

Program Correctness

Block 4

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Outline



Exercise 7.1: Powers

Exercise 7.2: Factorial

Exercise 7.8: Dijkstra's FUSC

Summing an array

Square Roo

Square Root (Similar to Exercise 7.3) Exercise 7.4: Square Root, Revisited

Integral Division

Exercise 7.5: Integral Division

Exercise 7.6: Variation on Integral Division Efficiency Comparison Exercises 7.5 and 7.

Exercise 7.1



```
\begin{array}{l} \textbf{const} \ n: \ \mathbb{N}; \\ \textbf{var} \ x, \ y: \ \mathbb{Z}; \\ \{P: \ \textbf{true}\} \\ T \\ \{Q: \ x=n^2 \land y=n^3\} \end{array}
```

- ▶ We are only allowed to multiply by 2 and 3, and use addition.
- ightharpoonup Use "replace a constant by a variable" to find J and B.

Exercise 7.1: Invariant and Guard



P: true

$$Q: x = n^2 \wedge y = n^3$$

- 0 We decide that we need a **while**-program: we are not allowed to use assignments x := n * n; y := n * x;
- 1 Choose an invariant J, and guard B such that $J \wedge \neg B \Rightarrow Q$.

Exercise 7.1: Invariant and Guard



P: true

$$Q: x = n^2 \wedge y = n^3$$

- 0 We decide that we need a **while**-program: we are not allowed to use assignments x := n * n; y := n * x;
- 1 Choose an invariant J, and guard B such that $J \land \neg B \Rightarrow Q$. We replace the constant n by the variable k:

$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$

 $B: k \neq n$

Clearly, J and $\neg B$ imply Q.

Exercise 7.1: Initialization and Variant



```
P: trueQ: x=n^2 \wedge y=n^3 \ J: x=k^2 \wedge y=k^3 \wedge 0 \leq k \leq n \ B: k 
eq n
```

2 **Initialization**: Find T_0 such that $\{P\}$ T_0 $\{J\}$.

```
\{P: \mathbf{true}\}\ (* calculus; n \in \mathbb{N} *)

\{0 = 0^2 \land 0 = 0^3 \land 0 \le 0 \le n\}

k := 0; x := 0; y := 0;

\{J: x = k^2 \land y = k^3 \land 0 \le k \le n\}
```

Exercise 7.1: Initialization and Variant



```
P: trueQ: x=n^2 \wedge y=n^3 \ J: x=k^2 \wedge y=k^3 \wedge 0 \leq k \leq n \ B: k 
eq n
```

2 **Initialization**: Find T_0 such that $\{P\}$ T_0 $\{J\}$.

```
\{P: 	extbf{true}\}\ (* calculus; n \in \mathbb{N}^*)\ \{0 = 0^2 \land 0 = 0^3 \land 0 \le 0 \le n\}\ k := 0; x := 0; y := 0;\ \{J: x = k^2 \land y = k^3 \land 0 \le k \le n\}
```

3 **Variant**: We take $vf = n - k \in \mathbb{Z}$. We must show $vf \ge 0$. Clearly, $J \wedge B \Rightarrow n - k \ge 0$ as J contains the conjunct $k \le n$.



$$P$$
: true

$$Q: x = n^2 \wedge y = n^3$$

$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$

 $B: k \neq n$

We can relate x, y, and k. Actually, we look into k+1 (why?):

$$(k+1)^3 = k^3 + 3k^2 + 3k + 1$$
$$= k^3 + 3(k^2 + k) + 1$$
$$\{y = k^3, x = k^2\}$$
$$= y + 3(x + k) + 1$$

Similarly:

$$(k+1)^2 = k^2 + 2k + 1$$

= $x + 2k + 1$

We shall use these equalities in the body of the loop.



$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$

 $B: k \ne n$

4 Body of the loop:
$$\{J \land B \land vf = V\}$$
 S $\{J \land vf < V\}$ $\{J \land B \land vf = V\}$

$$y := y + 3 * (x + k) + 1;$$

$$x := x + 2 * k + 1;$$

$$k := k + 1;$$

$$\{J \wedge vf < V\}$$



$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$
 $B: k \ne n$

4 Body of the loop:
$$\{J \land B \land vf = V\} S \{J \land vf < V\}$$

$$\begin{cases} J \wedge B \wedge vf = V \\ \{x = k^2 \wedge y = k^3 \wedge 0 \le k \le n \wedge k \ne n \wedge n - k = V \} \end{cases}$$

$$y := y + 3 * (x + k) + 1;$$

$$x := x + 2 * k + 1;$$

$$k := k + 1;$$

$$\{J \wedge vf < V\}$$



$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$
 $B: k \ne n$

4 Body of the loop:
$$\{J \land B \land vf = V\}$$
 S $\{J \land vf < V\}$ $\{J \land B \land vf = V\}$ $\{x = k^2 \land y = k^3 \land 0 \le k \le n \land k \ne n \land n - k = V\}$ (* prepare assignment to y : use $(k+1)^3 = y + 3(x+k) + 1$ *)

$$y := y + 3 * (x + k) + 1;$$

$$x := x + 2 * k + 1;$$

$$k := k + 1$$
;

$$\{J \wedge vf < V\}$$



$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$
 $B: k \ne n$

4 Body of the loop: $\{J \land B \land vf = V\}$ S $\{J \land vf < V\}$ $\{J \land B \land vf = V\}$ $\{x = k^2 \land y = k^3 \land 0 \le k \le n \land k \ne n \land n - k = V\}$ (* prepare assignment to y: use $(k+1)^3 = y + 3(x+k) + 1$ *) $\{x = k^2 \land y + 3(x+k) + 1 = (k+1)^3 \land 0 \le k < n \land n - k = V\}$ y := y + 3 * (x + k) + 1:

$$x := x + 2 * k + 1;$$

$$k:=k+1;$$

$$\{J \land vf < V\}$$



$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$
 $B: k \ne n$

4 Body of the loop: $\{J \land B \land vf = V\}$ S $\{J \land vf < V\}$ $\{J \land B \land vf = V\}$ $\{x = k^2 \land y = k^3 \land 0 \le k \le n \land k \ne n \land n - k = V\}$ (* prepare assignment to y: use $(k+1)^3 = y + 3(x+k) + 1$ *) $\{x = k^2 \land y + 3(x+k) + 1 = (k+1)^3 \land 0 \le k < n \land n - k = V\}$ y := y + 3 * (x + k) + 1; $\{x = k^2 \land y = (k+1)^3 \land 0 \le k < n \land n - k = V\}$

$$x := x + 2 * k + 1;$$

 $\{J \land vf < V\}$

$$k:=k+1;$$



$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$

 $B: k \ne n$

4 Body of the loop: $\{J \land B \land vf = V\}$ S $\{J \land vf < V\}$ $\{J \land B \land vf = V\}$ $\{x = k^2 \land y = k^3 \land 0 \le k \le n \land k \ne n \land n - k = V\}$ (* prepare assignment to y: use $(k+1)^3 = y + 3(x+k) + 1 *$) $\{x = k^2 \land y + 3(x+k) + 1 = (k+1)^3 \land 0 \le k < n \land n - k = V\}$ y := y + 3 * (x + k) + 1; $\{x = k^2 \land y = (k+1)^3 \land 0 \le k < n \land n - k = V\}$ (* prepare assignment to x: use $(k+1)^2 = x + 2k + 1 *$) x := x + 2 * k + 1:

$$k:=k+1;$$

 $\{J \land vf < V\}$



$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$

 $B: k \ne n$

4 Body of the loop: $\{J \land B \land vf = V\}$ S $\{J \land vf < V\}$ $\{J \land B \land vf = V\}$ $\{x = k^2 \land y = k^3 \land 0 \le k \le n \land k \ne n \land n - k = V\}$ (* prepare assignment to y: use $(k+1)^3 = y + 3(x+k) + 1$ *) $\{x = k^2 \land y + 3(x+k) + 1 = (k+1)^3 \land 0 \le k < n \land n - k = V\}$ y := y + 3 * (x + k) + 1; $\{x = k^2 \land y = (k+1)^3 \land 0 \le k < n \land n - k = V\}$ (* prepare assignment to x: use $(k+1)^2 = x + 2k + 1$ *) $\{x + 2k + 1 = (k+1)^2 \land y = (k+1)^3 \land 0 \le k < n \land n - k = V\}$ x := x + 2 * k + 1:

$$k := k + 1;$$

$$\{J \wedge vf < V\}$$



$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$

 $B: k \ne n$

4 Body of the loop: $\{J \land B \land vf = V\} S \{J \land vf < V\}$ $\{J \wedge B \wedge vf = V\}$ $\{x = k^2 \land y = k^3 \land 0 \le k \le n \land k \ne n \land n - k = V\}$ (* prepare assignment to y: use $(k+1)^3 = y + 3(x+k) + 1$ *) $\{x = k^2 \land y + 3(x+k) + 1 = (k+1)^3 \land 0 \le k \le n \land n-k = V\}$ u := u + 3 * (x + k) + 1: $\{x = k^2 \land y = (k+1)^3 \land 0 \le k \le n \land n-k = V\}$ (* prepare assignment to x: use $(k+1)^2 = x + 2k + 1$ *) $\{x+2k+1=(k+1)^2 \land y=(k+1)^3 \land 0 \le k < n \land n-k=V\}$ x := x + 2 * k + 1; $\{x = (k+1)^2 \land y = (k+1)^3 \land 0 \le k < n \land n-k = V\}$

$$\{J \wedge \mathit{vf} < V\}$$

k := k + 1:

 $\{J \land vf < V\}$



$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$

 $B: k \ne n$

```
4 Body of the loop: \{J \land B \land vf = V\} S \{J \land vf < V\}
     \{J \wedge B \wedge vf = V\}
    \{x = k^2 \land y = k^3 \land 0 \le k \le n \land k \ne n \land n - k = V\}
       (* prepare assignment to y: use (k+1)^3 = y + 3(x+k) + 1*)
    \{x = k^2 \land y + 3(x+k) + 1 = (k+1)^3 \land 0 \le k \le n \land n-k = V\}
  u := u + 3 * (x + k) + 1:
     \{x = k^2 \land y = (k+1)^3 \land 0 \le k \le n \land n-k = V\}
       (* prepare assignment to x: use (k+1)^2 = x + 2k + 1*)
     \{x+2k+1=(k+1)^2 \land y=(k+1)^3 \land 0 \le k < n \land n-k=V\}
  x := x + 2 * k + 1;
    \{x = (k+1)^2 \land y = (k+1)^3 \land 0 \le k < n \land n-k = V\}
       (* complete preparation for k := k + 1 *)
  k := k + 1;
```



$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$

 $B: k \ne n$

4 Body of the loop: $\{J \land B \land vf = V\} S \{J \land vf < V\}$ $\{J \wedge B \wedge vf = V\}$ $\{x=k^2 \land y=k^3 \land 0 \le k \le n \land k \ne n \land n-k=V\}$ (* prepare assignment to y: use $(k+1)^3 = y + 3(x+k) + 1$ *) $\{x = k^2 \land y + 3(x+k) + 1 = (k+1)^3 \land 0 \le k < n \land n-k = V\}$ u := u + 3 * (x + k) + 1: $\{x = k^2 \land y = (k+1)^3 \land 0 \le k \le n \land n-k = V\}$ (* prepare assignment to x: use $(k+1)^2 = x + 2k + 1$ *) $\{x+2k+1=(k+1)^2 \land y=(k+1)^3 \land 0 \le k < n \land n-k=V\}$ x := x + 2 * k + 1; $\{x = (k+1)^2 \land y = (k+1)^3 \land 0 \le k < n \land n-k = V\}$ (* complete preparation for k := k + 1 *) $\{x = (k+1)^2 \land y = (k+1)^3 \land 0 \le k+1 \le n \land n-(k+1) \le V\}$ k := k + 1;

$$\{J \land vf < V\}$$

 $\{J \land vf < V\}$



$$J: x = k^2 \wedge y = k^3 \wedge 0 \le k \le n$$
 $B: k \ne n$

4 Body of the loop: $\{J \land B \land vf = V\} S \{J \land vf < V\}$ $\{J \wedge B \wedge vf = V\}$ $\{x=k^2 \land y=k^3 \land 0 \le k \le n \land k \ne n \land n-k=V\}$ (* prepare assignment to y: use $(k+1)^3 = y + 3(x+k) + 1$ *) $\{x = k^2 \land y + 3(x+k) + 1 = (k+1)^3 \land 0 \le k \le n \land n-k = V\}$ u := u + 3 * (x + k) + 1: $\{x = k^2 \land y = (k+1)^3 \land 0 \le k < n \land n-k = V\}$ (* prepare assignment to x: use $(k+1)^2 = x + 2k + 1$ *) $\{x+2k+1=(k+1)^2 \land y=(k+1)^3 \land 0 \le k < n \land n-k=V\}$ x := x + 2 * k + 1; $\{x = (k+1)^2 \land y = (k+1)^3 \land 0 \le k < n \land n-k = V\}$ (* complete preparation for k := k + 1 *) $\{x = (k+1)^2 \land y = (k+1)^3 \land 0 \le k+1 \le n \land n-(k+1) \le V\}$ k := k + 1: $\{x = k^2 \land y = k^3 \land 0 \le k \le n \land n - k \le V\}$

Exercise 7.1: Conclusion



5 The command $\{P\}$ T_0 ; while B do S end $\{Q\}$ solves the problem:

```
const n : \mathbb{N}:
var x, y, k : \mathbb{Z};
  \{P: \mathsf{true}\}
k := 0:
x := 0:
y := 0;
  \{J: x = k^2 \land y = k^3 \land 0 < k < n\}
     (* vf = n - k *)
while k \neq n do
  y := y + 3 * x + 3 * k + 1;
  x := x + 2 * k + 1:
  k := k + 1;
end;
  \{Q: x = n^2 \wedge y = n^3\}
```

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Exercise 7.1: Powers

Exercise 7.2: Factorial

Exercise 7.8: Dijkstra's FUSC

Summing an array

Square Root

Square Root (Similar to Exercise 7.3) Exercise 7.4: Square Root, Revisited

Integral Division

Exercise 7.5: Integral Division

Exercise 7.6: Variation on Integral Division Efficiency Comparison Exercises 7.5 and 7.

Exercise 7.2: Factorial



```
egin{aligned} 	extsf{Var} & x, & n: & \mathbb{Z}; \ & \{P: & n \geq 0 \land X = n!\} \ & T \ & \{Q: & x = X\} \end{aligned}
```

Recall the heuristic generalization.

Exercise 7.2: Invariant and Guard



$$P: n \geq 0 \wedge X = n!$$
 $Q: x = X$

- 0 We assume that there is no function 'fact' available. We decide that we need a while-program.
- 1 Choose an invariant J, and guard B such that $J \land \neg B \Rightarrow Q$. We use the heuristic generalization.

$$egin{aligned} J: (x \cdot n! = X) \wedge n \geq 0 \ B: n
eq 0 \end{aligned}$$

By definition 0! = 1. Therefore, J and $\neg B$ imply Q.

