

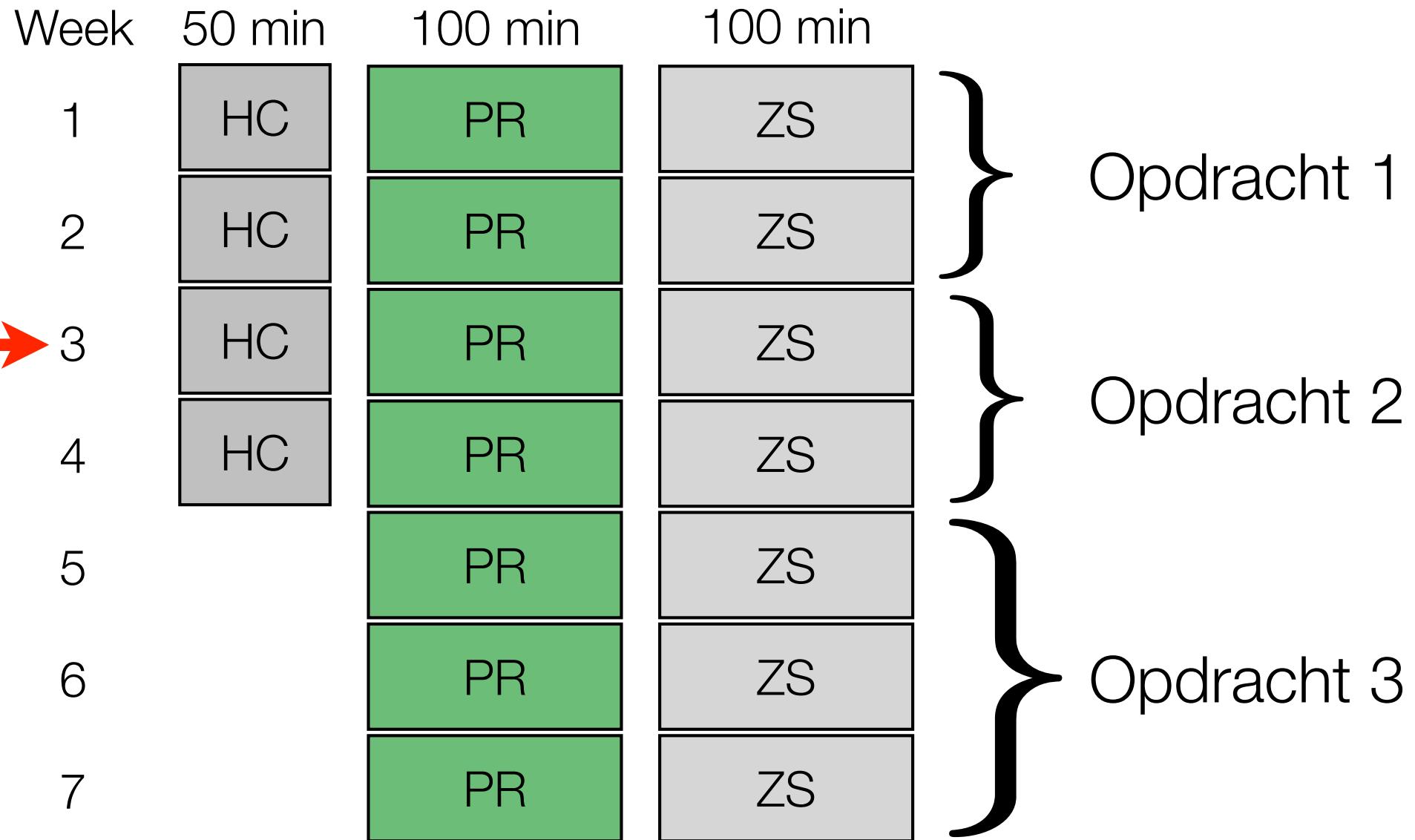


# Wiskunde 13/14

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September 17th 2012 version 1.0

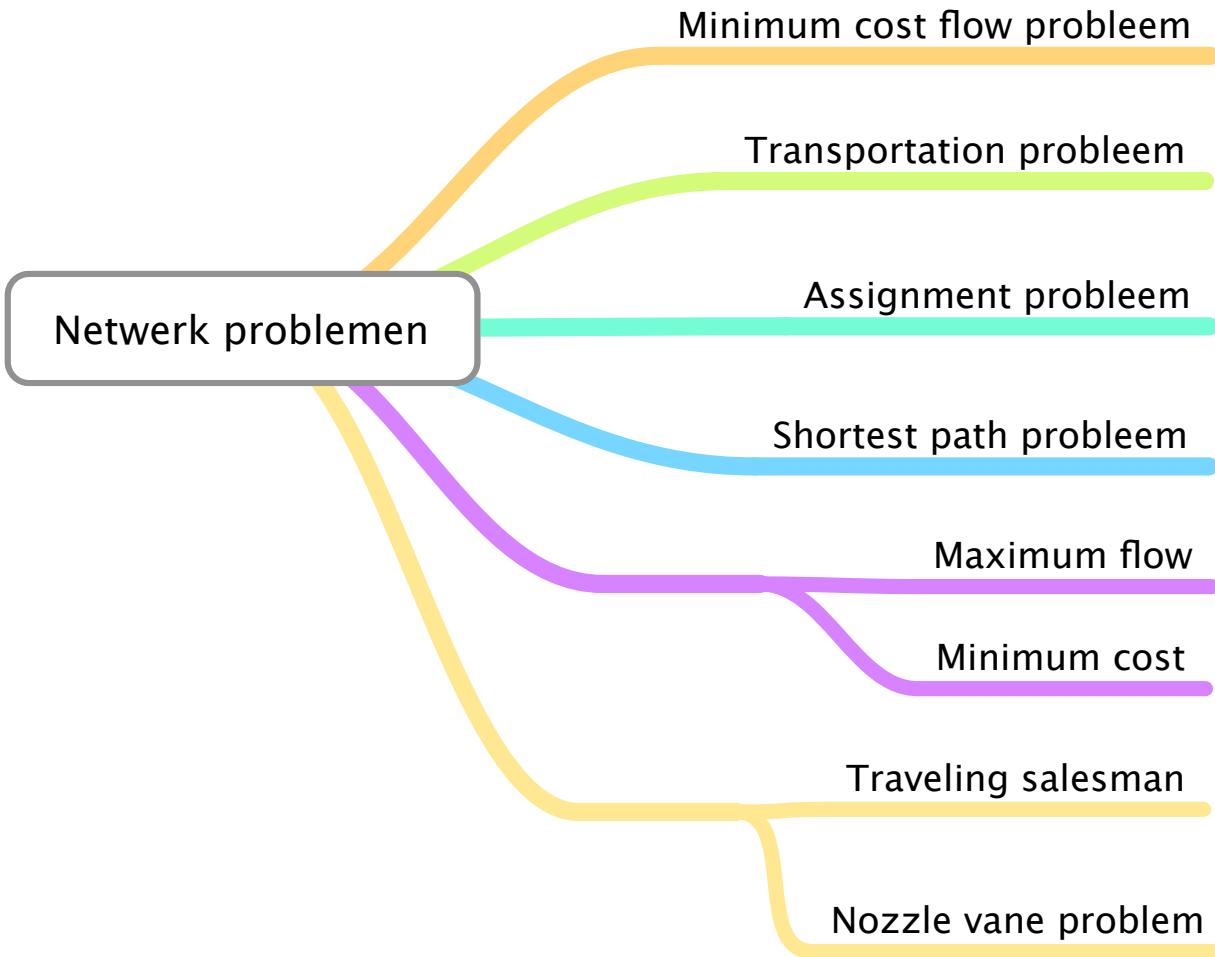
## Optimization : Les 3

# Organisatie van de cursus



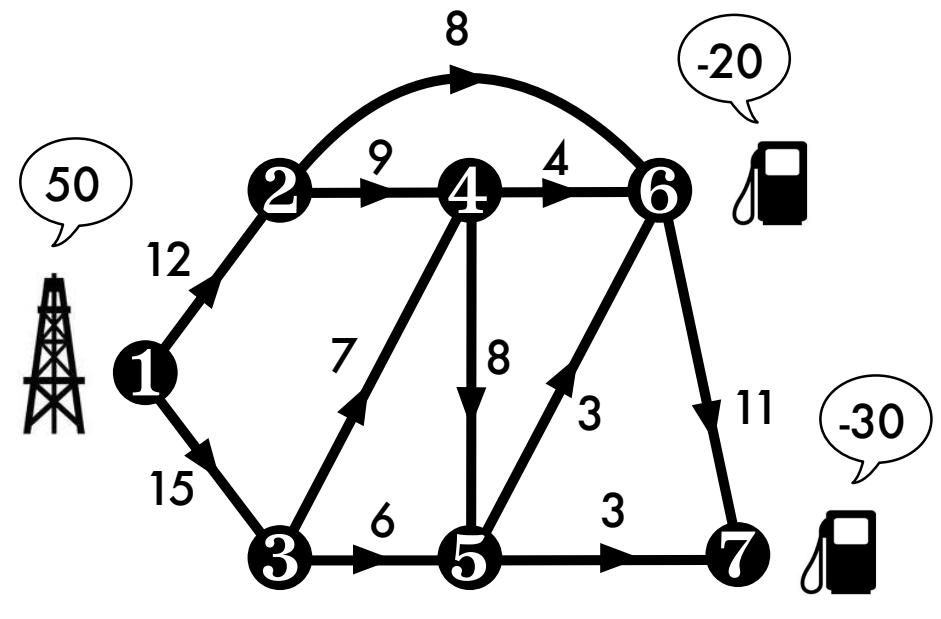
# Overzicht

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# Network problems : Minimum cost flow network

- Nodes : **1, 2, ..., 7**
- Arcs : **(1,2), (1,3), (2,4), ..., (6,7)**
- Cost coefficients :  $c_{1,2}, c_{2,3}, \dots, c_{6,7}$
- Flow :  $x_{1,2}, x_{2,3}, \dots, x_{6,7}$
- Sources : Node 1, Sinks : Nodes 6,7
- Constraints :
  - Flow  $\sum_j x_{i,j} - \sum_k x_{k,i} = b_i$ ,
  - Source  $S \equiv \sum_{\{i: b_i > 0\}} b_i$
  - Sink  $D \equiv - \sum_{\{i: b_i < 0\}} b_i$ .



