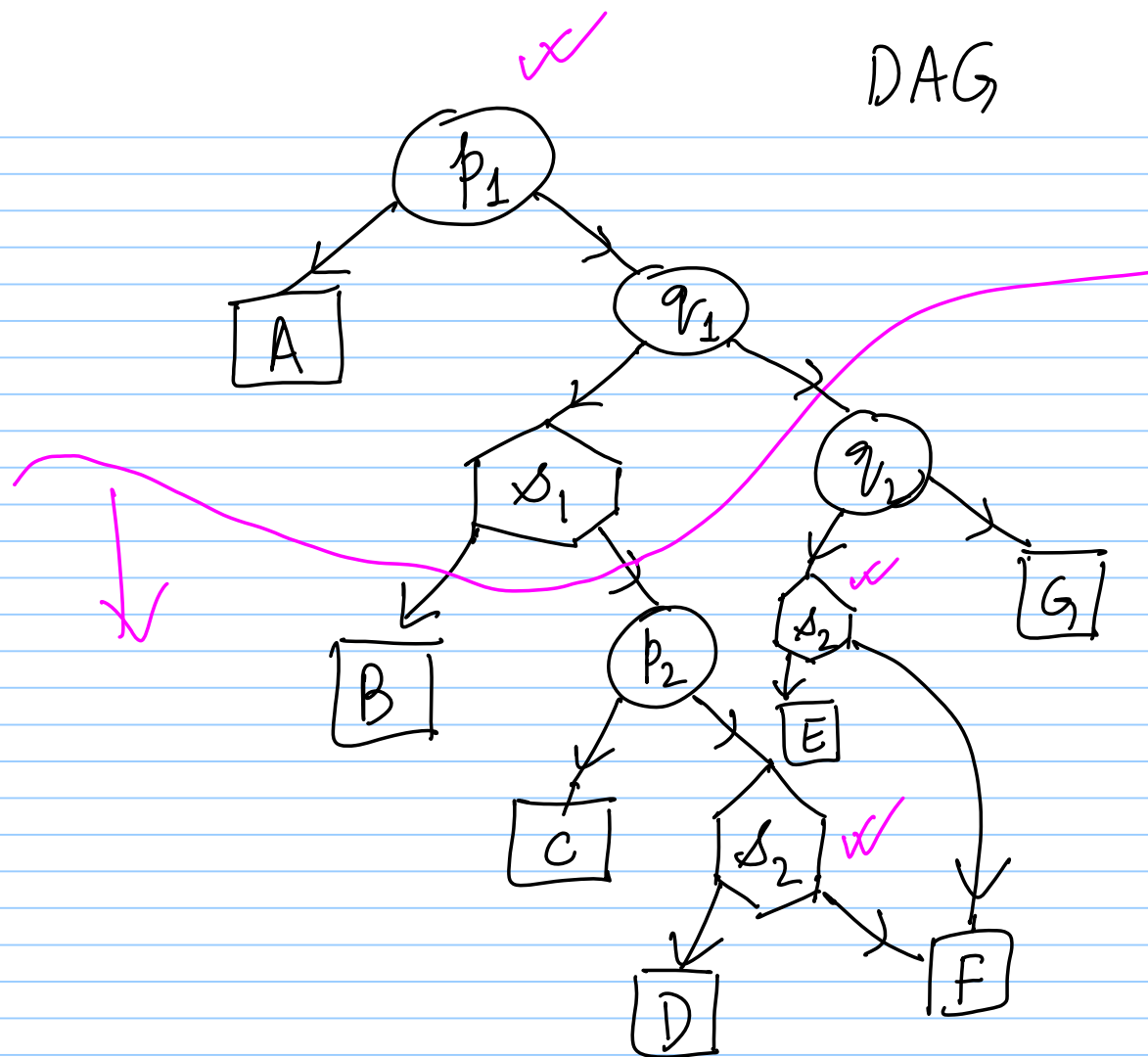
$$3n+1$$



~~$O(n)$~~



$$|\pi| = O(n)$$



Claim: Ignoring the time spent to locate the left endpoint of a segment, time that it takes to insert the  $i^{\text{th}}$  segment & update the trapezoidal map is  $O(k_i)$  where  $k_i$ : # newly created trapezoids.

Proof (idea): For each bullet path crossed, trim the bullet and create a new face.

$$E[k_i] = ? O(1)$$

Lemma: Consider the RIC of a trapezoidal map, and let

$k_i$  : # new trapezoids created when the  $i^{\text{th}}$  segment is added.

