

Outline

Last time: homotopies & equivalences (§9)

Today :

- A correction on joins of universes
- Contractible types & maps (§10)
 - Main result: eqivs as contr. maps
via: coherently invertible maps

A correction on joins of universes

In Lecture 4 I defined the join of $(U, \tau_U), (V, \tau_V)$ as the universe obtained by reflecting 4 type families

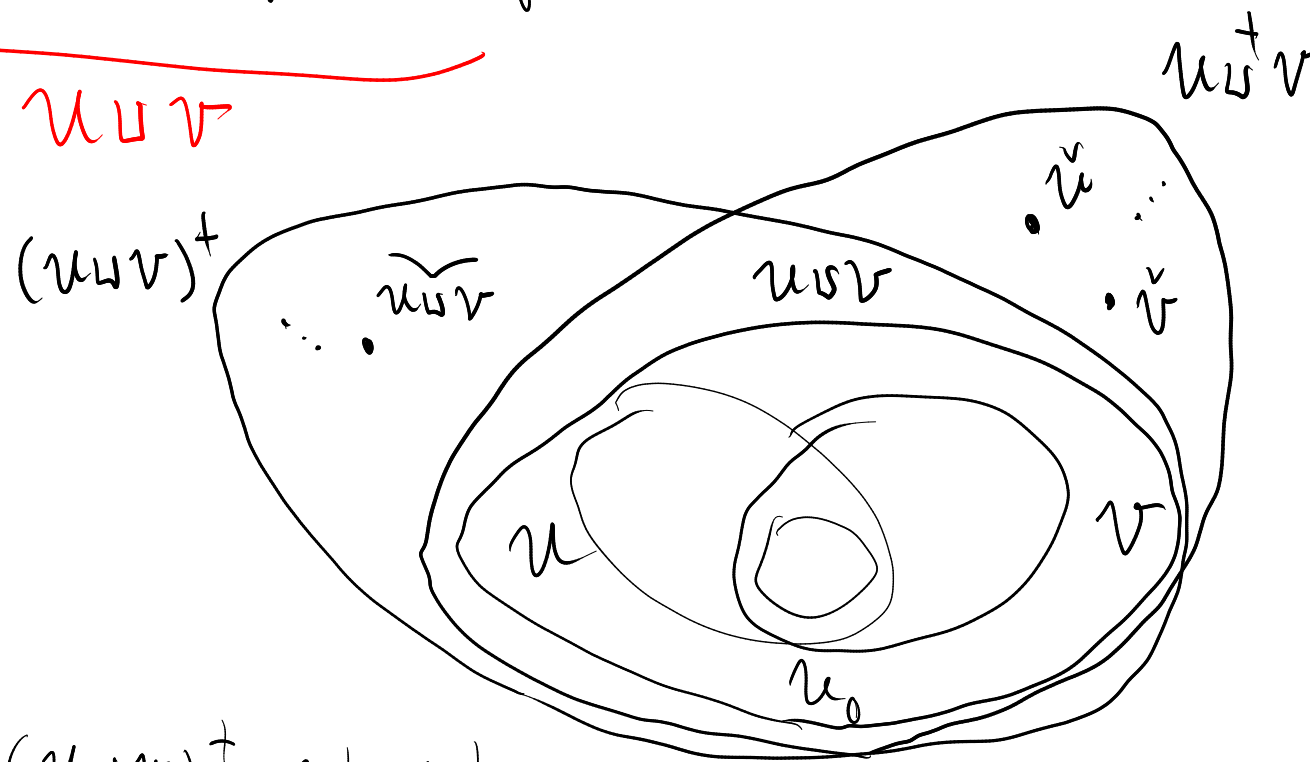
$$\left. \begin{array}{l} \cancel{\vdash U \text{ type}} \\ \cancel{\vdash V \text{ type}} \\ \underbrace{X:U \vdash \tau_U(X) \text{ type} \quad X:V \vdash \tau_V(X) \text{ type}} \end{array} \right\} \rightsquigarrow U \sqcup^+ V$$

- In Agda, we use universe polymorphism to give types to type formers, e.g.,

$$+ : U \rightarrow V \rightarrow U \sqcup V$$

Want $U \sqcup U \approx U$

\rightsquigarrow join semi-lattices w/ $+$, $(U \sqcup V)^+ = U^+ \sqcup V^+$ etc.



