

# Theoretical Computer Science - Compact

Janis Hutz  
<https://janishutz.com>

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*“Sie können also alle C Programme in Kanonischer Ordnung aufzählen. Sollten Sie dies tun. Wahrscheinlich nicht. Was aber zählt ist, sie **können** es tun”*

- Prof. Dr. Dennis Komm, 2025

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Compact Summary of the book Theoretische Informatik

by Prof. Dr. Juraj Hromkovic

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1 Introduction

This summary aims to provide a simple, easy to understand and short overview over the topics covered, with approaches for proofs, important theorems and lemmas, as well as definitions.

It does not aim to serve as a full replacement for the book or my main summary, but as a supplement to both of them.

It also lacks some formalism and is only intended to give some intuition, six pages are really not enough for a formal and complete overview of the topic.

As general recommendations, try to substitute possibly “weird” definitions in multiple choice to see a definition from the book.

## 2 Alphabets, Words, etc

### 2.2 Alphabets, Words, Languages

**Definition 2.1:** (*Alphabet*) Set  $\Sigma$ . Important alphabets:  $\Sigma_{\text{bool}}$ ,  $\Sigma_{\text{lat}}$  (all latin chars),  $\Sigma_{\text{Keyboard}}$  (all chars on keyboard),  $\Sigma_m$  ( $m$ -ary numbers)

**Definition 2.2:** (*Word*) Possibly empty (denoted  $\lambda$ ) sequences of characters from  $\Sigma$ .  $|w|$  is the length,  $\Sigma^*$  is the set of all words and  $\Sigma^+ = \Sigma^* - \{\lambda\}$

**Definition 2.3:** (*Konkatenation*)  $\text{Kon}(x, y) = xy$ , (so like string concat).  $(xy)^n$  is  $n$ -times repeated concat.

**Definition 2.4:** (*Reversal*)  $a^R$ , simply read the word backwards.

**Definition 2.6:** (*Prefix, Suffix, Subword*)  $v$  in  $w = vy$ ;  $s$  in  $w = xs$ ; Subword  $u$  in  $w = xuy$ ;  $x, y$  possibly  $\lambda$

**Definition 2.7:** (*Appearance*)  $|x|_a$  is the number of times  $a \in \Sigma$  appears in  $x$

**Definition 2.8:** (*Canonical ordering*) Ordered by length and then by first non-common letter:

$$u < v \iff |u| < |v| \vee (|u| = |v| \wedge u = x \cdot s_i \cdot u' \wedge v = x \cdot s_j \cdot v' \text{ for any } x, u', v' \in \Sigma^* \text{ and } i < j)$$

**Definition 2.9:** (*Language*)  $L \subseteq \Sigma^*$ , and we define  $L^C = \Sigma^* - L$  as the complement, with  $L_\emptyset$  being the empty language, whereas  $L_\lambda$  is the language with just the empty word in it.

**Concatenation:**  $L_1 \cdot L_2 = \{vw | v \in L_1 \wedge w \in L_2\}$  and  $L^{i+1} = L^i \cdot L \ \forall i \in \mathbb{N}$ .

**Cleen Star:**  $L^* = \bigcup_{i \in \mathbb{N}} L^i$  and  $L^+ = L \cdot L^*$

Of note is that there are irregular languages whose Cleen Star is regular, most notably, the language  $L = \{w \in \{0\}^* \mid |w| \text{ is prime}\}$ 's Cleen Star is regular, due to the fact that the prime factorization is regular

**Lemma 2.1:**  $L_1 L_2 \cup L_1 L_2 = L_1(L_2 \cup L_3)$  **Lemma 2.2:**  $L_1(K_2 \cap L_3) \subseteq L_1 L_2 \cap L_1 L_3$

For multiple choice questions, really think of how the sets would look to determine if they fulfill a requirement.

