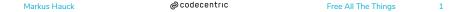
Free All The Things

Markus Hauck



Free All The Things

- well known: free monads
- maybe known: free applicatives
- free monoids
- free <you name it>

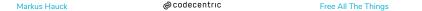


Goal Of This Talk

- how many of you wrote a Free X
- how many of you used Free...
 - Monad
 - Applicative
 - Functor
 - other?
- Goal: explain the technique behind "Free X" + example

The Road Ahead

- Demonstrate a recipe to "free" things
- Using: Free Monads, Applicatives, Functors
- New thing: Free Boolean Algebra + Example



What's The Problem

A free functor is left adjoint to a forgetful functor what's the problem?



What Is Free

A free "thing" **FreeA** on a type(class) A is a A and a function

def inject(x: A): FreeA

such that for any other "thing" B and a function

val $f: A \Rightarrow B$

there exists a unique homomorphism g such that

What Is Free

- still sounds complicated?
- · there is a recipe
 - AST
 - inject
 - interpreter
 - check laws



Why Free

- nice API using typeclass
- use Free X as if it was X
- program reified into datastructure
- structure can be analyzed/optimized
- one program many interpretations

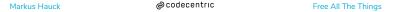


Disclaimer Before We Start

- deep embeddings / initial encoding / data structure representation
- not: finally tagless, optimization



Freeing The Monad



The Monad Typeclass

```
trait Monad[F[_]] {
    def pure[A](x: A): F[A]

def flatMap[A, B](fa: F[A])(f: A => F[B]): F[B]
}
```

Give Me The Laws

```
// Left identity
pure(a).flatMap(f) === f(a)

// Right identity
fa.flatMap(pure) === fa

// Associativity
fa.flatMap(f).flatMap(g) ===
fa.flatMap(a => f(a).flatMap(g))
```

Applying The Recipe

```
trait Monad[F[_]] {
   def pure[A](x: A): F[A]

def flatMap[A, B](fa: F[A])(f: A => F[B]): F[B]
}
```

- now comes our recipe
 - AST
 - inject
 - interpreter
 - check laws



Freeing The Monad

```
sealed abstract class Free[F[ ], A]
1
2
      final case class Pure[F[ ], A](a: A)
3
          extends Free[F, A]
4
5
      final case class FlatMap[F[], A, B](
6
          fa: Free[F, A],
7
          f: A \Rightarrow Free[F, B]
8
          extends Free[F, B]
9
10
      final case class Inject[F[_], A](fa: F[A])
11
          extends Free[F, A]
12
```

Markus Hauck @ codecentric Free All The Things 13

Freeing The Monad

```
implicit def freeMonad[F[_], A]: Monad[Free[F, ?]] =
new Monad[Free[F, ?]] {
    def pure[A](x: A): Free[F, A] = Pure(x)

def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] =
FlatMap(fa, f)
}
```

14

Interpreter

```
def runFree[F[_], M[_]: Monad, A](nat: F ~> M)(
    free: Free[F, A]): M[A] = free match {
    case Pure(x) => Monad[M].pure(x)

    case Inject(fa) => nat(fa)

    case FlatMap(fa, f) =>
    Monad[M].flatMap(runFree(nat)(fa))(x =>
    runFree(nat)(f(x)))
}
```

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What about the laws?

```
// The associativity law
  fa.flatMap(f).flatMap(g) ===
     fa.flatMap(fa, a \Rightarrow f(a).flatMap(q))
  val exp1 = FlatMap(FlatMap(fa, f), g)
  val exp2 = FlatMap(fa, (a: Int) => FlatMap(f(a), q))
3
  exp1 != exp2
```

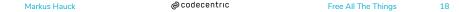
What about the laws?





The Laws

- actually, we don't satisfy them
- programmer: after interpretation it's no longer visible
- mathematician: that's not the free monad!
- tradeoff: during construction vs during interpretation



The Right Free Monad

- common transformation: associate flatMap's to the right
- avoids having to rebuild the tree repeatedly during construction
- how: during construction time



Transforming Free Monads

```
def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] = fa match {
    case Pure(x) => f(x) // Left identity
    case Inject(fa) => FlatMap(Inject(fa), f)
    case FlatMap(ga, g) => // Associativity
    FlatMap(ga, (a: Any) => FlatMap(g(a), f))
}
```

Transforming Free Monads

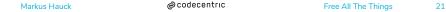
```
def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] = fa match {
    case Pure(x) => f(x) // Left identity
    case Inject(fa) => FlatMap(Inject(fa), f)
    case FlatMap(ga, g) => // Associativity
    FlatMap(ga, (a: Any) => FlatMap(g(a), f))
}
```

Transforming Free Monads

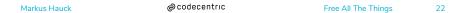
```
def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] = fa match {
    case Pure(x) => f(x) // Left identity
    case Inject(fa) => FlatMap(Inject(fa), f)
    case FlatMap(ga, g) => // Associativity
    FlatMap(ga, (a: Any) => FlatMap(g(a), f))
}
```

We Freed Monads

- DSL with monadic expressiveness
- context sensitive, branching, loops, fancy control flow
- familiarity with monadic style for DSL
- big drawback: interpreter has limited possibilities



Freeing The Applicative



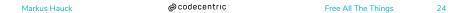
Freeing The Applicative

- free monads are great, but also limited
- we can't analyze the programs
- how about a smaller abstraction?



