

1.)

```
75     def swap(self):
76         if self.Header is None or self.Current.Next is None:
77             return -1
78         else:
79             x = self.Current.Data
80             self.Current.Data = self.Current.Next.Data
81             self.removeCurrentNext()
82             self.insertCurrentNext(x)
83
84             return 0
```

2.)

- a.) A Python Dictionary is an Array object.
- b.) `dic.update({key,value})`
- c.) `dic["key"] = "val2"`
- d.) The for loop will be executed for however many key:value pairs are in the dictionary

3.)

```
def inOrderRank(self, num):
    def create_rank_list(self):
        rank_list = []
        def __visit__(n):
            if (n != None):
                __visit__(n.Left)
                rank_list.append(n.Data)
                __visit__(n.Right)

        __visit__(self.Root)
        return rank_list

    x = create_rank_list()

    for i in range(len(x)):
        if x[i] == num:
            return i
    return -1
```

4.)

```

180 def compare_arrays(A, B):
181     def check(arr1, arr2):
182         ans = []
183         for i in range(len(arr1)):
184             count = 0
185             for j in range(len(arr2)):
186                 if arr1[i] == arr2[j]:
187                     count = count + 1
188             if count == 0:
189                 ans.append(arr1[i])
190         return ans
191
192     C1 = check(A,B)
193     C2 = check(B,A)
194     return C1 + C2
195
196 A = [20, 40, 70, 30, 10, 80, 50, 90, 60]
197 B = [35, 45, 55, 60, 50, 40]
198
199 print(compare_arrays(A,B))

```

a.)

b.) C1:

i = 0 and ans = [20]

i = 1 and ans = [20]

i = 2 and ans = [20, 70]

i = 3 and ans = [20, 70, 30]

i = 4 and ans = [20, 70, 30, 10]

i = 5 and ans = [20, 70, 30, 10, 80]

i = 6 and ans = [20, 70, 30, 10, 80]

i = 7 and ans = [20, 70, 30, 10, 80, 90]

i = 8 and ans = [20, 70, 30, 10, 80, 90]

C2:

i = 0 and ans = [35]

i = 1 and ans = [35, 45]

i = 2 and ans = [35, 45, 55]

i = 3 and ans = [35, 45, 55]

i = 4 and ans = [35, 45, 55]

i = 5 and ans = [35, 45, 55]

C1 + C2:

[20, 70, 30, 10, 80, 90, 35, 45, 55]

c.)

180

181

182 c1

183 c2 \* n-1

184 c3 \* n-1

185 c4 \* (m-1)\*m/2

186 c5 \* (m-1)\*m/2

187 c6 \* (m-1)\*m/2

188 c7 \* n-1

189 c8 \* n-1

190 c9

191

192 c10

193 c11

193 c12

$$T(n) = C_1 + C_2(n-1) + C_3(n-1) + C_4\left(\frac{(n-1)(n)}{2}\right) + C_5\left(\frac{(n-1)n}{2}\right) + C_6\left(\frac{(n-1)n}{2}\right) + C_7(n-1) + C_9 + C_{10} + C_{11} + C_{12}$$

$$T(n) = C_{13} + C_{14}(n-1) + C_{15}\left(\frac{(n-1)(n)}{2}\right)$$

$$T(n) = C_{13} + C_{14}n - C_{14} + \frac{C_{15}}{2}(n^2 - n)$$

$$T(n) = C_{16} + C_{14}n + C_{17}n^2 - C_{17}n \quad \text{and} \quad \text{canceling out } C_{17}n^2 \text{ and } C_{17}n$$

$$T(n) = C_{16} + C_{14}n + C_{17}n^2 - C_{17}n \leq C_{16} + C_{14}n + C_{17}n^2$$

$$T(n) \leq C_{16} + C_{14}n + C_{17}n^2 \leq (C_{16} + C_{14} + C_{17})nm^2 \leq C_{18}nm^2$$

$$\boxed{T(n) = O(nm^2)}$$

5.)