### 2.4 Kolmogorov-Complexity

**Definition 2.17:** (*Kolmogorov-Complexity*)  $K(x)$  for  $x \in (\Sigma_{\text{bool}})^*$  is the minimum of all binary lengths of Pascal programs that output  $x$ , where the Program doesn't have to compile, i.e. we can describe processes informally

**Lemma 2.4:** For each word  $x$  exists constant  $d$  s.t.  $K(x) \leq |x| + d$ , for which we can use a program that simply includes a `write(x)` command

**Definition 2.18:** (*Of natural number*)  $K(n) = K(\text{Bin}(x))$  with  $|\text{Bin}(x)| = \lceil \log_2(x+1) \rceil$

**Lemma 2.5:** For each  $n \in \mathbb{N} \exists w_n \in (\Sigma_{\text{bool}})^n$  s.t.  $K(w_n) \geq |w_n| = n$ , i.e. exists a non-compressible word.

**Theorem 2.1:** Kolmogorov-Complexity doesn't depend on programming language. It only differs in constant

**Definition 2.19:** (*Randomness*)  $x \in (\Sigma_{\text{bool}})^*$  random if  $K(x) \geq |x|$ , thus for  $n \in \mathbb{N}$ ,  $K(n) \geq \lceil \log_2(n+1) \rceil - 1$

**Theorem 2.3:** (*Prime number*)  $\lim_{n \rightarrow \infty} \frac{\text{Prime}(n)}{\frac{n}{\ln(n)}} = 1$  with  $\text{Prime}(n)$  the number of prime numbers on  $[0, n] \subseteq \mathbb{N}$

**Proofs** Proofs in which we need to show a lower bound for Kolmogorov-Complexity (almost) always work as follows: Assume for contradiction that there are no words with  $K(w) > f$  for all  $w \in W$ . We count the number  $m$  of words in  $W$  and the number  $n$  of programs of length  $\leq f$  ( $f$  being the given, lower bound). We will have  $m - n > 0$ , which means, there are more different words than there are Programs with Kolmogorov-Complexity  $\leq f$ , which is a contradiction to our assumption.

There are  $\lfloor \frac{n}{k} \rfloor + 1$  numbers divisible by  $k$  in the set  $\{0, 1, \dots, n\}$ .

#### Laws of logarithm

- $\log_a(x) + \log_a(y) = \log_a(x \cdot y)$
- $y \log_a(x) = \log_a(x^y)$
- $\log_a(1) = 0$
- $\log_a(x) - \log_a(y) = \log_a(x \div y)$
- $\log_a(x) = \frac{\ln(x)}{\ln(a)}$

### 3 Finite Automata

#### 3.2 Representation

We can note the automata using graphical notation similar to graphs or as a series of instructions like this:

```
select input = a1 goto i1
      ⋮
input = ak goto ik
```

**Definition 3.1:** (*Finite Automaton*)  $A = (Q, \Sigma, \delta, q_0, F)$  with

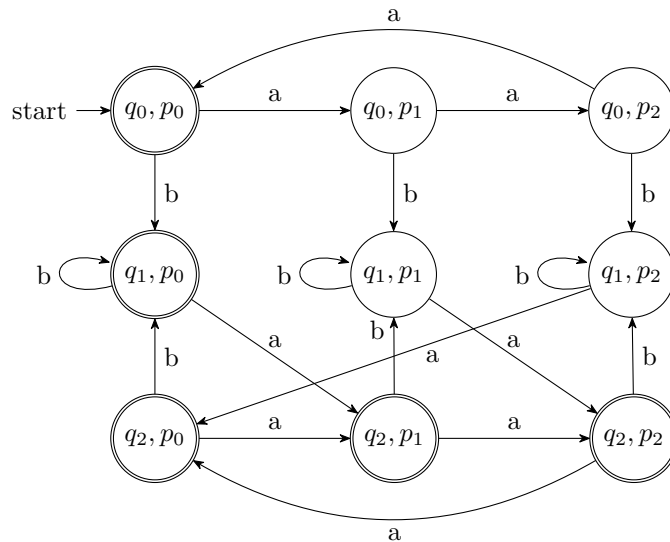
- $Q$  set of states
- $\Sigma$  input alphabet
- $\delta(q, a) = p$  transition from  $q$  on reading  $a$  to  $p$
- $q_0$  initial state
- $F \subseteq Q$  accepting states
- $\mathcal{L}_{EA}$  regular languages (accepted by FA)