Contractible types

We use contractibility to capture uniqueness / ^{unique} existence

\rightsquigarrow (look up talk by Emily Riehl - Topos Colloquium)

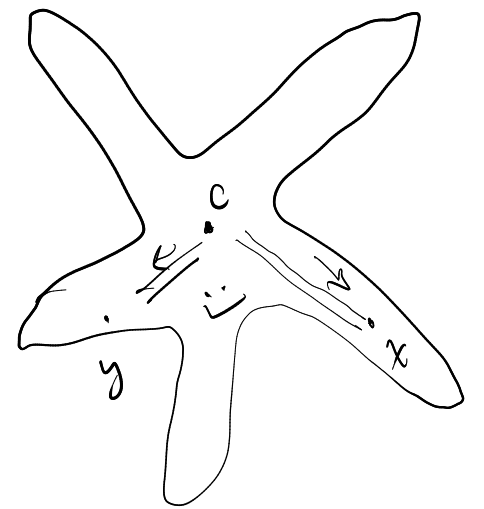
Def For any type A , let $\text{is-contr}(A) \coloneqq \sum_{c:A} \prod_{x:A} c = x$

(props as types interpretation of

"there is a $c:A$ st. every $x:A$ is equal to c ")

c is called the center of contraction

$C : \text{const}_c \sim \text{id}_A$ is called the contraction
($A \rightarrow A$)



exercise:

If A is contr., then
 $x \stackrel{A}{=} y$ is as well.

Example 1 The unit type $\mathbb{1}$ is contr. w/ $c := *$, $C := \prod_{x:\mathbb{1}} * = x$
 (1-induction)

Example 2 For a type A , w/ $a:A$,
 the Σ -type $\sum_{x:A} a = x$ is contr. w/ center
 $(a, \text{refl}(a))$

Observation A is contr. iff $! = \text{const}_*: A \rightarrow \mathbb{1}$ is an equiv.

" \rightarrow " Have $c, C: \text{const}_c \sim \text{id}_A$. Then let $g := \text{const}_c: \mathbb{1} \rightarrow A$

$$g \circ ! \doteq \text{const}_c \stackrel{C}{\sim} \text{id}_A \quad \checkmark$$

$$\& \quad ! \circ g \doteq \text{const}_* \stackrel{!}{\sim} \text{id}_{\mathbb{1}} \quad \checkmark$$

" \leftarrow "

Have $g: \mathbb{1} \rightarrow A$ inverse of $!$, let $c := g *$, for any $x:A$

$$\text{so } c \doteq (g \circ !)(x) = x.$$

Singleton induction

Since contr. types are equiv. to $\mathbb{1}$, they have all the str. props as $\mathbb{1}$.

Def Let A be a type w/ $a: A$. We say A satisfies singleton induction if for every type family B over A the eval. map at a :

$$\text{ev-pt}_a : \left(\prod_{x:A} B(x) \right) \longrightarrow B(a) \quad \text{has a section, i.e.,}$$

We have

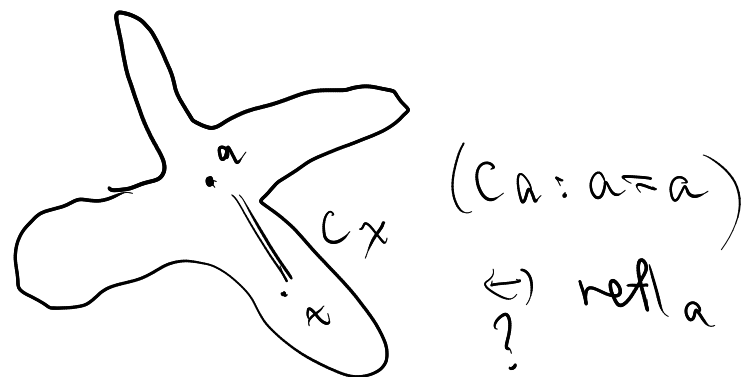
$$\text{ind-sing}_a : B(a) \longrightarrow \prod_{x:A} B(x)$$

$$\text{comp-sing}_a : \text{ev-pt}_a \circ \text{ind-sing}_a \sim \text{id}_{B(a)}$$

