

Databases Autumn 2025

Hand-In Exercise 2

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Total Points	

Task	Points

Task 1

Given excerpt of the relational schema plus added relations and integrety constrain befor the first bullet point.

```
1 ===== Entity =====
2 CREATE TABLE Lecturer (
3     LecturerID    INT PRIMARY KEY,
4     FirstName     VARCHAR(255),
5     LastName      VARCHAR(255),
6     Title         VARCHAR(255)
7 );
8
9 CREATE TABLE Lecture (
10    Title          VARCHAR(255) PRIMARY KEY,
11    CreditPoints   INT,
12    SemesterWeekHours INT,
13    LecturerID     INT,
14    FOREIGN KEY (LecturerID) REFERENCES Lecturer (LecturerID)
15 );
16
17 CREATE TABLE Exercise (
18     ExID          INT PRIMARY KEY,
19     No            INT,
20     Semester      VARCHAR(255),
21     LectureTitle  VARCHAR(255),
22     FOREIGN KEY (LectureTitle) REFERENCES Lecture (Title)
23 );
24
25 CREATE TABLE Author (
26     AuthorID      INT PRIMARY KEY,
27     LastName      VARCHAR(255),
28     FirstName     VARCHAR(255),
29     Title         VARCHAR(255)
30 );
31
32 CREATE TABLE Task (
33     TaskID        INT PRIMARY KEY,
34     Points         INT,
35     Difficulty    INT,
36     Text           VARCHAR(65535),
37     AuthorID       INT,
38     FOREIGN KEY (AuthorID) REFERENCES Author (AuthorID)
39 );
40
41 ===== Relationships =====
```

```
42 | CREATE TABLE Contains (
43 |     ExID      INT,
44 |     TaskID    INT,
45 |     PRIMARY KEY (ExID, TaskID),
46 |     FOREIGN KEY (ExID) REFERENCES Exercise (ExID),
47 |     FOREIGN KEY (TaskID) REFERENCES Task (TaskID),
48 |     Sequence   INT
49 | );
50 |
51 | CREATE TABLE Consists_of (
52 |     SuperTaskID INT,
53 |     SubTaskID   INT,
54 |     PRIMARY KEY (SuperTaskID, SubTaskID),
55 |     FOREIGN KEY (SuperTaskID) REFERENCES Task (TaskID),
56 |     FOREIGN KEY (SubTaskID) REFERENCES Task (TaskID),
57 |     Sequence    INT
58 | );
59 |
60 | ===== Integrity =====
61 | CREATE ASSERTION contains_only_super_tasks
62 |     CHECK ( NOT EXISTS (
63 |         SELECT *
64 |         FROM Contains c
65 |         WHERE EXISTS (
66 |             SELECT *
67 |             FROM Consists_of co
68 |             WHERE co.SubTaskID = c.TaskID
69 |         )
70 |     );
71 |
72 | CREATE ASSERTION consists_of_non_recursive
73 |     CHECK ( NOT EXISTS (
74 |         SELECT *
75 |         FROM Consists_of x
76 |         WHERE EXISTS (
77 |             SELECT *
78 |             FROM Consists_of y
79 |             WHERE y.SuperTaskID = x.SubTaskID
80 |         )
81 |     );
```

- *The title of the lecture has to be unique and may not be altered if any exercise is available for the lecture.*

The first part of the point is already enforced bc the title of the lecture is a PRIMARY KEY and therefore must be unique. The second part can be enforced with this trigger.

```
1 CREATE TRIGGER BlockTitleChange
2 BEFORE UPDATE OF Title ON Lecture
3 REFERENCING OLD AS lOld NEW AS lNew
4 WHEN (lNew.Title != lOld.Title AND
5      EXISTS (SELECT 1
6          FROM Exercise e
7          WHERE e.LectureTitle = lOld.Title))
8 (ROLLBACK WORK);
```