Recall

- we follow the same pattern
- AST
- inject
- interpreter
- check laws



The Applicative Typeclass

```
trait Applicative[F[_]] {
   def pure[A](x: A): F[A]

def ap[A, B](fab: F[A => B], fa: F[A]): F[B]
}
```

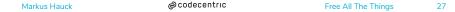
25

AST for FreeApplicative

```
sealed abstract class FreeAp[F[ ], A]
2
      final case class Pure [F], A (a: A)
3
          extends FreeAp[F. A]
4
5
      final case class Ap[F[ ], A, B](
6
          fab: FreeAp[F, A \Rightarrow B],
7
          fa: FreeAp[F, A])
8
          extends FreeAp[F, B]
9
10
      final case class Inject[F[\ ], A](fa: F[A])
11
          extends FreeAp[F, A]
12
```

Laws

```
1 // identity
   Ap(Pure(identity), v) === v
3
   // composition
   Ap(Ap(Ap(Pure(.compose), u), v), w) ===
     Ap(u, Ap(v, w))
  // homomorphism
   Ap(Pure(f), Pure(x)) === Pure(f(x))
10
   // interchange
   Ap(u, Pure(y)) === Ap(Pure((v)), u)
12
```



Don't Forget The Laws

```
def ap[A, B](fab: FreeAp[F, A \Rightarrow B],
1
                    fa: FreeAp[F, A]): FreeAp[F, B] =
2
        (fab, fa) match {
3
          case (Pure(f), Pure(x)) =>
4
            Pure(f(x)) // homomorphism
5
          case (u. Pure(v)) =>
6
            Ap(Pure((f: A \Rightarrow B) \Rightarrow f(y)), u) // interchange)
7
          case ( , ) => Ap(fab, fa)
8
```

Don't Forget The Laws

```
def ap[A, B](fab: FreeAp[F, A \Rightarrow B],
1
                    fa: FreeAp[F, A]): FreeAp[F, B] =
2
        (fab, fa) match {
3
          case (Pure(f), Pure(x)) =>
4
            Pure(f(x)) // homomorphism
5
          case (u. Pure(v)) =>
6
            Ap(Pure((f: A \Rightarrow B) \Rightarrow f(y)), u) // interchange)
7
          case ( , ) => Ap(fab, fa)
8
```

Running FreeApplicatives

```
def runFreeAp[F[\_], M[\_]: Applicative, A](
1
         nat: F \sim M)(free: FreeAp[F, A]): M[A] =
2
       free match {
3
         case Pure(x) => Applicative[M].pure(x)
4
         case Inject(fa) => nat(fa)
5
         case Ap(fab, fa) =>
6
           Applicative[M]
7
             .ap(runFreeAp(nat)(fab), runFreeAp(nat)(fa))
8
```

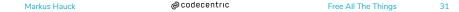
We Freed Applicatives

- DSL with applicative expressiveness
- context insensitive
- pure computation over effectful arguments
- more freedom during interpretation



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Freeing The Functor



And Once Again

- AST
- inject
- interpreter
- check laws



The Functor Typeclass

```
trait Functor[F[_]] {
    def map[A, B](fa: F[A])(f: A => B): F[B]
}
```

```
sealed abstract class FreeFunctor[F[_], A]

case class Fmap[F[_], X, A](fa: F[X])(f: X => A)
    extends FreeFunctor[F, A]

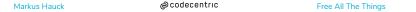
case class Inject[F[_], A](fa: F[A])
    extends FreeFunctor[F, A]
```



```
sealed abstract class FreeFunctor[F[_], A]

case class Fmap[F[_], X, A](fa: F[X])(f: X => A)
    extends FreeFunctor[F, A]

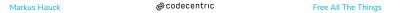
def inject[F[_], A](value: F[A]) =
    Fmap(value)(identity)
```



Clean Code Police



• only one subclass?



```
sealed abstract class Fmap[F[_], A] {
1
       type X
2
       def fa: F[X]
       def f: X => A
4
5
6
     def inject[F[], A](v: F[A]) = new Fmap[F, A] {
7
       type X = A
8
       def fa = v
9
       def f = identity
10
11
```

```
sealed abstract class Coyoneda [F[], A]
1
       type X
2
       def fa: F[X]
       def f: X => A
6
     def inject[F[], A](v: F[A]) = new Coyoneda[F, A] {
7
       type X = A
8
       def fa = v
9
       def f = identity
10
11
```



Free Functor Instance

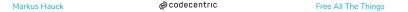
```
implicit def covoFun[F[ ]]: Functor[Covoneda[F, ?]] =
1
        new Functor[Covoneda[F, ?]] {
2
          def map[A, B](coyo: Coyoneda[F, A])(
3
              q: A \Rightarrow B: Covoneda[F, B] =
4
            new Coyoneda[F, B] {
5
              type X = coyo.X
6
              def fa = coyo.fa
7
              def f = q.compose(coyo.f)
8
9
10
```