5.)

$$a) T(n) = 4T\left(\frac{n}{2}\right) + n^3$$

$$a=4 \quad b=2 \quad d=3$$

$$a=4 < b^d = 2^3 = 8$$

$$T(n) = O(n^3)$$

$$b.) T(n) = 4T\left(\frac{n}{2}\right) + n^2$$

$$a=4 \quad b=2 \quad d=2$$

$$a=4 = b^d = 2^2 = 4$$

$$T(n) = O(n^2 \lg(n))$$

$$c.) T(n) = 4T\left(\frac{n}{2}\right) + 1$$

$$a=4 \quad b=2 \quad d=0$$

$$a=4 > b^d = 1$$

$$T(n) = O(n^{\log_2 4}) = O(n^2)$$

6.)



b.) a) graph =  $\begin{bmatrix} [0, 1, 0, 1, 0, 0, 0] \\ [1, 0, 1, 0, 0, 0, 1] \\ [1, 1, 0, 0, 0, 0, 0] \\ [0, 0, 0, 0, 1, 1, 0] \\ [0, 0, 0, 0, 0, 1, 0] \\ [0, 1, 0, 0, 0, 0, 0] \\ [0, 0, 1, 0, 0, 1, 0] \end{bmatrix}$

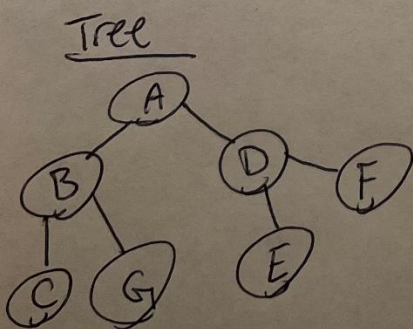
b.) Adjacency List

A	BD
B	ACG
C	AB
D	EF
E	F
F	B
G	CF

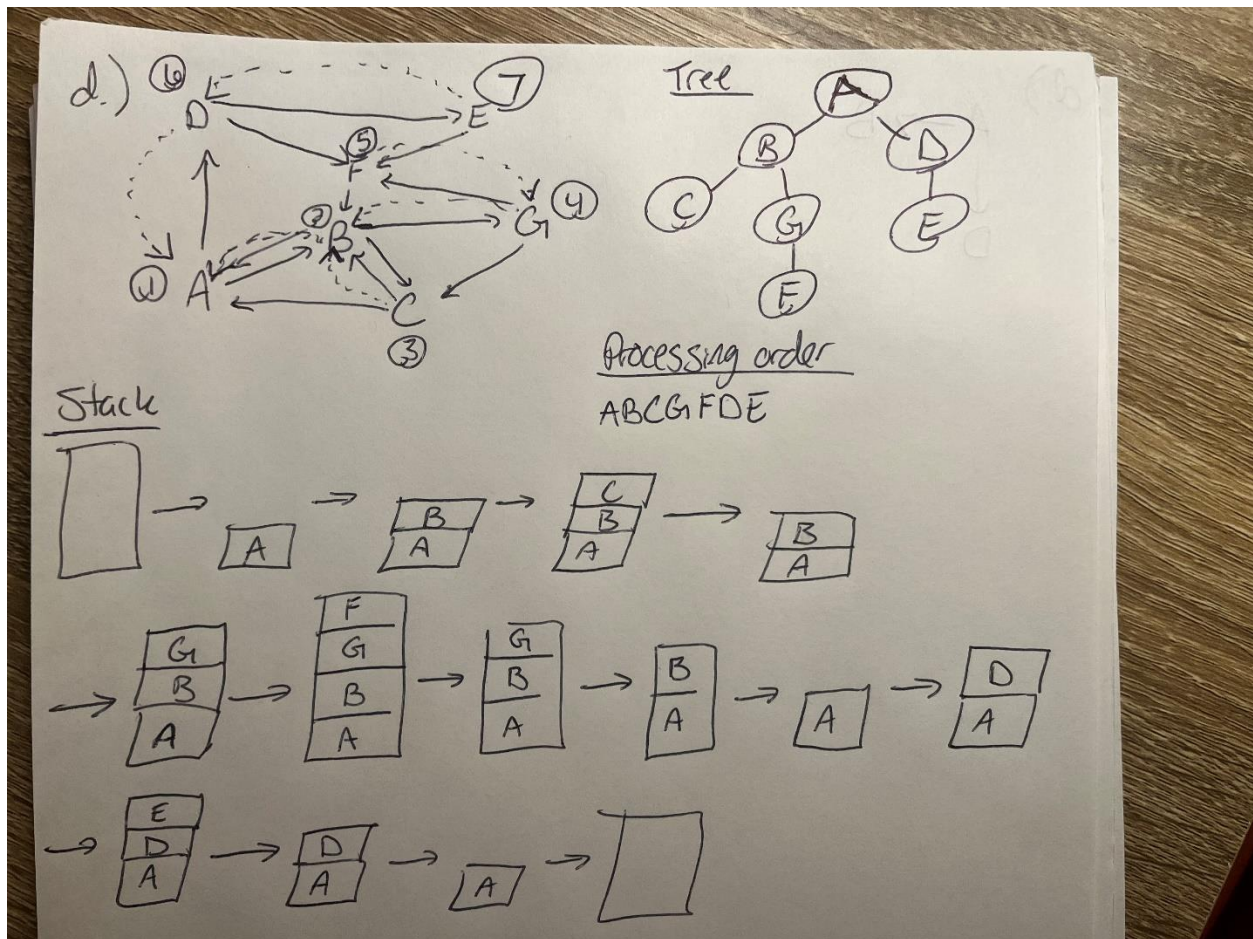
c.) Queue

[ ]  
 [A]  
 [AB]  
 [ABD]  
 [BD]  
 [BDC]  
 [BDCG]  
 [ACG]  
 [DCGE]  
 [DCGEF]

[CGEF]  
 [GEF]  
 [EF]  
 [F]  
 [ ]



Processing List  
 ABDCGEF



7.)

a.)

```

205 def find_max(A, right):
206     print("Params: ", A, right)
207     if right == 0:
208         print("Return: ", A[right])
209         return A[right]
210     else:
211         x = find_max(A, right - 1)
212         print(A[right], " compared to ", x)
213         if A[right] > x:
214             print("Return: ", A[right])
215             return A[right]
216         else:
217             print("Return: ", x)
218             return x
219
220 A = [17, 62, 49, 73, 26, 51]
221 find_max(A, 5)

```

b.)

Params: [17, 62, 49, 73, 26, 51] 5

Params: [17, 62, 49, 73, 26, 51] 4

