Initially: EBP and ESP point to the top and bottom of the bar stack frame.

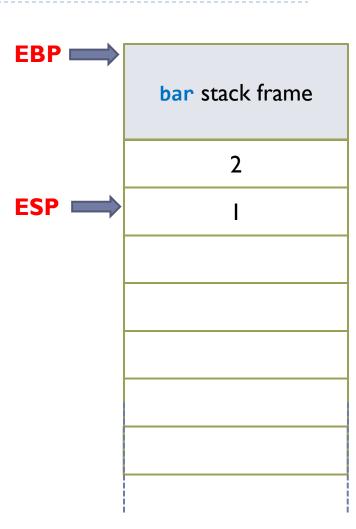
```
EBP •
              bar stack frame
ESP I
```

```
void bar( ) {
  foo(1, 2);
}
int foo(int x, int y){
  int z = x + y;
  return z;
}
```

## Step I: Push function parameters to the stack.

- Function parameters are stored in reverse order.
- ESP is updated to denote the lowest stack location due to the push operation.

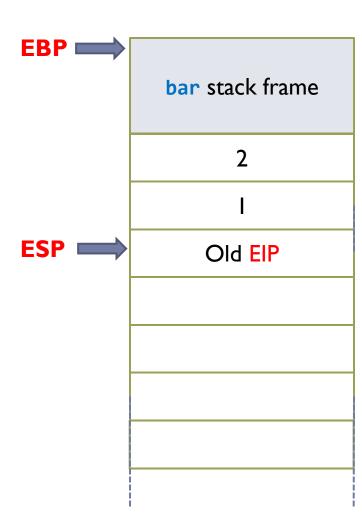
```
void bar( ) {
  foo(1, 2);
}
int foo(int x, int y){
  int z = x + y;
  return z;
}
```



# Step 2: Push the current instruction pointer (EIP) to the stack.

- This is the return address in function bar after we finish function foo.
- ESP is updated to denote the lowest stack location due to the push operation.

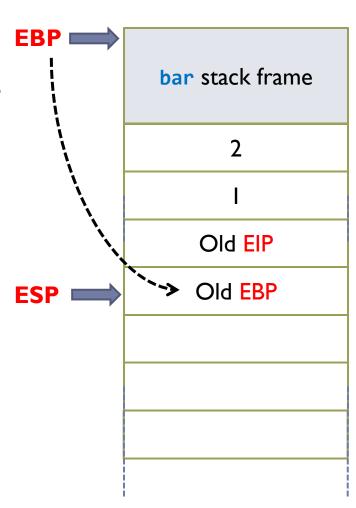
```
void bar( ) {
  foo(1, 2);
}
int foo(int x, int y){
  int z = x + y;
  return z;
}
```



### Step 3: Push the EBP of function bar to the stack.

- This can help restore the top of function bar stack frame when we finish function foo.
- ESP is updated to denote the lowest stack location due to the push operation.

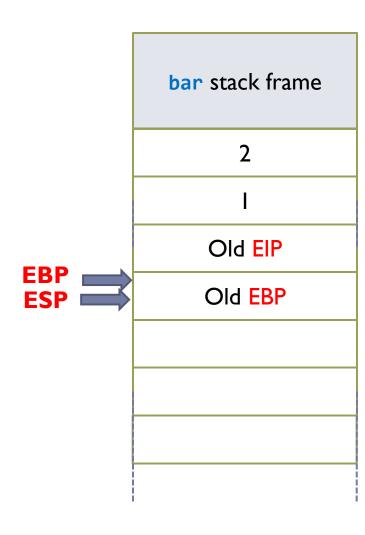
```
void bar() {
  foo(1, 2);
}
int foo(int x, int y){
  int z = x + y;
  return z;
}
```



## Step 4: Adjust EBP for function foo stack frame.

Move EBP to ESP of bar stack frame

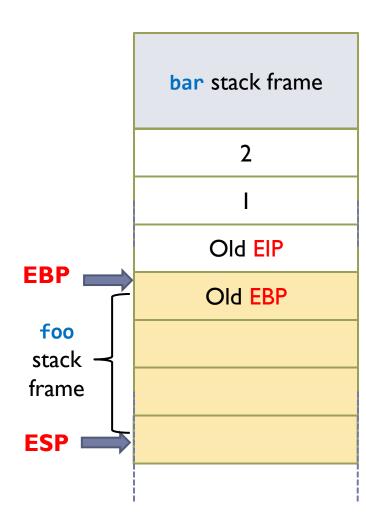
```
void bar( ) {
   foo(1, 2);
}
int foo(int x, int y){
   int z = x + y;
   return z;
}
```



### Step 5: Adjust ESP for function foo stack frame.

- Move ESP to some location below to create a new stack frame for function foo
- The stack space for function foo is precalculated based on the source code. It is used for storing the local variables and intermediate results.

```
void bar() {
  foo(1, 2);
}
int foo(int x, int y){
  int z = x + y;
  return z;
}
```



# Step 6: Execute function foo within its stack frame.

The returned result will be stored in the register EAX.

```
void bar( ) {
  foo(1, 2);
}
int foo(int x, int y){
  int z = x + y;
  return z;
}
```