$\hat{\delta}(q_0, w) = p$  is the end state reached when we process word  $w$  from state  $q_0$ , and  $(q, w) \mid_M^* (p, \lambda)$  is the formal definition, with  $\mid_M^*$  representing any number of steps  $\mid_M$  executed (transitive hull).

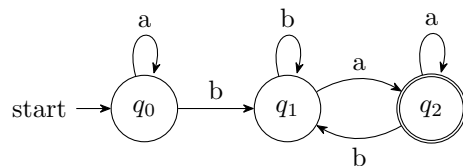
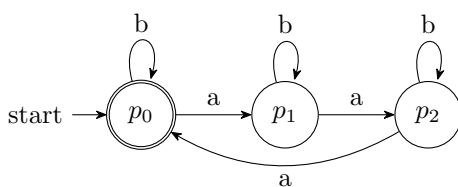
The class  $\text{Cl}[q_i]$  represents all possible words for which the FA is in this state. Be cautious when defining them, make sure that no extra words from other classes could appear in the current class, if this is not intended.

Sometimes, we need to combine two (or more) FA to form one larger one. We can do this easily with product automata. To create one from two automata  $M_1$  (states  $q_i$ ) and  $M_2$  (states  $p_j$ ) we do the following steps:

1. Write down the states as tuples of the form  $(q_i, p_j)$  (i.e. form a grid by writing down one of the automata vertically and the other horizontally)
2. From each state, the automata on the horizontal axis decides for the input symbol if we move left or right, whereas the automata on the vertical axis decides if we move up or down.



For the automata



- (a) Module to compute  $|w|_b \equiv |w| \pmod{3}$ . States  $q \in Q_a$       (b) Module to compute  $w$  contains sub.  $ba$  and ends in  $a$ . States  $p \in Q_b$

Figure 3.1: Graphical representation of the Finite Automaton of Task 9 in 2025

### 3.4 Proofs of nonexistence

We have three approaches to prove non-regularity of words. Below is an informal guide as to how to do proofs using each of the methods and possible pitfalls.

For all of them start by assuming that  $L$  is regular.

#### Lemma 3.3

##### Regular words

##### Lemma 3.3

Let  $A$  be a FA over  $\Sigma$  and let  $x \neq y \in \Sigma^*$ , such that  $\hat{\delta}_A(q_0, x) = \hat{\delta}_A(q_0, y)$ . Then for each  $z \in \Sigma^*$  there exists an  $r \in Q$ , such that  $xz, yz \in \text{Cl}[r]$ , and we thus have

$$xz \in L(A) \iff yz \in L(A)$$

1. Pick a FA  $A$  over  $\Sigma$  and say that  $L(A) = L$
  2. Pick  $|Q| + 1$  words  $x$  such that  $xy = w \in L$  with  $|y| > 0$ .
  3. State that via pigeonhole principle there exists w.l.o.g  $i < j \in \{1, \dots, |Q| + 1\}$ , s.t.  $\hat{\delta}_A(q_0, x_i) = \hat{\delta}_A(q_0, x_j)$ .
  4. Build contradiction by picking  $z$  such that  $x_i z \in L$ .
  5. Then, if  $z$  was picked properly, since  $i < j$ , we have that  $x_j z \notin L$ , since the lengths do not match
- That is a contradiction, which concludes our proof

#### Pumping Lemma

##### Pumping-Lemma für reguläre Sprachen

##### Lemma 3.4

Let  $L$  be regular. Then there exists a constant  $n_0 \in \mathbb{N}$ , such that each word  $w \in \Sigma^*$  with  $|w| \geq n_0$  can be decomposed into  $w = yxz$ , with