Exercise 7.2: Initialization and Variant



$$egin{aligned} P: & n \geq 0 \wedge X = n! \ Q: x = X \ & J: x \cdot n! = X \wedge n \geq 0 \ & B: n
eq 0 \end{aligned}$$

2 Initialization: Find a command T_0 such that $\{P\}$ T_0 $\{J\}$

```
\{P: X = n! \land n \geq 0\}
(* \ calculus \ *)
\{X = 1 \cdot n! \land n \geq 0\}
x := 1;
\{J: \ x \cdot n! = X \land n \geq 0\}
```

3 Variant function: $vf \in \mathbb{Z}$ and $J \wedge B \Rightarrow vf \geq 0$ Clearly, n must decrease until n=0. We choose $vf=n \in \mathbb{N}$. Because J contains the conjunct $n \geq 0$, we have that $J \wedge B \Rightarrow vf > 0$ holds trivially.



$$J:x\cdot n!=X\wedge n\geq 0$$
 $B:n
eq 0$ $vf=n$

4 Body of the loop:
$$\{J \wedge B \wedge vf = V\}$$
 S $\{J \wedge vf < V\}$ $\{J \wedge B \wedge vf = V\}$

$$\{J \wedge vf < V\}$$



$$egin{aligned} J:x\cdot n! &= X\wedge n \geq 0\ B:n
eq 0 \end{aligned}$$

$$egin{array}{ll} \{J \wedge B \wedge v f = V\} \ \{x \cdot n! = X \ \wedge \ n \geq 0 \ \wedge \ n
eq 0 \wedge n = V\} \end{array}$$

$$\{J \wedge vf < V\}$$



$$egin{aligned} J:x\cdot n! &= X\wedge n \geq 0\ B:n
eq 0 \end{aligned}$$

$$egin{aligned} \{J \wedge B \wedge v & f = V \} \ \{x \cdot n! = X \ \wedge \ n \geq 0 \ \wedge \ n
eq 0 \wedge n = V \} \ & (* \ n = V > 0 \Rightarrow n! = n \cdot (n-1)! \wedge n - 1 \geq 0 \wedge n - 1 < V \ ^*) \end{aligned}$$

$$\{J \wedge vf < V\}$$



$$egin{aligned} J:x\cdot n! &= X\wedge n \geq 0\ B:n
eq 0\ vf &= n \end{aligned}$$

$$egin{aligned} \{J \wedge B \wedge vf &= V\} \ \{x \cdot n! &= X \ \wedge \ n \geq 0 \ \wedge \ n \neq 0 \wedge n = V\} \ &\text{(* } n = V > 0 \Rightarrow n! = n \cdot (n-1)! \wedge n - 1 \geq 0 \wedge n - 1 < V \ *) \ \{x \cdot n \cdot (n-1)! &= X \ \wedge \ n - 1 \geq 0 \ \wedge \ n - 1 < V\} \end{aligned}$$

$$\{J \wedge vf < V\}$$



$$egin{aligned} J:x\cdot n! &= X\wedge n \geq 0\ B:n
eq 0\ vf &= n \end{aligned}$$

4 Body of the loop: $\{J \land B \land vf = V\} \ S \ \{J \land vf < V\}$

```
egin{aligned} \{J \wedge B \wedge vf &= V\} \ \{x \cdot n! &= X \wedge n \geq 0 \wedge n \neq 0 \wedge n = V\} \ (* n &= V > 0 \Rightarrow n! &= n \cdot (n-1)! \wedge n - 1 \geq 0 \wedge n - 1 < V \ \{x \cdot n \cdot (n-1)! &= X \wedge n - 1 \geq 0 \wedge n - 1 < V\} \ x &:= x * n; \end{aligned}
```

x = x + n

$$\{J \wedge vf < V\}$$



$$egin{aligned} J:x\cdot n! &= X\wedge n \geq 0\ B:n
eq 0\ vf &= n \end{aligned}$$

$$\{J \wedge B \wedge vf = V\}$$
 $\{x \cdot n! = X \wedge n \ge 0 \wedge n \ne 0 \wedge n = V\}$
 $(* n = V > 0 \Rightarrow n! = n \cdot (n-1)! \wedge n - 1 \ge 0 \wedge n - 1 < V *)$
 $\{x \cdot n \cdot (n-1)! = X \wedge n - 1 \ge 0 \wedge n - 1 < V\}$
 $x := x * n;$
 $\{x \cdot (n-1)! = X \wedge n - 1 > 0 \wedge n - 1 < V\}$

$$\{J \wedge vf < V\}$$



```
egin{aligned} J:x\cdot n! &= X\wedge n \geq 0\ B:n 
eq 0\ vf &= n \end{aligned}
```

```
4 Body of the loop: \{J \land B \land vf = V\} S \{J \land vf < V\}
  \{J \wedge B \wedge vf = V\}
  \{x \cdot n! = X \land n > 0 \land n \neq 0 \land n = V\}
     (*n = V > 0 \Rightarrow n! = n \cdot (n-1)! \land n-1 > 0 \land n-1 < V *)
  \{x \cdot n \cdot (n-1)! = X \land n-1 > 0 \land n-1 < V\}
x := x * n:
  \{x \cdot (n-1)! = X \land n-1 > 0 \land n-1 < V\}
n := n - 1:
  \{J \wedge vf < V\}
```



```
egin{aligned} J:x\cdot n! &= X\wedge n \geq 0\ B:n 
eq 0\ vf &= n \end{aligned}
```

```
4 Body of the loop: \{J \land B \land vf = V\} S \{J \land vf < V\}
  \{J \wedge B \wedge vf = V\}
  \{x \cdot n! = X \land n > 0 \land n \neq 0 \land n = V\}
     (*n = V > 0 \Rightarrow n! = n \cdot (n-1)! \land n-1 > 0 \land n-1 < V *)
  \{x \cdot n \cdot (n-1)! = X \land n-1 > 0 \land n-1 < V\}
x := x * n:
  \{x \cdot (n-1)! = X \land n-1 > 0 \land n-1 < V\}
n := n - 1:
   \{x \cdot n! = X \wedge n > 0 \wedge n < V\}
  \{J \wedge vf < V\}
```

Exercise 7.2: Conclusion



5 The command $\{P\}$ T_0 ; while B do S end $\{Q\}$ solves the problem:

```
var x, n : \mathbb{Z};
  \{P: X = n! \land n > 0\}
x := 1:
  \{J: x \cdot n! = X \wedge n > 0\}
    (* vf = n *)
while n \neq 0 do
  x := x * n;
  n := n - 1;
end:
\{Q: x = X\}
```

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Exercise 7.8: Dijkstra's FUSC

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Square Roof

Square Root (Similar to Exercise 7.3) Exercise 7.4: Square Root, Revisited

Integral Division

Exercise 7.5: Integral Division

Exercise 7.6: Variation on Integral Division Efficiency Comparison Exercises 7.5 and 7.

Exercise 7.8: FUSC



Dijkstra's FUSC function (check EWD 570):

$$f(0) = 0 \ f(1) = 1 \ f(2 \cdot n) = f(n) \ f(2 \cdot n + 1) = f(n) + f(n + 1)$$

Exercise 7.8: FUSC



Dijkstra's FUSC function (check EWD 570):

$$f(0) = 0$$
 $f(1) = 1$
 $f(2 \cdot n) = f(n)$
 $f(2 \cdot n + 1) = f(n) + f(n + 1)$

For instance:

$$f(2) = f(2 \cdot 1) = f(1) = 1$$



Dijkstra's FUSC function (check EWD 570):

$$f(0) = 0 \ f(1) = 1 \ f(2 \cdot n) = f(n) \ f(2 \cdot n + 1) = f(n) + f(n + 1)$$

$$f(2) = f(2 \cdot 1) = f(1) = 1$$

 $f(3) = f(2 \cdot 1 + 1) = f(1) + f(1 + 1) = 1 + 1 = 2$



Dijkstra's FUSC function (check EWD 570):

$$f(0) = 0 \ f(1) = 1 \ f(2 \cdot n) = f(n) \ f(2 \cdot n + 1) = f(n) + f(n + 1)$$

$$f(2) = f(2 \cdot 1) = f(1) = 1$$

 $f(3) = f(2 \cdot 1 + 1) = f(1) + f(1 + 1) = 1 + 1 = 2$
 $f(4) = f(2 \cdot 2) = f(2) = 1$



Dijkstra's FUSC function (check EWD 570):

$$f(0) = 0 \ f(1) = 1 \ f(2 \cdot n) = f(n) \ f(2 \cdot n + 1) = f(n) + f(n + 1)$$

$$f(2) = f(2 \cdot 1) = f(1) = 1$$

 $f(3) = f(2 \cdot 1 + 1) = f(1) + f(1 + 1) = 1 + 1 = 2$
 $f(4) = f(2 \cdot 2) = f(2) = 1$
 $f(5) = ??$



Dijkstra's FUSC function (check EWD 570):

$$f(0) = 0 \ f(1) = 1 \ f(2 \cdot n) = f(n) \ f(2 \cdot n + 1) = f(n) + f(n + 1)$$

$$f(2) = f(2 \cdot 1) = f(1) = 1$$

$$f(3) = f(2 \cdot 1 + 1) = f(1) + f(1 + 1) = 1 + 1 = 2$$

$$f(4) = f(2 \cdot 2) = f(2) = 1$$

$$f(5) = f(2) + f(3) = 3$$



Dijkstra's FUSC function (check EWD 570):

$$f(0) = 0 \ f(1) = 1 \ f(2 \cdot n) = f(n) \ f(2 \cdot n + 1) = f(n) + f(n + 1)$$

Also notice:

$$f(1) = 1$$

= 0 + 1
= $f(0) + f(0 + 1)$
= $f(2 \cdot 0 + 1)$



Dijkstra's FUSC function:

$$egin{aligned} f(0) &= 0 \ f(1) &= 1 \ f(2 \cdot n) &= f(n) \ f(2 \cdot n+1) &= f(n) + f(n+1) \end{aligned}$$

We consider the specification:

```
egin{aligned} 	extsf{var} & n, \ x: \ \mathbb{Z}; \ & \{P: \ n \geq 0 \land Z = f(n)\} \ T \ & \{Q: \ Z = x\} \end{aligned}
```

Hint:

Use $Z = y \cdot f(n) + x \cdot f(n+1)$ in the invariant.



$$P: n \geq 0 \wedge Z = f(n)$$

$$Q:Z=x$$

- 0 We decide that we need a **while**-program: We only have a recurrence for f(n), so we need iteration.
- 1 Choose an invariant J, and guard B such that $J \land \neg B \Rightarrow Q$. Following the hint, we choose:

$$egin{aligned} J:n \geq 0 \wedge (Z=y \cdot f(n) + x \cdot f(n+1)) \ B:n
eq 0 \end{aligned}$$

We have:

$$egin{aligned} J \wedge
eg B \Rightarrow Z &= y \cdot f(0) + x \cdot f(0+1) \ Z &= y \cdot 0 + x \cdot 1 \ Z &= x \end{aligned}$$

Exercise 7.8: Initialization and Variant



$$egin{aligned} f(0) &= 0 & P: n \geq 0 \wedge Z = f(n) \ f(1) &= 1 & Q: Z = x \ f(2 \cdot n) &= f(n) & J: n \geq 0 \wedge Z = y \cdot f(n) + x \cdot f(n+1) \ f(2 \cdot n + 1) &= f(n) + f(n+1) & B: n
eq 0 \end{aligned}$$

2 Initialization: Find a command T_0 such that $\{P\}$ T_0 $\{J\}$

```
egin{aligned} \{P: \ n \geq 0 \land Z = f(n)\} \ & 	ext{(* calculus; } f(n) 	ext{ is defined for } n \geq 0 	ext{ *)} \ & \{n \geq 0 \land Z = 1 \cdot f(n) + 0 \cdot f(n+1)\} \ y := 1; \ x := 0; \ & \{J: \ n \geq 0 \land Z = y \cdot f(n) + x \cdot f(n+1)\} \end{aligned}
```

3 Variant function: $vf \in \mathbb{Z}$ and $J \wedge B \Rightarrow vf \geq 0$. We choose $vf = n \in \mathbb{Z}$ and so $J \wedge B \Rightarrow vf \geq 0$, because J contains the conjunct n > 0.



By observing the inductive part of the definition of f:

$$f(2\cdot n)=f(n) \ f(2\cdot n+1)=f(n)+f(n+1)$$

We infer that we should work towards a command of the form:

```
\{J:\ n\geq 0 \land Z=y\cdot f(n)+x\cdot f(n+1)\ \land\ B:\ n
eq 0\ \land\ n=V\}
while n \neq 0 do
 if n \mod 2 = 0 then
    S_1: (* Do something if n is even *)
  else
    S_2: (* Do something else if n is odd *)
 end:
 S_3: (* Modify n *)
end:
 \{J \wedge vf < V\}
```



if
$$n \mod 2 = 0$$
 then
$$\{n \mod 2 = 0 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\}$$

$$\{0 < n = V \land Z = y \cdot f(n \ \mathsf{div} \ 2) + x \cdot f(n \ \mathsf{div} \ 2 + 1)\}$$
 else



```
if n \mod 2 = 0 then \{n \mod 2 = 0 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\} (* (n \mod 2 = 0) \Rightarrow n = 2(n \operatorname{div} 2) + n \operatorname{mod} 2 = 2(n \operatorname{div} 2) *)
```

```
\{0 < n = V \land Z = y \cdot f(n \ \mathsf{div} \ 2) + x \cdot f(n \ \mathsf{div} \ 2 + 1)\} else
```



```
if n \mod 2 = 0 then
```

```
 \begin{cases} n \bmod 2 = 0 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1) \} \\ \binom{*}{n \bmod 2} = 0) \Rightarrow n = 2(n \bmod 2) + n \bmod 2 = 2(n \bmod 2) * ) \\ \binom{*}{n \bmod 2} = 0 \land 0 < n = V \land Z = y \cdot f(2(n \bmod 2)) + x \cdot f(2(n \bmod 2) + 1) \} \end{cases}
```

```
\{0 < n = V \land Z = y \cdot f(n \ \mathsf{div} \ 2) + x \cdot f(n \ \mathsf{div} \ 2 + 1)\} else
```



```
if n \bmod 2 = 0 then \{n \bmod 2 = 0 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\}\ (* (n \bmod 2 = 0) \Rightarrow n = 2(n \operatorname{div} 2) + n \operatorname{mod} 2 = 2(n \operatorname{div} 2)*) \{n \bmod 2 = 0 \land 0 < n = V \land Z = y \cdot f(2(n \operatorname{div} 2)) + x \cdot f(2(n \operatorname{div} 2) + 1)\}\ (* logic; definition f(n); n > 0 \land (n \operatorname{mod} 2 = 0) \Rightarrow n \geq 2*)
```

```
\{0 < n = V \land Z = y \cdot f(n \ \mathsf{div} \ 2) + x \cdot f(n \ \mathsf{div} \ 2 + 1)\} else
```