- *For a lecture, no more than 10 credit points may be awarded.*

This part can be enforced with this assertion.

```
1 CREATE ASSERTION no_more_than_10_credits
2     CHECK (NOT EXISTS (
3         SELECT * FROM Lecture
4         WHERE CreditPoints > 10))
5     NOT DEFERRABLE;
```

- *Lecturers may give multiple lectures.*

This part is beeing enforced with this.

```
1 CREATE TABLE Lecture (
2     Title           VARCHAR(255) PRIMARY KEY,
3     CreditPoints    INTEGER,
4     SemesterWeekHours INTEGER,
5     LecturerID      INTEGER,
6     --> FOREIGN KEY (LecturerID) REFERENCES Lecturer (LecturerID)
7 );
```

- *A lecture may include several exercises. An exercise always belongs to exactly one lecture.*

The first part is already enforced via the foreign key while for the second part we have to add the NOT NULL so we guarantee that an exersise must be in one lecture.

```
1 CREATE TABLE Exercise (
2     ExID            INT PRIMARY KEY,
3     No              INT,
4     Semester        VARCHAR(255),
5     --> LectureTitle VARCHAR(255) NOT NULL,
6     --> FOREIGN KEY (LectureTitle) REFERENCES Lecture (Title)
7 );
```

- Before a new author is entered into the system, it should be checked that no other author with the same first name, last name and title is present.
This part is being enforced with this trigger.

```
1 CREATE TRIGGER BlockDuplicateAuthor
2 BEFORE INSERT ON Author
3 REFERENCING NEW AS aNew
4 WHEN (EXISTS (SELECT 1
5     FROM Author a
6     WHERE a.FirstName = aNew.FirstName
7         AND a.LastName = aNew.LastName
8         AND a.Title = aNew.Title))
9 (ROLLBACK WORK);
```

Task 2 Foreign Keys

2a) Foreign Keys

For the relation `Lecture`, the following foreign key is defined in SQL:

```
1 CREATE TABLE Lecture (
2     ...
3     FOREIGN KEY (fk_lecturer) REFERENCES Lecturer(pk_lecturer)
4 );
```

In the following, this foreign key is simulated using (i) an assertion and (ii) triggers.

(i) Assertion (static integrity constraint)

An `ASSERTION` ensures that no tuple in `Lecture` references a non-existent `Lecturer`. We formulate this as a double `NOT EXISTS` statement that expresses the same semantics as a foreign key.

```
1 CREATE ASSERTION FK_Lecture_Lecturer_OK
2 CHECK (
3     NOT EXISTS (
4         SELECT 1
5         FROM Lecture L
6         WHERE L.fk_lecturer IS NOT NULL
7         AND NOT EXISTS (
8             SELECT 1
9             FROM Lecturer R
10            WHERE R.pk_lecturer = L.fk_lecturer
11        )
12    )
13 ) DEFERRABLE INITIALLY IMMEDIATE;
```

This assertion guarantees that for every non-NULL value in `Lecture.fk_lecturer`, a matching `Lecturer.pk_lecturer` exists. The condition is checked globally and can be deferred to commit time using the DEFERRABLE clause.

(ii) **Triggers (dynamic integrity constraint)**

The same referential integrity can be dynamically enforced by triggers (E-C-A pattern):

- **Child-side:** Prevent inserting or updating a `Lecture` with a non-existent `fk_lecturer`.
- **Parent-side:** Prevent deleting or updating a `Lecturer` that is still referenced by a `Lecture`.

Child-side triggers (on Lecture)

```
1 CREATE TRIGGER Lecture_FK_Check_Ins
2 BEFORE INSERT ON Lecture
3 WHEN (NEW.fk_lecturer IS NOT NULL AND
4     NOT EXISTS (SELECT 1 FROM Lecturer R
5                 WHERE R.pk_lecturer = NEW.fk_lecturer))
6 ( ROLLBACK WORK );
7
8 CREATE TRIGGER Lecture_FK_Check_Upd
9 BEFORE UPDATE OF fk_lecturer ON Lecture
10 REFERENCING NEW AS Lnew
11 WHEN (Lnew.fk_lecturer IS NOT NULL AND
12     NOT EXISTS (SELECT 1 FROM Lecturer R
13                 WHERE R.pk_lecturer = Lnew.fk_lecturer))
14 ( ROLLBACK WORK );
```

