Graph Theory

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1 Independence, matching, covers

1.1 Definitions

Let $G = (V, E)$ be a graph.

Definition 1.1.1. A set $S \subseteq V$ is an ***independent set*** if $G[S]$ contains no edges. The ***independence number*** $\alpha(G)$ is the maximum possible cardinality of an independent set.

Definition 1.1.2. A set $T \subseteq V$ is a ***vertex cover*** if

$$\forall e \in E.\ T \cap e \neq \emptyset.$$

The ***vertex cover number*** $\beta(G)$ is the minimum possible cardinality of a vertex cover.

Definition 1.1.3. A set $M \subseteq E$ is a ***matching*** if

$$\forall e, f \in M.\ e \neq f \Rightarrow e \cap f = \emptyset.$$

The ***matching number*** $\alpha'(G)$ is the maximum possible cardinality of a matching.

Definition 1.1.4. A set $C \subseteq E$ is an ***edge cover*** if every vertex of G is covered by at least one edge from C. If $\delta(G) \geq 1$, we define the ***edge cover number*** $\beta'(G)$ as the minimum possible cardinality of an edge cover.

Lemma 1.1.5. Let G be a graph. The following holds:

- $\alpha(G) + \beta(G) = n(G)$.
- $\alpha'(G) \leq \beta(G)$.
- $\alpha(G) \leq \beta'(G)$.
- If $\delta(G) \geq 1$, then $\alpha'(G) \leq \frac{n}{2} \leq \beta(G)$.

Theorem 1.1.6 (Gallai)**.** If $\delta(G) \geq 1$, then $\alpha'(G) + \beta'(G) = n(G)$.

1.2 Matchings

Definition 1.2.1. Let M be a matching. A path is an ***M-alternating path*** if the edges along the path alternate between M and $\overline{M} = E \setminus M$.

Definition 1.2.2. An M-alternating path is called an ***M-augmenting path*** if both ends of the path are uncovered by M.

Proposition 1.2.3. Let G be a graph and M a matching. If there exists an M-augmenting path P, then M is not a maximum matching.

Proof. We can construct a bigger matching $M' = M \triangle E(P)$, where $\triangle$ is the disjunctive union. $\square$

Theorem 1.2.4 (König)**.** Let G be a bipartite graph. Then the following holds:

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(a) $\alpha'(G) = \beta(G)$.

(a) If M is a matching witn no M-augmenting path, then M is a maximum matching.

Remark 1.2.5. There also exist graphs with $\alpha'(G) = \beta(G)$ that are not bipartite.

Corollary 1.2.6. If G is a bipartite graph, then $\alpha(G) = \beta'(G)$

Proof. We have

$$\alpha(G) = n(G) - \beta(G) = n(G) - \alpha'(G) = \beta'(G),$$

where the latter two equalities are due to König's and Gallai's theorem respectively. □

Definition 1.2.7. Let G be a bipartite graph with partite classes A and B. ***Hall's condition*** (HC) holds for A, if

$$\forall S \subseteq A.\ |S| \leq |N(S)|\,, \tag{HC}$$

where $N(S)$ is the neighbourhood of S.

Theorem 1.2.8 (Hall)**.** Let $G = (A \cup B, E)$ be a bipartite graph. A matching that covers A exists if and only if Hall's condition holds for A.

Definition 1.2.9. A matching M is a ***perfect matching*** if it covers all vertices.

Corollary 1.2.10. Let G be a bipartite graph. There exists a perfect matching of G if and only if $|A| = |B|$ and A satisfies Hall's condition.

Definition 1.2.11. Let $S \subseteq A$. The ***deficiency*** of S is defined as $\mathrm{def}(S) = |S| - |N(S)|$.

Theorem 1.2.12. Let $G = (A \cup B, E)$ be a bipartite graph and M a matching. Then at most

$$|A| - \max_{S \subseteq A}(\mathrm{def}(S))$$

vertices of A are covered.

Theorem 1.2.13. If G is a regular bipartite graph, then G has a perfect matching.

Theorem 1.2.14. Let M be a matching in G. There exists an M-augmenting path in G if and only of M is not maximum.

The Blossom algorithm is a known algorithm in $O(n\sqrt{n})$ that finds M-augmenting paths (in polynomial time). It also provides a maximum matching, and it allows us to find $\alpha'(G)$ and $\beta'(G)$ in polynomial time.

Let $o(G)$ be the number of odd components ($|V(C)| \equiv 1 \pmod 2$) in G.

Theorem 1.2.15. Tutte A graph G has a perfect matching if and only if

$$\forall S \subseteq V(G).\ |S| \geq o(G \setminus S). \tag{TC}$$

The condition (TC) is called ***Tutte's condition***.

Remark 1.2.16. In bipartite graphs (TC) implies (HC).

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Theorem 1.2.17 (Berge-Tutte formula)**.** A maximum matching in a graph G leaves exactly

$$\max_{S \subseteq V(G)} \left(o(G \setminus S) - |S|\right).$$

vertices uncovered. Equivalently,

$$\alpha'(G) = \frac{1}{2}\left(n - \max_{S \subseteq V(G)} \{o(G \setminus S) - |S|\}\right).$$

1.3 Factors

Definition 1.3.1. A ***factor*** is a spanning subgraph (subgraph that contains all vertices). A k***-factor*** is a k-regular spanning subgraph.

Theorem 1.3.2 (Petersen)**.** If G is a cubic graph witha at most one bridge, then G has a 1-factor.