Then  $E[k_i] = O(1)$ , where the expectation is taken over all permutations of the segments

Proof: Let  $\mathcal{T}_i$ : trapezoidal map after  $i$  insertions.  $|\mathcal{T}_i| = O(i)$

Let's say that a trapezoid  $\Delta$  depends on a segment  $\overline{s}$ ,

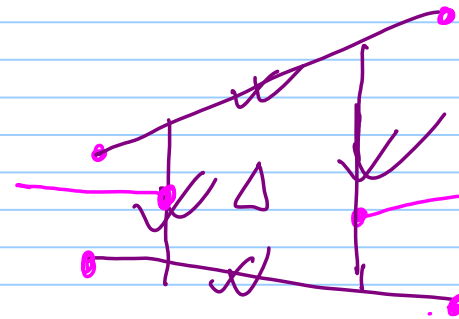
if  $\overline{s}$  would have caused  $\Delta$  to be created, had  $\overline{s}$  been added last.



$$\delta(\Delta, \delta) = \begin{cases} 1 & \text{if } \Delta \text{ depends on } \delta \\ 0 & \text{o.w.} \end{cases} \quad || \quad S_i = \{\delta_1, \dots, \delta_i\}$$

$$E[R_i] = \frac{1}{i} \sum_{\delta \in S_i} \sum_{\Delta \in \mathcal{T}_i} \delta(\Delta, \delta)$$

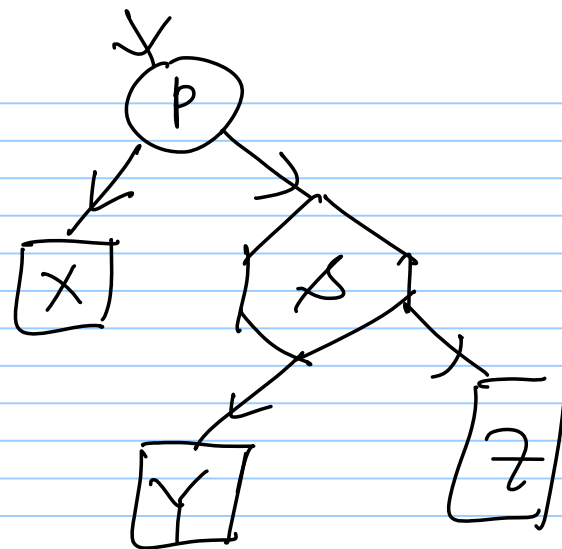
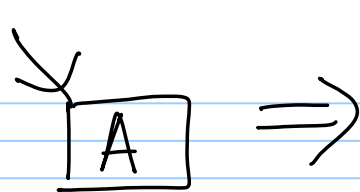
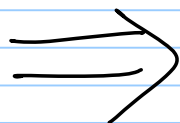
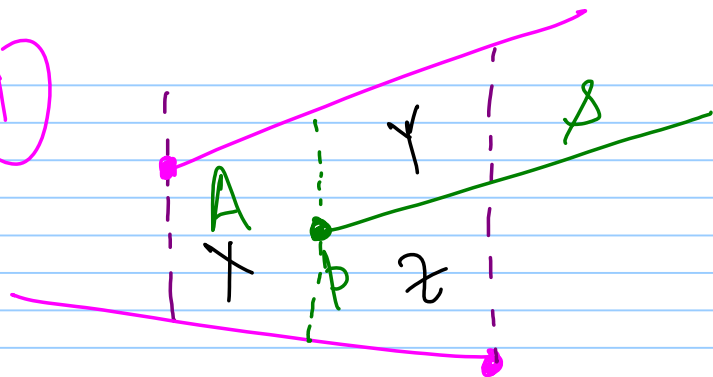
$$= \frac{1}{i} \sum_{\Delta \in \mathcal{T}_i} \sum_{\delta \in S_i} \delta(\Delta, \delta)$$

