



# Worksheet 5 (Solved)

HoTTEST Summer School 2022

The HoTTEST TAs

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## 1 (★)

Consider a function  $f : A \rightarrow B$ . Recall that a *retraction* of  $f$  is a function  $g : B \rightarrow A$  such that  $g \circ f \sim \text{id}_A$ . Construct a function

$$\text{retr}(f) \rightarrow \left( \prod_{a, a' : A} f(a) = f(a') \rightarrow a = a' \right).$$

This means that if  $f$  has a retraction, then it is an injection.

**Let  $\gamma : \text{retr}(f)$ . By  $\Sigma$ -induction, we may take  $\gamma \doteq \langle g, r \rangle$ , where  $g : B \rightarrow A$  and  $r : g \circ f \sim \text{id}_A$ . Let  $a, a' : A$  and  $p : f(a) = f(a')$ . We must find a term of type  $a = a'$ . By the action of  $g$  on  $p$ , we have that  $g(f(a)) = g(f(a'))$ . We also have that**

$$\begin{array}{ll} a = g(f(a)) & (r(a)^{-1}) \\ g(f(a')) = a'. & (r(a')) \end{array}$$

**Finally, by transitivity of equality, we have an identity**

$$r(a)^{-1} \cdot \text{ap}_g(p) \cdot r(a') : a = a'.$$

**2**   **(★★)**

Consider a commuting triangle

$$\begin{array}{ccc} A & \xrightarrow{h} & B \\ & \searrow f & \swarrow g \\ & X & \end{array} .$$

1. Suppose that  $h$  has a section  $s : B \rightarrow A$ . Prove that  $f \circ s \sim g$  and that  $\mathbf{sec}(f) \leftrightarrow \mathbf{sec}(g)$ .
2. Suppose that  $g$  has a retraction  $r : X \rightarrow B$ . Prove that  $r \circ f \sim h$  and that  $\mathbf{retr}(f) \leftrightarrow \mathbf{retr}(h)$ .
3. Prove that if any two of  $f$ ,  $g$ , and  $h$  are equivalences, then so is the third.
4. Prove that any retraction or section of an equivalence is itself an equivalence.

1. First, note that

$$f(s(b)) = g(h(s(b))) = g(b)$$

for all  $b : B$ . This means that  $f \circ s \sim g$ .

Next, suppose that  $f$  has a section  $s_f : X \rightarrow A$ . Then

$$g(h(s_f(x))) = f(s_f(x)) = x$$

for all  $x : X$ . Hence  $h \circ s_f$  is a section of  $g$ .

Conversely, suppose that  $g$  has a section  $s_g : X \rightarrow B$ . Then

$$f(s(s_g(x))) = g(s_g(x)) = x$$

for all  $x : X$ . Hence  $s \circ s_g$  is a section of  $f$ .

2. First, note that

$$r(f(a)) = r(g(h(a))) = h(a)$$

for all  $a : A$ . This means that  $r \circ f \sim h$ .

Next, suppose that  $f$  has a retraction  $r_f : X \rightarrow A$ . Then

$$r_f(g(h(a))) = r_f(f(a)) = a$$

for all  $a : A$ . Hence  $r_f \circ g$  is a retraction of  $h$ .

Conversely, suppose that  $h$  has a retraction  $r_h : B \rightarrow A$ . Then

$$r_h(r(f(a))) = r_h(h(a)) = a$$

for all  $a : A$ . Hence  $r_h \circ r$  is a retraction of  $f$ .

3. We have three cases to consider. For any equivalence  $e$ , let  $r_e$  and  $s_e$  denote its retraction and section, respectively.

- Suppose that  $h$  and  $g$  are equivalences. Then  $f$  has a section by part (1) and a retraction by part (2). Thus, it's an equivalence.
- Suppose that  $f$  and  $h$  are equivalences. Part (1) implies that  $g$  has a section. Moreover, since  $f \circ s_h \sim g$  by part (1),

$$h(r_f(g(b))) = h(r_f(f(s_h(b)))) = h(s_h(b)) = b$$

for all  $b : B$ . Thus,  $h \circ r_f$  is a retraction of  $g$ , so that  $g$  is an equivalence.

- Suppose that  $f$  and  $g$  are equivalences. Part (2) implies that  $h$  has a retraction. Moreover, since  $r_g \circ f \sim h$  by part (2),

$$h(s_f(g(b))) = r_g(f(s_f(g(b)))) = r_g(g(b)) = b$$

for all  $b : B$ . Thus,  $s_f \circ g$  is a section of  $h$ , so that  $h$  is an equivalence.

4. Notice a retraction  $r_e$  and a section  $s_e$  of an equivalence  $e : N \rightarrow M$  are commuting triangles  $\text{id}_N \sim r_e \circ e$  and  $\text{id}_M \sim e \circ s_e$ , respectively.

### 3 (★★)

Consider the type `Bool`, generated by

`true` : `Bool`  
`false` : `Bool`.

Define the type family `Eq-bool` : `Bool`  $\rightarrow$  `Bool`  $\rightarrow$   $\mathcal{U}_0$  by

`Eq-bool`(`true`, `true`) :=  $\mathbb{1}$   
`Eq-bool`(`true`, `false`) :=  $\emptyset$   
`Eq-bool`(`false`, `false`) :=  $\mathbb{1}$   
`Eq-bool`(`false`, `true`) :=  $\emptyset$ .

For every  $b, b' : \text{Bool}$ , define  $\varphi_{b,b'} : (b = b') \rightarrow \text{Eq-bool}(b, b')$  by path induction. Prove that  $\varphi_{b,b'}$  is an equivalence.

