

Worksheet 1 (Solved)

HoTTEST Summer School 2022

The HoTTEST TAS 4th July 2022

1 (*)

State the introduction and elimination rules for

- $1. \times \text{-types}$
- $2. \rightarrow$ -types
- 3. \prod -types

$$\frac{\Gamma \vdash a : A \quad \Gamma \vdash b : B}{\Gamma \vdash (a,b) : A \times B} \qquad \frac{\Gamma \vdash x : A \times B}{\Gamma \vdash \mathbf{pr}_{1}(x) : A} \qquad \frac{\Gamma \vdash x : A \times B}{\Gamma \vdash \mathbf{pr}_{2}(x) : B}$$

$$\frac{\Gamma, x : A \vdash b : B}{\Gamma \vdash \lambda(x : A).b : A \to B} \qquad \frac{\Gamma \vdash f : A \to B \quad \Gamma \vdash a : A}{\Gamma \vdash f(x) : B}$$

$$\frac{\Gamma, x : A \vdash b : B(x)}{\Gamma \vdash \lambda(x : A).b : \prod_{x : A} B(x)} \qquad \frac{\Gamma \vdash f : \prod_{x : A} B(x) \quad \Gamma \vdash a : A}{\Gamma \vdash f(a) : B(a)}$$

2 (*)

Fill in this proof tree:

3 (**)

Write a proof tree ending with a term of type $A \times B \to B \times A$ in the empty context.

4 (**)

For problems 2 and 3, what is the *logical* content of the proof tree? That is, under the "types are theorem" interpretation of Curry-Howard, what theorems have we proven?

Let A and B be propositions. Then problem 2 proves the tautology

$$A \to B \to A \wedge B$$

and problem 3 proves the tautology

$$A \wedge B \to B \wedge A$$

Next, what is the *computational* content of the proof tree? That is, under the "programs are proofs" interpretation of Curry-Howard, what programs have we written?

Problem 2 implements the function that takes in a and outputs a function taking in b and outputting (a,b).

Problem 3 implements the function that takes in the tuple (a,b) as input, and outputs the tuple (b,a)

$$\mathbf{5} \quad (\star \star \star)$$

Define the **swap** function $\sigma_{A,B}$ of type

$$\sigma_{A,B}: \left(\prod_{x:A}\prod_{y:B}C(x,y)\right) \to \left(\prod_{y:B}\prod_{x:A}C(x,y)\right)$$

and show that $\sigma_{B,A} \circ \sigma_{A,B}$ is (definitionally) equal to the identity.

$$\sigma_{A,B} \doteq \lambda \left(f: \prod_{x:A} \prod_{y:B} C(x,y) \right) . \lambda(y:B) . \lambda(x:A) . f(x,y)$$

By our computation rules:

$$(\sigma_{B,A} \circ \sigma_{A,B})(g) \doteq \sigma_{B,A}(\sigma_{A,B}(g))$$

$$\doteq \sigma_{B,A} ((\lambda f.\lambda(y:B).\lambda(x:A).f(x,y)) \quad (g))$$

$$\doteq \sigma_{B,A}(\lambda(y:B).\lambda(x:A).g(x,y))$$

$$\doteq (\lambda f.\lambda(x':A).\lambda(y':B).f(y',x')) \quad (\lambda(y:B).\lambda(x:A).g(x,y))$$

$$\doteq \lambda(x':A).\lambda(y':B). \quad ((\lambda(y:B).\lambda(x:A).g(x,y)) \quad (y',x'))$$

$$\doteq \lambda(x':A).\lambda(y':B).g(x',y')$$

$$\doteq g$$

We first substitute the definition of $\sigma_{A,B}$, then apply this definition to g. Then we expand the definition of $\sigma_{B,A}$, and apply this definition to the previous result. Finally, we apply our $\lambda(y:B).\lambda(x:A).g(x,y)$ term to the inputs (y',x') to recover something which is g up to α -equivalence. Since every step was a definitional equality we see $(\sigma_{B,A}\circ\sigma_{A,B})(g)\doteq g$ for each input g, so that $\sigma_{B,A}\circ\sigma_{A,B}$ is definitionally the identity.