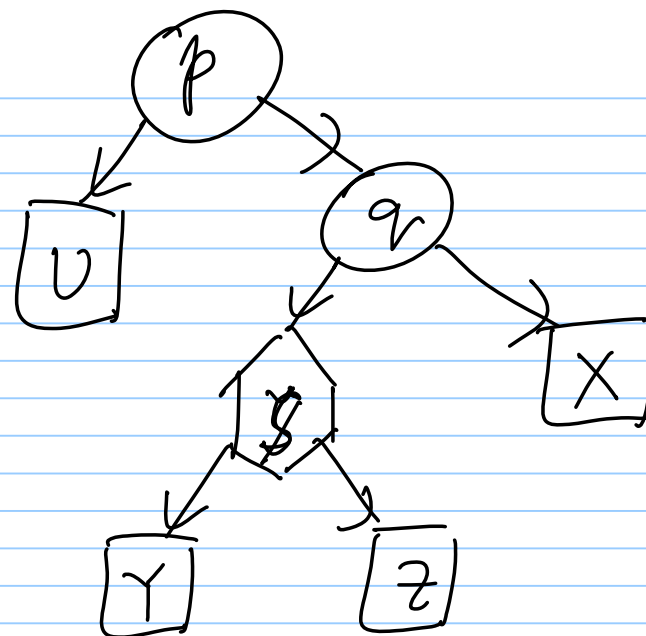
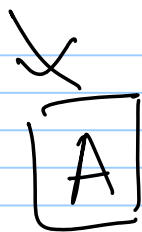
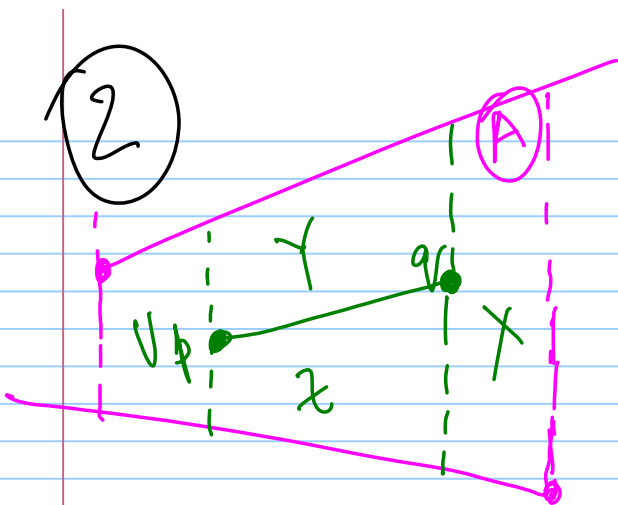


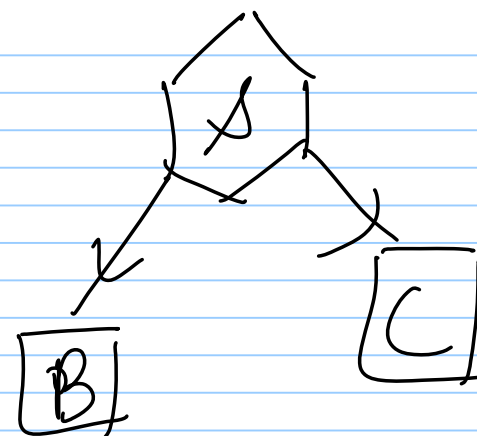
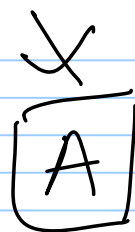
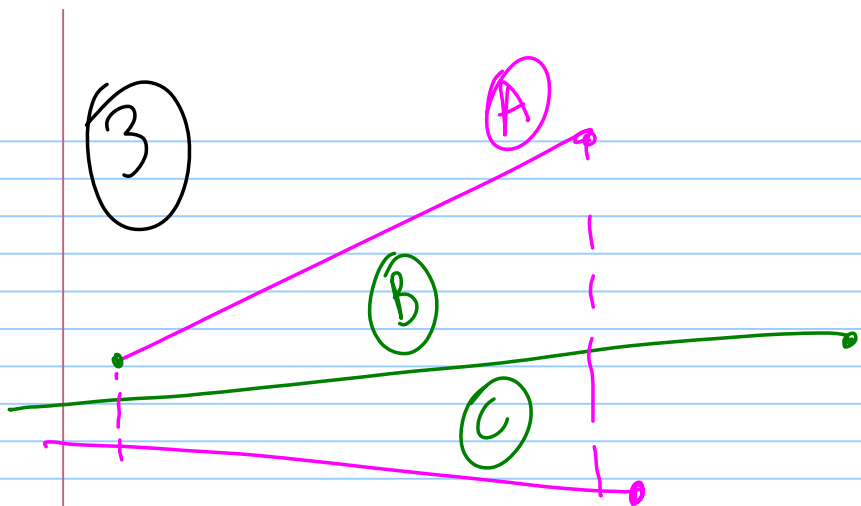
$$= \frac{1}{i} \sum_{\Delta \in \tau_i} 4 = \frac{1}{i} 4 \cdot |\tau_i| = O(1)$$

Q: How big can the data structure grow?  
 $\rightarrow O(n)$

①







Query Analysis: Consider a query point  $q$  that has been fixed by an adversary w/o the knowledge of the random order in which we insert the lines.

$P_i$ : prob that the trapezoid that contains  $q$  changes as a result of the  $i^{\text{th}}$  insertion.

Expected length of  $q_i$ 's search path in the final structure  
is at most  $\sum_{i=1}^n 3 P_i \leq 3 \sum_{i=1}^n \frac{4}{i}$

$$P_i = ? \frac{4}{i}$$

Argue using backward analysis

$$\approx 12 \log_2 n$$

CAVEAT !!

$$\Pr(|X(q) - E[X(q)]|) \leq ?$$