else



```
\begin{array}{l} \textbf{if } n \ \textbf{mod} \ 2 = 0 \ \textbf{then} \\ \{n \ \textbf{mod} \ 2 = 0 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\} \\ \text{(* } (n \ \textbf{mod} \ 2 = 0) \Rightarrow n = 2(n \ \textbf{div} \ 2) + n \ \textbf{mod} \ 2 = 2(n \ \textbf{div} \ 2) \ *) \\ \{n \ \textbf{mod} \ 2 = 0 \land 0 < n = V \land Z = y \cdot f(2(n \ \textbf{div} \ 2)) + x \cdot f(2(n \ \textbf{div} \ 2) + 1)\} \\ \text{(* } logic; definition } f(n); n > 0 \land (n \ \textbf{mod} \ 2 = 0) \Rightarrow n \geq 2 \ *) \\ \{0 < n = V \land Z = y \cdot f(n \ \textbf{div} \ 2) + x \cdot (f(n \ \textbf{div} \ 2) + f(n \ \textbf{div} \ 2 + 1))\} \end{array}
```



```
\begin{array}{l} \textbf{if } n \ \textbf{mod} \ 2 = 0 \ \textbf{then} \\ \{n \ \textbf{mod} \ 2 = 0 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\} \\ \text{(* } (n \ \textbf{mod} \ 2 = 0) \Rightarrow n = 2(n \ \textbf{div} \ 2) + n \ \textbf{mod} \ 2 = 2(n \ \textbf{div} \ 2) *) \\ \{n \ \textbf{mod} \ 2 = 0 \land 0 < n = V \land Z = y \cdot f(2(n \ \textbf{div} \ 2)) + x \cdot f(2(n \ \textbf{div} \ 2) + 1)\} \\ \text{(* logic; definition } f(n); n > 0 \land (n \ \textbf{mod} \ 2 = 0) \Rightarrow n \geq 2 *) \\ \{0 < n = V \land Z = y \cdot f(n \ \textbf{div} \ 2) + x \cdot (f(n \ \textbf{div} \ 2) + f(n \ \textbf{div} \ 2 + 1))\} \\ \text{(* calculus (common factor } f(n \ \textbf{div} \ 2)) *) \\ \\ \{0 < n = V \land Z = y \cdot f(n \ \textbf{div} \ 2) + x \cdot f(n \ \textbf{div} \ 2 + 1)\} \\ \textbf{else} \end{array}
```



```
 \begin{aligned} & \text{if } n \text{ mod } 2 = 0 \text{ then} \\ & \{ n \text{ mod } 2 = 0 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1) \} \\ & (* (n \text{ mod } 2 = 0) \Rightarrow n = 2(n \text{ div } 2) + n \text{ mod } 2 = 2(n \text{ div } 2) *) \\ & \{ n \text{ mod } 2 = 0 \land 0 < n = V \land Z = y \cdot f(2(n \text{ div } 2)) + x \cdot f(2(n \text{ div } 2) + 1) \} \\ & (* \text{ logic; definition } f(n); n > 0 \land (n \text{ mod } 2 = 0) \Rightarrow n \geq 2 *) \\ & \{ 0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot (f(n \text{ div } 2) + f(n \text{ div } 2 + 1)) \} \\ & (* \text{ calculus (common factor } f(n \text{ div } 2)) *) \\ & \{ 0 < n = V \land Z = (x + y) \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1) \} \end{aligned}
```



```
\begin{array}{l} \textbf{if } n \ \textbf{mod } 2 = 0 \ \textbf{then} \\ & \{ n \ \textbf{mod } 2 = 0 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1) \} \\ & (* \ (n \ \textbf{mod } 2 = 0) \Rightarrow n = 2(n \ \textbf{div } 2) + n \ \textbf{mod } 2 = 2(n \ \textbf{div } 2) *) \\ & \{ n \ \textbf{mod } 2 = 0 \land 0 < n = V \land Z = y \cdot f(2(n \ \textbf{div } 2)) + x \cdot f(2(n \ \textbf{div } 2) + 1) \} \\ & (* \ logic; \ definition \ f(n); \ n > 0 \land (n \ \textbf{mod } 2 = 0) \Rightarrow n \geq 2 \ *) \\ & \{ 0 < n = V \land Z = y \cdot f(n \ \textbf{div } 2) + x \cdot (f(n \ \textbf{div } 2) + f(n \ \textbf{div } 2 + 1)) \} \\ & (* \ calculus \ (common \ factor \ f(n \ \textbf{div } 2)) \ *) \\ & \{ 0 < n = V \land Z = (x + y) \cdot f(n \ \textbf{div } 2) + x \cdot f(n \ \textbf{div } 2 + 1) \} \\ & y := x + y; \\ & \{ 0 < n = V \land Z = y \cdot f(n \ \textbf{div } 2) + x \cdot f(n \ \textbf{div } 2 + 1) \} \\ & \textbf{else} \end{array}
```



if $n \mod 2 = 0$ then

```
\begin{array}{l} y := x + y; \ (\text{* see previous slide *}) \\ \{0 < n = V \land Z = y \cdot f(n \ \text{div } 2) + x \cdot f(n \ \text{div } 2 + 1)\} \\ \text{else} \\ \{n \ \text{mod} \ 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n + 1)\} \end{array}
```

$$\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}$$
 end;



```
\begin{array}{l} \textbf{if } n \ \textbf{mod} \ 2 = 0 \ \ \textbf{then} \\ y := x + y; \ ( * \textit{see previous slide} \ * ) \\ \{ 0 < n = V \land Z = y \cdot f(n \ \textbf{div} \ 2) + x \cdot f(n \ \textbf{div} \ 2 + 1) \} \\ \textbf{else} \\ \{ n \ \textbf{mod} \ 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n + 1) \} \\ ( * \textit{logic; similarly as before: } n = 2(n \ \textbf{div} \ 2) + 1 \ * ) \end{array}
```

```
\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\} end;
```



```
\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\} end;
```



```
\begin{array}{l} \textbf{if } n \ \textbf{mod} \ 2 = 0 \ \textbf{ then} \\ y := x + y; \ ( * \ \textit{see previous slide} \ * ) \\ \{ 0 < n = V \land Z = y \cdot f(n \ \textbf{div} \ 2) + x \cdot f(n \ \textbf{div} \ 2 + 1) \} \\ \textbf{else} \\ \{ n \ \textbf{mod} \ 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n + 1) \} \\ \ ( * \ \textit{logic}; \textit{similarly as before:} \ n = 2(n \ \textbf{div} \ 2) + 1 \ * ) \\ \{ 0 < n = V \land Z = y \cdot f(2(n \ \textbf{div} \ 2) + 1) + x \cdot f(2(n \ \textbf{div} \ 2) + 2) \} \\ \ ( * \ \textit{calculus} \ * ) \end{array}
```

```
\{0 < n = V \land Z = y \cdot f(n \ \mathsf{div} \ 2) + x \cdot f(n \ \mathsf{div} \ 2 + 1)\} end;
```



```
\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\} end;
```



```
 \begin{aligned} & \text{if } n \text{ mod } 2 = 0 \text{ then} \\ & y := x + y; \text{ (* see previous slide *)} \\ & \{ 0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1) \} \end{aligned} \\ & \text{else} \\ & \{ n \text{ mod } 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n + 1) \} \\ & \text{ (* logic; similarly as before: } n = 2(n \text{ div } 2) + 1 \text{ *)} \\ & \{ 0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2) + 2) \} \\ & \text{ (* calculus *)} \\ & \{ 0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2 + 1)) \} \\ & \text{ (* definition } f(n) \text{ expanded twice *)} \end{aligned}
```

```
\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\} end;
```



```
if n \bmod 2 = 0 then y := x + y; (* see previous slide *) \{0 < n = V \land Z = y \cdot f(n \operatorname{div} 2) + x \cdot f(n \operatorname{div} 2 + 1)\} else \{n \bmod 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n + 1)\} (* logic; similarly as before: n = 2(n \operatorname{div} 2) + 1*) \{0 < n = V \land Z = y \cdot f(2(n \operatorname{div} 2) + 1) + x \cdot f(2(n \operatorname{div} 2) + 2)\} (* calculus *) \{0 < n = V \land Z = y \cdot f(2(n \operatorname{div} 2) + 1) + x \cdot f(2(n \operatorname{div} 2 + 1))\} (* definition f(n) expanded twice *) \{0 < n = V \land Z = y \cdot (f(n \operatorname{div} 2) + f(n \operatorname{div} 2 + 1)) + x \cdot f(n \operatorname{div} 2 + 1)\}
```

```
\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\} end;
```



```
if n \mod 2 = 0 then
 y := x + y; (* see previous slide *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
else
   \{n \bmod 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\}
    (* logic; similarly as before: n = 2(n \operatorname{div} 2) + 1 *)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2) + 2)\}
    (* calculus *)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2 + 1))\}
    (* definition f(n) expanded twice *)
   \{0 < n = V \land Z = y \cdot (f(n \text{ div } 2) + f(n \text{ div } 2 + 1)) + x \cdot f(n \text{ div } 2 + 1)\}
    (* calculus (common factor f(n \operatorname{div} 2 + 1) *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
end:
```



```
if n \mod 2 = 0 then
 y := x + y; (* see previous slide *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
else
   \{n \bmod 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\}
    (* logic; similarly as before: n = 2(n \operatorname{div} 2) + 1*)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2) + 2)\}
    (* calculus *)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2 + 1))\}
   (* definition f(n) expanded twice *)
   \{0 < n = V \land Z = y \cdot (f(n \text{ div } 2) + f(n \text{ div } 2 + 1)) + x \cdot f(n \text{ div } 2 + 1)\}
    (* calculus (common factor f(n \operatorname{div} 2 + 1) *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + (x + y) \cdot f(n \text{ div } 2 + 1)\}
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
end:
```



```
if n \mod 2 = 0 then
 y := x + y; (* see previous slide *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
else
   \{n \bmod 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\}
    (* logic; similarly as before: n = 2(n \operatorname{div} 2) + 1*)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2) + 2)\}
    (* calculus *)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2 + 1))\}
   (* definition f(n) expanded twice *)
   \{0 < n = V \land Z = y \cdot (f(n \text{ div } 2) + f(n \text{ div } 2 + 1)) + x \cdot f(n \text{ div } 2 + 1)\}
    (* calculus (common factor f(n \operatorname{div} 2 + 1) *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + (x + y) \cdot f(n \text{ div } 2 + 1)\}
 x := x + y:
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
end:
```



```
if n \mod 2 = 0 then
 y := x + y; (* see previous slide *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
else
   \{n \bmod 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\}
    (* logic; similarly as before: n = 2(n \operatorname{div} 2) + 1*)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2) + 2)\}
    (* calculus *)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2 + 1))\}
   (* definition f(n) expanded twice *)
   \{0 < n = V \land Z = y \cdot (f(n \text{ div } 2) + f(n \text{ div } 2 + 1)) + x \cdot f(n \text{ div } 2 + 1)\}
    (* calculus (common factor f(n \operatorname{div} 2 + 1) *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + (x + y) \cdot f(n \text{ div } 2 + 1)\}
 x := x + y:
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
end; (* collect branches *)
\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
```



```
if n \mod 2 = 0 then
 y := x + y; (* see previous slide *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
else
   \{n \bmod 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\}
    (* logic; similarly as before: n = 2(n \operatorname{div} 2) + 1*)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2) + 2)\}
   (* calculus *)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2 + 1))\}
   (* definition f(n) expanded twice *)
   \{0 < n = V \land Z = y \cdot (f(n \text{ div } 2) + f(n \text{ div } 2 + 1)) + x \cdot f(n \text{ div } 2 + 1)\}
    (* calculus (common factor f(n \operatorname{div} 2 + 1) *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + (x + y) \cdot f(n \text{ div } 2 + 1)\}
 x := x + y:
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
end; (* collect branches *)
\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
 (* 0 < n = V \Rightarrow 0 < n \text{ div } 2 < V *)
```



```
if n \mod 2 = 0 then
 y := x + y; (* see previous slide *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
else
   \{n \bmod 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\}
    (* logic; similarly as before: n = 2(n \operatorname{div} 2) + 1*)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2) + 2)\}
   (* calculus *)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2 + 1))\}
   (* definition f(n) expanded twice *)
   \{0 < n = V \land Z = y \cdot (f(n \text{ div } 2) + f(n \text{ div } 2 + 1)) + x \cdot f(n \text{ div } 2 + 1)\}
    (* calculus (common factor f(n \operatorname{div} 2 + 1) *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + (x + y) \cdot f(n \text{ div } 2 + 1)\}
 x := x + y:
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
end; (* collect branches *)
\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
(* 0 < n = V \Rightarrow 0 < n \text{ div } 2 < V *)
\{0 < n \text{ div } 2 < V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
```



```
if n \mod 2 = 0 then
 y := x + y; (* see previous slide *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
else
   \{n \bmod 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\}
    (* logic; similarly as before: n = 2(n \operatorname{div} 2) + 1*)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2) + 2)\}
    (* calculus *)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2 + 1))\}
   (* definition f(n) expanded twice *)
   \{0 < n = V \land Z = y \cdot (f(n \text{ div } 2) + f(n \text{ div } 2 + 1)) + x \cdot f(n \text{ div } 2 + 1)\}
    (* calculus (common factor f(n \operatorname{div} 2 + 1) *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + (x + y) \cdot f(n \text{ div } 2 + 1)\}
 x := x + y:
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
end; (* collect branches *)
\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
(* 0 < n = V \Rightarrow 0 < n \text{ div } 2 < V *)
\{0 < n \text{ div } 2 < V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
n := n \operatorname{div} 2;
```



```
if n \mod 2 = 0 then
 y := x + y; (* see previous slide *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
else
   \{n \bmod 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\}
    (* logic; similarly as before: n = 2(n \operatorname{div} 2) + 1*)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2) + 2)\}
   (* calculus *)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2 + 1))\}
   (* definition f(n) expanded twice *)
   \{0 < n = V \land Z = y \cdot (f(n \text{ div } 2) + f(n \text{ div } 2 + 1)) + x \cdot f(n \text{ div } 2 + 1)\}
    (* calculus (common factor f(n \operatorname{div} 2 + 1) *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + (x + y) \cdot f(n \text{ div } 2 + 1)\}
 x := x + y;
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
end; (* collect branches *)
\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
(* 0 < n = V \Rightarrow 0 < n \text{ div } 2 < V *)
\{0 < n \text{ div } 2 < V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
n := n \operatorname{div} 2;
\{0 < n < V \land Z = y \cdot f(n) + x \cdot f(n+1)\}\
```



```
if n \mod 2 = 0 then
 y := x + y; (* see previous slide *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
else
   \{n \bmod 2 = 1 \land 0 < n = V \land Z = y \cdot f(n) + x \cdot f(n+1)\}
    (* logic; similarly as before: n = 2(n \operatorname{div} 2) + 1*)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2) + 2)\}
   (* calculus *)
   \{0 < n = V \land Z = y \cdot f(2(n \text{ div } 2) + 1) + x \cdot f(2(n \text{ div } 2 + 1))\}
   (* definition f(n) expanded twice *)
   \{0 < n = V \land Z = y \cdot (f(n \text{ div } 2) + f(n \text{ div } 2 + 1)) + x \cdot f(n \text{ div } 2 + 1)\}
    (* calculus (common factor f(n \operatorname{div} 2 + 1) *)
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + (x + y) \cdot f(n \text{ div } 2 + 1)\}
 x := x + y;
   \{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
end; (* collect branches *)
\{0 < n = V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
(* 0 < n = V \Rightarrow 0 < n \text{ div } 2 < V *)
\{0 < n \text{ div } 2 < V \land Z = y \cdot f(n \text{ div } 2) + x \cdot f(n \text{ div } 2 + 1)\}
n := n \operatorname{div} 2;
\{0 < n < V \land Z = y \cdot f(n) + x \cdot f(n+1)\}\
\{J \wedge vf < V\}
```

Exercise 7.8: Conclusion



5 The following command solves the problem:

```
var n, x, y : \mathbb{Z};
  \{P: n > 0 \land Z = f(n)\}
y := 1;
x := 0:
  \{J: n > 0 \land Z = y \cdot f(n) + x \cdot f(n+1)\}
    (* vf = n *)
while n \neq 0 do
  if n \mod 2 = 0 then
        y := x + y;
   else
        x := x + y;
  end:
  n:=n \operatorname{div} 2;
end;
  \{Q: x = Z\}
```

Outline



Exercise 7.1: Powers

Exercise 7.2: Factorial

Exercise 7.8: Dijkstra's FUSC

Summing an array

Square Root

Square Root (Similar to Exercise 7.3) Exercise 7.4: Square Root, Revisited

Integral Division

Exercise 7.5: Integral Division

Exercise 7.6: Variation on Integral Division Efficiency Comparison Exercises 7.5 and 7.