Thm A is contr. iff we have $a:A$ s.t. A sat. singleton ind. (a)

Thm A is contr. iff we have $a \in A$ s.t. A sat. singleton $\text{ind.}(a)$

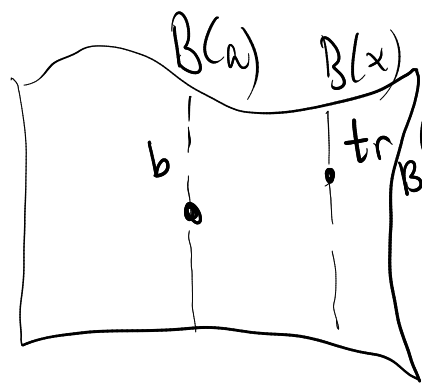
" \rightarrow " We have (a, C) : let $C'(x) := (C_a)^{-1} \cdot C_x$



$$C'(a) = C_a^{-1} \cdot C_a = \text{refl}_a.$$

WLOG assume $C_a = \text{refl}_a$.

Let now B be given

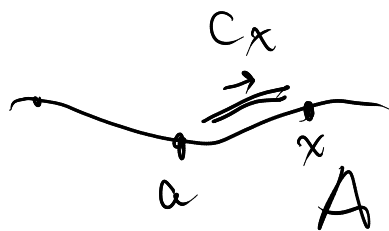


Def. $s : B(a) \rightarrow \prod_{x \in A} B(x)$

$$s(b, x) := \text{tr}_B(Cx, b)$$

Then for any b , $s(b, a) = \text{tr}_B(C_a, b)$

$$= \text{tr}_B(\text{refl}_a, b) = b \quad \checkmark$$



"←"

Suppose A w/ $a \in A$ sat. sing. ind.

of course, pick a as center.

to show $\prod_{x \in A} a = x$

by singleton ind. applied to $B := x$. $a = x$

it suffices to give $_ : B(a) \doteq a = a$
 \uparrow
 refl_a

□

Fibers & contractible map

Def Let $f: A \rightarrow B$, $b: B$. The fiber of f at b is

$$\text{fib}_f(b) := \sum_{x:A} f x = b$$

(prop as types version of

" b is in the image of f "

Obs For $p: a =_A a'$, we have

B

$$q: f a = b \quad \text{tr}_{x: f x = b}(p, q) =$$

$$(\text{ap}_f p)^{-1} \cdot q$$

/ pre-image of f at b .

(pf by path. ind)

Cor For all $(a, q), (a', q') : \text{fib}_f(b)$ we have

$$(a, q) = (a', q') \simeq \sum_{p: a = a'} \underbrace{q = \text{ap}_f p \cdot q'}_{(\text{ap}_f p)^{-1} \cdot q = q'}$$

Ind. by path ind. w/ "refl"

$$(\text{ap}_f p)^{-1} \cdot q \stackrel{12}{=} q'$$



Def A map $f: A \rightarrow B$ is contr. if all its fibers are,
 i.e., we have $\text{elt in } \prod_{y:B} \text{is-contr}(\text{fib}_f y)$

Thm Any contr. map is an equiv.

Pf From centers of contr. get $\prod_{y:B} \text{fib}_f y = \prod_{y:B} \sum_{x:A} f x = y$

by tt-choice, we get $g: B \rightarrow A$ & $G: f \circ g \sim \text{id}_B$ (sect. of f)

We want $g \circ f \sim \text{id}_A$, i.e., for all $x:A$, $(g \circ f)(x) = x$.

$f \circ g \circ f \stackrel{G \circ f}{\sim} f$ we get $p: f g f(x) = f(x)$, $(g f(x), p): \text{fib}_f(fx)$

also $(x, \text{refl}_{fx}): \text{fib}_f(fx)$, so $(g f(x), p) = (x, \text{refl}_{fx})$

so get $g f(x) = x$, as desired \square .