# Network problems : Minimum cost flow network

$$\text{minimize} \quad z = 12x_{1,2} + 15x_{1,3} + 9x_{2,4} + 8x_{1,2} + 7x_{3,4} + \\ x_{3,5} + 8x_{4,5} + 4x_{4,6} + 3x_{5,7} + 11x_{6,7}$$

subject to

$$x_{1,2} + x_{1,3} = 50$$

$$x_{2,4} + x_{2,6} - x_{1,2} = 0$$

$$x_{3,4} + x_{3,5} - x_{1,3} = 0$$

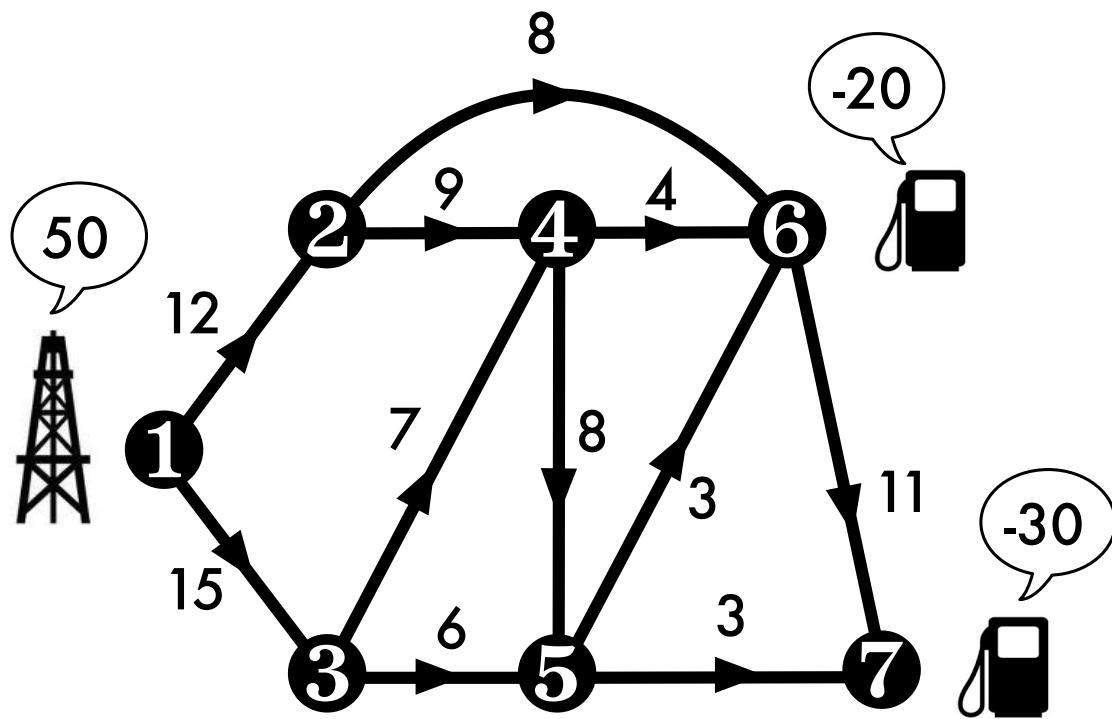
$$x_{4,5} + x_{4,6} - x_{2,4} - x_{3,4} = 0$$

$$x_{5,6} + x_{5,7} - x_{3,5} - x_{4,5} = 0$$

$$x_{6,7} - x_{2,6} - x_{4,6} - x_{5,6} = -20$$

$$-x_{5,7} - x_{6,7} = -30$$

$$0 \leq x \leq 30$$



# Network problems : Minimum cost flow network

The screenshot shows the Wolfram Mathematica interface with the title bar "Network flow problem example HC 3 Equations". The menu bar includes "Wolfram Mathematica FOR STUDENTS", "Demonstrations", "MathWorld", "Student Forum", and "Help". The main workspace displays the following Mathematica code:

```
In[59]:= ClearAll["Global`*"]

In[60]:= z = 12 x12 + 15 x13 + 9 x24 + 8 x26 + 7 x34 + 6 x35 + 8 x45 + 4 x46 + 5 x56 + 3 x57 + 11 x67;
c1 = x12 + x13 == 50;
c2 = x24 + x26 - x12 == 0 ;
c3 = x34 + x35 - x13 == 0 ;
c4 = x45 + x46 - x24 - x34 == 0 ;
c5 = x56 + x57 - x35 - x45 == 0 ;
c6 = x67 - x26 - x46 - x56 == -20 ;
c7 = -x57 - x67 == -30;
c8 = x12 >= 0 && x13 >= 0 && x24 >= 0 && x26 >= 0 && x34 >= 0 && x35 >= 0 &&
    x45 >= 0 && x46 >= 0 && x56 >= 0 && x57 >= 0 && x67 >= 0 ;
c9 = x12 <= 30 && x13 <= 30 && x24 <= 30 && x26 <= 30 && x34 <= 30 && x35 <= 30 &&
    x45 <= 30 && x46 <= 30 && x56 <= 30 && x57 <= 30 && x67 <= 30;
var = {x12, x13, x24, x26, x34, x35, x45, x46, x56, x57, x67};
Minimize[{z, c1 && c2 && c3 && c4 && c5 && c6 && c7 && c8 && c9}, var, Reals]

Out[71]= {1120, {x12 -> 20, x13 -> 30, x24 -> 0, x26 -> 20,
    x34 -> 0, x35 -> 30, x45 -> 0, x46 -> 0, x56 -> 0, x57 -> 30, x67 -> 0}}
```

The code defines variables and constraints for a minimum cost flow problem, then uses the `Minimize` function to find the solution.

# Network problems : Minimum cost flow network

- Formuleer als matrix vergelijking :

minimize     $z = c^T x$   
subject to     $Ax = b$   
                   $\ell \leq x \leq u,$

- Los op met Mathematica :

The screenshot shows a Mathematica notebook interface. The title bar reads "Network flow problem example HC 3 Matrix". The notebook contains the following code:

```
In[13]:= ClearAll["Global`*"]
In[14]:= c = {12, 15, 9, 8, 7, 6, 8, 4, 5, 3, 11};
           b = {50, 0, 0, 0, 0, -20, -30};

A = {{1, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0},
      {-1, 0, 1, 1, 0, 0, 0, 0, 0, 0, 0, 0},
      {0, -1, 0, 0, 1, 1, 0, 0, 0, 0, 0, 0},
      {0, 0, -1, 0, -1, 0, 1, 1, 0, 0, 0, 0},
      {0, 0, 0, 0, 0, -1, -1, 0, 1, 1, 0, 0},
      {0, 0, 0, -1, 0, 0, 0, -1, -1, 0, 1, 0},
      {0, 0, 0, 0, 0, 0, 0, 0, 0, -1, -1, 0}};

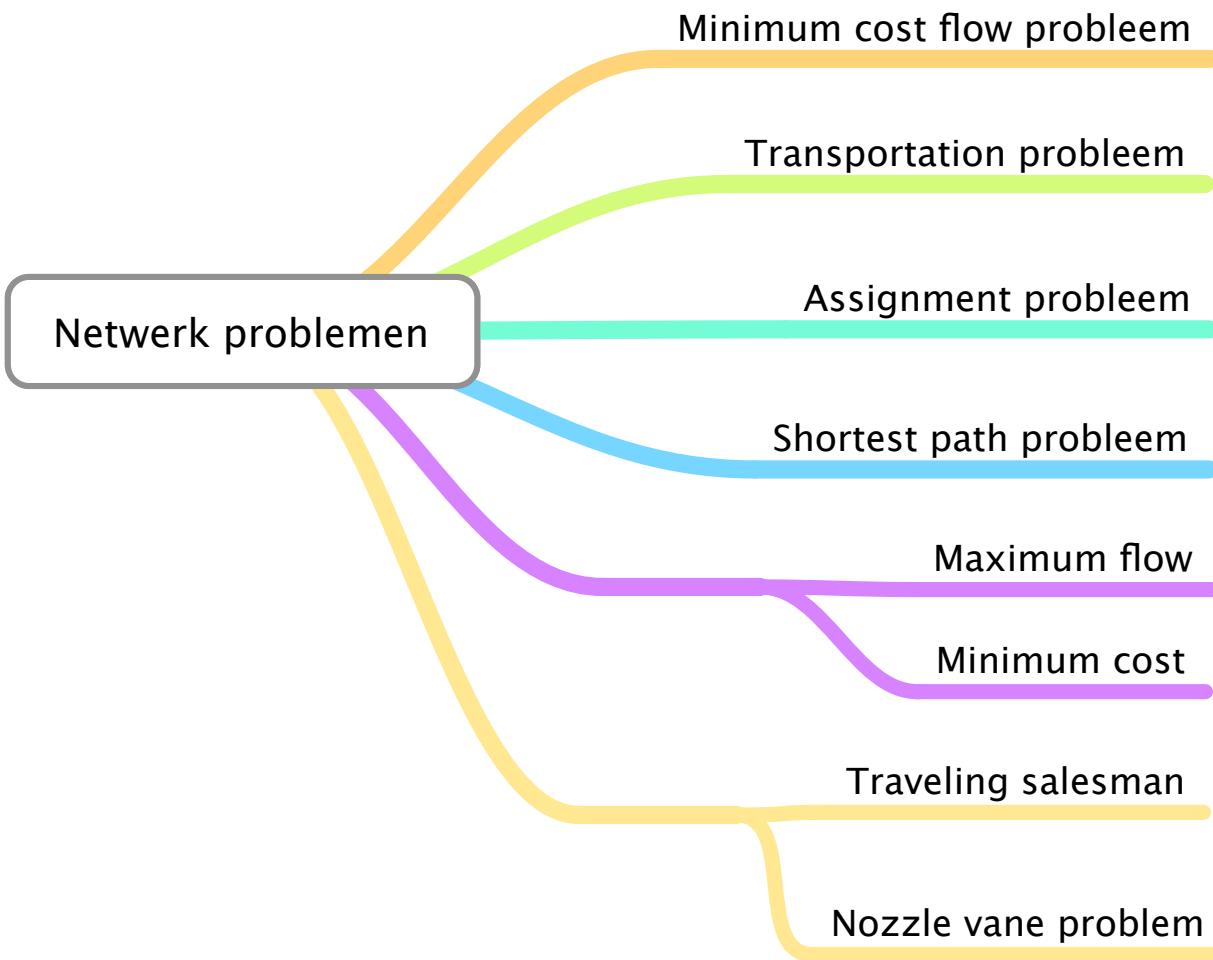
In[35]:= x = LinearProgramming[c, A, b]
           z = {c}.Transpose[{x}];
           Flatten[z, 1][[1]]
Out[35]= {20, 30, 0, 20, 0, 30, 0, 0, 0, 30, 0}
Out[36]= 1120
```