Example: Summing an array



```
\begin{array}{l} \textbf{const} \ n: \ \mathbb{N}, \ a: \ \textbf{array} \ [0..n) \ \textbf{of} \ \mathbb{Z}; \\ \textbf{var} \ x: \ \mathbb{Z}; \\ \{P: \ \textbf{true}\} \\ S \\ \{Q: \ x = \Sigma(a[i] \mid i: i \in [0..n))\} \end{array}
```

▶ To reduce the size of our formulas we write, for $0 \le k \le n$:

$$S(k) = \Sigma(a[i] \mid i:i \in [0..k))$$

▶ This way, we rewrite the postcondition: Q: x = S(n).

Summing an array: Recurrence



We consider a recurrence relation for $S(k) = \Sigma(a[i] \mid i : i \in [0..k))$. In this case:

$$S(0) = 0 \ 0 \leq k < n \Rightarrow S(k+1) = a[k] + S(k)$$

Summing an array: Recurrence



We consider a recurrence relation for $S(k) = \Sigma(a[i] \mid i : i \in [0..k))$. In this case:

$$egin{aligned} S(0) &= 0 \ 0 &\leq k < n \Rightarrow & S(k+1) = a[k] + S(k) \end{aligned}$$

Justification:

- lt is clear that S(0) = 0, since the domain of the sum is empty.
- For $0 \le k < n$, we compute S(k+1):

$$S(k+1) = (* definition *)$$
 $\Sigma(a[i] \mid i:i \in [0..k+1))$
 $= (* split domain: i = k \lor i < k *)$
 $a[k] + \Sigma(a[i] \mid i:i \in [0..k))$
 $= (* definition *)$
 $a[k] + S(k)$

Summing an array: Invariant and Guard



P: true

$$Q:x=S(n)$$

- 0 We expect to add values iteratively: we need a **while**-program.
- 1 Choose an invariant J and guard B such that $J \land \neg B \Rightarrow Q$. We obtain J by using "replacing a constant by a variable":
 - (1) Replace n in Q by variable k and
 - (2) Include the domain condition $0 \le k \le n$:

$$J:0\leq k\leq n\wedge x=S(k)$$

$$B: k \neq n$$

Summing an array: Invariant and Guard



P: true

$$Q: x = S(n)$$

- 0 We expect to add values iteratively: we need a **while**-program.
- 1 Choose an invariant J and guard B such that $J \land \neg B \Rightarrow Q$. We obtain J by using "replacing a constant by a variable":
 - (1) Replace n in Q by variable k and
 - (2) Include the domain condition $0 \le k \le n$:

$$J: 0 \leq k \leq n \wedge x = S(k)$$

 $B: k \neq n$

The proof obligation holds:

$$J \wedge \neg B \equiv 0 \leq k \leq n \wedge x = S(k) \wedge k = n$$

 \Rightarrow (* substitute $k = n$; logic *)
 $Q: x = S(n)$

Summing an array: Initialization & Variant



P: true

 $J: 0 \leq k \leq n \wedge x = S(k)$

 $B: k \neq n$

2 Initialization: Find a command T_0 such that $\{P\}$ T_0 $\{J\}$.

Summing an array: Initialization & Variant



```
P: 	extbf{true} \ J: 0 \leq k \leq n \wedge x = S(k) \ B: k 
eq n
```

2 Initialization: Find a command T_0 such that $\{P\}$ T_0 $\{J\}$. Since S(0)=0, it suffices to choose k:=0; x:=0; $\{P: \ \, \ \, \ \, \}$ $(*n\in\mathbb{N};S(0)=0\ ^*)$ $\{0\le 0\le n\land 0=S(0)\}$ k:=0; $\{0\le k\le n\land 0=S(k)\}$ x:=0; $\{J:\ 0< k< n\land x=S(k)\}$

Summing an array: Initialization & Variant



```
P: trueJ:0\leq k\leq n \wedge x=S(k)B:k
eq n
```

2 Initialization: Find a command T_0 such that $\{P\}$ T_0 $\{J\}$. Since S(0)=0, it suffices to choose k:=0; x:=0; $\{P: \mbox{ true}\}$ $(*n\in\mathbb{N};S(0)=0*)$ $\{0\le 0\le n\land 0=S(0)\}$ k:=0; $\{0\le k\le n\land 0=S(k)\}$ x:=0; $\{J:\ 0\le k\le n\land x=S(k)\}$

3 Variant function: Choose a $vf \in \mathbb{Z}$ and prove $J \wedge B \Rightarrow vf \geq 0$. Since initially k=0 and $B: k \neq n$, we must increase k. We choose $vf = n - k \in \mathbb{Z}$. Clearly, $J \wedge B \Rightarrow J \Rightarrow k < n \equiv vf > 0$.



$${J \wedge B \wedge vf = V}$$



```
\{J \wedge B \wedge vf = V\}
(* definitions J, B, and vf *)
\{0 \leq k \leq n \wedge x = S(k) \wedge k \neq n \wedge n - k = V\}
```



```
 \begin{cases} J \wedge B \wedge vf = V \\ \text{(* definitions } J, B, \text{ and } vf \text{ *)} \end{cases} \\ \{0 \leq k \leq n \wedge x = S(k) \wedge k \neq n \wedge n - k = V \} \\ \text{(* calculus; } k < n; \text{ prepare } k := k + 1; \text{ use recurrence *)} \\ \{0 \leq k + 1 \leq n \wedge x + a[k] = S(k + 1) \wedge n - (k + 1) < V \}
```



```
\{J \wedge B \wedge vf = V\}

(* definitions J, B, and vf *)

\{0 \le k \le n \wedge x = S(k) \wedge k \ne n \wedge n - k = V\}

(* calculus; k < n; prepare k := k + 1; use recurrence *)

\{0 \le k + 1 \le n \wedge x + a[k] = S(k + 1) \wedge n - (k + 1) < V\}

x := x + a[k];
```



```
\{J \wedge B \wedge vf = V\}

(* definitions J, B, and vf *)

\{0 \le k \le n \wedge x = S(k) \wedge k \ne n \wedge n - k = V\}

(* calculus; k < n; prepare k := k + 1; use recurrence *)

\{0 \le k + 1 \le n \wedge x + a[k] = S(k + 1) \wedge n - (k + 1) < V\}

x := x + a[k];

\{0 \le k + 1 \le n \wedge x = S(k + 1) \wedge n - (k + 1) < V\}
```







```
\{J \wedge B \wedge vf = V\}
    (* definitions J. B. and vf *)
  \{0 < k < n \land x = S(k) \land k \neq n \land n - k = V\}
     (* calculus; k < n; prepare k := k + 1; use recurrence *)
  \{0 < k+1 < n \land x + a[k] = S(k+1) \land n - (k+1) < V\}
x := x + a[k];
  \{0 < k+1 < n \land x = S(k+1) \land n - (k+1) < V\}
k := k + 1:
  \{0 < k < n \land x = S(k) \land n - k < V\}
    (* definitions J, and vf *)
  \{J \wedge vf < V\}
```

Summing an array: Conclusion



5 We conclude that $\{P\}$ T_0 ; while B do S end $\{Q\}$ solves the problem:

```
const n : \mathbb{N}, a : \operatorname{array} [0..n) of \mathbb{Z};
var x:\mathbb{Z}:
   \{P: \mathsf{true}\}
k := 0:
x := 0:
   \{J: 0 < k < n \land x = S(k)\}
     (* vf = n - k *)
while k \neq n do
   x := x + a[k];
   k := k + 1:
end:
   \{Q: x = \Sigma(a[i] \mid i: i \in [0..n))\}
```

Outline



Exercise 7.1: Powers

Exercise 7.2: Factorial

Exercise 7.8: Dijkstra's FUSC

Summing an array

Square Root

Square Root (Similar to Exercise 7.3) Exercise 7.4: Square Root, Revisited

Integral Division

Exercise 7.5: Integral Division

Exercise 7.6: Variation on Integral Division Efficiency Comparison Exercises 7.5 and 7.

Exercise 7.3



```
\begin{aligned} & \textbf{const } x: \ \mathbb{N}; \\ & \textbf{var } y: \ \mathbb{Z}; \\ & \{P: \ \textbf{true}\} \\ & T \\ & \{Q: \ y \geq 0 \land y^2 \leq x < (y+1)^2\} \end{aligned}
```

- 0 We assume that there is no function '**sqrt**' available, and decide that we need a **while**-program.
- 1 Choose an invariant J and guard B such that $J \land \neg B \Rightarrow Q$. We use the heuristic split conjuncts.

$$egin{aligned} J: y \geq 0 \wedge y^2 \leq x \ B: (y+1)^2 \leq x \end{aligned}$$

Clearly, $J \wedge \neg B \equiv Q$.

Exercise 7.3: Initialization and Variant



2 Initialization: Find a command T_0 such that

$$\{P: \text{ true}\}\ T_0\ \{J:\ y\geq 0 \land y^2\leq x\}$$

We have:

```
\{P: 	extbf{true}\}
(*x \in \mathbb{N} *)
\{0 \geq 0 \wedge 0^2 \leq x\}
y:=0;
\{J: y \geq 0 \wedge y^2 \leq x\}
```

3 Variant function: $vf = x - y^2 \in \mathbb{Z}$ J contains the conjunct $y^2 \le x$, so trivially $J \wedge B \Rightarrow vf \ge 0$.



$$egin{aligned} J: y &\geq 0 \wedge y^2 \leq x \ B: (y+1)^2 \leq x \ vf &= x-y^2 \end{aligned}$$

4 Body of the loop:
$$\{J \land B \land vf = V\}$$
 S $\{J \land vf < V\}$ $\{J \land B \land vf = V\}$

$$\{J \wedge vf < V\}$$



$$J:y\geq 0 \wedge y^2 \leq x \ B:(y+1)^2 \leq x \ vf=x-y^2$$

4 Body of the loop:
$$\{J \land B \land vf = V\}$$
 S $\{J \land vf < V\}$
$$\{J \land B \land vf = V\}$$

$$\{y \ge 0 \land y^2 \le x \land (y+1)^2 \le x \land x - y^2 = V\}$$

$$\{J \wedge vf < V\}$$



$$egin{aligned} J: y &\geq 0 \wedge y^2 \leq x \ B: (y+1)^2 \leq x \ vf &= x-y^2 \end{aligned}$$

4 Body of the loop:
$$\{J \land B \land vf = V\}$$
 $S \{J \land vf < V\}$
 $\{J \land B \land vf = V\}$
 $\{y \ge 0 \land y^2 \le x \land (y+1)^2 \le x \land x - y^2 = V\}$
(* logic; prepare $y := y + 1$ *)

$$\{J \wedge vf < V\}$$



$$J:y\geq 0 \wedge y^2 \leq x \ B:(y+1)^2 \leq x \ vf=x-y^2$$

4 Body of the loop:
$$\{J \land B \land vf = V\}$$
 $S \{J \land vf < V\}$
 $\{J \land B \land vf = V\}$
 $\{y \ge 0 \land y^2 \le x \land (y+1)^2 \le x \land x - y^2 = V\}$
(* logic; prepare $y := y+1$ *)
 $\{y+1 \ge 0 \land (y+1)^2 \le x \land x - (y+1)^2 < V\}$

$$\{J \wedge vf < V\}$$



$$J:y\geq 0 \wedge y^2 \leq x \ B:(y+1)^2 \leq x \ vf=x-y^2$$

4 Body of the loop:
$$\{J \land B \land vf = V\}$$
 S $\{J \land vf < V\}$ $\{J \land B \land vf = V\}$ $\{y \ge 0 \land y^2 \le x \land (y+1)^2 \le x \land x - y^2 = V\}$ (* logic; prepare $y := y+1$ *) $\{y+1 \ge 0 \land (y+1)^2 \le x \land x - (y+1)^2 < V\}$ $y := y+1$; $\{J \land vf < V\}$



$$egin{aligned} J: y &\geq 0 \wedge y^2 \leq x \ B: (y+1)^2 \leq x \ vf &= x-y^2 \end{aligned}$$

$$\begin{array}{l} \text{4 Body of the loop: } \{J \wedge B \wedge vf = V\} \; S \; \{J \wedge vf < V\} \\ \{J \wedge B \wedge vf = V\} \\ \{y \geq 0 \; \wedge \; y^2 \leq x \; \wedge \; (y+1)^2 \leq x \; \wedge \; x - y^2 = V\} \\ \text{(* logic; prepare } y := y+1 \; *) \\ \{y+1 \geq 0 \; \wedge \; (y+1)^2 \leq x \; \wedge \; x - (y+1)^2 < V\} \\ y := y+1; \\ \{y \geq 0 \wedge y^2 \leq x \wedge x - y^2 < V\} \\ \{J \wedge vf < V\} \end{array}$$

Exercise 7.3: Conclusion



```
5 We conclude \{P\} T_0; while B do S end \{Q\}:
           const x : \mathbb{N};
           var y: \mathbb{Z};
           u := 0:
              \{J: y > 0 \land y^2 < x\}
                (* vf = x - y^2 *)
           while (y + 1) * (y + 1) < x do
              y := y + 1:
           end:
              \{Q: y > 0 \land y^2 < x < (y+1)^2\}
```

Exercise 7.4: Same Spec, New Invariant



```
\begin{aligned} & \textbf{const } x: \ \mathbb{N}; \\ & \textbf{var } y: \ \mathbb{Z}; \\ & \{P: \ \textbf{true}\} \\ & T \\ & \{Q: \ y \geq 0 \land y^2 \leq x < (y+1)^2\} \end{aligned}
```

- We decide that we need a while-program:We assume that there is no function sqrt available.
- 1 Choose an invariant J, and guard B such that $J \wedge \neg B \Rightarrow Q$.

Exercise 7.4: Same Spec, New Invariant



```
\begin{aligned} & \textbf{const } x: \ \mathbb{N}; \\ & \textbf{var } y: \ \mathbb{Z}; \\ & \{P: \ \textbf{true}\} \\ & T \\ & \{Q: \ y \geq 0 \land y^2 \leq x < (y+1)^2\} \end{aligned}
```

- We decide that we need a while-program:We assume that there is no function sqrt available.
- 1 Choose an invariant J, and guard B such that $J \wedge \neg B \Rightarrow Q$. Before the invariant was

$$y \ge 0 \wedge y^2 \le x$$

This time, we use the heuristic replace expression by variable, with z in place of y + 1:

Exercise 7.4: Same Spec, New Invariant



```
\begin{aligned} & \textbf{const } x: \ \mathbb{N}; \\ & \textbf{var } y: \ \mathbb{Z}; \\ & \{P: \ \textbf{true}\} \\ & T \\ & \{Q: \ y \geq 0 \land y^2 \leq x < (y+1)^2\} \end{aligned}
```

- 0 We decide that we need a while-program:We assume that there is no function sqrt available.
- 1 Choose an invariant J, and guard B such that $J \land \neg B \Rightarrow Q$. Before the invariant was

$$y > 0 \wedge y^2 < x$$

This time, we use the heuristic replace expression by variable, with z in place of y + 1:

$$egin{aligned} J: 0 \leq y < z \ \land \ y^2 \leq x < z^2 \ B: z
eq y + 1 \end{aligned}$$

Clearly,
$$J \wedge \neg B \Rightarrow Q$$
.

