# Databases Autumn 2025 Hand-In Exercise 1

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Total Points

Task	Points	
	I	

# Task 1

Given excerpt of the relational schema plus added relations and integrety constrain befor the first bullet point.

```
= Entity =
   CREATE TABLE Lecturer (
            LecturerID
                        INT PRIMARY KEY,
            FirstName
                         VARCHAR(255),
            LastName
                         VARCHAR(255),
5
            Title
                         VARCHAR(255)
   );
   CREATE TABLE Lecture (
            Title
                                VARCHAR(255) PRIMARY KEY,
10
            CreditPoints
                                 INT,
11
            SemesterWeekHours
                                INT,
12
            LecturerID
                                 INT,
13
           FOREIGN KEY (LecturerID) REFERENCES Lecturer (LecturerID)
14
   );
15
16
   CREATE TABLE Exercise (
17
            ExID
                         INT PRIMARY KEY,
18
           No
                         INT,
19
                         VARCHAR(255),
            Semester
20
            Lecture Title VARCHAR(255),
21
           FOREIGN KEY (Lecture Title) REFERENCES Lecture (Title)
22
   );
23
24
   CREATE TABLE Author (
25
            AuthorID
                        INT PRIMARY KEY,
26
            LastName
                       VARCHAR(255),
27
            FirstName
                       VARCHAR(255),
28
            Title
                       VARCHAR(255)
29
   );
30
31
   CREATE TABLE Task (
32
            TaskID
                        INT PRIMARY KEY,
33
            Points
                        INT,
            Difficulty INT,
35
            Text
                       VARCHAR(65535),
            AuthorID
                        INT,
37
           FOREIGN KEY (AuthorID) REFERENCES Author (AuthorID)
38
   );
39
40
        Relationships ====
41
```

```
CREATE TABLE Contains (
            ExID
                        INT,
43
            TaskID
                       INT,
44
            PRIMARY KEY (ExID, TaskID),
45
            FOREIGN KEY (ExID) REFERENCES Exercise (ExID),
46
            FOREIGN KEY (TaskID) REFERENCES Task (TaskID),
47
            Sequence INT
48
   );
49
50
   CREATE TABLE Consists of (
51
            SuperTaskID INT,
52
            {\bf SubTaskID}
                          INT,
53
            PRIMARY KEY (SuperTaskID, SubTaskID),
54
            FOREIGN KEY (SuperTaskID) REFERENCES Task (TaskID),
55
            FOREIGN KEY (SubTaskID)
                                          REFERENCES Task (TaskID),
56
            Sequence
                          INT
57
   );
58
59
        = Integrity =
60
   CREATE ASSERTION contains only super tasks
61
     CHECK ( NOT EXISTS (
62
       SELECT *
63
       FROM Contains c
64
       WHERE EXISTS (
65
          SELECT *
         FROM Consists of co
67
          \textbf{WHERE} \ co.\,SubTaskID \ = \ c.\,TaskID 
69
     ) );
70
71
   CREATE ASSERTION consists of non recursive
72
     CHECK ( NOT EXISTS (
73
       \mathbf{SELECT} \ *
74
       FROM Consists of x
75
       WHERE EXISTS (
76
          SELECT *
77
         FROM Consists of y
78
         WHERE y \cdot SuperTaskID = x \cdot SubTaskID
79
80
     ) );
```

• The title of the lecture has to be unique and may not be altered if any exercise is available for the lecture.

The first part of the point is already enforced by the title of the lecture is a PRIMARY KEY and therefore must be unique. The second part can be enforced with this addition.

• For a lecture, no more than 10 credit points may be awarded. This part can be enforced with this assertion.

```
CREATE ASSERTION no_more_than_10_credits
CHECK (NOT EXISTS (

SELECT * FROM Lecture

WHERE CreditPoints > 10))

NOT DEFERRABLE;
```

• Lecturers may give multiple lectures.

This part is beeing enforced with this.

```
CREATE TABLE Lecture (
Title VARCHAR(255) PRIMARY KEY,
CreditPoints INTEGER,
SemesterWeekHours INTEGER,
LecturerID INTEGER,
FOREIGN KEY (LecturerID) REFERENCES Lecturer (LecturerID)
);
```

• A lecture may include several exercises. An exercise always belongs to exactly one lecture. The first part is already enforced via the foreign key while for the second part we have to add the NOT NULL so we guaratee that e exersise must be in one lecture.

• Before a new author is entered into the system, it should be checked that no other author with the same first name, last name and title is present.