- (i)  $|yx| \leq n_0$
- (ii)  $|x| \geq 1$

- (iii) For  $X = \{yx^kz \mid k \in \mathbb{N}\}$  either  $X \subseteq L$  or  $X \cap L = \emptyset$  applies

1. State that according to Lemma 3.4 there exists a constant  $n_0$  such that  $|w| \geq n_0$ .
2. Choose a word  $w \in L$  that is sufficiently long to enable a sensible decomposition for the next step.
3. Choose a decomposition, such that  $|yx| = n_0$  (makes it quite easy later). Specify  $y$  and  $x$  in such a way that for  $|y| = l$  and  $|x| = m$  we have  $l + m \leq n_0$
4. According to Lemma 3.4 (ii),  $m \geq 1$  and thus  $|x| \geq 1$ . Fix  $z$  to be the suffix of  $w = yxz$
5. Then according to Lemma 3.4 (iii), fill in for  $X = \{yx^kz \mid k \in \mathbb{N}\}$  we have  $X \subseteq L$ .
6. This will lead to a contradiction commonly when setting  $k = 0$ , as for a language like  $0^n 1^n$ , we have  $0^{(n_0-m)+km} 1^{n_0}$  as the word (with  $n_0 - m = l$ ), which for  $k = 0$  is  $u = 0^{n_0-m} 1^{n_0}$  and since  $m \geq 1$ ,  $u \notin L$  and thus by Lemma 3.4,  $X \cap L = \emptyset$ , but that is also not true, as the intersection is not empty (for  $k = 1$ )

#### Kolmogorov Complexity

1. We first need to choose an  $x$  such that  $L_x = \{y \mid xy \in L\}$ . If not immediately apparent, choosing  $x = a^{\alpha+1}$  for  $a \in \Sigma$  and  $\alpha$  being the exponent of the exponent of the words in the language after a variable rename. For example, for  $\{0^{n^2+2n} \mid n \in \mathbb{N}\}$ ,  $\alpha(m) = m^2 + 2m$ . Another common way to do this is for languages of the form  $\{a^n b^n \mid n \in \mathbb{N}\}$  to use  $x = a^m$  and  $L_{0^m} = \{y \mid 0^m y \in L\} = \{0^j 1^{m+j} \mid j \in \mathbb{N}\}$ .
2. Find the first word  $y_1 \in L_x$ . In the first example, this word would be  $y_1 = 0^{(m+1)^2 \cdot 2(m+1) - m^2 \cdot 2m+1}$ , or in general  $a^{\alpha(m+1) - \alpha(m)+1}$ . For the second example, the word would be  $y_1 = 1^m$ , i.e. with  $j = 0$
3. According to Theorem 3.1, there exists constant  $c$  such that  $K(y_k) \leq \lceil \log_2(k+1) \rceil + c$ . We often choose  $k = 1$ , so we have  $K(y_1) \leq \lceil \log_2(1+1) \rceil + c = 1 + c$  and with  $d = 1 + c$ ,  $K(y_1) \leq d$
4. This however leads to a contradiction, since the number of programs with length  $\leq d$  is at most  $2^d$  and thus finite, but our set  $L_x$  is infinite.

### Minimum number of states

To show that a language needs *at least*  $n$  states, use Lemma 3.3 and  $n$  words. We thus again do a proof by contradiction:

1. Assume that there exists FA with  $|Q| < n$ . We now choose  $n$  words (as short as possible), as we would for non-regularity proofs using Lemma 3.3 (i.e. find some prefixes). It is usually beneficial to choose prefixes with  $|w|$  small (consider just one letter,  $\lambda$ , then two and more letter words). An “easy” way to find the prefixes is to construct a finite automaton and then picking a prefix from each class
2. Construct a table for the suffixes using the  $n$  chosen words such that one of the words at entry  $x_{ij}$  is in the language and the other is not. ( $n \times n$  matrix, see below in example)
3. Conclude that we have reached a contradiction as every field  $x_{ij}$  contains a suffix such that one of the two words is in the language and the other one is not.