- Coherently invertible maps

Def $f: A \rightarrow B$ is coh. inv if we have

$$\left. \begin{array}{l} g: B \rightarrow A \\ G: f \circ g \sim \text{id}_B \\ H: g \circ f \sim \text{id}_A \\ K: G \cdot f \sim f \cdot H \quad (\text{as } \text{htpics } f \circ g \circ f \sim f) \end{array} \right\} \begin{array}{l} \text{has-inv}(f) \\ \text{is-coh-inv}(f) \end{array}$$

Thm $\text{is-coh-inv}(f) \rightarrow \text{is-contr}(f)$ (i.e., $\prod_{y:B} \text{is-contr}(\text{fib}_f y)$)

center at $y: (g y, G y)$ to show:

$$\prod_{y:B} \prod_{x:A} \prod_{q:fx=y} (g y, G y) = (x, q) \quad \text{By path ind. it suffices to give}$$

By '=' in fib's, give
for $x:A$

$$Hx: gf(x)=x, \quad G(fx) = \text{ap}_f(Hx) \cdot \text{refl}_{fx} \\ Kx \implies (f \cdot H)(x)$$

$$\prod_{x:A} (gf(x), G(fx)) = (x, \text{refl}_{fx})$$

□

Final goal $\text{has-inv}(f) \rightarrow \text{is-coh-inv}(f)$

Tool: naturality squares of htpy .

Def $f, g: A \rightarrow B$, $H: f \sim g$

$P: x \xrightarrow[A]{=} y$

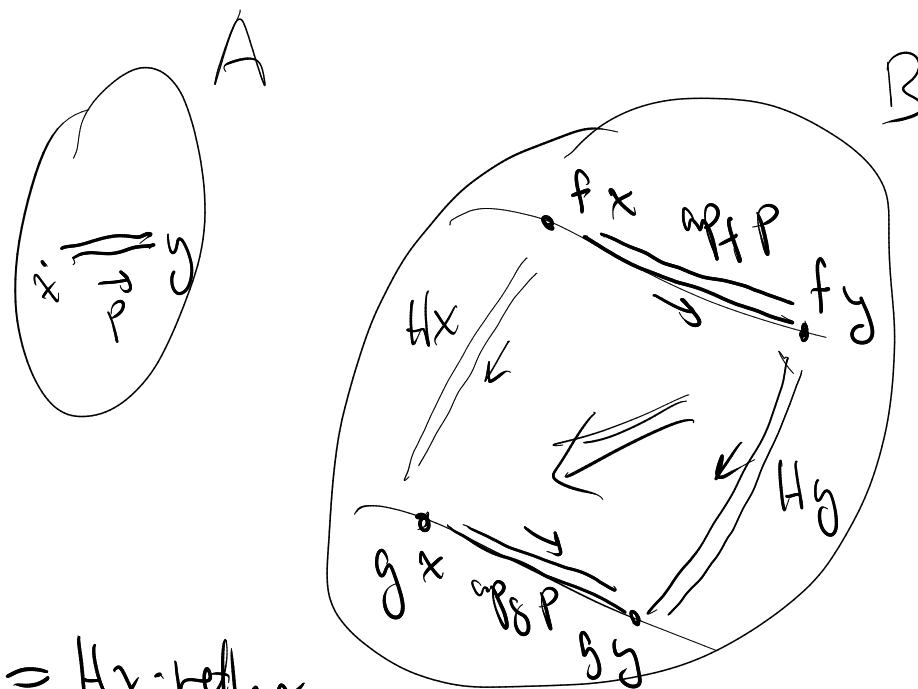
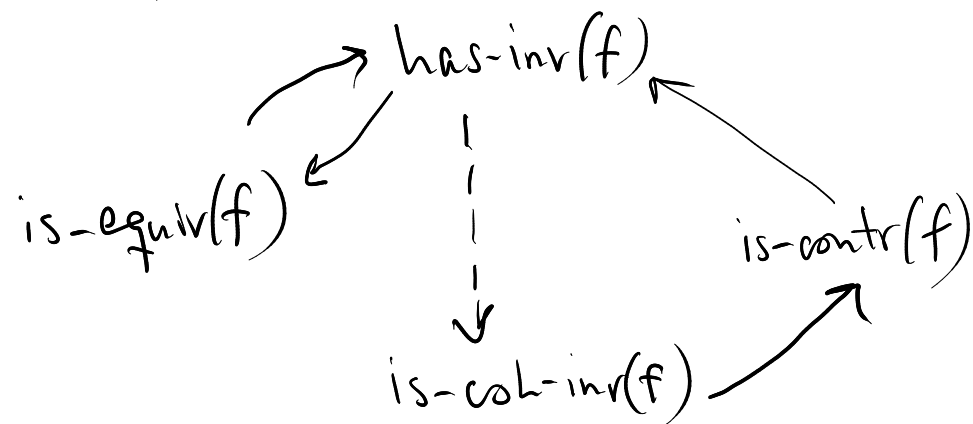
$\text{nat-htpy}(H, P): \text{ap}_f P \circ Hx = Hx \cdot \text{ap}_g P$