Theorem 1.3.3 (Petersen)**.** If G is a cubic graph and all cut edges lie on the same path then G has a 1-factor.

Proof. Omitted. □

Example 1.3.4. This is the smallest example of a cubic graph with no 1-factor.

Theorem 1.3.5 (Petersen)**.** Every bridgeless cubic graph decomposes into a 1-factor and 2-factor.

Theorem 1.3.6. If G is a k-regular graph and k is even, then G has a 2-factor.

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2 Connectivity

2.1 k-connectivity

Definition 2.1.1. The ***connectivity number*** $\kappa(G)$ is the minimum number of vertices in $S \subseteq V(G)$ such that $G - S$ is disconnected or contains only one vertex.

Remark 2.1.2. The latter conditions handles the case of complete graphs.

Definition 2.1.3. A graph G is k-connected if $\kappa(G) \geq k$.

Alternatively: G is k-connected if the removal of at most $k-1$ vertices always results in a connected graph with at most two vertices.

Proposition 2.1.4.

1. $\kappa(G) \leq \delta(G)$
2. $\kappa(G) \leq \beta(G)$
3. $\kappa(G) \leq n(G) - 2$

Theorem 2.1.5. The minimum number of edges in a k-connected graph of order $n > k > 2$ is $\lceil nk/2 \rceil$.

Proof. Using ***Harary graphs*** $H_{n,k}$. □

Definition 2.1.6. A set $F \subseteq E(G)$ is a ***disconnecting set*** if $E - F$ is disconnected

Definition 2.1.7. Let $\emptyset \neq A \subsetneq V(G)$ and let $E(A, \overline{A})$ be the set oe edges between A and $\overline{A}$. Then $E(A, \overline{A})$ is an ***edge cut***.

Clearly, an edge cut is a disconnecting set. Similarly, a minimal disconnecting set is an edge cut.

Definition 2.1.8. The ***edge-connectivity numver*** of G, $\kappa'(G)$ is the minimum number of edges in a disconnecting set.

Definition 2.1.9. A graph is ***k-edge-connected*** if the removal of less than k edges always leaves a connected graph.

Equivalently, a graph G is k-edge-connected if $k \leq \kappa'(G)$.

Theorem 2.1.10. Let G be a simple graph with $n(G) \geq 2$. Then $\kappa(G) \leq \kappa'(G) \leq \delta(G)$.

Corollary 2.1.11. Let G be a graph.

- If G is k-connected, then G is k-edge-connected.
- The minimum number of edges in a k-edge-connected graph on $n > k \geq 2$ vertices is $\lceil kn/2 \rceil$.

Theorem 2.1.12 (Whitney)**.** A graph G is 2-connected if and only if for every pair of distinct vertices $u, v \in V(G)$ there exist two internally disjoint u, v-paths.

Lemma 2.1.13 (Expansion)**.** If G is a k-connected and we add a new vertex v and k incident edges to the graph, we obtain a k-connected graph.

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A Exercises

Exercise sheet 1

1. For each of the following graphs G, determine $\alpha(G)$, $\alpha'(G)$, $\beta(G)$, and $\beta'(G)$:

   (a) $G = P_n$

   (b) $G = C_n$

2. Prove that $\frac{\beta(G)}{2} \leq \alpha'(G) \leq \beta(G)$ for any graph G.

3. (From lectures) Let G be a bipartite graph with bipartition $\{A, B\}$. The deficiency of the subset $U \subseteq A$ is defined as:

$$\mathrm{def}_G(U) := |U| - |N_G(U)|,$$

   and the deficiency of G is defined as:

$$\mathrm{def}(G) := \max_{U \subseteq A} \mathrm{def}_G(U).$$

   Prove that $\alpha'(G) = |A| - \mathrm{def}(G)$.

4. Let $F = \{F_1, \ldots, F_n\}$ be a family of subsets of a set Y. Prove there are distinct elements $a_1, \ldots, a_n$ such that $a_i \in F_i$ if and only if $|F| \leq |\cup_{S \in F} S|$ for all $F \subseteq S$. (Such a set is called a system of distinct representatives for F.)

5. (Extra) Let Γ be a finite group and $H \leq \Gamma$ with index n. Let $L = \{L_i\}_{i=1}^n$ be the left cosets of H and $R = \{R_i\}_{i=1}^n$ be the right cosets of H. Prove that there is a subset $\{h_1, \ldots, h_n\} \subseteq G$ such that $L = \{h_i H\}_{i=1}^n$ and $R = \{H h_i\}_{i=1}^n$.

Exercise sheet 2

1. For each of the following graphs G, determine the number of maximum matchings:

   (a) $G = K_n$

   (b) $G = K_{a,b}, a \leq b$

2. Let G be a graph such that $\delta(G) \geq 1$, with maximal matching M and minimal edge cover C. Prove the following equivalences:

   (a) M is a maximum matching if and only if M is contained in a minimum edge cover.

   (b) C is a minimum edge cover if and only if C contains a maximum matching.

3. Use deficiency (see sheet 1) to prove the König-Egerváry theorem. (Hint: find a matching and a vertex cover of the same size using a subset of maximum deficiency).

4. (From lectures) Let G be a bipartite graph with bipartition $\{A, B\}$ such that $|A| = |B|$. Prove that for A, Hall's condition holds if and only if Tutte's condition holds.

5. Prove Theorem 1.2.17 (Berge-Tutte formula).

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