**Example 3.1:** Let  $L = \{x1y \mid x \in (\Sigma_{\text{bool}})^*, y \in \{0,1\}^2\}$ . Show that any FA that accepts  $L$  needs at least four states.

Assume for contradiction that there exists EA  $A = (Q, \Sigma_{\text{bool}}, \delta_A, q_0, F)$  with  $|Q| < 4$ . Let's take the 4 words 00, 01, 10, 11. Then according to Lemma 3.3, there needs to exist a  $z$  such that  $xz \in L(A) \iff yz \in L(A)$  with  $\hat{\delta}_A(q_0, x) = \hat{\delta}_A(q_0, y)$  for  $x, y \in \{00, 01, 10, 11\}$ .

This however is a contradiction, as we can find a  $z$  for each of the pairs  $(x, y)$ , such that  $xz \in L(A)$ , but  $yz \notin L(A)$ . See for reference the below table (it contains suffixes  $z$  fulfilling prior condition):

	00	01	10	11
00	-	00	0	0
01		-	0	0
10			-	00
11				-

Thus, all four words have to lay in pairwise distinct states and we thus need at least 4 states to detect this language.

### 3.5 Non-determinism

The most notable differences between deterministic and non-deterministic FA is that the transition function is different:  $\delta : Q \times \Sigma \rightarrow \mathcal{P}(Q)$ . I.e., there can be any number of transitions for one symbol of  $\Sigma$  for each state. This is (in graphical notation) represented by arrows that have the same label going to different nodes.

It is also possible for there to not be a transition function for a certain element of the input alphabet. In that case, regardless of state, the NFA rejects, as it “gets stuck” in a state and can't finish processing.

Additionally, the NFA accepts  $x$  if it has at least one accepting calculation on  $x$ .

**Theorem 3.2:** For every NFA  $M$  there exists a FA  $A$  such that  $L(M) = L(A)$ . They are then called *equivalent*

**Potenzmengenkonstruktion** States are now sets of states of the NFA in which the NFA could be in after processing the preceding input elements and we have a special state called  $q_{\text{trash}}$ .

For each state, the set of states  $P = \hat{\delta}(q_0, z)$  for  $|z| = n$  represents all possible states that the NFA could be in after doing the first  $n$  calculations.

Correspondingly, we add new states if there is no other state that is in the same branch of the calculation tree  $\mathcal{B}_M(x)$ . So, in other words, we execute BFS on the calculation tree.

## 4 Turing Machines

### 4.3 Representation

Turing machines are much more capable than FA and NFA. A full definition of them can be found in the book on pages 96 - 98 (= pages 110 - 112 in the PDF).

For example, to detect a recursive language like  $\{0^n 1^n \mid n \in \mathbb{N}\}$  we simply replace the left and rightmost symbol with a different one and repeat until we only have the new symbol, at which point we accept, or there are no more 0s or 1s, at which point we reject.

The Turing Machines have an accepting  $q_{\text{accept}}$  and a rejecting state  $q_{\text{reject}}$  and a configuration is an element of  $\{\{\epsilon\} \cdot \Gamma^* \cdot Q \cdot \Gamma^+ \cup Q \cdot \{\epsilon\} \cdot \Gamma^+\}$  with  $\cdot$  being the concatenation and  $\epsilon$  the marker of the start of the band.

### 4.4 Multi-tape TM and Church's Thesis

$k$ -Tape Turing machines have  $k$  extra tapes that can be written to and read from, called memory tapes. They *cannot* write to the input tape. Initially the memory tapes are empty and we are in state  $q_0$ . All read/write-heads of the memory tapes can move in either direction, granted they have not reached the far left end, marked with  $\epsilon$ .

As with normal TMs, the Turing Machine  $M$  accepts  $w$  if and only if  $M$  reaches the state  $q_{\text{accept}}$  and rejects if it does not terminate or reaches the state  $q_{\text{reject}}$ .

