Sorting

CS61B Spring'15 Discussion 11

Some Interesting Properties

▶ In-Place Sort:

- Keeps sorted items in original array (destructive)
- No equivalent for linked list-based input

Stable Sort:

 Keeps elements with equal keys in same relative order in output as in input

Algorithm

- Partition array into unsorted portion *U* and sorted portion *S*.
- For each item x in U: swap x with left neighbor until left neighbor is $\leq x$.

Intuition: Insert *x* into correct position in *S*.

 7
 3
 9
 5

7 3 9 5

3 7 9 5

3 7 9 5

3 7 5 9

3 5 7 9

Question 1(a): Insertion Sort.106 351 214 873 615 172 333 564

Question 1(a): Insertion Sort.

```
      106
      351
      214
      873
      615
      172
      333
      564

      106
      351
      214
      873
      615
      172
      333
      564

      106
      351
      214
      873
      615
      172
      333
      564

      106
      214
      351
      873
      615
      172
      333
      564

      106
      214
      351
      615
      873
      172
      333
      564

      106
      214
      351
      615
      873
      172
      333
      564

      106
      172
      214
      351
      615
      873
      333
      564

      106
      172
      214
      333
      351
      615
      873
      564

      106
      172
      214
      333
      351
      564
      615
      873
      564

      106
      172
      214
      333
      351
      564
      615
      873
      564
```

Question 3.

- Worst-Case Runtime:
- Best-Case Runtime:
- ▶ In-Place?
- Stable?

Question 3.

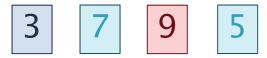
- Worst-Case Runtime: $\Theta(N^2)$
- ▶ Best–Case Runtime: $\Theta(N)$
- In-Place? ✓ Yes
- Stable? ✓ Yes

Algorithm

For each position i:

- Find smallest element x from i to end.
- Swap x and arr[i].





3 5 9 7

3 5 7 9

3 5 7 9

Question 1(b): Selection Sort.106 351 214 873 615 172 333 564

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Question 3.

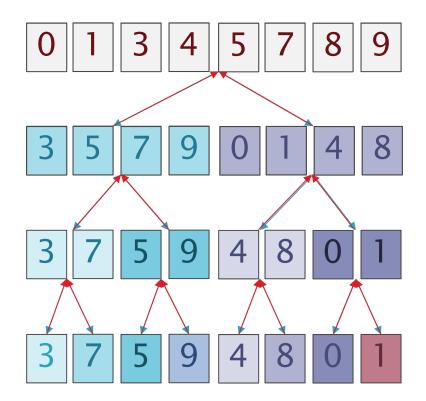
- Worst-Case Runtime:
- Best-Case Runtime:
- ▶ In-Place?
- Stable?

Question 3.

- Worst-Case Runtime: $\Theta(n^2)$
- ▶ Best–Case Runtime: $\Theta(n^2)$
- In-Place? ✓ Yes
- Stable? ✓ Yes

Algorithm

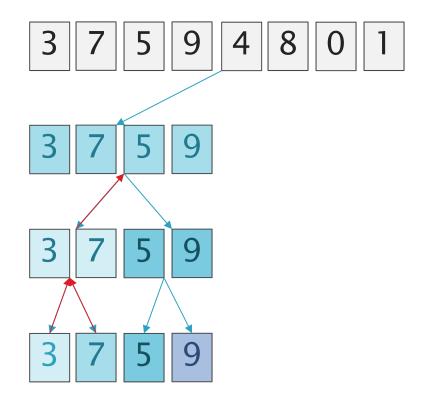
- Given N items, split into left half and right half.
- Mergesort left half.
- Mergesort right half.
- Merge the sorted halves together.



How? See Q4: MergeTwo.

Algorithm

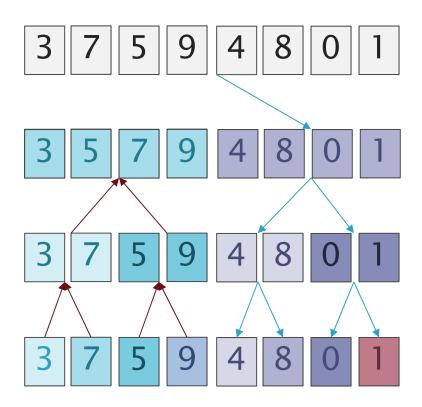
- Given N items, split into left half and right half.
- Mergesort left half.
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How? See Q4: MergeTwo.

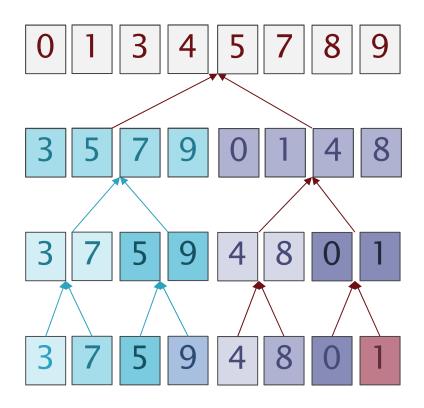
Algorithm

- Given N items, split into left half and right half.
- Mergesort left half.
- Mergesort right half.
- Merge the sorted halves together.



