

Weighted Finite Automata and Noncommutative Rational Series

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1 Weighted finite automata, linear representations, rational series

1.1 Weighted finite automata

Definition 1.1.1. Let K be a semiring and A an alphabet.

- (1) A **weighted (finite) automaton** (WFA) with weights in K is a tuple (Q, I, E, T) consisting of a finite set Q of **states**, and maps $I: Q \rightarrow K$ (**initial weights**), $E: Q \times A \times Q \rightarrow K$ (**transition function**), $T: Q \rightarrow K$ (**terminal weights**).
- (2) A triple (p, a, q) with $E(p, a, q) \neq 0$ is an **edge/transition** with **label** a , **starting state** p , **ending state** q and **weight** $E(p, a, q) \in K$.

- (3) A **path/run** is a sequence of edges $c = (q, a_1, q_1)(q_1, a_2, p_2) \dots (q_{n-1}, a_n, q_n)$. Its **weight** is

$$E(q, a_1, q_1) \cdot E(q_1, a_2, p_2) \cdots E(q_{n-1}, a_n, q_n)$$

and its **label** is $w = a_1 a_2 \dots a_n \in A^*$.

- (4) The **behaviour** of \mathcal{A} is the series $[[\mathcal{A}]] \in K \langle\langle A \rangle\rangle$ defined by

$$([[\mathcal{A}]], w) = \sum_{q_0, \dots, q_n \in Q} I(q_0) E(q_0, a_1, q_1) E(q_1, a_2, p_2) \cdots E(q_{n-1}, a_n, q_n) T(q_n),$$

where $w = a_1 \cdots a_n$, $a_i \in A$.

Definition 1.1.2. Terminology:

1. A state q is **initial** if $I(q) \neq 0$ and **terminal** if $T(q) \neq 0$.
2. A **successful run/accepting run** is a run from an initial state to a terminal state.

1.2 Linear representation, recognizable series

We wish to represent the data in a weighted finite automaton using adjacency matrices.

Definition 1.2.1. Let K be a semiring and A an alphabet.

- (1) A series $S \in K \langle\langle A \rangle\rangle$ is **(K-)recognizable** if there exist $n \geq 0$, $\lambda \in K^{1 \times b}$, $\gamma \in K^{n \times 1}$, and a monoid morphism $\mu: A^* \rightarrow K^{d \times d}$ such that for every $w \in A^*$ we have

$$(S, w) = \lambda \mu(w) \gamma.$$

- (2) The triple (λ, μ, γ) is a **linear representation** of S with **dimension** n .

Proposition 1.2.2. A series $S \in K \langle\langle A \rangle\rangle$ is recognizable if and only if there exists a WFA \mathcal{A} such that $S = [[\mathcal{A}]]$.

Proof. Suppose $S = [[\mathcal{A}]]$ with $\mathcal{A} = (Q, I, E, T)$. Without loss of generality let $Q = \{1, \dots, n\}$. Let

$$\lambda := \begin{bmatrix} I(1) & \dots & I(n) \end{bmatrix}, \quad \gamma := \begin{bmatrix} T(1) & \vdots & T(n) \end{bmatrix}^\top,$$

and for every $a \in A$ let

$$\mu(a) := \begin{bmatrix} E(1, a, 1) & \dots & E(1, a, n) \\ \vdots & & \vdots \\ E(1, a, n) & \dots & E(n, a, n) \end{bmatrix}.$$

This extends to a morphism

$$\mu: A^* \rightarrow K^{d \times d},$$

$$\mu(a_1 \cdots a_n) = \mu(a_1)\mu(a_2) \cdots \mu(a_n).$$

Then, for $p, q \in Q$, $w = a_1 \cdots a_n$, we have

$$\begin{aligned} u(w)_{p,q} &= [\mu(a_1) \cdots \mu(a_n)]_{p,q} \\ &= \sum_{p_1, \dots, p_{n-1}=1}^n \mu(a_1)_{p,p_1} \mu(a_2)_{p_1,p_2} \cdots \mu(a_n)_{p_{n-1},q} \end{aligned}$$

and

$$\begin{aligned} \lambda \mu(w) \gamma &= \sum_{p,p_1, \dots, p_{n-1}, q=1}^n \lambda_p \mu(a_1)_{p,p_1} \mu(a_2)_{p_1,p_2} \cdots \mu(a_n)_{p_{n-1},q} \gamma_q \\ &= \sum_{p,p_1, \dots, p_{n-1}, q=1}^n I(p) E(p, a_1, p_1) E(p_1, a_2, p_2) \cdots E(p_{n-1}, a_n, q) T(q) \\ &= ([[\mathcal{A}]], w). \end{aligned}$$

Conversly, let (λ, μ, γ) be a linear representation recognizing S . Let $Q := \{1, \dots, n\}$, $I(p) := \lambda_p$, $T(q) := \gamma_q$, $E(p, a, q) := \mu(a)_{p,q}$. The computation above shows that S is the behaviour of (Q, I, E, T) . \square

1.3 Model-theoretic characterization

Lemma 1.3.1. Let K be a semiring and A an alphabet.

- (1) For $x \in A^*$, the map $S \mapsto x^{-1}S$ is a K -module morphism.
- (2) For every $x, y \in A^*$ and for every $S \in K \langle\langle A \rangle\rangle$ we have $(xy)^{-1}S = y^{-1}(x^{-1}S)$.

Proof. Left as an exercise. \square

Definition 1.3.2. A submodule $M \subseteq K \langle\langle A \rangle\rangle$ is **stable** if for every $S \in M$ and every $x \in A^*$ we have $x^{-1}S \subseteq M$ (equivalently for every $a \in A$ we have $a^{-1}M \subseteq M$).

Theorem 1.3.3. A series $S \in K \langle\langle A \rangle\rangle$ is recognizable if and only if there exists a stable finitely generated left K -submodule $M \subseteq K \langle\langle A \rangle\rangle$ such that $S \in M$.

Proof. Next lecture. \square

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