**Lemma 4.1:** There exists an equivalent 1-Tape-TM for every TM.

**Lemma 4.2:** There exists an equivalent TM for each Multi-tape TM.

Church's Thesis states that the Turing Machines are a formalization of the term "Algorithm". It is the only axiom specific to Computer Science.

All the words that can be accepted by a Turing Machine are elements of  $\mathcal{L}_{RE}$  and are called *recursively enumerable*.

### 4.5 Non-Deterministic Turing Machines

The same ideas as with NFA apply here. The transition function also maps into the power set:

$$\delta : (Q - \{q_{\text{accept}}, q_{\text{reject}}\}) \times \Gamma \rightarrow \mathcal{P}(Q \times \Gamma \times \{L, R, N\})$$

Again, when constructing a normal TM from a NTM (which is not required at the Midterm, or any other exam for that matter in this course), we again apply BFS to the NTM's calculation tree.

**Theorem 4.2:** For an NTM  $M$  exists a TM  $A$  s.t.  $L(M) = L(A)$  and if  $M$  doesn't contain infinite calculations on words of  $(L(M))^C$ , then  $A$  always stops.

## 5 Computability

### 5.2 Diagonalization

The *set of binary encodings of all TMs* is denoted  $\text{KodTM}$  and  $\text{KodTM} \subseteq (\Sigma_{\text{bool}})^*$  and the upper bound of the cardinality is  $|(\Sigma_{\text{bool}})^*|$ , as there are infinitely many TMs.

Below is a list of countable objects. They all have corresponding Lemmas in the script, but omitted here:

- $\Sigma^*$  for any  $\Sigma$
- $\text{KodTM}$
- $\mathbb{N} \times \mathbb{N}$
- $\mathbb{Q}^+$

The following objects are uncountable:  $[0, 1]$ ,  $\mathbb{R}$ ,  $\mathcal{P}((\Sigma_{\text{bool}})^*)$

**Corollary 5.1:**  $|\text{KodTM}| < |\mathcal{P}((\Sigma_{\text{bool}})^*)|$  and thus there exist infinitely many not recursively enumerable languages over  $\Sigma_{\text{bool}}$

#### Proof of $L$ (not) recursively enumerable

Proving that a language *is* recursively enumerable is as difficult as providing a Turing Machine that accepts it.

Proving that a language is *not* recursively enumerable is likely easier. For it, let  $d_{ij} = 1 \iff M_i$  accepts  $w_j$ .

**Example 5.1:** Assume towards contradiction that  $L_{\text{diag}} \in \mathcal{L}_{RE}$ . Let

$$\begin{aligned} L_{\text{diag}} &= \{w \in (\Sigma_{\text{bool}})^* \mid w = w_i \text{ for an } i \in \mathbb{N} - \{0\} \text{ and } M_i \text{ does not accept } w_i\} \\ &= \{w \in (\Sigma_{\text{bool}})^* \mid w = w_i \text{ for an } i \in \mathbb{N} - \{0\} \text{ and } d_{ii} = 0\} \end{aligned}$$

Thus assume that,  $L_{\text{diag}} = L(M)$  for a Turing Machine  $M$ . Since  $M$  is a Turing Machine in the canonical ordering of all Turing Machines, so there exists an  $i \in \mathbb{N} - \{0\}$ , such that  $M = M_i$ .

This however leads to a contradiction, as  $w_i \in L_{\text{diag}} \iff d_{ii} = 0 \iff w_i \notin L(M_i)$ .

In other words,  $w_i$  is in  $L_{\text{diag}}$  if and only if  $w_i$  is not in  $L(M_i)$ , which contradicts our statement above, in which we assumed that  $L_{\text{diag}} \in \mathcal{L}_{RE}$

In other, more different, words,  $w_i$  being in  $L_{\text{diag}}$  implies (from the definition) that  $d_{ii} = 0$ , which from its definition implies that  $w_i \notin L(M_i)$

**Theorem 5.3:**  $L_{\text{diag}} \notin \mathcal{L}_{RE}$