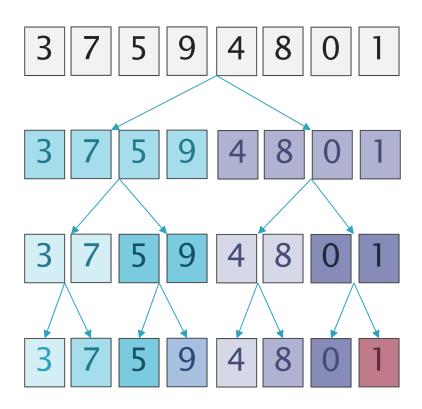
Algorithm

- Given N items, split into left half and right half.
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Algorithm

- Given N items, split into left half and right half.
- Mergesort left half.
- Mergesort right half.
- Merge the sorted halves together.



Question 1(c): Merge Sort.106 351 214 873 615 172 333 564

Question 1(c): Merge Sort.

```
      106
      351
      214
      873
      615
      172
      333
      564

      106
      351
      214
      873
      615
      172
      333
      564

      106
      351
      214
      873
      615
      172
      333
      564

      106
      351
      214
      873
      615
      172
      333
      564

      106
      351
      214
      873
      172
      615
      333
      564

      106
      214
      351
      873
      172
      333
      564
      615

      106
      172
      214
      333
      351
      564
      615
      873
```

Question 3.

- Worst-Case Runtime:
- Best-Case Runtime:
- ▶ In-Place?
- Stable?

Question 3.

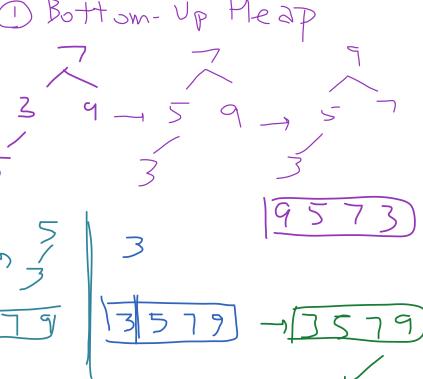
- ▶ Worst–Case Runtime: $\Theta(n \log n)$
- ▶ Best-Case Runtime: $\Theta(n \log n)$
- ▶ In-Place? **×** No
- Stable? ✓ Yes

Algorithm

- Create max-heap with bottomUpHeap().
- Repeatedly deleteMax() and place element at end of array.

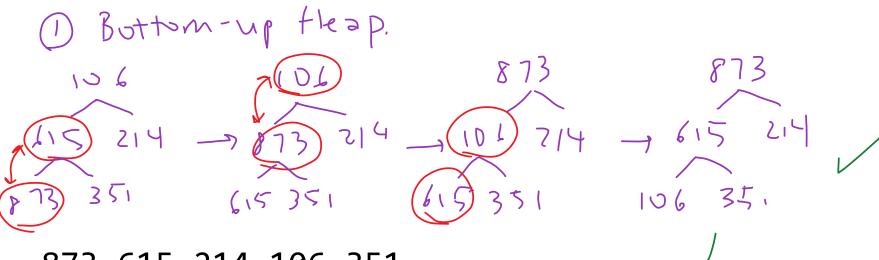
- 1. Start with first internal node and do reverse level-order traversal.
- 2. Bubble down to fix heap.

Sort: 7 3 9 5



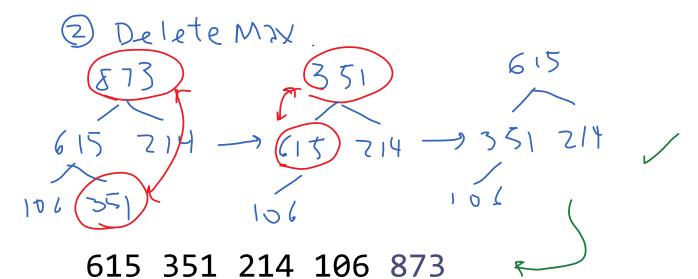
Question 1(d): Heap Sort.106 615 214 873 351

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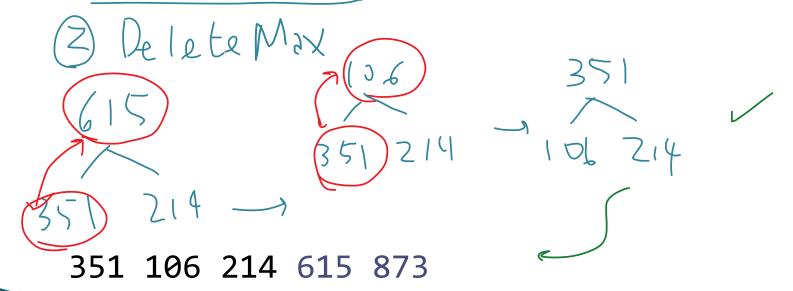


873 615 214 106 351

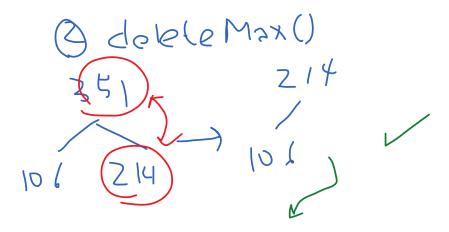
Question 1(d): Heap Sort.873 615 214 106 351



Question 1(d): Heap Sort.615 351 214 106 873



Question 1(d): Heap Sort.351 106 214 615 873



214 106 351 615 873

Question 1(d): Heap Sort.214 106 351 615 873

Question 3.