Defn is by path ind.: it suffices

to give $\text{nat-htpy}(H, \text{refl}_x) : \text{refl}_{fx} \cdot Hx = Hx \cdot \text{refl}_{gx}$

✓

□



Special case $h: A \rightarrow A$, $H: h \sim \text{id}_A$

$$x: A, \quad Hx: hx = x$$

$$\text{get } \text{ap}_h(Hx) = H(hx).$$

$$\begin{array}{ccc} hh(x) & \xrightarrow{\text{ap}_h(Hx)} & hx \\ H(hx) \parallel \downarrow & & \downarrow \parallel Hx \\ hx & \xrightarrow{Hx} & x \end{array}$$

Now Assume f has inv. $g: B \rightarrow A$, $G: fg \sim \text{id}_B$
 $f: A \rightarrow B$ $H: gf \sim \text{id}_A$

We improve G to $G': fg \sim \text{id}_B$ w/ $K: G' \cdot f \sim f \cdot H$.

$$\text{Let } y: B, \quad G'y := \left(fg(y) \xrightarrow{(G(fg(y)))^{-1}} fgfg(y) \xrightarrow{\text{ap}_f(H(gy))} fg(y) \xrightarrow{Gy} y \right)$$

Let $x: A$ show

$$\begin{array}{ccc} fgfgf(x) & \xrightarrow{\text{ap}_f(H(gf(x)))} & fgf(x) \\ G(fgf(x)) \parallel \downarrow & & \downarrow \parallel G(fx) \\ fgf(x) & \xrightarrow{\text{ap}_f(Hx)} & f(x) \end{array} \quad \begin{array}{c} \text{use} \\ H(gf(x)) \\ \parallel \\ \text{ap}_{gf}(Hx) \\ gf(x) \stackrel{Hx}{=} x \end{array} \quad \begin{array}{ccc} fgfgf(x) & \xrightarrow{\text{ap}_{fgf}(Hx)} & fgf(x) \\ G(fgf(x)) \parallel \downarrow & \text{nat-htpy}(G, f, Hx) & \downarrow \parallel G(fx) \\ fgf(x) & \xrightarrow{\text{ap}_f(Hx)} & f(x) \end{array}$$