**Let's construct an inverse  $\psi_{b,b'}$  of  $\varphi_{b,b'}$  by double induction on `Bool`:**

$\psi_{\text{true}, \text{true}} : \mathbb{1} \rightarrow (\text{true} = \text{true})$   
 $\psi(*) := \text{refl}_{\text{true}}$

$\psi_{\text{true}, \text{false}} : \emptyset \rightarrow (\text{true} = \text{false})$   
 $\psi(x) := \emptyset\text{-ind}(x)$

$\psi_{\text{false}, \text{false}} : \mathbb{1} \rightarrow (\text{false} = \text{false})$   
 $\psi(*) := \text{refl}_{\text{false}}$

$\psi_{\text{false}, \text{true}} : \emptyset \rightarrow (\text{false} = \text{true})$   
 $\psi(x) := \emptyset\text{-ind}(x)$ .

**By double induction on `Bool`, it's straightforward to check that  $\varphi_{b,b'} \circ \psi_{b,b'} \sim \text{id}_{\text{Eq-bool}(b,b')}$ . Moreover, by path induction on  $b = b'$ , it's straightforward to check that  $\psi_{b,b'} \circ \varphi_{b,b'} \sim \text{id}_{b=b'}$ .**

It is easy to show that  $\neg(\mathbb{1} = \emptyset)$ . As a consequence, we can prove that  $\neg(b = \text{neg-bool}(b))$  for every  $b : \text{Bool}$ .

## 4 (★★)

Prove that for all  $b : \text{Bool}$ ,

$$\neg \text{is-equiv}(\text{const}_b).$$

Also, prove that

$$\text{Bool} \not\cong \mathbb{1}.$$

To prove that  $\neg \text{is-equiv}(\text{const}_b)$ , suppose that  $\text{const}_b$  is an equivalence. In particular, it has a retraction, hence is injective by Problem 1. Since

$$\text{const}_b(\text{true}) \doteq b \doteq \text{const}_b(\text{false}),$$

it follows that  $\text{true} = \text{false}$ . But we know that

$$\text{true} \neq \text{neg-bool}(\text{true}),$$

where  $\text{neg-bool}(\text{true}) \doteq \text{false}$ . This gives us an element of  $\emptyset$ , as required.

Likewise, to prove that  $\text{Bool} \not\cong \mathbb{1}$ , suppose that we have an equivalence  $e : \text{Bool} \rightarrow \mathbb{1}$  with retraction  $r_e : \mathbb{1} \rightarrow \text{Bool}$ . By induction on  $\mathbb{1}$ , it's easy to check that  $r_e(x) = r_e(*)$  for all  $x : \mathbb{1}$ . Then

$$\text{true} = r_e(e(\text{true})) = r_e(*) = r_e(e(\text{false})) = \text{false}.$$

Again, this gives us an element of  $\emptyset$ .

**5**    **(★)**

Let  $A$  be a type and  $B$  be a type family over  $A$ . For each  $x, y : A$ , construct an inverse of the function

$$\text{inv}_{x,y} : (x = y) \rightarrow (y = x).$$

Further, for each  $p : x = y$ , construct an inverse of the function

$$\text{tr}_B(p) : B(x) \rightarrow B(y).$$

**We may define these inverses by path induction on  $x = y$ . Specifically, define  $\text{inv}_{x,y}^{-1} : (y = x) \rightarrow (x = y)$  by**

$$\text{inv}_{x,x}^{-1}(\text{refl}_x) := \text{refl}_x$$

**and define  $\text{tr}_B(p)^{-1} : B(y) \rightarrow B(x)$  by**

$$\text{tr}_B(\text{refl}_x)^{-1} := \text{id}_{B(x)}.$$

**Our definition of  $\text{tr}_B(p)^{-1}$  gives us a homotopy**

$$\text{tr}_B(p)^{-1} \sim \text{tr}_B(p^{-1})$$

**of functions  $B(y) \rightarrow B(x)$ .**

## 6 (★)

Let  $f, g : A \rightarrow B$  and  $H : f \sim g$ . Prove that  $\text{is-equiv}(f) \leftrightarrow \text{is-equiv}(g)$ .

To define a function  $\text{is-equiv}(f) \rightarrow \text{is-equiv}(g)$ , let  $\langle h, t_h, k, t_k \rangle : \underbrace{\text{sec}(f) \times \text{retr}(f)}_{\text{is-equiv}(f)}$ .

Note that for all  $b : B$  and  $a : A$ ,

$$\begin{aligned} g(h(b)) &= f(h(b)) & (H(h(b))^{-1}) \\ &= b & (t_h(b)) \end{aligned}$$

$$\begin{aligned} k(g(a)) &= k(f(a)) & (\text{ap}_k(H(a)^{-1})) \\ &= a. & (t_k(a)) \end{aligned}$$

Thus,  $h$  is a section of  $g$ , and  $k$  is a retraction of  $g$ . Now we may take

$$\langle h, \lambda b. H(h(b))^{-1} \cdot t_h(b), k, \lambda a. \text{ap}_k(H(a)^{-1}) \cdot t_k(a) \rangle : \underbrace{\text{sec}(g) \times \text{retr}(g)}_{\text{is-equiv}(g)}$$

as our element of  $\text{is-equiv}(g)$ . The function in the other direction is similar.

## 7 (★★)

Suppose that  $e, e' : A \rightarrow B$  are equivalences and that  $H : e \sim e'$ . Let  $s$  and  $s'$  denote the sections of  $e$  and  $e'$ , respectively. Prove that  $s$  and  $s'$  are homotopic.

Recall that any section of an equivalence is also a retraction of it. Therefore, we see that

$$s(b) = s'(e'(s(b))) = s'(e(s(b))) = s'(b)$$

for all  $b : B$ .