- Worst-Case Runtime:
- Best-Case Runtime:
- ▶ In-Place?
- Stable?

Question 3.

- Worst-Case Runtime: $\Theta(n \log n)$
- ▶ Best–Case Runtime: $\Theta(n)$
- In-Place? ✓ Yes
- Stable? * No Unless you add a secondary key (a timestamp) to break ties.

Merge sort isn't always better!

Question 1(e). Give an example of a situation when using insertion sort is more efficient than using merge sort.

Merge sort isn't always better!

- Question 1(e). Give an example of a situation when using insertion sort is more efficient than using merge sort.
- Insertion sort outperforms merge sort for lists that are "mostly sorted".
 - If list has only a few elements out of place
 - If all elements are within k positions of their proper place and $k < \log N$

2(a) Sorting II

Which sorting algorithm?

```
12 7 8 4 10 2 5 34 14
2 4 5 7 8 12 10 34 14
```

2(a) Sorting II

Which sorting algorithm?

```
12 7 8 4 10 2 5 34 14
2 4 5 7 8 12 10 34 14
```

Selection Sort.

```
      12
      7
      8
      4
      10
      2
      5
      34
      14

      2
      7
      8
      4
      10
      12
      5
      34
      14

      2
      4
      8
      7
      10
      12
      5
      34
      14

      2
      4
      5
      7
      10
      12
      8
      34
      14

      2
      4
      5
      7
      8
      12
      10
      34
      14
```

2(b) Sorting II

Which sorting algorithm?

23 45 12 4 65 34 20 43

12 23 45 4 65 34 20 43

2(b) Sorting II

Which sorting algorithm?

```
23 45 12 4 65 34 20 43
```

12 23 45 4 65 34 20 43

Insertion Sort.

```
23 45 12 4 65 34 20 43
```

12 23 45 4 65 34 20 43

2(c) Sorting II

Which sorting algorithm?

45 23 5 65 34 3 76 25

23 45 5 65 3 34 25 76

5 23 45 65 3 25 34 76

2(c) Sorting II

Which sorting algorithm?

```
45 23 5 65 34 3 76 25 23 45 5 65 3 34 25 76 5 23 45 65 3 25 34 76
```

Merge Sort.

```
      45
      23
      5
      65
      34
      3
      76
      25

      23
      45
      5
      65
      3
      34
      25
      76

      5
      23
      45
      65
      3
      25
      34
      76
```

2(d) Sorting II

Which sorting algorithm?

12 32 14 34 17 38 23 11

12 14 17 32 34 38 23 11

2(d) Sorting II

Which sorting algorithm?

```
12 32 14 34 17 38 23 11
12 14 17 32 34 38 23 11
```

Insertion Sort.

```
12 32 14 34 17 38 23 11
12 14 32 34 17 38 23 11
12 14 32 17 34 38 23 11
12 14 17 32 34 38 23 11
```

Suppose you are given two sorted arrays of ints. Fill in the method mergeTwo to return a new array containing all of the elements of both arrays in sorted order. Duplicates are allowed (if an element appears s times in a and t times in b, then it should appear s+t times in the returned array.

```
public static int[] mergeTwo(int[] a, int[] b) {
   // YOUR CODE HERE
}
```

- Create new array for result.
- Initialize three counters (for indices)
- Merge while both arrays have unmerged elements.
- Array b is done, so append rest of a.
- Array a is done, so append rest of b.

```
public static int[] mergeTwo(int[] a, int[] b) {
    int[] merged = new int[a.length + b.length];
    int aIndex = 0; // Current index in a.
    int bIndex = 0; // Current index in b.
    int mergedIndex = 0; // Current index in merged.
   while (aIndex < a.length && bIndex < b.length) {</pre>
        if (a[aIndex] < b[bIndex]) {</pre>
            merged[mergedIndex] = a[aIndex];
            aIndex++;
        } else {
            merged[mergedIndex] = b[bIndex];
            bIndex++;
        mergedIndex++;
                                                  // Continued on next slide.
```

```
while (aIndex < a.length) {
    merged[mergedIndex] = a[aIndex];
    aIndex++;
    mergedIndex++;
}

while (bIndex < b.length) {
    merged[mergedIndex] = b[bIndex];
    bIndex++;
    mergedIndex++;
}</pre>
```