Weighted Finite Automata and Noncommutative Rational Series

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1 WFA, Linear Representations, Rational Series

1.1 Weighted Finite Automata

Definition 1.1.1. Let K be a semiring and A an alphabet.

- (1) A weighted (finite) automaton (WFA) with weights in K is a tuple (Q, I, E, T) consisting of a finite set Q of states, and maps $I: Q \to K$ (initial weights), $E: Q \times A \times Q \to K$ (transition function), $T: Q \to K$ (terminal weights).
- (2) A triple (p, a, q) with $E(p, a, q) \neq 0$ is an **edge/transition** with **label** a, **starting state** p, **ending state** q and **weight** $E(p, a, q) \in K$.
- (3) A **path/run** is a sequence of edges $c = (q, a_1, q_1)(q_1, a_2, p_2) \dots (q_{n-1}, a, q_n)$. Its **weight** is

$$E(q, a_1, q_1) \cdot E(q_1, a_2, p_2) \cdots E(q_{n-1}, a_n, q_n)$$

and its label is $w = a_1 a_2 \dots a_n \in A^*$.

(4) The **behaviour** of \mathcal{A} is the series $[[\mathcal{A}]] \in K \langle \langle A \rangle \rangle$ defined by

$$([[\mathcal{A}]], w) = \sum_{q_0, \dots, q_n \in Q} I(q_0) E(q, a_1, q_1) E(q_1, a_2, p_2) \cdots E(q_{n-1}, a_n, q_n) T(q_n),$$

where $w = a_1 \cdots a_n, a_i \in A$.

Definition 1.1.2. Terminology:

- 1. A state q is **initial** if $I(q) \neq 0$ and **terminal** if $T(q) \neq 0$.
- 2. A *successful run/accepting run* is a run from an initial state to a terminal state.

1.2 Linear representation, recognizable series

Definition 1.2.1. Let K be a semiring and A an alphabet.

(1) A series $S \in K \langle \langle A \rangle \rangle$ is **(K-)recognizable** if there exist $n \geq 0$, $\lambda \in K^{1 \times b}$, $\gamma \in K^{n \times 1}$, and a monoid morphism $\mu \colon A^* \to K^{d \times d}$ such that for every $w \in A^*$ we have

$$(S, w) = \lambda \mu(w) \gamma.$$

Proposition 1.2.2. A series $S \in K\langle\langle A \rangle\rangle$ is recognizable if and only if there exists a WFA \mathcal{A} such that $S = [[\mathcal{A}]]$.