The notebook also displays the output of the code at the bottom:

```
100% ▶
```

# Overzicht : Transportation probleem

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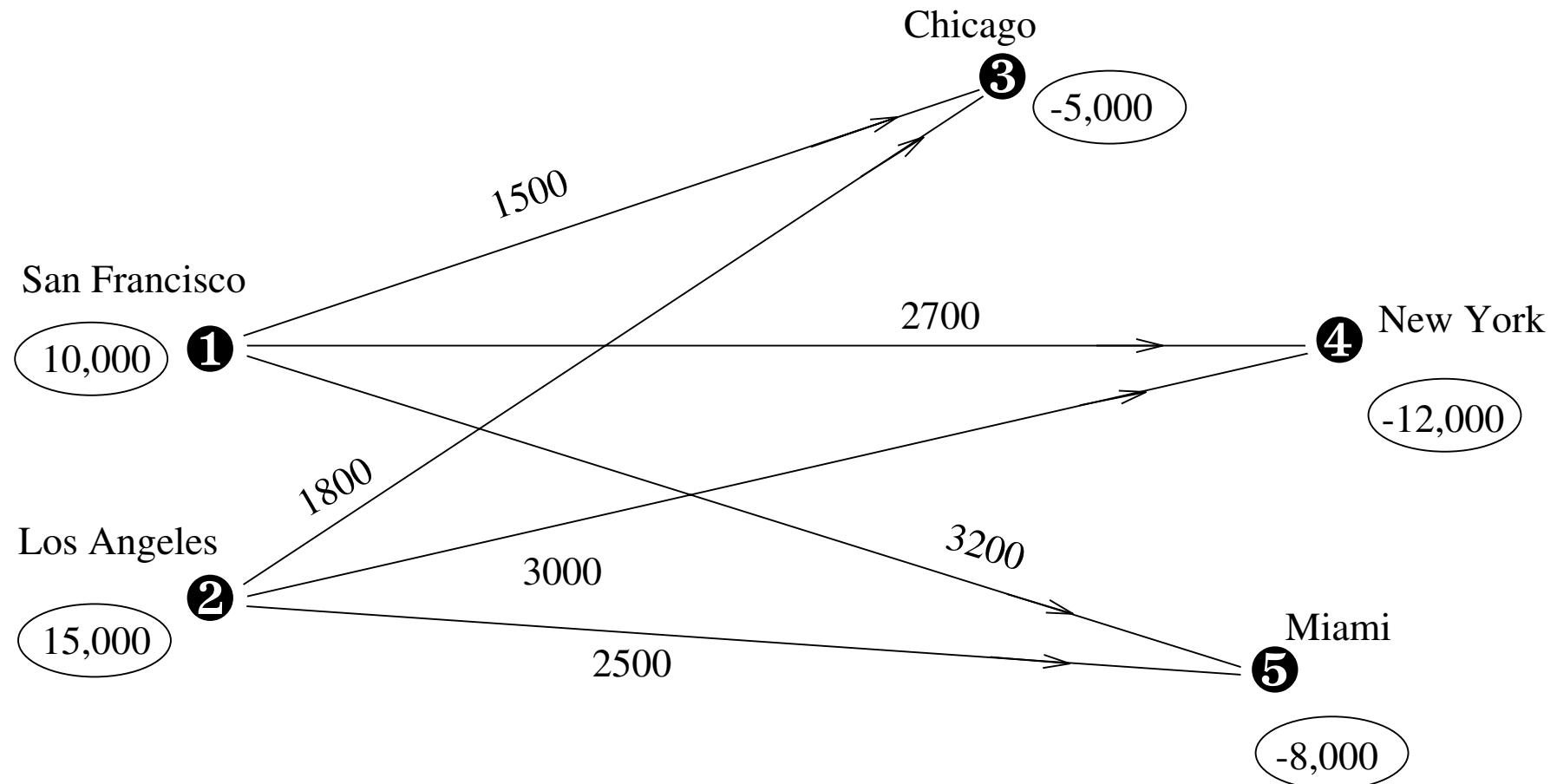
# Network problems : Transportation probleem

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- Een transportation problem modelleert de verplaatsing van goederen van leveranciers naar klanten. Aan verplaatsingen zijn kosten verbonden
- Elke node is of een *source* of een *sink*
- Elke arc gaat van een *source node* naar een *sink node*
- *De flow conservation constraints* zijn :

$$\begin{aligned} \text{• voor een source met } b_i > 0, \quad \sum_j x_{i,j} = b_i \\ \text{• voor een sink met } b_i < 0. \quad - \sum_k x_{k,i} = b_i \end{aligned}$$

# Network problems : Transportation probleem



# Network problems : Transportation probleem

The screenshot shows the Wolfram Mathematica interface for students. The title bar reads "Transportation problem example HC 3 Equations". The menu bar includes "Wolfram Mathematica FOR STUDENTS", "Demonstrations", "MathWorld", "Student Forum", and "Help". The main workspace displays the following Mathematica code:

```
In[95]:= ClearAll["Global`*"]

In[105]:= z = 1500 x13 + 1800 x23 + 2700 x14 + 3000 x24 + 3200 x15 + 2500 x25;
c1 = x13 + x14 + x15 == 10000;
c2 = x23 + x24 + x25 == 15000;
c3 = -x13 - x23 == -5000;
c4 = -x14 - x24 == -12000;
c5 = -x15 - x25 == -8000;
c6 = x13 >= 0 && x14 >= 0 && x15 >= 0 && x23 >= 0 && x24 >= 0 && x25 >= 0;
var = {x13, x14, x15, x23, x24, x25};
Minimize[{z, c1 && c2 && c3 && c4 && c5 && c6}, var, Reals]

Out[113]= {62000000, {x13 -> 0, x14 -> 10000, x15 -> 0, x23 -> 5000, x24 -> 2000, x25 -> 8000}}
```

The code defines variables and constraints for a transportation problem, then uses the Minimize function to find the optimal solution.

# Network problems : Transportation probleem

The screenshot shows a Mathematica notebook interface with the title "Transportation problem example HC 3 Matrix". The notebook contains the following code:

```
In[114]:= ClearAll["Global`*"]

In[115]:= c = {1500, 2700, 3200, 1800, 3000, 2500};
           b = {10 000, 15 000, -5000, -12 000, -8000};

A = 
$$\begin{pmatrix} 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 \\ -1 & 0 & 0 & -1 & 0 & 0 \\ 0 & -1 & 0 & 0 & -1 & 0 \\ 0 & 0 & -1 & 0 & 0 & -1 \end{pmatrix};$$


In[118]:= x = LinearProgramming[c, A, b]
           z = {c}.Transpose[{x}];
           Flatten[z, 1][[1]]

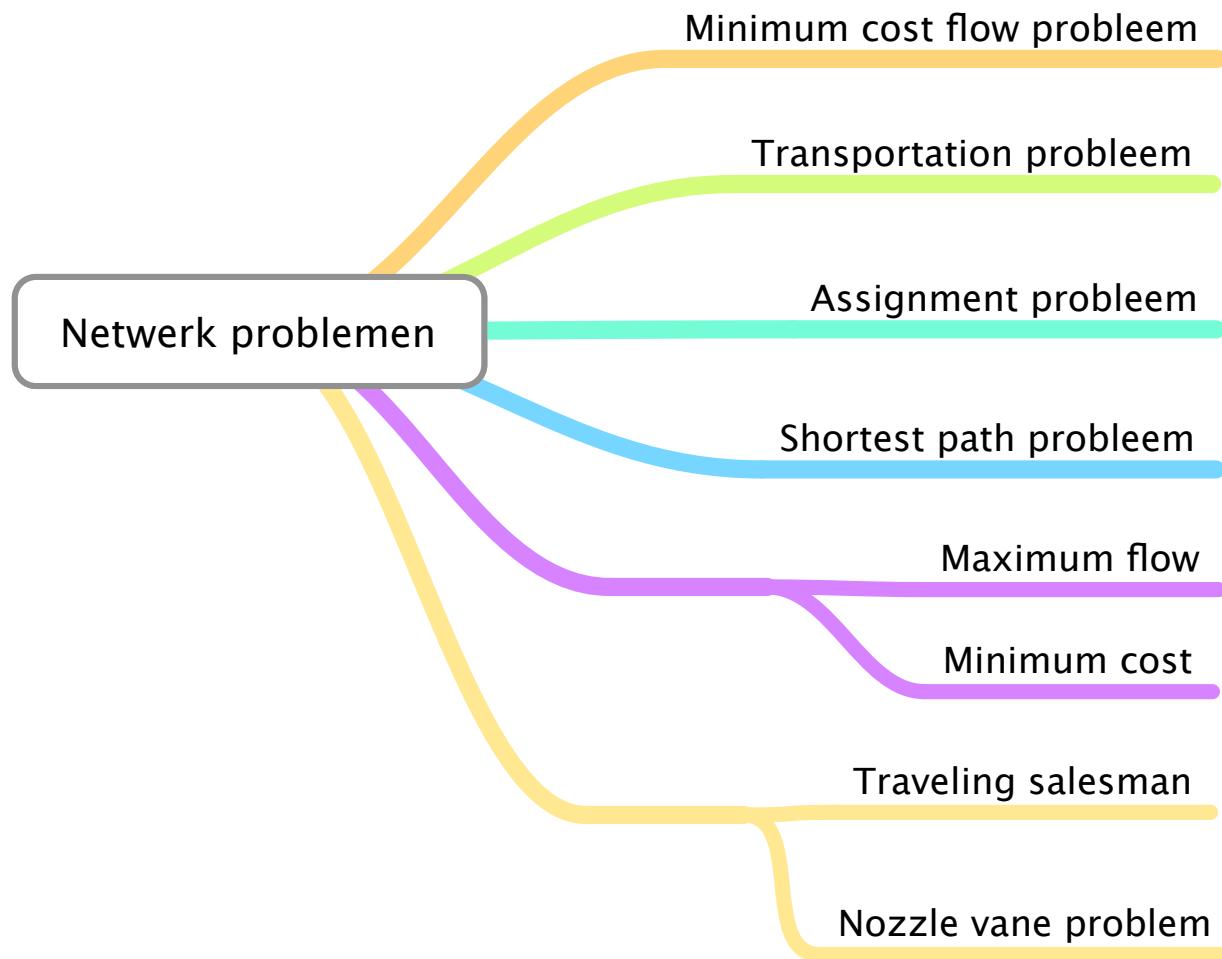
Out[118]= {0, 10 000, 0, 5000, 2000, 8000}

Out[119]= 62 000 000
```

The notebook also includes a zoom control at the bottom left and a status bar at the bottom right indicating "100%".