Exercise 7.4: Initialization and Variant



```
P: true J: 0 \leq y < z \wedge y^2 \leq x < z^2 B: z 
eq y + 1
```

2 Initialization: Find a command T_0 such that $\{P\}$ T_0 $\{J\}$.

```
\{P: 	extbf{true}\}\ (* calculus; x \in \mathbb{N}^*)\ \{0 \le 0 < x + 1 \land 0^2 \le x < (x+1)^2\}\ y := 0;\ z := x+1;\ \{J: \ 0 \le y < z \land y^2 \le x < z^2\}
```

3 Variant function:

Exercise 7.4: Initialization and Variant



```
P: true J: 0 \leq y < z \wedge y^2 \leq x < z^2 B: z 
eq y+1
```

2 Initialization: Find a command T_0 such that $\{P\}$ T_0 $\{J\}$.

```
\{P: 	extbf{true}\}\ (* calculus; x \in \mathbb{N} *)\ \{0 \le 0 < x + 1 \land 0^2 \le x < (x+1)^2\}\ y := 0;\ z := x + 1;\ \{J: \ 0 < y < z \land y^2 < x < z^2\}
```

3 Variant function: We choose $vf = z - y \in \mathbb{Z}$ J contains the conjunct $0 \le y < z$, so $J \land B \Rightarrow vf \ge 0$ holds. The body of the loop will narrow down the interval [y, z).



$$\{J \wedge B \wedge vf = V\}$$



$$\begin{cases} J \wedge B \wedge vf = V \\ \{ 0 \leq y < z \ \wedge \ y^2 \leq x < z^2 \ \wedge \ y + 1 \neq z \ \wedge \ z - y = V \} \end{cases}$$





```
 \begin{cases} J \wedge B \wedge vf = V \\ \{0 \leq y < z \ \wedge \ y^2 \leq x < z^2 \ \wedge \ y+1 \neq z \ \wedge \ z-y = V \} \\ \text{(* First, } y+1 < z. \ \textit{Then } y+1 < z \equiv y+2 \leq z \Rightarrow y < (y+z) \ \textit{div } 2 < z \ ^*) \\ \{0 < y < (y+z) \ \textit{div } 2 < z \ \wedge \ y^2 < x < z^2 \wedge z-y = V \} \end{cases}
```





```
 \begin{cases} J \land B \land vf = V \} \\ \{0 \leq y < z \land y^2 \leq x < z^2 \land y + 1 \neq z \land z - y = V \} \\ \text{(* First, } y + 1 < z. \text{ Then } y + 1 < z \equiv y + 2 \leq z \Rightarrow y < (y + z) \text{ div } 2 < z \text{ *}) \\ \{0 \leq y < (y + z) \text{ div } 2 < z \land y^2 \leq x < z^2 \land z - y = V \} \\ m := (y + z) \text{ div } 2; \\ \{0 \leq y < m < z \land y^2 \leq x < z^2 \land z - y = V \} \end{cases}
```



```
 \begin{cases} J \wedge B \wedge vf = V \} \\ \{0 \leq y < z \wedge y^2 \leq x < z^2 \wedge y + 1 \neq z \wedge z - y = V \} \\ \text{(* First, } y + 1 < z. \text{ Then } y + 1 < z \equiv y + 2 \leq z \Rightarrow y < (y + z) \text{ div } 2 < z \text{ *)} \\ \{0 \leq y < (y + z) \text{ div } 2 < z \wedge y^2 \leq x < z^2 \wedge z - y = V \} \\ m := (y + z) \text{ div } 2; \\ \{0 \leq y < m < z \wedge y^2 \leq x < z^2 \wedge z - y = V \} \\ \text{if } m * m \leq x \text{ then}
```

else

$$\{J \land vf < V\}$$



```
 \begin{cases} J \wedge B \wedge vf = V \} \\ \{0 \leq y < z \wedge y^2 \leq x < z^2 \wedge y + 1 \neq z \wedge z - y = V \} \\ \text{$(^*\textit{First}, y + 1 < z. \textit{Then } y + 1 < z \equiv y + 2 \leq z \Rightarrow y < (y + z) \textit{ div } 2 < z ^*)$} \\ \{0 \leq y < (y + z) \textit{ div } 2 < z \wedge y^2 \leq x < z^2 \wedge z - y = V \} \\ m := (y + z) \textit{ div } 2; \\ \{0 \leq y < m < z \wedge y^2 \leq x < z^2 \wedge z - y = V \} \\ \textit{ if } m * m \leq x \textit{ then } \\ \{m^2 \leq x \wedge 0 \leq y < m < z \wedge y^2 \leq x < z^2 \wedge z - y = V \}
```

else

$$\{m^2 > x \land 0 \le y < m < z \land y^2 \le x < z^2 \land z - y = V\}$$

$$\{J \wedge vf < V\}$$



else

$$\{m^2 > x \land 0 \le y < m < z \land y^2 \le x < z^2 \land z - y = V\}$$

$$\{J \wedge vf < V\}$$



```
 \begin{cases} J \wedge B \wedge vf = V \} \\ \{0 \leq y < z \wedge y^2 \leq x < z^2 \wedge y + 1 \neq z \wedge z - y = V \} \\ \text{(* First, } y + 1 < z. \text{ Then } y + 1 < z \equiv y + 2 \leq z \Rightarrow y < (y + z) \text{ div } 2 < z \text{ *)} \\ \{0 \leq y < (y + z) \text{ div } 2 < z \wedge y^2 \leq x < z^2 \wedge z - y = V \} \\ m := (y + z) \text{ div } 2; \\ \{0 \leq y < m < z \wedge y^2 \leq x < z^2 \wedge z - y = V \} \\ \text{if } m * m \leq x \text{ then} \\ \{m^2 \leq x \wedge 0 \leq y < m < z \wedge y^2 \leq x < z^2 \wedge z - y = V \} \\ \text{(* logic, combine conjuncts 1 and 3; calculus (prepare update to y) *)} \\ \{0 \leq m < z \wedge m^2 \leq x < z^2 \wedge z - m < V \}
```

else

$$\{m^2 > x \land 0 \le y < m < z \land y^2 \le x < z^2 \land z - y = V\}$$

$$\{J \wedge vf < V\}$$



```
\{J \wedge B \wedge vf = V\}
  \{0 < y < z \land y^2 < x < z^2 \land y + 1 \neq z \land z - y = V\}
    (* First, y + 1 < z. Then y + 1 < z \equiv y + 2 < z \Rightarrow y < (y + z) div 2 < z^*)
  \{0 \le y < (y+z) \text{ div } 2 < z \land y^2 \le x < z^2 \land z - y = V\}
m := (y + z) \text{ div } 2;
  \{0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
if m * m < x then
     \{m^2 < x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
      (* logic, combine conjuncts 1 and 3; calculus (prepare update to y) *)
     \{0 < m < z \land m^2 < x < z^2 \land z - m < V\}
  u := m:
else
     \{m^2 > x \land 0 \le y \le m \le z \land y^2 \le x \le z^2 \land z - y = V\}
```

$$\{J \wedge vf < V\}$$



```
\{J \wedge B \wedge vf = V\}
  \{0 < y < z \land y^2 < x < z^2 \land y + 1 \neq z \land z - y = V\}
     (* First, y + 1 < z. Then y + 1 < z \equiv y + 2 < z \Rightarrow y < (y + z) div <math>2 < z^*)
  \{0 \le y < (y+z) \text{ div } 2 < z \land y^2 \le x < z^2 \land z - y = V\}
m := (y + z) \text{ div } 2;
  \{0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
if m * m < x then
     \{m^2 < x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
       (* logic, combine conjuncts 1 and 3; calculus (prepare update to y) *)
     \{0 < m < z \land m^2 < x < z^2 \land z - m < V\}
  u := m:
     \{0 < y < z \land y^2 < x < z^2 \land z - y < V\}
else
     \{m^2 > x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
```

$$\{J \wedge vf < V\}$$



```
\{J \wedge B \wedge vf = V\}
  \{0 < y < z \land y^2 < x < z^2 \land y + 1 \neq z \land z - y = V\}
    (* First, y + 1 < z. Then y + 1 < z \equiv y + 2 < z \Rightarrow y < (y + z) div 2 < z^*)
  \{0 < y < (y+z) \text{ div } 2 < z \land y^2 < x < z^2 \land z - y = V\}
m := (y + z) \text{ div } 2;
  \{0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
if m * m < x then
     \{m^2 < x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
      (* logic, combine conjuncts 1 and 3; calculus (prepare update to y) *)
     \{0 < m < z \land m^2 < x < z^2 \land z - m < V\}
  y := m;
     \{0 < y < z \land y^2 < x < z^2 \land z - y < V\}
else
     \{m^2 > x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
       (* logic, combine conjuncts 1 and 3; calculus (prepare update to z) *)
```

$$\{J \wedge vf < V\}$$



```
\{J \wedge B \wedge vf = V\}
  \{0 < y < z \land y^2 < x < z^2 \land y + 1 \neq z \land z - y = V\}
     (* First, y + 1 < z. Then y + 1 < z \equiv y + 2 < z \Rightarrow y < (y + z) div 2 < z^*)
  \{0 < y < (y+z) \text{ div } 2 < z \land y^2 < x < z^2 \land z - y = V\}
m := (y + z) \operatorname{div} 2;
  \{0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
if m * m < x then
     \{m^2 < x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
      (* logic, combine conjuncts 1 and 3; calculus (prepare update to y) *)
     \{0 < m < z \land m^2 < x < z^2 \land z - m < V\}
  y := m;
     \{0 < y < z \land y^2 < x < z^2 \land z - y < V\}
else
     \{m^2 > x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
       (* logic, combine conjuncts 1 and 3; calculus (prepare update to z) *)
     \{0 \le y \le m \land y^2 \le x \le m^2 \land m - y \le V\}
```

$$\{J \wedge vf < V\}$$



```
\{J \wedge B \wedge vf = V\}
  \{0 < y < z \land y^2 < x < z^2 \land y + 1 \neq z \land z - y = V\}
    (* First, y + 1 < z. Then y + 1 < z \equiv y + 2 < z \Rightarrow y < (y + z) div 2 < z^*)
  \{0 < y < (y+z) \text{ div } 2 < z \land y^2 < x < z^2 \land z - y = V\}
m := (y + z) \text{ div } 2;
  \{0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
if m * m < x then
     \{m^2 < x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
      (* logic, combine conjuncts 1 and 3; calculus (prepare update to y) *)
     \{0 < m < z \land m^2 < x < z^2 \land z - m < V\}
  y := m;
     \{0 < y < z \land y^2 < x < z^2 \land z - y < V\}
else
     \{m^2 > x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
       (* logic, combine conjuncts 1 and 3; calculus (prepare update to z) *)
     \{0 < y < m \land y^2 < x < m^2 \land m - y < V\}
  z := m:
```

$$\{J \wedge vf < V\}$$

 $\{J \wedge vf < V\}$



```
\{J \wedge B \wedge vf = V\}
  \{0 < y < z \land y^2 < x < z^2 \land y + 1 \neq z \land z - y = V\}
    (* First, y + 1 < z. Then y + 1 < z \equiv y + 2 < z \Rightarrow y < (y + z) div 2 < z^*)
  \{0 < y < (y+z) \text{ div } 2 < z \land y^2 < x < z^2 \land z - y = V\}
m := (y + z) \text{ div } 2;
  \{0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
if m * m < x then
     \{m^2 < x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
      (* logic, combine conjuncts 1 and 3; calculus (prepare update to y) *)
     \{0 < m < z \land m^2 < x < z^2 \land z - m < V\}
  y := m;
     \{0 < y < z \land y^2 < x < z^2 \land z - y < V\}
else
     \{m^2 > x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
       (* logic, combine conjuncts 1 and 3; calculus (prepare update to z) *)
     \{0 < y < m \land y^2 < x < m^2 \land m - y < V\}
  z := m;
    \{0 < y < z \land y^2 < x < z^2 \land z - y < V\}
end
```



```
\{J \wedge B \wedge vf = V\}
  \{0 < y < z \land y^2 < x < z^2 \land y + 1 \neq z \land z - y = V\}
    (* First, y + 1 < z. Then y + 1 < z \equiv y + 2 < z \Rightarrow y < (y + z) div 2 < z^*)
  \{0 < y < (y+z) \text{ div } 2 < z \land y^2 < x < z^2 \land z - y = V\}
m := (y + z) \text{ div } 2;
  \{0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
if m * m < x then
     \{m^2 < x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
      (* logic, combine conjuncts 1 and 3; calculus (prepare update to y) *)
     \{0 < m < z \land m^2 < x < z^2 \land z - m < V\}
  y := m;
     \{0 < y < z \land y^2 < x < z^2 \land z - y < V\}
else
     \{m^2 > x \land 0 < y < m < z \land y^2 < x < z^2 \land z - y = V\}
      (* logic, combine conjuncts 1 and 3; calculus (prepare update to z) *)
     \{0 < y < m \land y^2 < x < m^2 \land m - y < V\}
  z := m;
     \{0 < y < z \land y^2 < x < z^2 \land z - y < V\}
end (* collect branches *)
  \{0 \le y < z \land y^2 \le x < z^2 \land z - y < V\}
  \{J \wedge vf < V\}
```

Exercise 7.4: Conclusion



5 Conclude that $\{P\}$ T_0 ; while B do S end $\{Q\}$ solves the problem.

```
\{P: \mathsf{true}\}
y := 0;
z := x + 1:
  \{J: 0 \leq y < z \wedge y^2 \leq x < z^2\}
   (* vf = z - y *)
while y + 1 \neq z do
  m := (y + z) \text{ div } 2;
  if m * m < x then
     y := m;
   else
     z := m;
  end:
end:
  \{Q: y > 0 \land y^2 \le x < (y+1)^2\}
```



Exercise 7.3:

```
y:=0; while (y+1)*(y+1) \leq x do y:=y+1; end;
```

Exercise 7.4:

```
\begin{array}{l} y := 0; \\ z := x + 1; \\ \textbf{while} \ y + 1 \neq z \ \textbf{do} \\ m := (y + z) \ \textbf{div} \ 2; \\ \textbf{if} \ m * m \leq x \ \textbf{then} \\ y := m; \\ \textbf{else} \\ z := m; \\ \textbf{end}; \\ \textbf{end}; \end{array}
```



```
Exercise 7.3:
```

```
\begin{array}{l} y:=0;\\ \text{while } (y+1)*(y+1) \leq x \text{ do}\\ y:=y+1;\\ \text{end}; \end{array}
```

Exercise 7.4:

```
\begin{array}{l} y := 0; \\ z := x + 1; \\ \text{while } y + 1 \neq z \text{ do} \\ m := (y + z) \text{ div } 2; \\ \text{if } m * m \leq x \text{ then} \\ y := m; \\ \text{else} \\ z := m; \\ \text{end}; \\ \text{end}; \end{array}
```

Take x = 1000. Compare the number of iterations: 7.3 31 times.



```
Exercise 7.3:
```

```
\begin{array}{l} y:=0;\\ \text{while } (y+1)*(y+1) \leq x \text{ do}\\ y:=y+1;\\ \text{end}; \end{array}
```

Exercise 7.4:

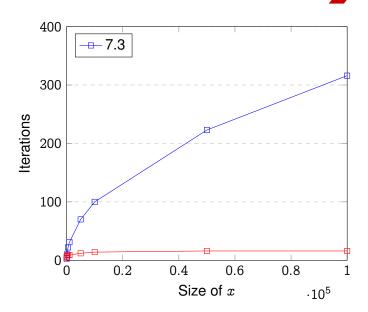
```
\begin{array}{l} y := 0; \\ z := x + 1; \\ \textbf{while} \ y + 1 \neq z \ \textbf{do} \\ m := (y + z) \ \textbf{div} \ 2; \\ \textbf{if} \ m * m \leq x \ \textbf{then} \\ y := m; \\ \textbf{else} \\ z := m; \\ \textbf{end}; \\ \textbf{end}; \end{array}
```

Take x = 1000. Compare the number of iterations:

7.3 31 times.