Parent-side triggers (on Lecturer)

```
1 CREATE TRIGGER Lecturer_Ref_Block_Del
2 BEFORE DELETE ON Lecturer
3 REFERENCING OLD AS Pold
4 WHEN (EXISTS (SELECT 1 FROM Lecture L
5                 WHERE L.fk_lecturer = Pold.pk_lecturer))
6 ( ROLLBACK WORK );
7
8 CREATE TRIGGER Lecturer_Ref_Block_Upd
9 BEFORE UPDATE OF pk_lecturer ON Lecturer
10 REFERENCING OLD AS Pold NEW AS Pnew
11 WHEN (EXISTS (SELECT 1 FROM Lecture L
12                 WHERE L.fk_lecturer = Pold.pk_lecturer))
13 ( ROLLBACK WORK );
```

- The **child-side** triggers prevent inserting or updating a lecture referencing a non-existent lecturer.
- The **parent-side** triggers prevent deleting or changing a lecturer ID if at least one lecture still refers to it.
- The `ROLLBACK WORK` aborts the transaction on violation (NO ACTION-like).

- 2b) Simulating FOREIGN KEY (FkExam) REFERENCES Exam(PkExam) ON DELETE SET NULL with triggers

We enforce referential integrity procedurally:

- **Child-side (Student).** Block INSERT/UPDATE that sets FkExam to a non-existent Exam.PkExam (allow NULL).
- **Parent-side (Exam) on DELETE.** Before deleting an Exam row, set Student.FkExam := NULL for all referencing students (ON DELETE SET NULL).
- **Parent-side (Exam) on UPDATE PkExam.** Default is NO ACTION; block if referenced.

Child-side checks (INSERT/UPDATE on Student)

```
1 CREATE TRIGGER Student_FK_Check_Ins
2 BEFORE INSERT ON Student
3 WHEN ( NEW.FkExam IS NOT NULL AND
4       NOT EXISTS (SELECT 1 FROM Exam E
5                     WHERE E.PkExam = NEW.FkExam) )
6   ( ROLLBACK WORK );
7
8 CREATE TRIGGER Student_FK_Check_Upd
9 BEFORE UPDATE OF FkExam ON Student
10 REFERENCING NEW AS Snew
11 WHEN ( Snew.FkExam IS NOT NULL AND
12       NOT EXISTS (SELECT 1 FROM Exam E
13                     WHERE E.PkExam = Snew.FkExam) )
14   ( ROLLBACK WORK );
```

Parent-side action (DELETE on Exam) — simulate ON DELETE SET NULL

```
1 CREATE TRIGGER Exam_Delete_SetNull
2 BEFORE DELETE ON Exam
3 REFERENCING OLD AS Eold
4 (
5   UPDATE Student
6   SET FkExam = NULL
7   WHERE FkExam = Eold.PkExam
8 );
```

Parent-side block (UPDATE of PkExam) — simulate ON UPDATE NO ACTION

```
1 CREATE TRIGGER Exam_Block_Update_PK
2 BEFORE UPDATE OF PkExam ON Exam
3 REFERENCING OLD AS Eold NEW AS Enew
4 WHEN ( EXISTS (SELECT 1 FROM Student S
5                     WHERE S.FkExam = Eold.PkExam) )
6   ( ROLLBACK WORK );
```

Notes.

- Ensure Student (FkExam) allows NULL; otherwise SET NULL would fail.
- The DELETE trigger clears references before parent removal.
- Using ROLLBACK WORK follows the lecture's "block on violation" pattern.

Task 3

Using Assertions or Triggers because integrity constraints span multiple relations and not only one (Create Table Statement).

- a) Using Assertions because constraints too general and not specific event.

e.g. no new orders when ...