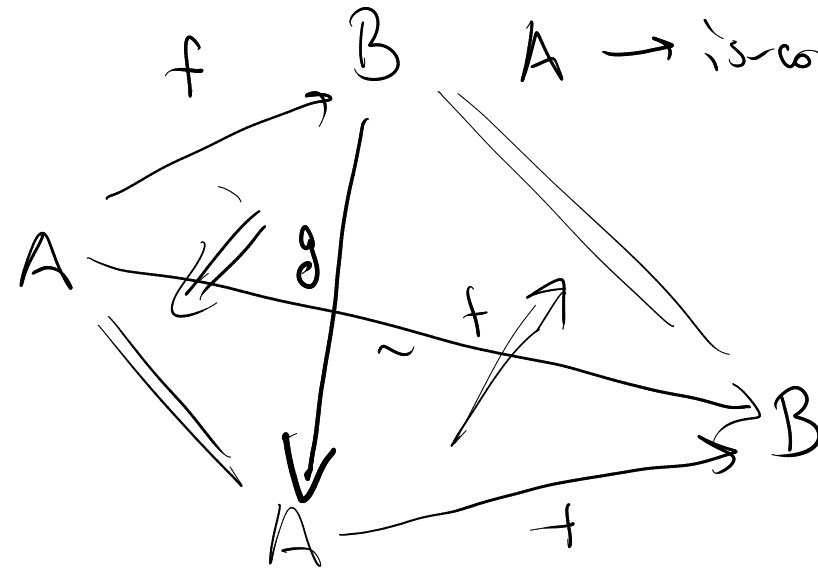
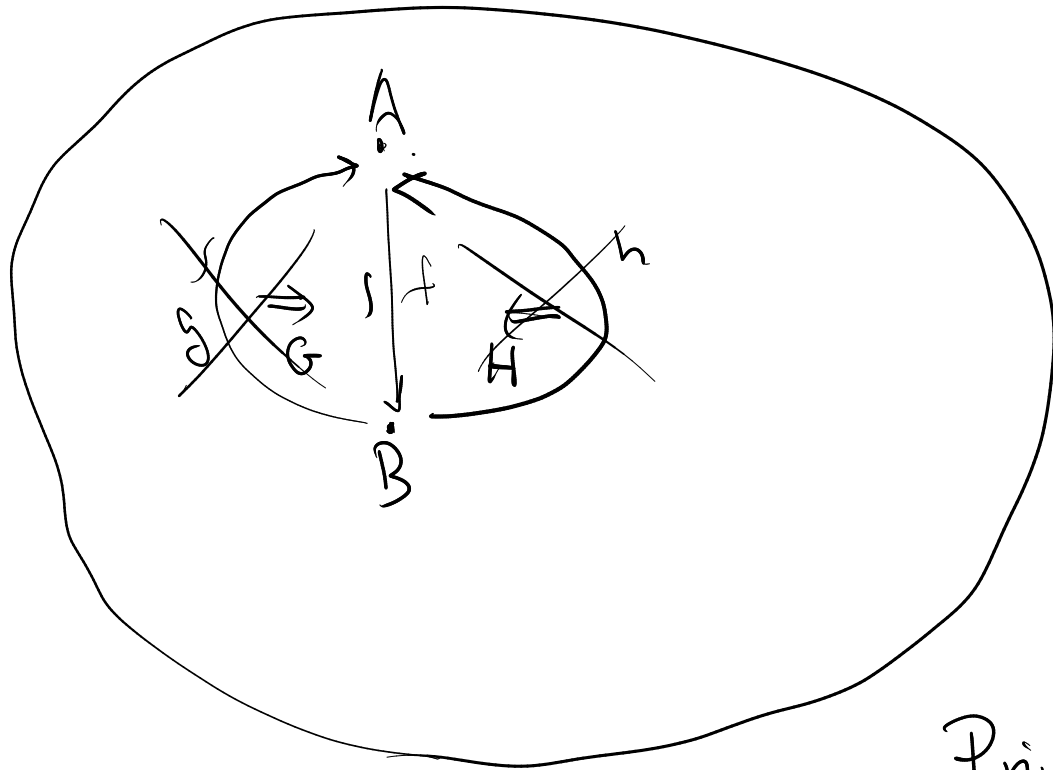
Geom intuition for why $\text{is-equiv}(f)$ & $\text{is-coh-inv}(f)$ are proposition

Assume f is an equiv.

$$(\text{is-prop}(A) := \prod_{x, y: A} x = y)$$



$$A \rightarrow \text{is-contr}(A)$$



Principle of contracting away.

If A is contr., B type fam/ A ,

$$B(c) \simeq \sum_{x:A} B(x), \quad \text{special case } \sum_{x:A} \sum_{p:a=x} B(x, p) \simeq B(a, \text{ref}_a)$$