# Overzicht : Assignment probleem

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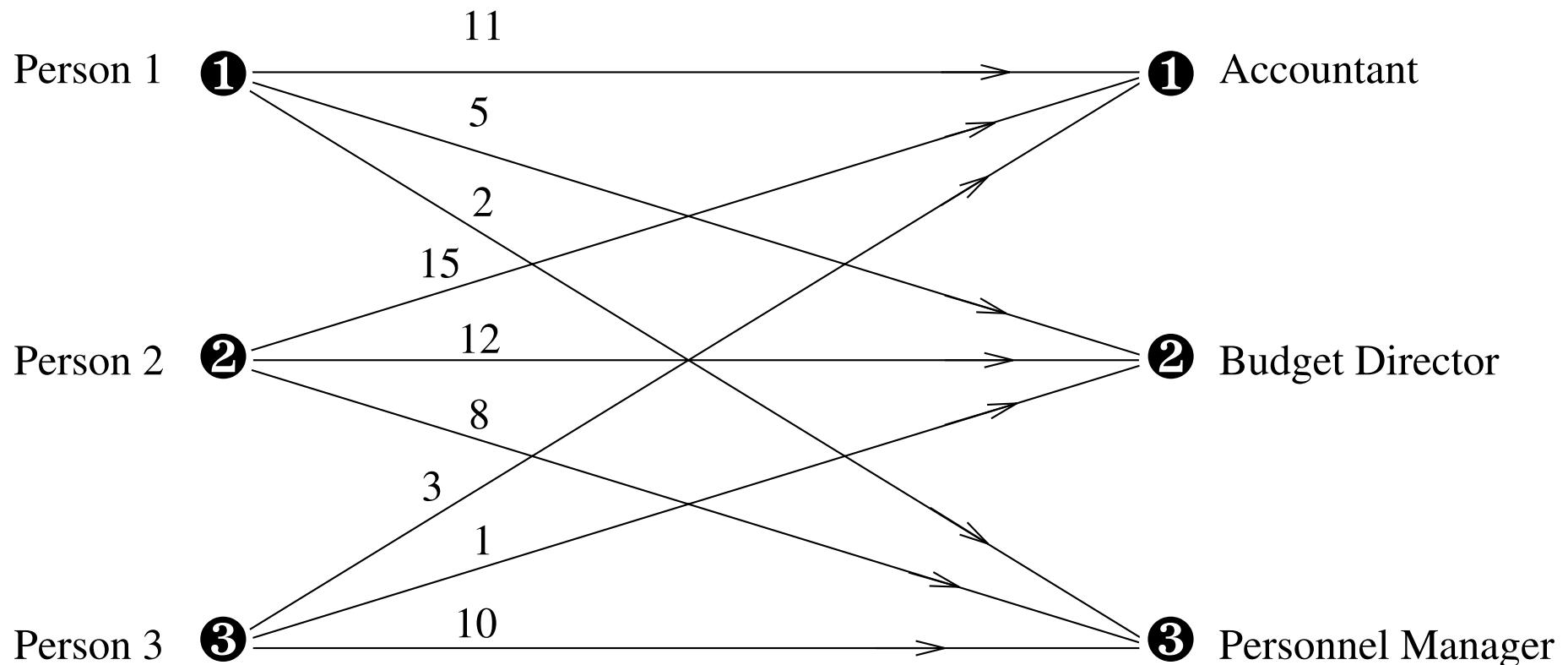
# Network problems : Assignment probleem

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- Persoon 1 :
  - Business degree
  - Just out of school
- Persoon 2 :
  - Business degree
  - 10 years of corporate experience
  - Management experience
- Persoon 3:
  - Management experience
  - Anthropology
- Functie 1: Accountant
- Functie 2: Budget Director
- Functie 3 : Personnel Manager

# Network problems : Assignment probleem

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# Network problems : Assignment probleem

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$$\begin{aligned} \text{maximize } z = & 11x_{1,1} + 5x_{1,2} + 2x_{1,3} + 15x_{2,1} + 12x_{2,2} \\ & + 8x_{2,3} + 3x_{3,1} + 1x_{3,2} + 10x_{3,3} \end{aligned}$$

subject to the constraints

$$x_{1,1} + x_{1,2} + x_{1,3} = 1$$

$$x_{2,1} + x_{2,2} + x_{2,3} = 1$$

$$x_{3,1} + x_{3,2} + x_{3,3} = 1$$

$$-x_{1,1} - x_{2,1} - x_{3,1} = -1$$

$$-x_{1,2} - x_{2,2} - x_{3,2} = -1$$

$$-x_{1,3} - x_{2,3} - x_{3,3} = -1$$

$$0 \leq x \leq 1.$$

# Network problems : Assignment probleem

The screenshot shows the Wolfram Mathematica interface with the title bar "Assignment problem example HC 3 Equations". The menu bar includes "Wolfram Mathematica FOR STUDENTS", "Demonstrations", "MathWorld", "Student Forum", and "Help". The main workspace displays the following Mathematica code:

```
In[154]:= ClearAll["Global`*"]

In[155]:= z = 11 x11 + 5 x12 + 2 x13 + 15 x21 + 12 x22 + 8 x23 + 3 x31 + 1 x32 + 10 x33;
c1 = x11 + x12 + x13 == 1;
c2 = x21 + x22 + x23 == 1;
c3 = x31 + x32 + x33 == 1;
c4 = -x11 - x21 - x31 == -1;
c5 = -x12 - x22 - x32 == -1;
c6 = -x13 - x23 - x33 == -1;
c6 = x11 >= 0 && x12 >= 0 && x13 >= 0 &&
      x21 >= 0 && x22 >= 0 && x23 >= 0 &&
      x31 >= 0 && x32 >= 0 && x33 >= 0;
var = {x11, x12, x13, x21, x22, x23, x31, x32, x33};
Minimize[{z, c1 && c2 && c3 && c4 && c5 && c6}, var, Reals]

Out[164]= {16, {x11 -> 0, x12 -> 1, x13 -> 0, x21 -> 0, x22 -> 0, x23 -> 1, x31 -> 1, x32 -> 0, x33 -> 0}}
```

# Network problems : Assignment probleem

Assignment problem example HC 3 Matrix

Wolfram Mathematica | FOR STUDENTS Demonstrations | MathWorld

```
In[137]:= ClearAll["Global`*"]

In[138]:= c = {11, 5, 2, 15, 12, 8, 3, 1, 10};
           b = {1, 1, 1, -1, -1, -1};

A = 
$$\begin{pmatrix} 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \\ -1 & 0 & 0 & -1 & 0 & 0 & -1 & 0 & 0 \\ 0 & -1 & 0 & 0 & -1 & 0 & 0 & -1 & 0 \\ 0 & 0 & -1 & 0 & 0 & -1 & 0 & 0 & -1 \end{pmatrix};$$

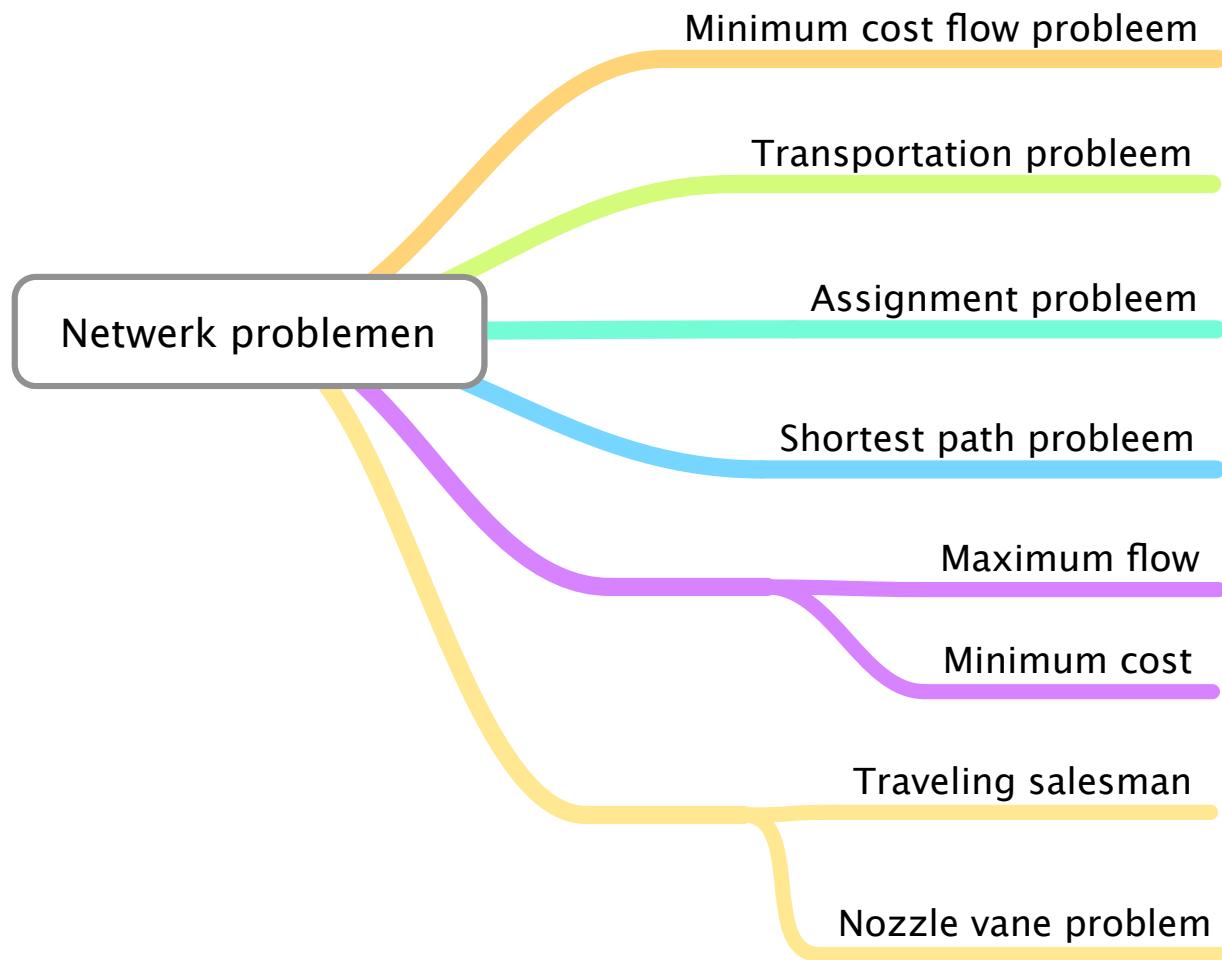

In[141]:= x = LinearProgramming[c, A, b]
           z = {c}.Transpose[{x}];
           Flatten[z, 1][[1]]

Out[141]= {0, 1, 0, 0, 0, 1, 1, 0, 0}

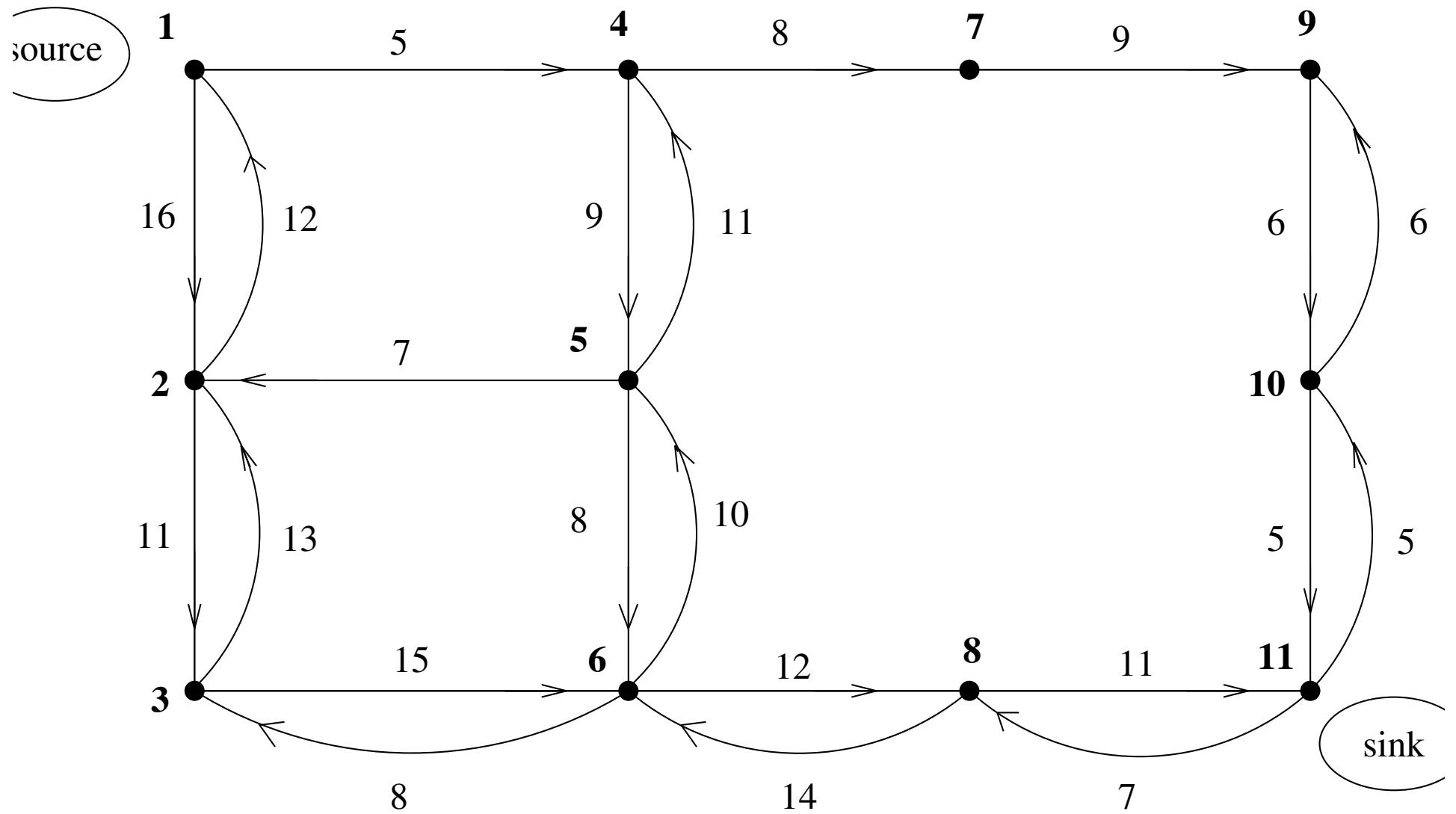
Out[142]= 16
```