This part is beeing enforced with this.

```
CREATE TABLE Author (
AuthorID INT PRIMARY KEY,

LastName VARCHAR(255),

FirstName VARCHAR(255),

Title VARCHAR(255),

VNIQUE (FirstName, LastName, Title)

VIIQUE (FirstName, LastName, Title)
```

# Task 2 Foreign Keys

#### 2a) Foreign Keys

For the relation Lecture, the following foreign key is defined in SQL:

```
CREATE TABLE Lecture (
...,
FOREIGN KEY (fk_lecturer) REFERENCES Lecturer(pk_lecturer)
);
```

In the following, this foreign key is simulated using (i) an assertion and (ii) triggers.

# (i) Assertion (static integrity constraint)

An ASSERTION ensures that no tuple in Lecture references a non-existent Lecturer. We formulate this as a double NOT EXISTS statement that expresses the same semantics as a foreign key.

```
CREATE ASSERTION FK_Lecture_Lecturer_OK
CHECK (
NOT EXISTS (
SELECT 1
FROM Lecture L
WHERE L. fk_lecturer IS NOT NULL
AND NOT EXISTS (
SELECT 1
FROM Lecturer R
WHERE R. pk_lecturer = L. fk_lecturer

WHERE R. pk_lecturer

DEFERRABLE INITIALLY IMMEDIATE;
```

This assertion guarantees that for every non-NULL value in Lecture.fk\_lecturer, a matching Lecturer.pk\_lecturer exists. The condition is checked globally and can be deferred to commit time using the DEFERRABLE clause.

#### (ii) Triggers (dynamic integrity constraint)

The same referential integrity can be dynamically enforced by triggers (E-C-A pattern):

- Child-side: Prevent inserting or updating a Lecture with a non-existent fk\_lecturer.
- Parent-side: Prevent deleting or updating a Lecturer that is still referenced by a Lecture.

```
Child-side triggers (on Lecture)
  CREATE TRIGGER Lecture FK Check Ins
  BEFORE INSERT ON Lecture
  WHEN (NEW. fk lecturer IS NOT NULL AND
        NOT EXISTS (SELECT 1 FROM Lecturer R
                     WHERE R. pk lecturer = NEW. fk lecturer))
   ( ROLLBACK WORK );
  CREATE TRIGGER Lecture FK Check Upd
  BEFORE UPDATE OF fk lecturer ON Lecture
  REFERENCING NEW AS Lnew
  WHEN (Lnew.fk lecturer IS NOT NULL AND
11
        NOT EXISTS (SELECT 1 FROM Lecturer R
12
                     WHERE R. pk lecturer = Lnew.fk lecturer))
13
   ( ROLLBACK WORK );
```

```
Parent-side triggers (on Lecturer)
  CREATE TRIGGER Lecturer Ref Block Del
  BEFORE DELETE ON Lecturer
  REFERENCING OLD AS Pold
  WHEN (EXISTS (SELECT 1 FROM Lecture L
                 WHERE L. fk lecturer = Pold.pk lecturer))
   ( ROLLBACK WORK );
  CREATE TRIGGER Lecturer Ref Block Upd
  BEFORE UPDATE OF pk lecturer ON Lecturer
  REFERENCING OLD AS Pold NEW AS Pnew
10
  WHEN (EXISTS (SELECT 1 FROM Lecture L
11
                 WHERE L. fk lecturer = Pold.pk lecturer))
12
   ( ROLLBACK WORK );
```

- The **child-side** triggers prevent inserting or updating a lecture referencing a non-existent lecturer.
- The **parent-side** triggers prevent deleting or changing a lecturer ID if at least one lecture still refers to it.
- The ROLLBACK WORK aborts the transaction on violation (NO ACTION-like).

# 2b) Simulating FOREIGN KEY (FkExam) REFERENCES Exam(PkExam) ON DELETE SET NULL with triggers

We enforce referential integrity procedurally:

- Child-side (Student). Block INSERT/UPDATE that sets FkExam to a non-existent Exam.PkExam (allow NULL).
- Parent-side (Exam) on DELETE. Before deleting an Exam row, set Student.FkExam := NULL for all referencing students (ON DELETE SET NULL).
- Parent-side (Exam) on UPDATE PkExam. Default is NO ACTION; block if referenced.

```
Child-side checks (INSERT/UPDATE on Student) F
  CREATE TRIGGER Student FK Check Ins
1
  BEFORE INSERT ON Student
2
  WHEN ( NEW.FkExam IS NOT NULL AND
         NOT EXISTS (SELECT 1 FROM Exam E
                      WHERE E. PkExam = NEW. FkExam)
   ( ROLLBACK WORK );
  CREATE TRIGGER Student FK Check Upd
  BEFORE UPDATE OF FkExam ON Student
  REFERENCING NEW AS Snew
10
  WHEN (Snew.FkExam IS NOT NULL AND
11
         NOT EXISTS (SELECT 1 FROM Exam E
12
                      WHERE E. PkExam = Snew . FkExam)
13
   ( ROLLBACK WORK );
```

```
Parent-side action (DELETE on Exam) — simulate ON DELETE SET NULL

CREATE TRIGGER Exam_Delete_SetNull

BEFORE DELETE ON Exam

REFERENCING OLD AS Eold

(
UPDATE Student

SET FkExam = NULL

WHERE FkExam = Eold.PkExam

);
```

```
Parent-side block (UPDATE of PkExam) — simulate ON UPDATE NO ACTION |

CREATE TRIGGER Exam_Block_Update_PK

BEFORE UPDATE OF PkExam ON Exam

REFERENCING OLD AS Eold NEW AS Enew

WHEN (EXISTS (SELECT 1 FROM Student S

WHERE S.FkExam = Eold.PkExam) )

(ROLLBACK WORK);
```

Notes.

- Ensure Student (FkExam) allows NULL; otherwise SET NULL would fail.
- The DELETE trigger clears references before parent removal.
- Using ROLLBACK WORK follows the lecture's "block on violation" pattern.

# Task 3

Using Assertions or Triggers because integrity constraints span multiple relations and not only one (Create Table Statement).

a) Using Assertions because constraints too general and not specific event. e.g. no new orders when ...

U2109: each User must have SUM(c.CreditScore) >= -1000 creditpoints over all credits.

```
CREATE ASSERTION CreditLimitU2109
        CHECK(NOT EXISTS(
            SELECT u.UID
            FROM User u
            JOIN Regulation r ON u.Country = r.Country
            JOIN Credit c ON u.UID = c.UID
            WHERE r.RegulationCode = 'U2109'
            GROUP BY u.UID
            HAVING SUM(c.CreditScore) < -1000))</pre>
        DEFFERABLE INITIALLY DEFERRED
<u>U2304:</u> each User must have SUM(c.CreditScore) >= (-0.1 * MAX(b.BetAmount))
        CREATE ASSERTION CreditLimitU2109
        CHECK(NOT EXISTS(
            SELECT u.UID
            FROM User u
            JOIN Regulation r ON u.Country = r.Country
            JOIN Credit c ON u.UID = c.UID
            JOIN Bet b ON u.UID = b.UID
            WHERE r.RegulationCode = 'U2304'
            GROUP BY u.UID
            HAVING SUM(c.CreditScore) < (-0.1 * MAX(b.BetAmount))))</pre>
        DEFFERABLE INITIALLY DEFERRED
```

#### b) Materialized View:

```
CREATE VIED DailyOnlineTime
REFRESH FORCE ON COMMIT
BUILD DEFERRED AS
SELECT u.UID, u.Name, u.Country, s.Date,
SUM(s.Duration) AS dailyTotalTime
FROM User u
JOIN Regulation r ON u.Country = r.Country
```

JOIN Session s ON u.UID = s.UID

JOIN Bet b ON u.UID = b.UID

WHERE r.RegulationCode = 'M5475'

AND u.Age BETWEEN 14 AND 18

GROUP BY u.Name, u.Country, s.Date;

c)

# Task 4

When we say something is **DEFERRABLE INITIALLY DEFERRED**, it means that the database will wait until the end of a transaction before checking certain rules, such as foreign key constraints. This is useful when two tables depend on each other. For example, a **Lecture** table and a **Lecture** table. A lecture record might point to the lecturer who gives it, while the lecturer record might also depend on the lecture. If we try to insert both of these in the same transaction, the database might raise an error because when we insert the first one, the other does not yet exist. If the rule is **NOT DEFERRABLE**, the database checks immediately, and the operation fails. If it is **DEFERRABLE INITIALLY IMMEDIATE**, we would need to manually tell the database to delay the check. However, if it is **DEFERRABLE INITIALLY DEFERRED**, the database automatically waits until the transaction ends to perform the checks. This allows both inserts to complete first, and the database only verifies the rules afterward, which helps prevent errors.