Assignment 1

BIOINFORMATICS (CIS 455)

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Problem 1-1 Jones & Pevzner, Problem 2.1

1.

We are going to write the pseudocode for an algorithm, that given a list of n numbers, returns the largest and smallest number in that list.

Algorithm 1 Finding largest and smallest numbers

```
1: Initialize max to L[0] and min to L[0]
 2: for n in number list L do
3:
      if n < min then
          min = n
4:
      end if
5:
      if n > max then
6:
7:
         max = n
      end if
8:
9: end for
10: Return max and min
```

The time complexity of my algorithm will be O(n) = n, and the running time will be 2(n-1) + 2 because for every number in the list that starts from the second position after initialization, we need to check if the number is greater than the current maximum, or smaller than the current minimim.

2.

Now we will implement an algorithm that performs only 3n/2 comparisons to find the smallest and largest numbers in the list. For this implementation we will use the tournament method described below (Divide & Conquer).

Algorithm 2 Divide & Conquer pseudocode

```
1: We have a list l with n integers l = (l_1, l_2, ... l_n)
 2: function FINDMAXMIN(list,alow,ahigh)
      if lengthlist = 1 then
 3:
 4:
          Maximum and minimum are the same
       return maximum and minimum
      end if
 5:
      if lengthlist = 2 then
 6:
          Compare both elements in the list to find the maximum and minimum
 7:
       return maximum and minimum
      end if
 8:
      Divide the array in half, and store the two halves in arrL and arrR
 9:
      Call findmaxmin to find the maximum and minimum of arrL and arrR
10:
      if Maximum of arrL or arrR is greater than the current max then
11:
          Update global maximum
12:
      end if
13:
      if Minimum of arrL or arrR is smaller than the current min then
14:
          Update global minimum
15:
      end ifreturn the global maximum and the global minimum
17: end function
18: In the end, we have the maximum and minimum of the given list of integers.
```

In this case, we will divide the initial list with n numbers into two lists of n/2 each, for eacg of this lists, we will need to find the maximum and minimum of this lists, which can be done in n/2 - 1 comparisons. So in total we will need 3n/2 - 2 comparisons to find the maximum and minimum of a list containing n numbers.

Problem 1-2 Jones & Pevzner, Problem 2.2

Now we will write the pseudocode for two algorithms that iterate over every index from (0,0,...,0) to $(n_1,n_2,...,n_d)$, one of them will be recursive and the other one iterative.

Iterative pseudocode:

Algorithm 3 Iterative pseudocode

```
    We have two indexes lists l = (0,0,...0) and n = (n<sub>1</sub>, n<sub>2</sub>,...n<sub>d</sub>).
    We will print all the possible combination of elements
    for each (i<sub>1</sub>, i<sub>2</sub>,..., i<sub>d</sub>) from (0,0,...,0) to (n<sub>1</sub>, n<sub>2</sub>,..., n<sub>d</sub>) do
    Return list i.
```

Recursive pseudocode:

5: end for

Algorithm 4 Recursive pseudocode

```
1: We have two indexes lists l = (0, 0, ...0) and n = (n_1, n_2, ...n_d).
     2: We will recursively build a tree to print all the possible combinations
     3: function Combinations(n,lev,index)
           while levid do
    4:
              Add n_i ndex to list list
                                           if level = 0 then
    5:
              Return array of 0's
Б:
8:
          if index = d then
9:
              Reset index
10:
          end if
11:
          if We reach the lenght of l and used all d then
12:
              Go down one level and reset index
13:
          end if
14:
          for n_i in n do
15:
              Return () Combinations(n,lev,i)
16:
17:
              Return the sum value
```

Problem 1-3 Jones & Pevzner, Problem 2.3

- Yes, $\log n = O(n)$, because O(n) is an upper bound for $\log n$, as it grows faster.
- No, $\log n = \Omega(n)$, because $\Omega(n)$ is not a lower bound for $\log n$, as $\log n$ grows significantly slower.
- No, $\log n = \Theta(n)$, because $\log n$ is not O(n) and $\Omega(n)$ at the same time.

Problem 1-4 Jones & Pevzner, Problem 2.17

Will the viruses eventually kill all the bacteria?

In order to find if the viruses end up killing all the bacteria in the Petri dish, we only need to prove that the growth rate of the viruses is higher than the growth rate of the bacteria.

In the first minute, the virus kills one bacteria, and produces another copy of himself, and all the remaining bacteria reproduce, making 2 viruses and 2(n-1) bacteria.

The number of viruses continue to double at each step, also, the number of bacteria double at each step too, but before the bacteria double, the viruses kill as many bacteria as the number of viruses there are at that moment. Then, the growth rate of bacteria is slower than the growth rate of the viruses, because is slowed by the number of viruses at each step. Thus, at some point, the bacteria will be exterminated by the viruses.

Algorithm design

Algorithm 5 Algorithm for calculating the number of steps

- 1: At the beginning we have $n_v = 1$ virus and $n_b = n$ bacteria
- 2: At each step, the number of viruses double, the bacteria double too, but n_v are killed by the virus too before they double.
- 3: while $n_b > n_v$ do
- 4: $n_b = 2 * (n_b n_v)$
- 5: $n_v = 2 * n_v$
- 6: steps + +
- 7: end while=0

Problem 1-5



Figure 1: DNA Sequence

5.A

The highest total score is 3599 with a max score of 435.

5.B

The description of the result with the highest score is: Cloning vector pAx-CALRL, complete sequence.

5.C

The query coverage of this result is: 99%

5.D

The DNA sequence in Jurassic Park is **fictional**, because the tool we used couldn't find a match without gaps in the database. The lowest fraction of gaps found was 12%, that's why the maximum Ident value is 88%.

Problem 1-6 Rosalind