# Overzicht : Shortest path probleem

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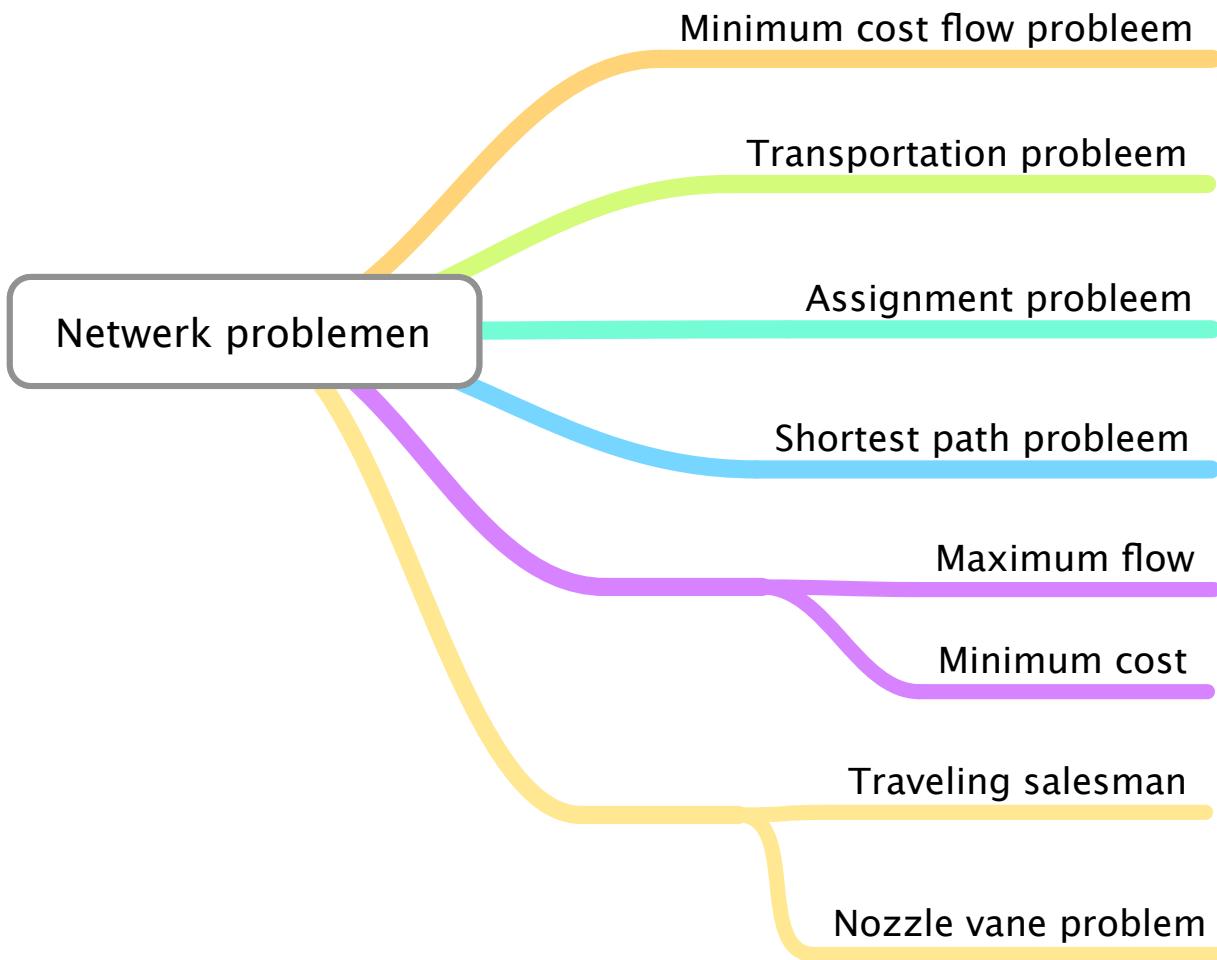


# Network problems : Shortest path probleem

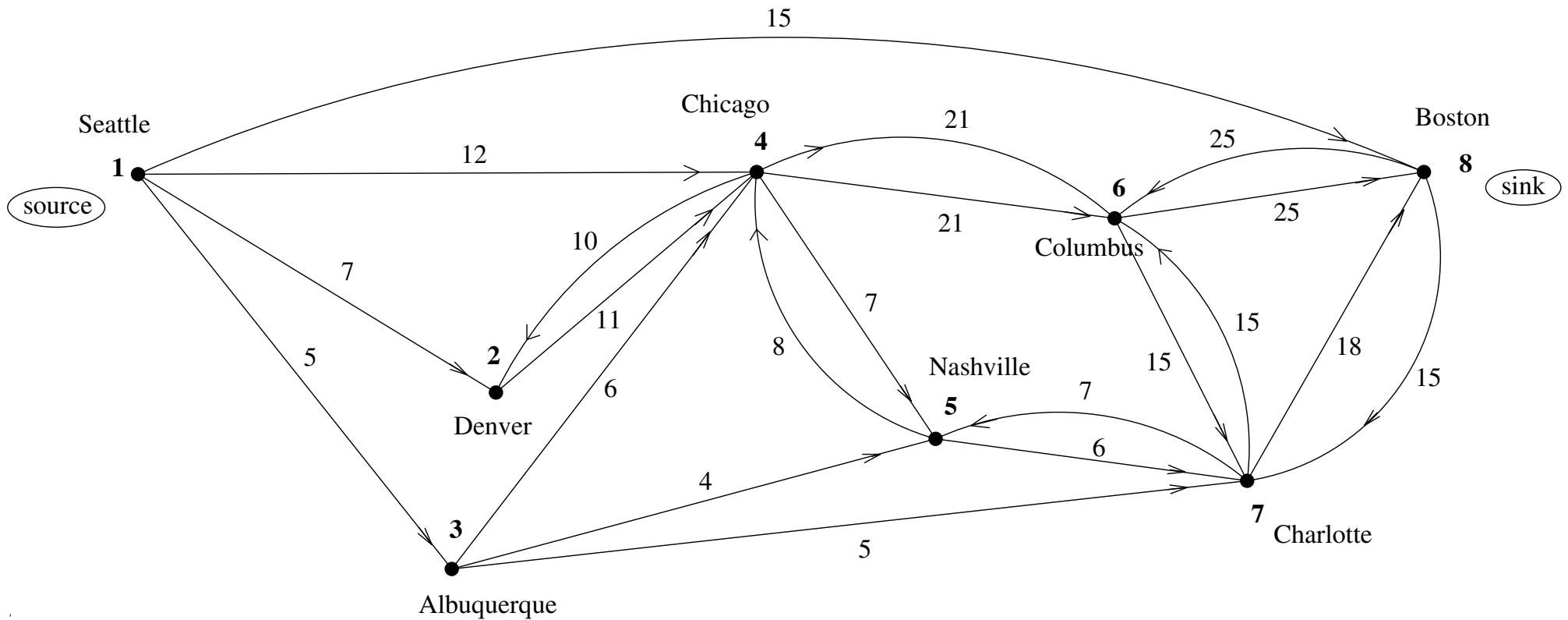


# Overzicht : Maximum flow/Minimum cost probleem

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# Network problems : Maximum flow



# Network problems : Maximum flow problem

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$$\underset{x, f}{\text{maximize}} \quad z = f$$

$$\begin{aligned} & \text{subject to} \quad \sum_j x_{1,j} - \sum_k x_{k,1} = f \\ & \quad \sum_j x_{i,j} - \sum_k x_{k,i} = 0, \quad i = 2, \dots, m-1 \\ & \quad \sum_j x_{m,j} - \sum_k x_{k,m} = -f \\ & \quad 0 \leq x_{i,j} \leq u_{i,j}. \end{aligned}$$

# Network problems : Maximum flow / Minimum cost

---

$$\underset{x, f}{\text{maximize}} \quad z = f$$

$$\begin{aligned} & \text{subject to} \quad \sum_j x_{1,j} - \sum_k x_{k,1} = f \\ & \quad \sum_j x_{i,j} - \sum_k x_{k,i} = 0, \quad i = 2, \dots, m-1 \\ & \quad \sum_j x_{m,j} - \sum_k x_{k,m} = -f \\ & \quad 0 \leq x_{i,j} \leq u_{i,j}. \end{aligned}$$

$$\underset{x}{\text{minimize}} \quad z = -x_{m,1}$$

$$\begin{aligned} & \text{subject to} \quad \sum_j x_{i,j} - \sum_k x_{k,i} = 0, \quad i = 1, \dots, m \\ & \quad 0 \leq x_{i,j} \leq u_{i,j}. \end{aligned}$$