7.4 9 times - the size of
$$[y,z)$$
 decreases in each iteration: $[0,1001) \rightarrow [0,500) \rightarrow [0,250) \rightarrow [0,125) \rightarrow [0,62) \rightarrow [31,62) \rightarrow [31,46) \rightarrow [31,38) \rightarrow [31,34) \rightarrow [31,32).$





Outline



Exercise 7.1: Powers

Exercise 7.2: Factorial

Exercise 7.8: Dijkstra's FUSC

Summing an array

Square Roo

Square Root (Similar to Exercise 7.3) Exercise 7.4: Square Root, Revisited

Integral Division

Exercise 7.5: Integral Division

Exercise 7.6: Variation on Integral Division Efficiency Comparison Exercises 7.5 and 7.6

Exercise 7.5: Invariant



```
\begin{array}{l} \textbf{const} \ y: \ \mathbb{N}^+; \\ \textbf{var} \ x, \ q: \ \mathbb{Z}; \\ \{P: \ x=X \land X \geq 0\} \\ T \\ \{Q: \ X=q \cdot y + x \ \land \ 0 \leq x < y\} \end{array}
```

This way, e.g., given X = 11 and y = 4, we obtain q = 2 and x = 3.

Exercise 7.5: Invariant



```
\begin{array}{l} \textbf{const} \ y: \ \mathbb{N}^+; \\ \textbf{var} \ x, \ q: \ \mathbb{Z}; \\ \{P: \ x=X \land X \geq 0\} \\ T \\ \{Q: \ X=q \cdot y + x \ \land \ 0 \leq x < y\} \end{array}
```

This way, e.g., given X = 11 and y = 4, we obtain q = 2 and x = 3.

- 0 We need a **while**-program: we cannot use **div**, **mod**, and multiplication, so we repeatedly use addition and subtraction.
- 1 Choose an invariant J, and guard B such that $J \land \neg B \Rightarrow Q$. We use the heuristic split conjuncts:

$$egin{aligned} J:X=q\cdot y+x\ \land\ 0\leq x\ B:y\leq x \end{aligned}$$

Clearly,
$$J \wedge \neg B \equiv Q$$
.

Exercise 7.5: Initialization and Variant



$$egin{aligned} P:x=X\wedge X &\geq 0\ Q:X=q\cdot y+x\wedge 0 &\leq x < y\ J:X=q\cdot y+x\wedge 0 &\leq x\ B:y &\leq x \end{aligned}$$

2 Initialization: Find a command T_0 such that $\{P\}$ T_0 $\{J\}$.

```
\{P: x = X \land X \geq 0\}
(* logic; calculus *)
\{X = 0 \cdot y + x \land 0 \leq x\}
q:=0;
\{J: X = q \cdot y + x \land 0 \leq x\}
```

3 Variant function:

Exercise 7.5: Initialization and Variant



$$egin{aligned} P:x=X\wedge X &\geq 0\ Q:X=q\cdot y+x\wedge 0 &\leq x < y\ J:X=q\cdot y+x\wedge 0 &\leq x\ B:y &\leq x \end{aligned}$$

2 Initialization: Find a command T_0 such that $\{P\}$ T_0 $\{J\}$.

$$\{P: x = X \land X \ge 0\}$$

 $(* logic; calculus *)$
 $\{X = 0 \cdot y + x \land 0 \le x\}$
 $q:=0;$
 $\{J: X = q \cdot y + x \land 0 \le x\}$

3 Variant function:

We need to decrease x until x < y, so we choose $vf = x \in \mathbb{Z}$. Since J contains $0 \le x$, it is trivial that $J \wedge B \Rightarrow vf \ge 0$.



$$\{J \wedge B \wedge vf = V\}$$

$$\{J \wedge vf < V\}$$



$$\left\{ \begin{matrix} J \wedge B \ \wedge \ vf = V \\ X = q \cdot y + x \ \wedge \ 0 \leq x \ \wedge \ y \leq x \ \wedge \ x = V \\ \end{matrix} \right\}$$

$$egin{aligned} x := x - y; \ & \ q := q + 1; \ & \ \{J \wedge vf < V\} \end{aligned}$$



```
egin{aligned} \{J \wedge B \ \wedge \ vf = V\} \ \{X = q \cdot y + x \ \wedge \ 0 \leq x \ \wedge \ y \leq x \ \wedge \ x = V\} \ (*\ y > 0; \textit{prepare}\ x := x - y; \textit{calculus}; \textit{logic}\ *) \end{aligned} x := x - y; q := q + 1; \{J \wedge vf < V\}
```



```
egin{aligned} \{J \wedge B \ \wedge \ vf = V\} \ \{X = q \cdot y + x \ \wedge \ 0 \leq x \ \wedge \ y \leq x \ \wedge \ x = V\} \ (*\ y > 0; \textit{prepare}\ x := x - y; \textit{calculus; logic}\ *) \ \{X = (q+1) \cdot y + x - y \ \wedge \ 0 \leq x - y \ \wedge \ x - y < V\} \ x := x - y; \ q := q+1; \ \{J \wedge vf < V\} \end{aligned}
```



```
egin{aligned} \{J \wedge B \ \wedge \ vf = V\} \ \{X = q \cdot y + x \ \wedge \ 0 \leq x \ \wedge \ y \leq x \ \wedge \ x = V\} \ (*\ y > 0; \textit{prepare}\ x := x - y; \textit{calculus; logic}\ *) \ \{X = (q+1) \cdot y + x - y \ \wedge \ 0 \leq x - y \ \wedge \ x - y < V\} \ x := x - y; \ q := q+1; \ \{J \wedge vf < V\} \end{aligned}
```



```
 \begin{cases} J \wedge B \ \wedge \ vf = V \} \\ \{X = q \cdot y + x \ \wedge \ 0 \leq x \ \wedge \ y \leq x \ \wedge \ x = V \} \\ \text{(* } y > 0 \text{; prepare } x := x - y \text{; calculus; logic *)} \\ \{X = (q+1) \cdot y + x - y \ \wedge \ 0 \leq x - y \ \wedge \ x - y < V \} \\ x := x - y; \\ \{X = (q+1) \cdot y + x \wedge 0 \leq x \wedge x < V \} \\ q := q+1; \\ \{J \wedge vf < V \}
```



```
 \begin{cases} J \wedge B \ \wedge \ vf = V \} \\ \{X = q \cdot y + x \ \wedge \ 0 \leq x \ \wedge \ y \leq x \ \wedge \ x = V \} \\ \text{(* } y > 0 \text{; prepare } x := x - y \text{; calculus; logic *)} \\ \{X = (q+1) \cdot y + x - y \ \wedge \ 0 \leq x - y \ \wedge \ x - y < V \} \\ x := x - y; \\ \{X = (q+1) \cdot y + x \wedge 0 \leq x \wedge x < V \} \\ q := q+1; \\ \{J \wedge vf < V \}
```



```
 \begin{cases} J \wedge B \ \wedge \ vf = V \} \\ \{X = q \cdot y + x \ \wedge \ 0 \leq x \ \wedge \ y \leq x \ \wedge \ x = V \} \\ \text{(* } y > 0 \text{; prepare } x := x - y \text{; calculus; logic *)} \\ \{X = (q+1) \cdot y + x - y \ \wedge \ 0 \leq x - y \ \wedge \ x - y < V \} \\ x := x - y; \\ \{X = (q+1) \cdot y + x \wedge 0 \leq x \wedge x < V \} \\ q := q+1; \\ \{X = q \cdot y + x \wedge 0 \leq x \wedge x < V \} \\ \{J \wedge vf < V \}
```

Exercise 7.5: Conclusion



5 Conclusion: We derived $\{P\}$ T_0 ; while B do S end $\{Q\}$

```
 \begin{cases} P: \ x = X \land X \geq 0 \rbrace \\ q := 0; \\ \{J: X = q \cdot y + x \land 0 \leq x \} \\ \quad (^* \ v\! f = x \ ^*) \end{cases}  while y \leq x do  x := x - y; \\ q := q + 1;  end;  \{Q: \ X = q \cdot y + x \land 0 \leq x < y \}
```

Exercise 7.6: Same Spec, New Invariant



```
\begin{array}{l} \textbf{const} \ y: \ \mathbb{N}^+; \\ \textbf{var} \ x, \ z, \ q, \ i: \ \mathbb{Z}; \\ \{P: \ x=X \land X \geq 0\} \\ T \\ \{Q: \ X=q\cdot y+x \land 0 \leq x < y\} \end{array}
```

Same exercise as 7.5, now with more hints.

Exercise 7.6: Same Spec, New Invariant



```
\begin{array}{l} \textbf{const} \ y: \ \mathbb{N}^+; \\ \textbf{var} \ x, \ z, \ q, \ i: \ \mathbb{Z}; \\ \{P: \ x=X \land X \geq 0\} \\ T \\ \{Q: \ X=q\cdot y+x \land 0 \leq x < y\} \end{array}
```

Same exercise as 7.5, now with more hints.

- ▶ We can use multiplication and division by 2.
- ightharpoonup Before in each iteration we had 'x := x y'. The invariant was

$$X = q \cdot y + x \wedge 0 \leq x$$

Exercise 7.6: Same Spec, New Invariant



```
\begin{array}{l} \textbf{const} \ y: \ \mathbb{N}^+; \\ \textbf{var} \ x, \ z, \ q, \ i: \ \mathbb{Z}; \\ \{P: \ x = X \wedge X \geq 0\} \\ T \\ \{Q: \ X = q \cdot y + x \wedge 0 \leq x < y\} \end{array}
```

Same exercise as 7.5, now with more hints.

- ▶ We can use multiplication and division by 2.
- ightharpoonup Before in each iteration we had 'x := x y'. The invariant was

$$X = q \cdot y + x \wedge 0 \leq x$$

Now the invariant involves a new variable z:

$$X = q \cdot z + x \wedge 0 \leq x < z \wedge z = 2^i \cdot y \wedge i \geq 0$$

What is the role of z?

Exercise 7.6: Initialization



$$P: x = X \wedge X \geq 0$$

$$J: X = q \cdot z + x \wedge 0 \leq x < z \wedge z = 2^i \cdot y \wedge i \geq 0$$

1 Choose a guard B such that $J \wedge \neg B \Rightarrow Q$. We choose $B: z \neq y$.

Exercise 7.6: Initialization



$$P: x = X \wedge X \geq 0$$
 $J: X = q \cdot z + x \wedge 0 \leq x < z \wedge z = 2^i \cdot y \wedge i \geq 0$

- 1 Choose a guard B such that $J \wedge \neg B \Rightarrow Q$. We choose $B: z \neq y$.
- 2 Initialization: We find a command T_0 such that $\{P\}$ T_0 $\{J\}$. We can easily initialize the first conjunct of J with q := 0.

```
egin{aligned} \{P: x = X \wedge X \geq 0\} \ T_1 \ \{P_0: x = X \wedge 0 \leq x < z \wedge i \geq 0 \wedge z = 2^i \cdot y\} \ (* \ calculus \ *) \ \{X = 0 \cdot z + x \wedge 0 \leq x < z \wedge i \geq 0 \wedge z = 2^i \cdot y\} \ q := 0; \ \{J: \ X = q \cdot z + x \wedge 0 \leq x < z \wedge i \geq 0 \wedge z = 2^i \cdot y\} \end{aligned}
```

But deriving T_1 requires more work: an auxiliary loop.

Exercise 7.6: Auxiliary Loop



We derive

$$egin{aligned} \{P: x = X \wedge X \geq 0\} \ T_1 \ \{P_0: x = X \wedge 0 \leq x < z \wedge i \geq 0 \wedge z = 2^i \cdot y\} \end{aligned}$$

For the invariant and guard we choose:

$$J_0: 0 \leq x = X \wedge i \geq 0 \wedge z = 2^i \cdot y \ B_0: z \leq x$$

Clearly $J_0 \wedge \neg B_0 \equiv P_0$. J_0 is easy to initialize (without proof): $z := y; \ i := 0;$ We choose the variant function $vf_0 = x - z \in \mathbb{Z}$. Since $B_0 \equiv vf_0 \geq 0$, it is trivial that $J_0 \wedge B_0 \Rightarrow vf_0 \geq 0$



$$\{J_0 \wedge B_0 \wedge vf_0 = V\}$$

$$\{J_0 \wedge vf_0 < V\}$$



$$\{J_0 \wedge B_0 \wedge v f_0 = V\}$$

 $\{0 \leq x = X \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge z \leq x \wedge x - z = V\}$

$$\{J_0 \wedge vf_0 < V\}$$



$$egin{aligned} \{J_0 \wedge B_0 \wedge v f_0 &= V \} \ \{0 \leq x = X \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge z \leq x \wedge x - z = V \} \ \textbf{(* } y > 0 \Rightarrow z = 2^i \cdot y > 0 ext{; prepare } z := 2 * z * \textbf{)} \ \{0 \leq x = X \wedge i + 1 \geq 0 \wedge 2 \cdot z = 2^{i+1} \cdot y \wedge x - 2 \cdot z < V \} \end{aligned}$$

$$\{J_0 \wedge vf_0 < V\}$$



```
egin{aligned} \{J_0 \wedge B_0 \wedge v f_0 &= V\} \ \{0 \leq x = X \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge z \leq x \wedge x - z = V\} \ (*\ y > 0 \Rightarrow z = 2^i \cdot y > 0 	ext{; prepare } z := 2*z*) \ \{0 \leq x = X \wedge i + 1 \geq 0 \wedge 2 \cdot z = 2^{i+1} \cdot y \wedge x - 2 \cdot z < V\} \ z := 2*z; \end{aligned}
```

$$\{J_0 \wedge vf_0 < V\}$$



```
egin{aligned} \{J_0 \wedge B_0 \wedge v f_0 &= V \} \ \{0 \leq x = X \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge z \leq x \wedge x - z = V \} \ (*\ y > 0 \Rightarrow z = 2^i \cdot y > 0; \textit{prepare}\ z := 2 * z *) \ \{0 \leq x = X \wedge i + 1 \geq 0 \wedge 2 \cdot z = 2^{i+1} \cdot y \wedge x - 2 \cdot z < V \} \ z := 2 * z; \ \{0 \leq x = X \wedge i + 1 \geq 0 \wedge z = 2^{i+1} \cdot y \wedge x - z < V \} \ \{J_0 \wedge v f_0 < V \} \end{aligned}
```



```
egin{aligned} \{J_0 \wedge B_0 \wedge v f_0 = V\} \ \{0 \leq x = X \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge z \leq x \wedge x - z = V\} \ (*y > 0 \Rightarrow z = 2^i \cdot y > 0; \textit{prepare } z := 2 * z *) \ \{0 \leq x = X \wedge i + 1 \geq 0 \wedge 2 \cdot z = 2^{i+1} \cdot y \wedge x - 2 \cdot z < V\} \ z := 2 * z; \ \{0 \leq x = X \wedge i + 1 \geq 0 \wedge z = 2^{i+1} \cdot y \wedge x - z < V\} \ i := i + 1; \ \{J_0 \wedge v f_0 < V\} \end{aligned}
```



```
egin{aligned} \{J_0 \wedge B_0 \wedge v f_0 &= V\} \ \{0 \leq x = X \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge z \leq x \wedge x - z = V\} \ (*\ y > 0 \Rightarrow z = 2^i \cdot y > 0; \textit{prepare } z := 2*z^*) \ \{0 \leq x = X \wedge i + 1 \geq 0 \wedge 2 \cdot z = 2^{i+1} \cdot y \wedge x - 2 \cdot z < V\} \ z := 2*z; \ \{0 \leq x = X \wedge i + 1 \geq 0 \wedge z = 2^{i+1} \cdot y \wedge x - z < V\} \ i := i+1; \ \{0 \leq x = X \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge z \leq x \wedge x - z < V\} \ \{J_0 \wedge v f_0 < V\} \end{aligned}
```

Exercise 7.6: Summing Up Initialization



We derived the following auxiliary loop for initialization:

```
egin{aligned} z &:= y; \ i &:= 0; \ \mathbf{while} \ z &\leq x \ \mathbf{do} \ z &:= 2 * z; \ i &:= i + 1; \ \mathbf{end}; \ q &:= 0; \end{aligned}
```

We may now return to design the main loop.

