| Ans 1  | :- Normalization is a process in DBMs to reduce reedundancy and dependency by organizing duta in        |
|--|---|
|  | tables. It divides large tables into smaller ones and links them using relationships.                   |
|  | Here are some normal forms:   |
|  | · 1st normal form (INF): Ensure that each columns   |
|  | contains atomic values and each has a unique value.  2nd normal form (2NF): A table is in 2NF, if it is |
| de entre de la companya del la companya de la companya del la companya de la comp | in INF and all non-key attributes are fully dependent   |
| National Conference on the Con | on the preimary key.  • 3rd Normal form (3NF): A table is in 3NF if it is in                            |
| and the second s | · 3 <sup>eq</sup> Noemal form (3NF): A table is in 3NF if it is in                                      |
|  | 2NF & there is no teansitive dependency i-e non-<br>ky atteibute are not dependent on other non-key     |
|  | afficientes.  |

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| Ans 2. | A primary key is a unique identifier for secords in a  |
|        | A primary key is a unique identifier for secords in a table. It must contain unique valves and cannot  |
|        | table. It must contain wayer races serves  |
|        | contain will value A table sun have song   |
|        | primary key, which can be a sirgle column ot a   |
|        | and in afting of columns   |
|        | combination of columns.  |
|        |  |
|        | It is important because of the following points.   |
|        | It is important because of the following points.  1) It ensures that each record in a table is unique. |
|        | 2) It have in establishing relationship between tables   |
|        | reten vyed with foreign keys.  3) It is essential for indexing, making searches and                    |
|        | 3) It is essential for indexing, making searches and   |
|        | rectrievals fastire.   |
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| -      |  |

| _ Anu 3 | 1) Atomicity: Ensures that a transaction is either fully                        |
|---------|---|
| -       | completed or fully rolled back. It prevents partial                             |
|         | apdates.  |
| -       | 2) Consistency: Guarantees that a transaction is either                         |
|         | fully will take the database from one valid                                     |
| _       | State to another ralid state, maintaining dela                                  |
|         | integrity.  |
|         | 3) Isolation: Ensures that transactions are executed                            |
|         | independently of each other. One transaction's                                  |
| -       | changes will not be visible to others until it                                  |
|         | is completed.   |
|         | 4) Durability: Ensures that once a transaction is                               |
|         | committed, its changes are permanent even in                                    |
|         | committed, its changes are permanent, even in<br>the event of a system failure. |
| (       |   |
|         | these proporties are critical to ensuring the correctness                       |
|         | and reliability of transactions in a database.                                  |
|         | J ' S WWalper Coll  |
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| Ans y. | A deadlock in DBMs occurs when the or more  |
|--------|---|
|        | transactions are weating for each other to release  |
|        | locks, and none of them can proceed for example,  |
|        |   |
|        |   |
|        | mails for resources 2, while Transaction B holds a  |
|        | Lock on Resources 2 and waite for Resources 1, both   |
|        | will be stuck.  |
|        |   |
|        | Deadlack can be prevented by:   |
|        | Deadlock Can be prevented by:  Deadlock Prevention: By using protocols that  avoid the conditions leading to deadlock, like |
|        | and the specific leading to dead out tite   |
|        | word the conditions reading to deact bear the   |
|        | resource ordering and timeouts.   |
|        | 2) Deadlock Detection: Regularly checking for deadlock  |
|        | eycles in the system and aborting one of the  |
|        | transaction to break the cycle.   |
|        | 3) Deadlock Avoidance: Using algorithms like the  |
|        | Banker's algorithm to theek the system State  |
|        | h O a resulting lasks arranged it designed  |
|        | before granting locks, envening it doesn't lead   |
|        | to a deadlock.  |
|        |   |
|        |   |
|        |   |

| - Ans 5: | A distributed database is one in which data          |
|----------|--|
|          | is stored across multiple Jocautions, but it appears |
|          | to user also single database.                        |
|          | ·  |
|          | Advantages:-   |
|          | 1) Reliability and Availabity: - If one site fails,  |
| _        | data base can continue functioning using other       |
| _        | sites.   |
|          | 2) Scalability: - It's easier to add new sites or    |
| _        | nodes without affecting the whole system.            |
|          | 3) Local control: - Each site can have control over  |
| _        | its data, providing flexibility and faster           |
| _        | Jocal queries.                                       |
|          |  |
| -        | Disadvantages:-                                      |
|          | 1) Complexity: - Managing a distributed system       |
|          | is more complex due to mored for coordination        |
|          | across multiple sites.                               |
|          | 2) Security Risks: - Since data is spread across     |
|          | locations, it increases the syrface area for         |
|          | potential security threats.                          |
|          | 3) Data Integrity: - keeping the data consistent     |
|          | across all sites can be chaddenging especially       |
|          |  |
|          | with concurrent transactions.                        |
|          |  |
|          |  |