# Network problems : Minimum cost flow probleem

```
WOLFRAM LANGUAGE | FOR STUDENTS
```

```
In[1]:= ClearAll["Global`*"]

z = -x81;

c1 = x12 + x13 + x14 + x18      - x81 == 0;
c2 = x24      - x12 - x42 == 0;
c3 = x34 + x35 + x37      - x13 == 0;
c4 = x45 + x46 + x42      - x14 - x24 - x34 - x54 - x64 == 0;
c5 = x54 + x57      - x45 - x35 - x75 == 0;
c6 = x64 + x67 + x68      - x46 - x76 - x86 == 0;
c7 = x78 + x76 + x75      - x37 - x57 - x67 - x87 == 0;
c8 = x86 + x87 + x81      - x18 - x68 - x78 == 0;

c9 = x12 ≥ 0 && x13 ≥ 0 && x14 ≥ 0 && x18 ≥ 0 && x81 ≥ 0 &&
     x24 ≥ 0 && x12 ≥ 0 && x42 ≥ 0 &&
     x34 ≥ 0 && x35 ≥ 0 && x37 ≥ 0 && x13 ≥ 0 &&
     x45 ≥ 0 && x46 ≥ 0 && x42 ≥ 0 && x14 ≥ 0 && x24 ≥ 0 && x34 ≥ 0 && x54 ≥ 0 && x64 ≥ 0 &&
     x54 ≥ 0 && x57 ≥ 0 && x45 ≥ 0 && x35 ≥ 0 && x75 ≥ 0 &&
     x64 ≥ 0 && x67 ≥ 0 && x68 ≥ 0 && x46 ≥ 0 && x76 ≥ 0 && x86 ≥ 0 &&
     x78 ≥ 0 && x76 ≥ 0 && x75 ≥ 0 && x37 ≥ 0 && x57 ≥ 0 && x67 ≥ 0 && x87 ≥ 0 &&
     x86 ≥ 0 && x87 ≥ 0 && x81 ≥ 0 && x18 ≥ 0 && x68 ≥ 0 && x78 ≥ 0;

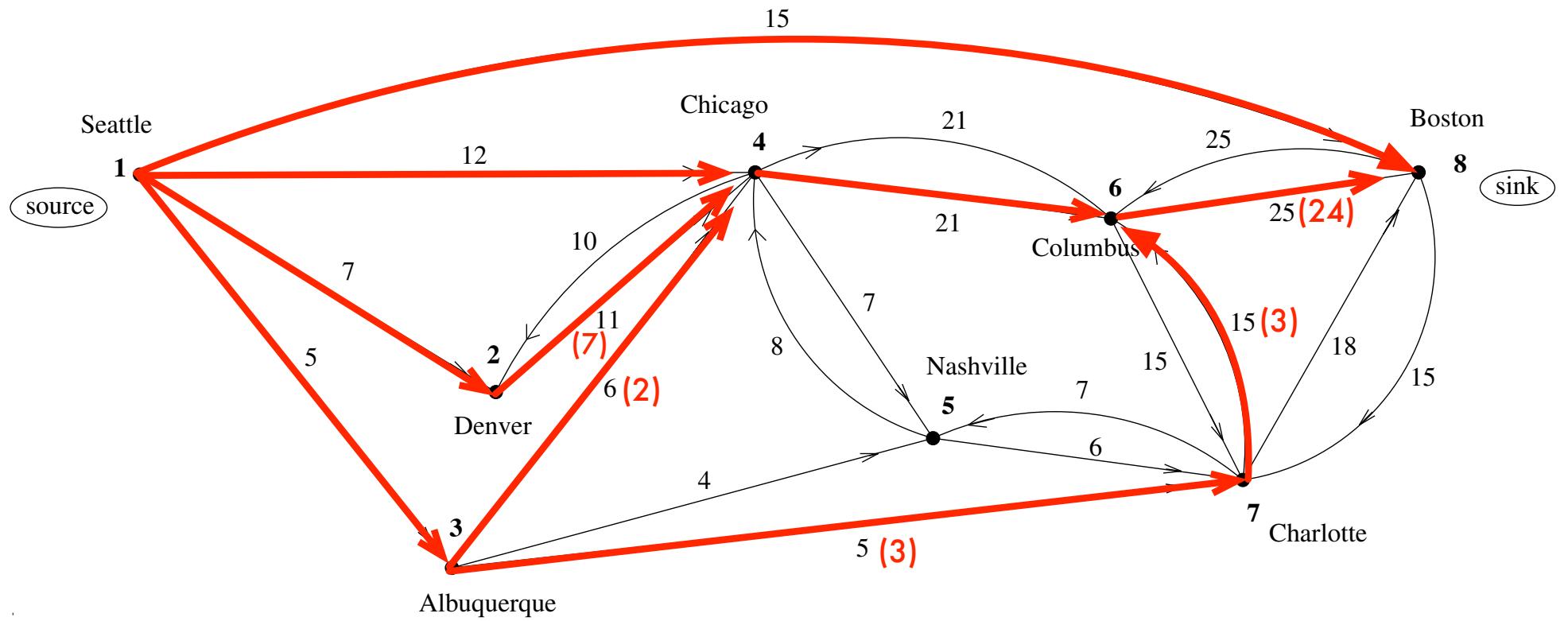
c10 = x12 ≤ 7 && x13 ≤ 5 && x14 ≤ 12 && x18 ≤ 15 &&
      x24 ≤ 11 && x12 ≤ 7 && x42 ≤ 10 &&
      x34 ≤ 6 && x35 ≤ 4 && x37 ≤ 5 && x13 ≤ 5 &&
      x45 ≤ 7 && x46 ≤ 21 && x42 ≤ 10 && x14 ≤ 12 && x24 ≤ 11 && x34 ≤ 6 && x54 ≤ 8 && x64 ≤ 21 &&
      x54 ≤ 8 && x57 ≤ 6 && x45 ≤ 7 && x35 ≤ 4 && x75 ≤ 7 &&
      x64 ≤ 21 && x67 ≤ 15 && x68 ≤ 25 && x46 ≤ 21 && x76 ≤ 15 && x86 ≤ 25 &&
      x78 ≤ 18 && x76 ≤ 15 && x75 ≤ 7 && x37 ≤ 5 && x57 ≤ 6 && x67 ≤ 15 && x87 ≤ 18 &&
      x86 ≤ 25 && x87 ≤ 15 && x18 ≤ 15 && x68 ≤ 25 && x78 ≤ 18;

var = {x12, x13, x14, x18, x42, x24, x34, x35, x37, x54, x64, x45, x46, x75, x57, x76, x86, x67, x68, x78, x87, x81};

Minimize[{z, c1 && c2 && c3 && c4 && c5 && c6 && c7 && c8 && c9 && c10}, var, Reals]

Out[14]= {-39, {x12 → 7, x13 → 5, x14 → 12, x18 → 15, x42 → 0, x24 → 7, x34 → 2, x35 → 0, x37 → 3, x54 → 0,
           x64 → 0, x45 → 0, x46 → 21, x75 → 0, x57 → 0, x76 → 3, x86 → 0, x67 → 0, x68 → 24, x78 → 0, x87 → 0, x81 → 39}}
```

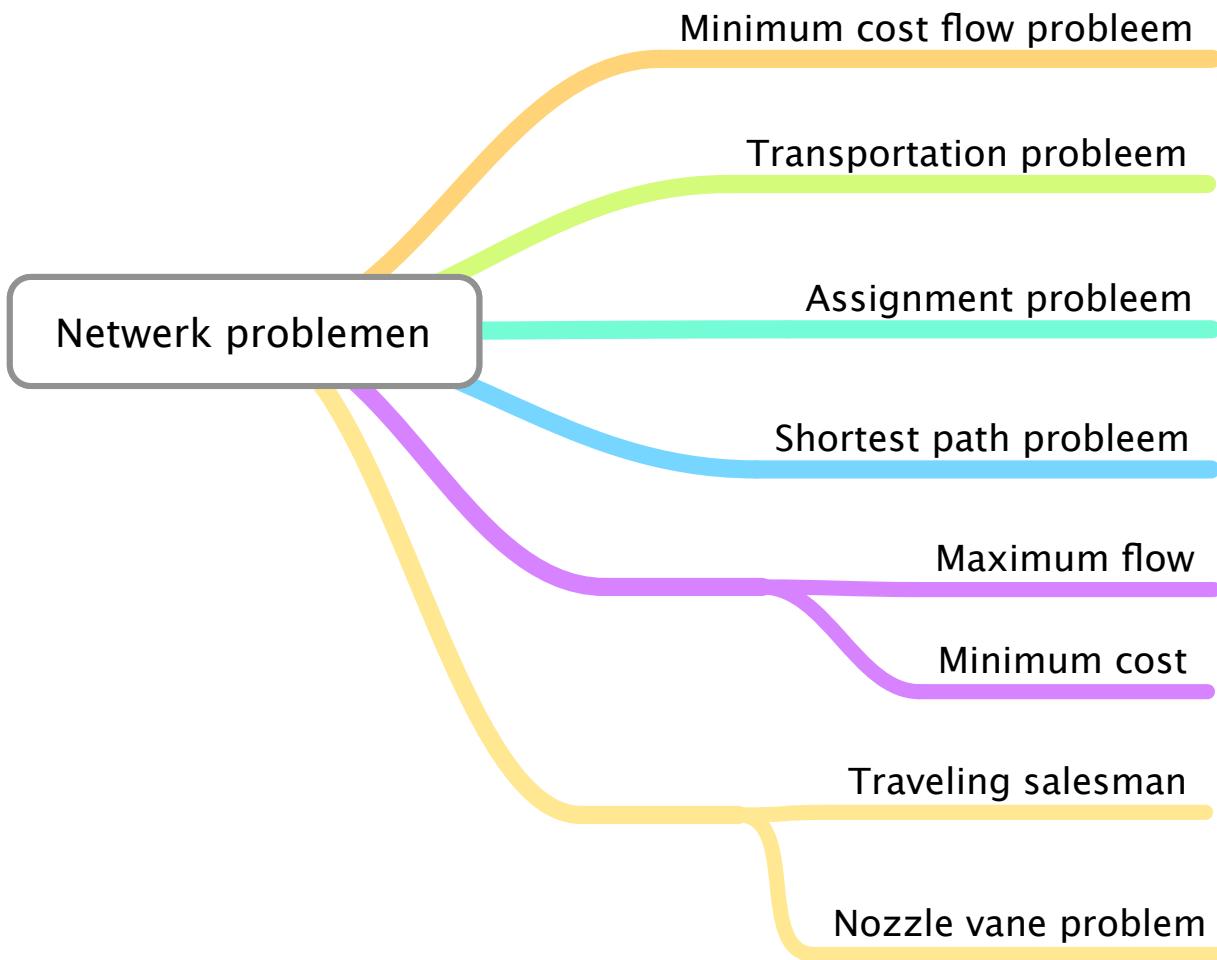
# Network problems : Maximum flow / minimum cost



```
{-39, {x12→7, x13→5, x14→12, x18→15, x42→0, x24→7, x34→2, x35→0, x37→3, x54→0, x64→0, x45→0, x46→21, x75→0, x57→0, x76→3, x86→0, x67→0, x68→24, x78→0, x87→0, x81→39}}
```