Exercise 7.6: Variant



Recall:

$$J: X = q \cdot z + x \wedge 0 \leq x < z \wedge i \geq 0 \wedge z = 2^i \cdot y$$

 $B: z \neq y$

3 Variant function:

We choose $\mathit{vf} = i \in \mathbb{Z}$.

Clearly, $J \wedge B \Rightarrow vf \geq 0$, since $i \geq 0$ is a conjunct of J.



$$\{J \wedge B \wedge vf = V\}$$

$$\{J \land vf < V\}$$



$$\begin{cases} J \wedge B \wedge vf = V \\ X = q \cdot z + x \wedge 0 \le x < z \wedge i \ge 0 \wedge z = 2^i \cdot y \wedge z \ne y \wedge i = V \end{cases}$$

$$\{J \land vf < V\}$$



```
 \begin{array}{l} \{J \wedge B \ \wedge \ vf = V\} \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < z \ \wedge \ i \geq 0 \ \wedge \ z = 2^i \cdot y \ \wedge \ z \neq y \ \wedge \ i = V\} \\ \text{(* prepare } i := i-1; z \neq y \ \wedge \ z = 2^i \cdot y \ \wedge \ i \geq 0 \Rightarrow i > 0 \ ^*)} \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < z \ \wedge \ i - 1 \geq 0 \ \wedge \ z = 2 \cdot 2^{i-1} \cdot y \ \wedge \ i - 1 < V\} \end{array}
```





```
 \begin{cases} J \land B \land vf = V \} \\ \{X = q \cdot z + x \land 0 \leq x < z \land i \geq 0 \land z = 2^i \cdot y \land z \neq y \land i = V \} \\ \text{(* prepare } i := i - 1; z \neq y \land z = 2^i \cdot y \land i \geq 0 \Rightarrow i > 0 \text{ *}) \\ \{X = q \cdot z + x \land 0 \leq x < z \land i - 1 \geq 0 \land z = 2 \cdot 2^{i - 1} \cdot y \land i - 1 < V \} \\ i := i - 1; \\ \{X = q \cdot z + x \land 0 \leq x < z \land i \geq 0 \land z = 2 \cdot 2^i \cdot y \land i < V \}
```



```
 \begin{cases} J \wedge B \ \wedge \ vf = V \} \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < z \ \wedge \ i \geq 0 \ \wedge \ z = 2^i \cdot y \ \wedge \ z \neq y \ \wedge \ i = V \} \\ \text{$(*\ prepare}\ i := i-1; z \neq y \ \wedge \ z = 2^i \cdot y \ \wedge \ i \geq 0 \Rightarrow i > 0\ ^*)$} \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < z \ \wedge \ i - 1 \geq 0 \ \wedge \ z = 2 \cdot 2^{i-1} \cdot y \ \wedge \ i - 1 < V \} \\ i := i-1; \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < z \ \wedge \ i \geq 0 \ \wedge \ z = 2 \cdot 2^i \cdot y \ \wedge \ i < V \} \\ \text{$(*\ z \ is \ even; } z = 2(z \ \text{div}\ 2); \ calculus\ ^*)$}
```



```
 \begin{cases} J \wedge B \ \wedge \ vf = V \} \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < z \ \wedge \ i \geq 0 \ \wedge \ z = 2^i \cdot y \ \wedge \ z \neq y \ \wedge \ i = V \} \\ \text{$(*\ prepare}\ i := i-1; z \neq y \ \wedge \ z = 2^i \cdot y \ \wedge \ i \geq 0 \Rightarrow i > 0\ ^*)$} \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < z \ \wedge \ i - 1 \geq 0 \ \wedge \ z = 2 \cdot 2^{i-1} \cdot y \ \wedge \ i - 1 < V \} \\ i := i-1; \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < z \ \wedge \ i \geq 0 \ \wedge \ z = 2 \cdot 2^i \cdot y \ \wedge \ i < V \} \\ \text{$(*\ z \ is \ even; \ z = 2(z \ \text{div}\ 2); \ calculus}\ ^*)$} \\ \{X = 2 \cdot q \cdot (z \ \text{div}\ 2) + x \wedge 0 \leq x < 2(z \ \text{div}\ 2) \wedge i \geq 0 \wedge z \ \text{div}\ 2 = 2^i \cdot y \wedge i < V \} \end{cases}
```



```
 \begin{cases} J \wedge B \wedge vf = V \} \\ \{X = q \cdot z + x \wedge 0 \leq x < z \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge z \neq y \wedge i = V \} \\ \text{$(*$ prepare $i := i - 1$; $z \neq y \wedge z = 2^i \cdot y \wedge i \geq 0 \Rightarrow i > 0$ *)$} \\ \{X = q \cdot z + x \wedge 0 \leq x < z \wedge i - 1 \geq 0 \wedge z = 2 \cdot 2^{i-1} \cdot y \wedge i - 1 < V \} \\ i := i - 1; \\ \{X = q \cdot z + x \wedge 0 \leq x < z \wedge i \geq 0 \wedge z = 2 \cdot 2^i \cdot y \wedge i < V \} \\ \text{$(*$ z is even; $z = 2(z \text{ div } 2)$; calculus *)$} \\ \{X = 2 \cdot q \cdot (z \text{ div } 2) + x \wedge 0 \leq x < 2(z \text{ div } 2) \wedge i \geq 0 \wedge z \text{ div } 2 = 2^i \cdot y \wedge i < V \} \\ z := z \text{ div } 2; \end{cases}
```



```
 \begin{cases} J \wedge B \wedge vf = V \} \\ \{X = q \cdot z + x \wedge 0 \leq x < z \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge z \neq y \wedge i = V \} \\ \text{$(^*$ prepare $i := i - 1$; $z \neq y \wedge z = 2^i \cdot y \wedge i \geq 0 \Rightarrow i > 0$^*$)} \\ \{X = q \cdot z + x \wedge 0 \leq x < z \wedge i - 1 \geq 0 \wedge z = 2 \cdot 2^{i - 1} \cdot y \wedge i - 1 < V \} \\ i := i - 1; \\ \{X = q \cdot z + x \wedge 0 \leq x < z \wedge i \geq 0 \wedge z = 2 \cdot 2^i \cdot y \wedge i < V \} \\ \text{$(^*$ z is even; $z = 2(z \text{ div } 2)$; calculus$^*$)} \\ \{X = 2 \cdot q \cdot (z \text{ div } 2) + x \wedge 0 \leq x < 2(z \text{ div } 2) \wedge i \geq 0 \wedge z \text{ div } 2 = 2^i \cdot y \wedge i < V \} \\ z := z \text{ div } 2; \\ \{X = 2 \cdot q \cdot z + x \wedge 0 \leq x < 2 \cdot z \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge i < V \}
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```
 \begin{cases} J \wedge B \wedge vf = V \} \\ \{X = q \cdot z + x \wedge 0 \leq x < z \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge z \neq y \wedge i = V \} \\ \text{$(*$ prepare $i := i-1$; $z \neq y \wedge z = 2^i \cdot y \wedge i \geq 0 \Rightarrow i > 0^*$)} \\ \{X = q \cdot z + x \wedge 0 \leq x < z \wedge i - 1 \geq 0 \wedge z = 2 \cdot 2^{i-1} \cdot y \wedge i - 1 < V \} \\ i := i-1$; \\ \{X = q \cdot z + x \wedge 0 \leq x < z \wedge i \geq 0 \wedge z = 2 \cdot 2^i \cdot y \wedge i < V \} \\ \text{$(*z$ is even; $z = 2(z$ div 2); calculus*)$} \\ \{X = 2 \cdot q \cdot (z \text{ div } 2) + x \wedge 0 \leq x < 2(z \text{ div } 2) \wedge i \geq 0 \wedge z \text{ div } 2 = 2^i \cdot y \wedge i < V \} \\ z := z \text{ div } 2\text{;} \\ \{X = 2 \cdot q \cdot z + x \wedge 0 \leq x < 2 \cdot z \wedge i \geq 0 \wedge z = 2^i \cdot y \wedge i < V \} \\ q := 2 * q; \end{cases}
```



```
 \begin{cases} J \wedge B \ \wedge \ vf = V \} \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < z \ \wedge \ i \geq 0 \ \wedge \ z = 2^i \cdot y \ \wedge \ z \neq y \ \wedge \ i = V \} \\ \text{(* prepare } i := i - 1; z \neq y \ \wedge \ z = 2^i \cdot y \ \wedge \ i \geq 0 \Rightarrow i > 0 \ ^*) \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < z \ \wedge \ i - 1 \geq 0 \ \wedge \ z = 2 \cdot 2^{i - 1} \cdot y \ \wedge \ i - 1 < V \} \\ i := i - 1; \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < z \ \wedge \ i \geq 0 \ \wedge \ z = 2 \cdot 2^i \cdot y \ \wedge \ i < V \} \\ \text{(* } z \text{ is even; } z = 2(z \text{ div } 2); \text{ calculus *}) \\ \{X = 2 \cdot q \cdot (z \text{ div } 2) + x \wedge 0 \leq x < 2(z \text{ div } 2) \wedge i \geq 0 \wedge z \text{ div } 2 = 2^i \cdot y \wedge i < V \} \\ z := z \text{ div } 2; \\ \{X = 2 \cdot q \cdot z + x \ \wedge \ 0 \leq x < 2 \cdot z \ \wedge \ i \geq 0 \ \wedge \ z = 2^i \cdot y \ \wedge \ i < V \} \\ q := 2 * q; \\ \{X = q \cdot z + x \ \wedge \ 0 \leq x < 2 \cdot z \ \wedge \ i \geq 0 \ \wedge \ z = 2^i \cdot y \ \wedge \ i < V \} \end{cases}
```



```
\{J \wedge B \wedge vf = V\}
   \{X = q \cdot z + x \land 0 \le x < z \land i \ge 0 \land z = 2^i \cdot y \land z \ne y \land i = V\}
     (* prepare i := i - 1; z \neq y \land z = 2^i \cdot y \land i > 0 \Rightarrow i > 0*)
   \{X = q \cdot z + x \land 0 < x < z \land i - 1 > 0 \land z = 2 \cdot 2^{i-1} \cdot y \land i - 1 < V\}
i := i - 1:
   {X = q \cdot z + x \land 0 < x < z \land i > 0 \land z = 2 \cdot 2^i \cdot y \land i < V}
     (* z is even; z = 2(\overline{z} \operatorname{div} 2); calculus *)
   \{X = 2 \cdot q \cdot (z \text{ div } 2) + x \land 0 \le x < 2(z \text{ div } 2) \land i \ge 0 \land z \text{ div } 2 = 2^i \cdot y \land i < V\}
z := z \operatorname{div} 2
   \{X = 2 \cdot q \cdot z + x \land 0 < x < 2 \cdot z \land i > 0 \land z = 2^i \cdot y \land i < V\}
q := 2 * q:
   {X = q \cdot z + x \land 0 < x < 2 \cdot z \land i > 0 \land z = 2^{i} \cdot y \land i < V}
if x < z then
      \{X = q \cdot z + x \land 0 \le x < 2 \cdot z \land x < z \land i > 0 \land z = 2^i \cdot y \land i < V\}
```

else

$$\{X = q \cdot z + x \land z < x < 2 \cdot z \land i > 0 \land z = 2^i \cdot y \land i < V\}$$

$$\{J \land vf < V\}$$



```
\{J \wedge B \wedge vf = V\}
   \{X = q \cdot z + x \land 0 \le x < z \land i > 0 \land z = 2^i \cdot y \land z \ne y \land i = V\}
      (* prepare i := i - 1; z \neq y \land z = 2^i \cdot y \land i > 0 \Rightarrow i > 0*)
   \{X = q \cdot z + x \land 0 \le x \le z \land i - 1 \ge 0 \land z = 2 \cdot 2^{i-1} \cdot y \land i - 1 \le V\}
i := i - 1:
   {X = q \cdot z + x \land 0 < x < z \land i > 0 \land z = 2 \cdot 2^{i} \cdot y \land i < V}
    (* z is even; z = 2(z \operatorname{div} 2); calculus *)
   \{X = 2 \cdot q \cdot (z \text{ div } 2) + x \land 0 \le x < 2(z \text{ div } 2) \land i \ge 0 \land z \text{ div } 2 = 2^i \cdot y \land i < V\}
z:=z \operatorname{div} 2:
   \{X = 2 \cdot q \cdot z + x \land 0 < x < 2 \cdot z \land i > 0 \land z = 2^{i} \cdot y \land i < V\}
q := 2 * q:
   \{X = q \cdot z + x \land 0 \le x < 2 \cdot z \land i \ge 0 \land z = 2^i \cdot y \land i < V\}
if x < z then
     \{X = q \cdot z + x \land 0 < x < 2 \cdot z \land x < z \land i > 0 \land z = 2^i \cdot y \land i < V\}
   skip:
else
      \{X = q \cdot z + x \land z < x < 2 \cdot z \land i > 0 \land z = 2^i \cdot y \land i < V\}
```