Params: [17, 62, 49, 73, 26, 51] 3

Params: [17, 62, 49, 73, 26, 51] 2

Params: [17, 62, 49, 73, 26, 51] 1

Params: [17, 62, 49, 73, 26, 51] 0

Return: 17

62 compared to 17

Return: 62

49 compared to 62

Return: 62

73 compared to 62

Return: 73

26 compared to 73

Return: 73

51 compared to 73



7.)

$$c.) T(n) = 1 + T(n-1), T(1) = 0$$

$$d.) T(n) = 1 + [1 + T(n-2)] = 2 + T(n-2)$$

$$T(n) = 1 + 1 + [1 + T(n-3)] = 3 + T(n-3)$$

$$T(n) = k + T(n-k)$$

$$T(1) = 0 \text{ when } n-k=1 \text{ or } k=n-1$$

$$T(n) = n-1 + T(1) = n-1 + 0 = n-1$$

$$T(n) = n-1 < n$$

$$T(n) = O(n)$$

8.)

a.) MergeSort is Big-Oh class  $n \lg(n)$  because the fundamental step of the algorithm is the comparisons in the Merge function. In the worst case the cost of Merge is  $n-1$  and Merge will be called however many times it takes to combine all subarrays into one final array by combining two at a time. QuickSort is in Big-Oh class  $n^2$  because in the worst case when the pivot is always on at the left or right index, the cost of the partition function can be  $n-1(n)/2$ .

b.) MergeSort average case is  $n \lg(n)$  because no algorithm will do more or less than that amount of work. QuickSort average case is  $n \lg(n)$  because the best case is  $n \lg(n)$  when the partition algorithm does  $n$  amount of work. The worst case for QuickSort is very unlikely, so the average case is best estimated using the best case.

c.) No answer. Ran out of time 😞

d.) No answer. Ran out of time 😞