# Overzicht : Traveling salesman

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# Network problems : Traveling salesman

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- Een verkoper bezoekt  $n$  steden op een ronde, de steden zijn genummerd als  $0, 1, \dots, n - 1$ .
- Elke stad wordt één keer bezocht en de reis eindigt bij het begin
- De afstand tussen de steden is gegeven als  $c_{ij}$
- De reis wordt beschreven door  $s_0 = 0, s_1, s_2, \dots, s_{n-1}$
- De variabele  $x_{ij}$  wordt 1 als stad  $j$  bezocht wordt na stad  $i$
- Er zijn veel combinatie mogelijk, voor  $n - 1$  steden zijn dit  $(n - 1)!$  routes

$$50! = 3.041 \times 10^{64}$$

# Network problems : Traveling salesman

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- De objective functie is

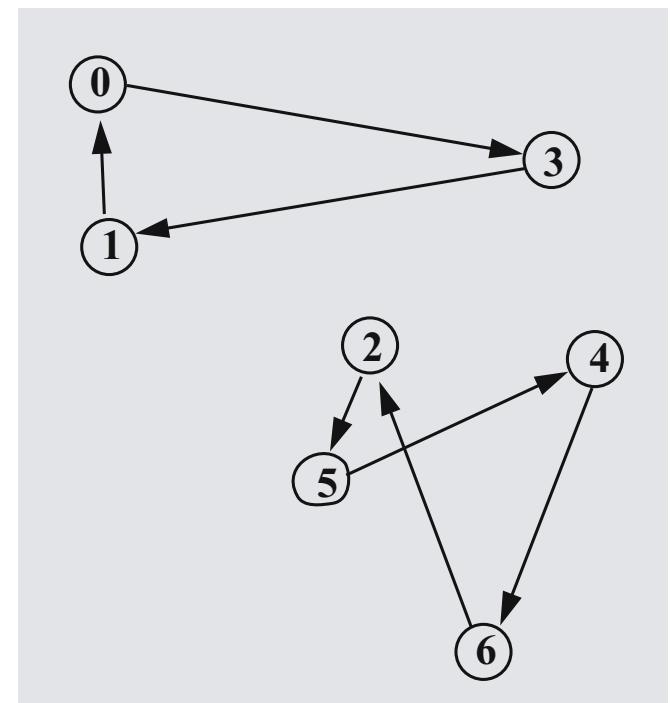
$$\text{minimize} \sum_i \sum_j c_{ij} x_{ij}$$

- met de constraints

$$\sum_j x_{ij} = 1, \quad i = 0, 1, \dots, n - 1$$

$$\sum_i x_{ij} = 1, \quad j = 0, 1, \dots, n - 1$$

- De constraints dwingen af dat de verkoper bij het bezoeken van elke stad afkomstig is uit één voorgaande stad en op reis gaat naar één volgende stad
- Er zijn echter meer constraints nodig...



# Network problems : Traveling salesman

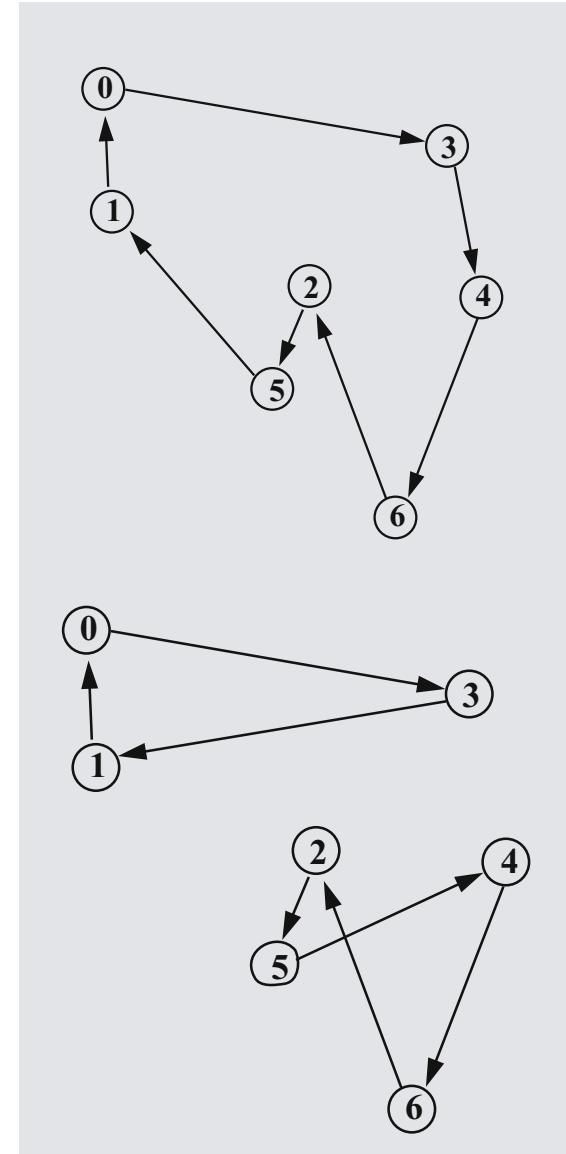
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- Er zijn constraints nodig die afdwingen dat de reis continu verloopt
  - Gegeven een reis  $s_0 = 0, s_1, s_2, \dots, s_{n-1}$
  - dan is  $t_i$  met  $i = 0, 1, \dots, n$  het volgorde nummer van stad  $i$
  - dus  $t_0 = 0, t_{s_1} = 1, t_{s_2} = 2$ , etc dus:  $t_{s_i} = i, i = 0, 1, \dots, n - 1$
  - het volgt dat  $t_j = t_i + 1$  als  $x_{ij} = 1$
  - en dus omdat  $t_i$  een geheel getal is volgt dat  $t_j \geq \begin{cases} t_i + 1 - n & \text{if } x_{ij} = 0, \\ t_i + 1 & \text{if } x_{ij} = 1. \end{cases}$
  - of  $t_j \geq t_i + 1 - n(1 - x_{ij}), \quad i \geq 0, j \geq 1, i \neq j$

# Network problems : Traveling salesman

- Case 1

$S$	$t_S = 0$	$(i, j)$	$x_{i,j} = 1$	$t_j \geq t_i + 1$	$t_j \geq t_i + 1$
$S_0 = 0$	$t_0 = 0$	$(0, 3)$	$x_{0,3} = 1$	$t_3 \geq t_0 + 1$	$1 \geq 0 + 1$
$S_1 = 3$	$t_3 = 1$	$(3, 4)$	$x_{3,4} = 1$	$t_4 \geq t_3 + 1$	$2 \geq 1 + 1$
$S_2 = 4$	$t_4 = 2$	$(4, 6)$	$x_{4,6} = 1$	$t_6 \geq t_4 + 1$	$3 \geq 2 + 1$
$S_3 = 6$	$t_6 = 3$	$(6, 2)$	$x_{6,2} = 1$	$t_2 \geq t_6 + 1$	$4 \geq 3 + 1$
$S_4 = 2$	$t_2 = 4$	$(2, 5)$	$x_{2,5} = 1$	$t_5 \geq t_2 + 1$	$5 \geq 4 + 1$
$S_5 = 5$	$t_5 = 5$	$(5, 1)$	$x_{5,1} = 1$	$t_1 \geq t_5 + 1$	$6 \geq 5 + 1$

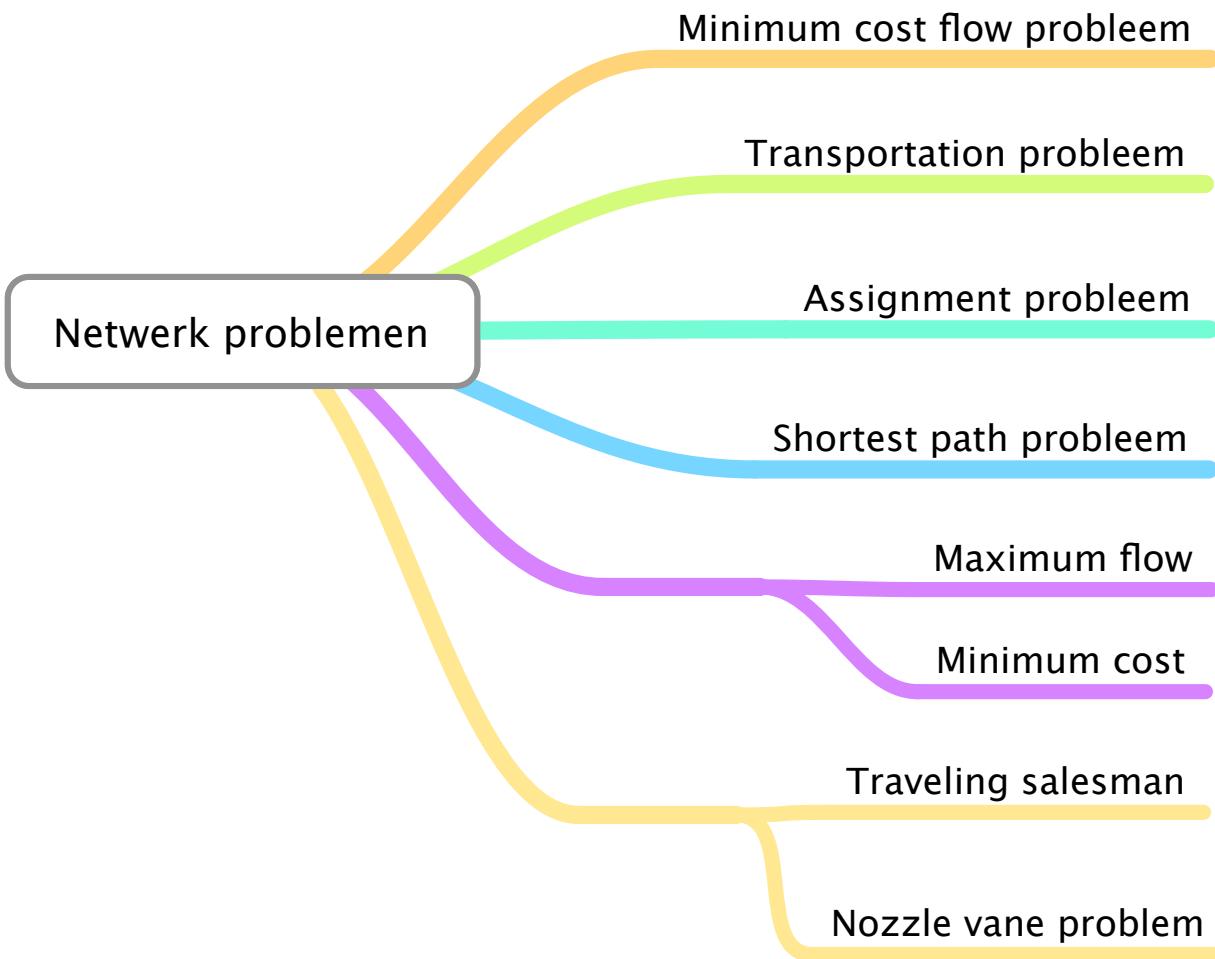


- Case 2

$S$	$t_S = 0$	$(i, j)$	$x_{i,j} = 1$	$t_j \geq t_i + 1$	$t_j \geq t_i + 1$
$S_0 = 2$	$t_2 = 0$	$(2, 5)$	$x_{2,5} = 1$	$t_5 \geq t_2 + 1$	$1 \geq 0 + 1$
$S_1 = 5$	$t_5 = 1$	$(5, 4)$	$x_{5,4} = 1$	$t_4 \geq t_5 + 1$	$2 \geq 1 + 1$
$S_2 = 4$	$t_4 = 2$	$(4, 6)$	$x_{4,6} = 1$	$t_6 \geq t_4 + 1$	$3 \geq 2 + 1$
$S_3 = 6$	$t_6 = 3$	$(6, 2)$	$x_{6,2} = 1$	$t_2 \geq t_6 + 1$	$0 \leq 3 + 1$