$$\{J \land vf < V\}$$



```
\{J \wedge B \wedge vf = V\}
   \{X = q \cdot z + x \land 0 \le x < z \land i > 0 \land z = 2^i \cdot y \land z \ne y \land i = V\}
      (* prepare i := i - 1; z \neq y \land z = 2^i \cdot y \land i > 0 \Rightarrow i > 0*)
  \{X = q \cdot z + x \land 0 < x < z \land i - 1 > 0 \land z = 2 \cdot 2^{i-1} \cdot y \land i - 1 < V\}
i := i - 1:
  \{X = q \cdot z + x \land 0 < x < z \land i > 0 \land z = 2 \cdot 2^i \cdot y \land i < V\}
    (* z is even; z = 2(\overline{z} \operatorname{div} 2); calculus *)
   \{X=2\cdot q\cdot (z\ \mathsf{div}\ 2)+x\wedge 0\leq x< 2(z\ \mathsf{div}\ 2)\wedge i\geq 0\wedge z\ \mathsf{div}\ 2=2^i\cdot y\wedge i< V\}
z:=z \operatorname{div} 2:
  {X = 2 \cdot q \cdot z + x \land 0 < x < 2 \cdot z \land i > 0 \land z = 2^{i} \cdot y \land i < V}
q := 2 * q:
  \{X = q \cdot z + x \land 0 \le x < 2 \cdot z \land i \ge 0 \land z = 2^i \cdot y \land i < V\}
if x < z then
      \{X = q \cdot z + x \land 0 < x < 2 \cdot z \land x < z \land i > 0 \land z = 2^i \cdot y \land i < V\}
  skip;
      \{J \land vf < V\}
else
      \{X = q \cdot z + x \land z < x < 2 \cdot z \land i > 0 \land z = 2^i \cdot y \land i < V\}
```

$$\{J \land vf < V\}$$



```
\{J \wedge B \wedge vf = V\}
   \{X = q \cdot z + x \land 0 \le x < z \land i > 0 \land z = 2^i \cdot y \land z \ne y \land i = V\}
      (* prepare i := i - 1; z \neq y \land z = 2^i \cdot y \land i > 0 \Rightarrow i > 0*)
   \{X = q \cdot z + x \land 0 \le x \le z \land i - 1 \ge 0 \land z = 2 \cdot 2^{i-1} \cdot y \land i - 1 \le V\}
i := i - 1:
   {X = q \cdot z + x \land 0 < x < z \land i > 0 \land z = 2 \cdot 2^{i} \cdot y \land i < V}
     (* z is even; z = 2(\overline{z} \operatorname{div} 2); calculus *)
   \{X = 2 \cdot q \cdot (z \text{ div } 2) + x \wedge 0 \leq x < 2(z \text{ div } 2) \wedge i \geq 0 \wedge z \text{ div } 2 = 2^i \cdot y \wedge i < V\}
z := z \operatorname{div} 2:
   \{X = 2 \cdot q \cdot z + x \land 0 < x < 2 \cdot z \land i > 0 \land z = 2^i \cdot y \land i < V\}
q := 2 * q;
  \{X = q \cdot z + x \land 0 \le x < 2 \cdot z \land i \ge 0 \land z = 2^i \cdot y \land i < V\}
if x < z then
      \{X = q \cdot z + x \ \land \ 0 \le x < 2 \cdot z \ \land \ x < z \ \land \ i \ge 0 \ \land \ z = 2^i \cdot y \ \land \ i < V\}
   skip;
      \{J \land vf < V\}
else
      \{X = q \cdot z + x \land z < x < 2 \cdot z \land i > 0 \land z = 2^{i} \cdot y \land i < V\}
      \{X = (q+1) \cdot z + x - z \land 0 < x - z < z \land i > 0 \land z = 2^i \cdot y \land i < V\}
```

```
end \{J \land vf < V\}
```



```
\{J \wedge B \wedge vf = V\}
   \{X = q \cdot z + x \land 0 \le x < z \land i > 0 \land z = 2^i \cdot y \land z \ne y \land i = V\}
      (* prepare i := i - 1; z \neq y \land z = 2^i \cdot y \land i > 0 \Rightarrow i > 0*)
   \{X = q \cdot z + x \land 0 \le x \le z \land i - 1 \ge 0 \land z = 2 \cdot 2^{i-1} \cdot y \land i - 1 \le V\}
i := i - 1:
   \{X = q \cdot z + x \land 0 < x < z \land i > 0 \land z = 2 \cdot 2^i \cdot y \land i < V\}
     (* z is even; z = 2(\overline{z} \operatorname{div} 2); calculus *)
   \{X = 2 \cdot q \cdot (z \text{ div } 2) + x \wedge 0 \leq x < 2(z \text{ div } 2) \wedge i \geq 0 \wedge z \text{ div } 2 = 2^i \cdot y \wedge i < V\}
z := z \operatorname{div} 2:
   \{X = 2 \cdot q \cdot z + x \land 0 < x < 2 \cdot z \land i > 0 \land z = 2^i \cdot y \land i < V\}
q := 2 * q;
  \{X = q \cdot z + x \land 0 \le x < 2 \cdot z \land i > 0 \land z = 2^i \cdot y \land i < V\}
if x < z then
      \{X = q \cdot z + x \ \land \ 0 \le x < 2 \cdot z \ \land \ x < z \ \land \ i \ge 0 \ \land \ z = 2^i \cdot y \ \land \ i < V\}
   skip;
      \{J \land vf < V\}
else
      \{X = q \cdot z + x \land z < x < 2 \cdot z \land i > 0 \land z = 2^{i} \cdot y \land i < V\}
      \{X = (q+1) \cdot z + x - z \land 0 \le x - z < z \land i \ge 0 \land z = 2^i \cdot y \land i < V\}
   a := a + 1:
```

end $\{J \land vf < V\}$



```
\{J \wedge B \wedge vf = V\}
   \{X = q \cdot z + x \land 0 \le x \le z \land i \ge 0 \land z = 2^i \cdot y \land z \ne y \land i = V\}
      (* prepare i := i - 1; z \neq y \land z = 2^i \cdot y \land i > 0 \Rightarrow i > 0*)
   \{X = q \cdot z + x \ \land \ 0 \le x < z \ \land \ i - 1 \ge 0 \ \land \ z = 2 \cdot 2^{i - 1} \cdot y \ \land \ i - 1 < V\}
i := i - 1:
   {X = q \cdot z + x \land 0 < x < z \land i > 0 \land z = 2 \cdot 2^{i} \cdot y \land i < V}
    (* z is even; z = 2(\overline{z} \operatorname{div} 2); calculus *)
   \{X = 2 \cdot q \cdot (z \text{ div } 2) + x \wedge 0 \leq x < 2(z \text{ div } 2) \wedge i \geq 0 \wedge z \text{ div } 2 = 2^i \cdot y \wedge i < V\}
z := z \operatorname{div} 2:
   \{X = 2 \cdot q \cdot z + x \land 0 < x < 2 \cdot z \land i > 0 \land z = 2^i \cdot y \land i < V\}
q := 2 * q;
  \{X = q \cdot z + x \land 0 \le x < 2 \cdot z \land i \ge 0 \land z = 2^i \cdot y \land i < V\}
if x < z then
      \{X = q \cdot z + x \land 0 < x < 2 \cdot z \land x < z \land i > 0 \land z = 2^i \cdot y \land i < V\}
   skip;
      \{J \land vf < V\}
else
      \{X = q \cdot z + x \land z < x < 2 \cdot z \land i > 0 \land z = 2^{i} \cdot y \land i < V\}
      \{X = (q+1) \cdot z + x - z \land 0 \le x - z \le z \land i \ge 0 \land z = 2^i \cdot y \land i \le V\}
   q := q + 1;
      \{X = q \cdot z + x - z \land 0 \le x - z < z \land i \ge 0 \land z = 2^i \cdot y \land i < V\}
```

$$\{J \land vf < V\}$$



```
\{J \wedge B \wedge vf = V\}
   \{X = q \cdot z + x \land 0 \le x \le z \land i \ge 0 \land z = 2^i \cdot y \land z \ne y \land i = V\}
      (* prepare i := i - 1; z \neq y \land z = 2^i \cdot y \land i > 0 \Rightarrow i > 0*)
   \{X = q \cdot z + x \ \land \ 0 \le x < z \ \land \ i - 1 \ge 0 \ \land \ z = 2 \cdot 2^{i - 1} \cdot y \ \land \ i - 1 < V\}
i := i - 1:
   {X = q \cdot z + x \land 0 < x < z \land i > 0 \land z = 2 \cdot 2^{i} \cdot y \land i < V}
    (* z is even; z = 2(\overline{z} \operatorname{div} 2); calculus *)
   \{X = 2 \cdot q \cdot (z \text{ div } 2) + x \land 0 \le x \le 2(z \text{ div } 2) \land i \ge 0 \land z \text{ div } 2 = 2^i \cdot y \land i \le V\}
z := z \operatorname{div} 2:
   \{X = 2 \cdot q \cdot z + x \land 0 < x < 2 \cdot z \land i > 0 \land z = 2^i \cdot y \land i < V\}
q := 2 * q;
  \{X = q \cdot z + x \land 0 \le x < 2 \cdot z \land i \ge 0 \land z = 2^i \cdot y \land i < V\}
if x < z then
      \{X = q \cdot z + x \land 0 < x < 2 \cdot z \land x < z \land i > 0 \land z = 2^i \cdot y \land i < V\}
   skip;
      \{J \land vf < V\}
else
      \{X = q \cdot z + x \land z < x < 2 \cdot z \land i > 0 \land z = 2^{i} \cdot y \land i < V\}
      \{X = (q+1) \cdot z + x - z \land 0 \le x - z \le z \land i \ge 0 \land z = 2^i \cdot y \land i \le V\}
   q := q + 1;
      \{X = q \cdot z + x - z \land 0 < x - z < z \land i > 0 \land z = 2^i \cdot y \land i < V\}
   x := x - z:
```

$$\{J \land vf < V\}$$

 $\{J \land vf < V\}$



```
\{J \wedge B \wedge vf = V\}
   \{X = q \cdot z + x \land 0 \le x \le z \land i \ge 0 \land z = 2^i \cdot y \land z \ne y \land i = V\}
      (* prepare i := i - 1; z \neq y \land z = 2^i \cdot y \land i > 0 \Rightarrow i > 0*)
  \{X = q \cdot z + x \ \land \ 0 \le x < z \ \land \ i - 1 \ge 0 \ \land \ z = 2 \cdot 2^{i - 1} \cdot y \ \land \ i - 1 < V\}
i := i - 1:
  {X = q \cdot z + x \land 0 < x < z \land i > 0 \land z = 2 \cdot 2^{i} \cdot y \land i < V}
    (* z is even; z = 2(\overline{z} \operatorname{div} 2); calculus *)
  \{X = 2 \cdot q \cdot (z \text{ div } 2) + x \land 0 \le x \le 2(z \text{ div } 2) \land i \ge 0 \land z \text{ div } 2 = 2^i \cdot y \land i \le V\}
z := z \operatorname{div} 2:
  \{X = 2 \cdot q \cdot z + x \land 0 < x < 2 \cdot z \land i > 0 \land z = 2^i \cdot y \land i < V\}
q := 2 * q;
  \{X = q \cdot z + x \land 0 \le x < 2 \cdot z \land i \ge 0 \land z = 2^i \cdot y \land i < V\}
if x < z then
      \{X = q \cdot z + x \land 0 < x < 2 \cdot z \land x < z \land i > 0 \land z = 2^i \cdot y \land i < V\}
  skip;
      \{J \land vf < V\}
else
      \{X = q \cdot z + x \land z < x < 2 \cdot z \land i > 0 \land z = 2^{i} \cdot y \land i < V\}
      \{X = (q+1) \cdot z + x - z \land 0 \le x - z \le z \land i \ge 0 \land z = 2^i \cdot y \land i \le V\}
   q := q + 1;
      \{X = q \cdot z + x - z \land 0 \le x - z < z \land i \ge 0 \land z = 2^i \cdot y \land i < V\}
   x := x - z:
      \{J \land vf < V\}
end (* collect branches *)
```

Exercise 7.6: Conclusion



5 Conclusion: We derived the following program fragment.

```
\{P: x = X \wedge X \geq 0\}
z := y:
i := 0:
while z < x do
  z := 2 * z;
  i := i + 1:
end:
a := 0;
while z \neq y do
  i := i - 1:
  z:=z \operatorname{div} 2;
  q := 2 * q;
  if x > z then
     q := q + 1;
     x := x - z
  end:
end:
  \{Q: X = q \cdot y + x \land 0 \le x < y\}
```

Above, i is a ghost variable: we only use it to make the proof easier.

Exercise 7.6: Conclusion



5 Conclusion: We derived the following program fragment.

```
\{P: x = X \land X \ge 0\}
z := y;
i := 0:
while z < x do
  z := 2 * z:
  i := i + 1:
end:
q := 0:
while z \neq y do
  i := i - 1;
  z:=z \operatorname{div} 2:
  q := 2 * q;
  if x > z then
     q := q + 1;
     x := x - z
  end:
end:
  \{Q: X = q \cdot y + x \land 0 < x < y\}
```

Above, i is a ghost variable: we only use it to make the proof easier. We can remove all assignments to i.

Exercise 7.6: Conclusion



5 Conclusion: We derived the following program fragment.

```
{P: x = X \land X > 0}
z := y;
while z \leq x do
  z := 2 * z;
end:
q := 0;
while z \neq y do
  z := z \operatorname{div} 2;
  q := 2 * q;
  if x > z then
     q := q + 1;
     x := x - z
  end;
end:
  \{Q: X = q \cdot y + x \land 0 \le x < y\}
```

The final program.



Exercise 7.5:

```
\begin{array}{l} q:=0;\\ \textbf{while}\ y\leq x\ \textbf{do}\\ x:=x-y;\\ q:=q+1;\\ \textbf{end}; \end{array}
```



Exercise 7.5:

```
q:=0; while y\leq x do x:=x-y; q:=q+1; end:
```

Exercise 7.6:

```
z := y;
while z \leq x do
  z := 2 * z;
end:
q := 0:
while z \neq y do
  z:=z \operatorname{div} 2:
  q := 2 * q;
  if x > z then
     q := q + 1;
     x := x - z
  end;
end;
```



The program in 7.6 is generally faster than the one in 7.5.



The program in 7.6 is generally faster than the one in 7.5.

Consider some sample values: x = 864, y = 23.

► The program in 7.5 returns q = 37 (because 864 = 23 * 37 + 13) That is, the assignment q := q + 1 is executed 37 times.



The program in 7.6 is generally faster than the one in 7.5.

Consider some sample values: x = 864, y = 23.

- ► The program in 7.5 returns q = 37 (because 864 = 23 * 37 + 13) That is, the assignment q := q + 1 is executed 37 times.
- The analysis for the program in 7.6 is more complicated. The auxiliary loop is executed n times, where n is the smallest value for which $2^n \cdot y > x$. That is, this loop runs at most $n = \lfloor \log_2(x \text{ div } y) \rfloor = \lfloor \log_2(q) \rfloor$.
 - That is, this loop runs at most $n = \lfloor \log_2(x \text{ div } y) \rfloor = \lfloor \log_2(q) \rfloor$ For q = 37, we have n = 5 iterations.
- After the auxiliary loop, we have $y \le z$ and z div $2 \le x < z$. The guard of the **if** is true in the first iteration and so q = 1. In each following iteration q is doubled (at least). Since $2^6 = 64 > 37$, this loop is executed at most 6 times. For q = 37, the total amount of iterations is approximately 10.

In terms of complexity analysis (not in this course): Linear time (7.5) vs logarithmic (7.6).



The End

- ► This week:
 - Exercises 7.1, 7.2, and 7.8 // Square root (Exercises 7.3 and 7.4) // Integral division (Exercises 7.5 and 7.6)
- Next week: Chapter 8. Read in advance!