

Information flow Libraries

- privacy concerns in software systems
 - code must not leak sensitive data to the internet
- ↳ Restrict data propagation (**Information flow**)

Haskell Library: Mandatory Access Control (MAC)

What is IFC?

- Specify how information may propagate in the system:
 - "Sensitive inputs may not flow"
- Track data flows
- Detect data leaks (and suppress)

Building IFC is hard

We need DSLs, custom oracles (operating systems, browsers, etc.), compilers, ...

This is easier in Haskell!

- Haskell pure abstractions can directly express IFC analysis (no side-effects!)
- No need for custom compilers just an API for Haskell.

Several Approaches

- MAC : Static
- HLIO : Hybrid
- LIO : Dynamic

Challenge: Write a MAC

program that leaks but still obeys MAC properties

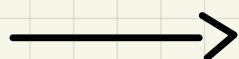
Haskell's type system restricts where I/O is allowed

These libraries prevent leaks by:

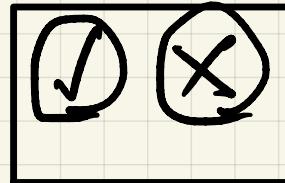
- Untrusted code may perform I/O only through IFC library. There is an IO monad we can only do IO when wrapped in.
- Wrapped with security types.

Example

secret password



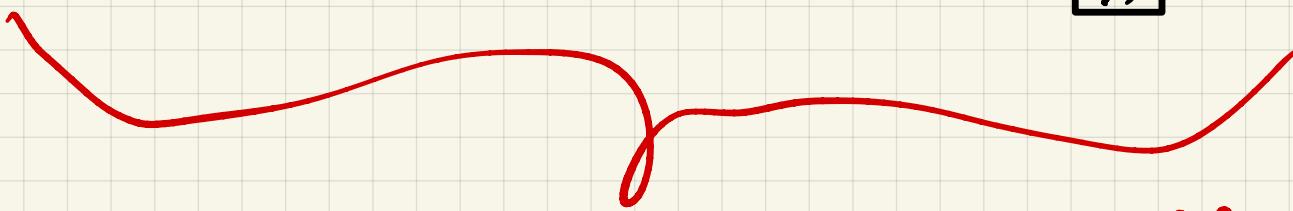
Passwd APP



Leaked password DB



Attack Server



isWeakPwd :: String → Bool (No side effects)



↑
No queries!

isWeakPwd :: String → IO Bool (Can leak!)

- Public observable outputs can leak password
- IO might be used incorrectly

Sol. IFC libs disallow public outputs when data is in scope.

MAC

- Multi-parameter type classes : Reuse type system
- Safe Haskell : untrusted code can't cheat the ts
- Small ~200 LOC

How to specify security policies?

- Security labels data L and data H
- "Can-flow-to" type class represents order between labels:
 $\text{class } L \sqsubseteq L'$ where
give class instances for allowed information flows:
 $\text{instance } L \sqsubseteq L \text{ where}$
 $\text{instance } L \sqsubseteq H \text{ where}$
 $\text{instance } H \sqsubseteq H \text{ where}$
- \sqsubseteq : flow is reflexive and transitive (antisymmetric)

Label data

- Define ADT for labeled data
`newtype Labeled l a = Labeled a`
- These types explicitly assign labels to data
`Password :: Labeled H String`

- ADT for secure computations
 $\text{newty}\ \text{pe}\ \text{MAC}\ l\ a = \text{MAC}(\text{IO}\ a)$
 instance Monad (MAC L) where...
- Encapsulate IO actions that don't leak
- Handle data at security level L
 $\text{readPwdFile} :: \text{MAC}\ H\ \text{String}$
- Only trusted code can run secure computations

How do we not leak?

- Follows MAC principles:
 - 1) No read-up: IO actions may not read data at higher levels
 - 2) No write-down: IO actions may not write data at lower levels

We have safety from the module system in which for:

$$\begin{aligned} \text{newty}\ \text{pe}\ \text{Labeled}\ L\ a &= \text{Labeled}\ a \\ \text{label} :: L \subseteq h \Rightarrow a \rightarrow \text{MAC}\ L(\text{Labeled}\ h\ a) \\ \text{unlabel} :: L \subseteq h \Rightarrow \text{Labeled}\ L\ a \rightarrow \text{MAC}\ h\ a \end{aligned}$$

- You cannot call unlabel when it's unsafe and the password is safe behind a wrapper.

$$\begin{aligned} \text{ToLabeled} :: L \subseteq h \Rightarrow \\ \text{MAC}\ h\ a \\ \rightarrow \text{MAC}(L(\text{Labeled}\ h\ a)) \end{aligned}$$