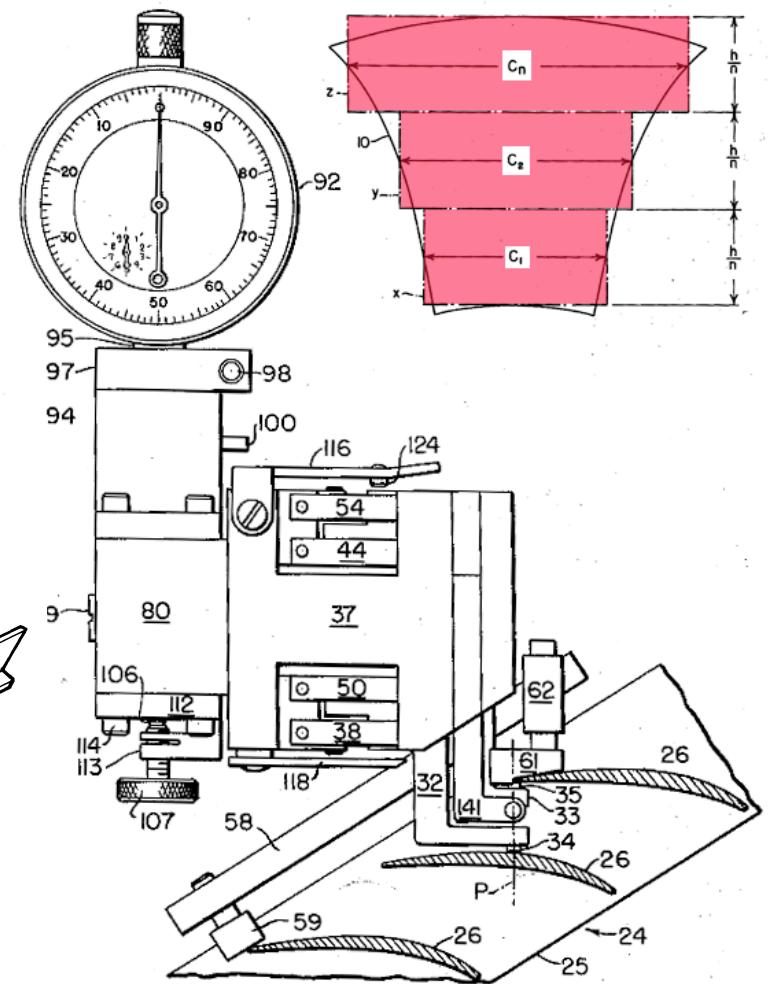
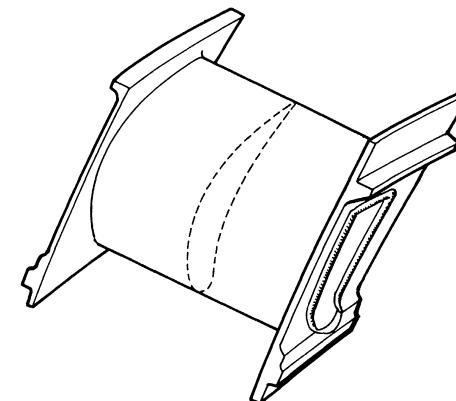
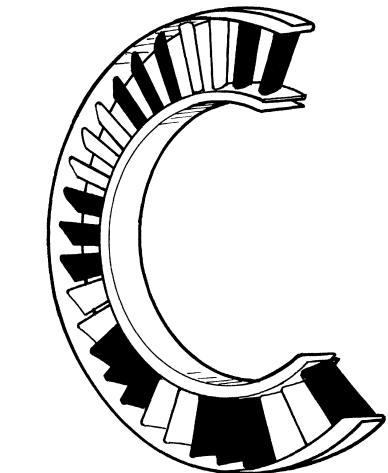
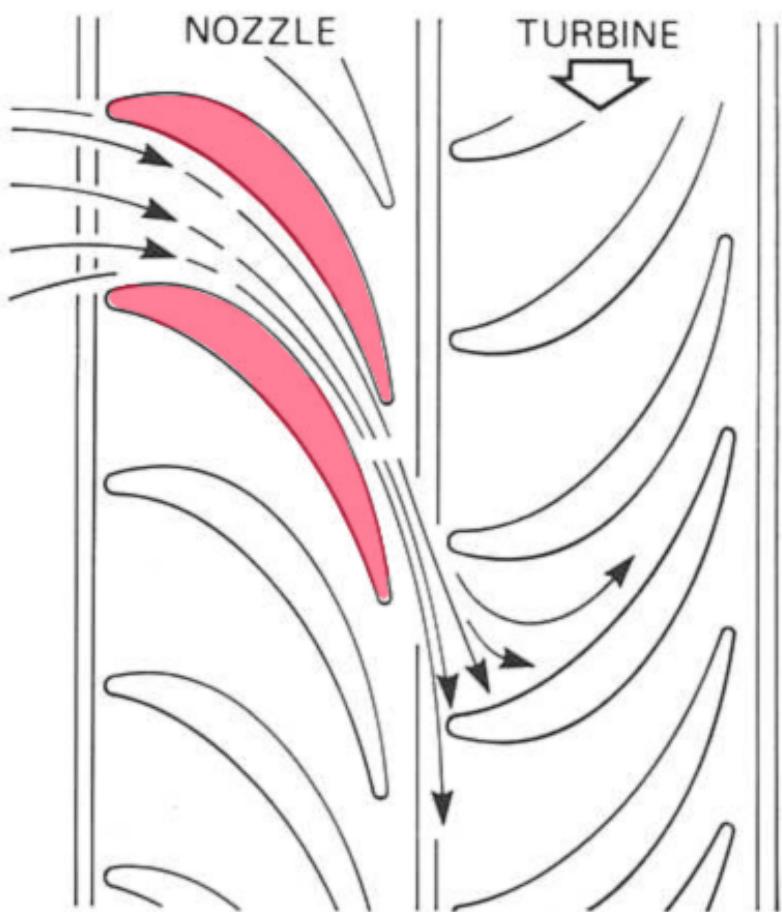
# Overzicht : Nozzle vane probleem

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# Nozzle vane problem

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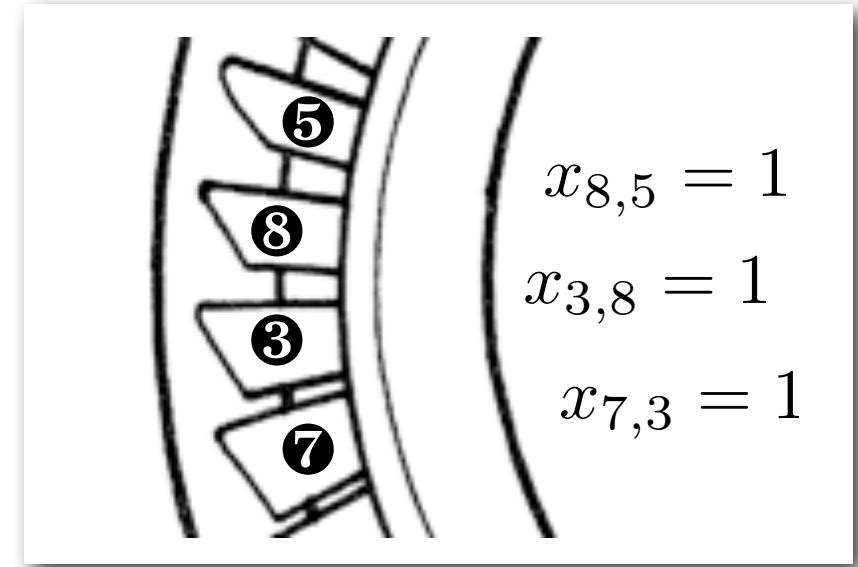
# Nozzle vane problem : Oppervlakte

- Een nozzle vane  $j$  die direct rechts (met de klok mee) van nozzle vane  $i$  geplaatst wordt vormt een doorgang met een oppervlakte van  $c_{ij}$
- Het totale oppervlak van het doorstroom opening is :

$$T = \sum_{i} \sum_{j} c_{ij} x_{ij},$$

where

$$x_{ij} = \begin{cases} 1 & \text{if vane } j \text{ is placed} \\ & \text{clockwise-adjacent to vane } i, \text{ and} \\ 0 & \text{otherwise.} \end{cases}$$



# Nozzle vane problem : Oppervlakte

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- Met een meetinstrument worden twee waarden gemeten

$$c_{ij} = A_i + B_j \quad \text{for all } i \text{ and } j.$$

- De ideale nozzle vane guide plaatsing wordt bereikt als alle  $c_{ij}$  gelijk zijn
  - Er geldt dan dat  $c_{ij} = c_{kl}$  voor alle  $i, j$  en  $k, l$  als  $x_{ij} = 1$  en  $x_{kl} = 1$

- De het totale doorstroom oppervlak wordt gegeven door

$$T = \sum_{i=1}^n (A_i + B_i)$$

- Het gemiddelde oppervlak per nozzle kanaal is

$$\bar{d} = \frac{\sum_i (A_i + B_i)}{n}$$

- Waarmee het ideale kanaal gegeven wordt door

$$c_{ij} - \bar{d} = 0 \quad \text{for all } i \text{ and } j \text{ with } x_{ij} = 1.$$

# Nozzle vane problem : Probleemstelling

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## Problem P1

Minimize

$$U(X) = \sum_{i} \sum_{j} (\bar{d} - (A_i + B_j))^2 x_{ij}$$

subject to

$$\sum_j x_{ij} = 1 \quad \text{for all } i,$$

$$\sum_i x_{ij} = 1 \quad \text{for all } j,$$

$$x_{ij} = \{0, 1\} \quad \text{for all } i \text{ and } j,$$

no subtours.

## Problem P2

Minimize

$$h(X) = \sum_{i} \sum_{j} (A_i B_j) x_{ij}$$

subject to

$$\sum_j x_{ij} = 1 \quad \text{for all } i,$$

$$\sum_i x_{ij} = 1 \quad \text{for all } j,$$

$$x_{ij} = \{0, 1\} \quad \text{for all } i \text{ and } j,$$

no subtours.

# Vragen?

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