U2109: each User must have $\text{SUM}(c.\text{CreditScore}) \geq -1000$ creditpoints over all credits.

```
CREATE ASSERTION CreditLimitU2109
CHECK(NOT EXISTS(
    SELECT u.UID
    FROM User u
    JOIN Regulation r ON u.Country = r.Country
    JOIN Credit c ON u.UID = c.UID
    WHERE r.RegulationCode = 'U2109'
    GROUP BY u.UID
    HAVING SUM(c.CreditScore) < -1000))
DEFERRABLE INITIALLY DEFERRED
```

U2304: each User must have $\text{SUM}(c.\text{CreditScore}) \geq (-0.1 * \text{MAX}(b.\text{BetAmount}))$

```
CREATE ASSERTION CreditLimitU2109
CHECK(NOT EXISTS(
    SELECT u.UID
    FROM User u
    JOIN Regulation r ON u.Country = r.Country
    JOIN Credit c ON u.UID = c.UID
    JOIN Bet b ON u.UID = b.UID
    WHERE r.RegulationCode = 'U2304'
    GROUP BY u.UID
    HAVING SUM(c.CreditScore) < (-0.1 * MAX(b.BetAmount)))
DEFERRABLE INITIALLY DEFERRED
```

- b) Materialized View:

```
CREATE VIEW DailyOnlineTime
REFRESH FORCE ON COMMIT
```

```
BUILD DEFERRED AS
    SELECT u.UID, u.Name, u.Country, s.Date,
           SUM(s.Duration) AS dailyTotalTime
      FROM User u
     JOIN Regulation r ON u.Country = r.Country
     JOIN Session s ON u.UID = s.UID
     JOIN Bet b ON u.UID = b.UID
    WHERE r.RegulationCode = 'M5475'
      AND u.Age BETWEEN 14 AND 18
   GROUP BY u.Name, u.Country, s.Date;
```

c) Refresh Materliazed View on Demand:

Using additionally a trigger to "manually" trigger the refresh of the view. First changing the view to "ON DEMAND". Second creating a trigger that refreshes the view when event happens.

```
CREATE VIEW DailyOnlineTime
REFRESH FORCE ON DEMAND
BUILD DEFERRED AS
    SELECT u.UID, u.Name, u.Country, s.Date,
           SUM(s.Duration) AS dailyTotalTime
      FROM User u
     JOIN Regulation r ON u.Country = r.Country
     JOIN Session s ON u.UID = s.UID
     JOIN Bet b ON u.UID = b.UID
    WHERE r.RegulationCode = 'M5475'
      AND u.Age BETWEEN 14 AND 18
   GROUP BY u.Name, u.Country, s.Date;
```

```
CREATE TRIGGER refreshView_onM4575_change
REFRESH FORCE ON DEMAND
AFTER INSERT OR DELETE OR UPDATE OF RegulationCode, Country
ON Regulation
REFERENCING OLD AS OldReg NEW AS NewReg
WHEN((NewReg.RegulationCode = 'M5475')
      OR (OldReg.RegulationCode = 'M5475'))
  (CALL DBMS_MVIEW.REFRESH('DailyOnlineTime'));
```

Used own notation of view-refresh because no example or definition with refresh on slides.

Task 4

When we say something is **DEFERRABLE INITIALLY DEFERRED**, it means that the database will wait until the end of a transaction before checking certain rules, such as foreign key

constraints. This is useful when two tables depend on each other. For example, a `Lecture` table and a `Lecturer` table. A lecture record might point to the lecturer who gives it, while the lecturer record might also depend on the lecture. If we try to insert both of these in the same transaction, the database might raise an error because when we insert the first one, the other does not yet exist. If the rule is **NOT DEFERRABLE**, the database checks immediately, and the operation fails. If it is **DEFERRABLE INITIALLY IMMEDIATE**, we would need to manually tell the database to delay the check. However, if it is **DEFERRABLE INITIALLY DEFERRED**, the database automatically waits until the transaction ends to perform the checks. This allows both inserts to complete first, and the database only verifies the rules afterward, which helps prevent errors.