Free Functor Instance

```
implicit def coyoFun[F[_]]: Functor[Coyoneda[F, ?]] =
1
        new Functor[Covoneda[F, ?]] {
2
          def map[A, B](coyo: Coyoneda[F, A])(
3
              q: A \Rightarrow B: Covoneda[F, B] =
4
            new Coyoneda[F, B] {
5
              type X = coyo.X
6
              def fa = coyo.fa
7
              def f = q.compose(coyo.f)
8
9
10
```

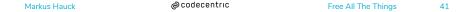
Free Functor Interpreter

```
def runCoyo[F[_]: Functor, A](
coyo: Coyoneda[F, A]): F[A] =
Functor[F].map(coyo.fa)(coyo.f)
```



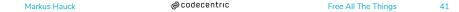
We Freed Functors

- DSL with hmm functorial expressiveness?
- map fusion! (functor law)
- boring interpreter, though



We Freed Functors

- DSL with hmm functorial expressiveness?
- map fusion! (functor law)
- boring interpreter, though
- still fun!

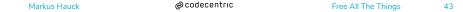


Freeing The Monoid



The Monoid Typeclass

```
trait Monoid[A] {
def empty: A
def combine(x: A, y: A): A
}
```



The Free Monoid — First Try

```
sealed abstract class FreeMonoid[+A]

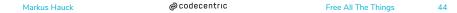
case object Empty extends FreeMonoid[Nothing]

case class Inject[A](x: A) extends FreeMonoid[A]

case class Combine[A](x: FreeMonoid[A],

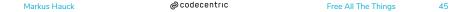
y: FreeMonoid[A])

extends FreeMonoid[A]
```



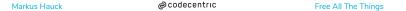
The Laws

```
1  // left identity
2  empty |+| x === x
3
4  // right identity
5  x |+| empty === x
6
7  // associativity
8  1 |+| (2 |+| 3) === (1 |+| 2) |+| 3
```



The Laws and Free Monoid

- let's try to enforce those laws in our structure
- goal: correct by construction
- arbitrary decision: associate left vs right



Fixing Associativity

```
sealed trait NotCombine[+A]
1
2
      sealed abstract class FreeMonoid[+A]
3
4
      case object Empty
5
          extends FreeMonoid[Nothing]
6
          with NotCombine[Nothing]
8
      case class Inject[A](x: A)
9
          extends FreeMonoid[A]
10
          with NotCombine[A]
11
12
      case class Combine[A](x: NotCombine[A],
13
                             v: FreeMonoid[A])
14
          extends FreeMonoid[A]
15
```

The Problem With Neutral Elements

- get rid completely? not possible
- limit ourselves to a single element
- restrict Combine to have only real values on the left side
- goal: minimal canonical structure



Minimizing Structure — Extract Inject

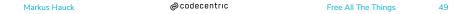
```
case class Inject[A](x: A)

sealed abstract class FreeMonoid[+A]

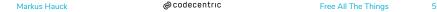
case object Empty extends FreeMonoid[Nothing]

case class Combine[A](x: Inject[A], y: FreeMonoid[A])

extends FreeMonoid[A]
```



Minimizing Structure — Remove Inject

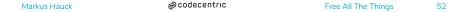


The Monoid Instance

```
implicit def monoid[A]: Monoid[FreeMonoid[A]] =
1
       new Monoid[FreeMonoid[A]] {
2
          override def empty = Empty
3
          override def combine(
4
              x: FreeMonoid[A],
5
              y: FreeMonoid[A]): FreeMonoid[A] = x match {
6
            case Empty
7
                                => V
            case Combine(h, t) => Combine(h, combine(t, y))
8
9
10
```

Minimizing Structure — List

```
sealed abstract class List[+A]
2
     case object Nil extends List[Nothing]
3
4
     case class Cons[A](head: A, tail: List[A])
5
         extends List[A]
6
```

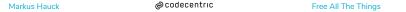


Now That We Can Free Anything



What should we free?

Monads

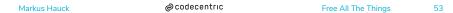


Now That We Can Free Anything



What should we free?

Monads Applicatives



Now That We Can Free Anything



What should we free?

Monads Applicatives Functors

Now That We Can Free Anything



What should we free?

Monads Applicatives Functors Monoids

Now That We Can Free Anything



What should we free?

Monads Applicatives Functors Monoids Semigroups

Now That We Can Free Anything

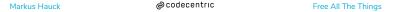


What should we free?

Monads Applicatives Functors Monoids Semigroups Groups

Credit Where It's Due

- Once upon a time: https://engineering.wingify.com/posts/Free-objects/
- use free boolean algebra to define DSL for event predicates
- credits to Chris Stucchio (@stucchio)



Let's Free A Boolean Algebra

- DSL: and, or, not, true, false
- we know what to do, so let's go!
- AST
- inject
- interpreter
- check laws



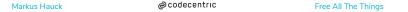
Boolean Algebras

```
trait BoolAlgebra[A] {
  def tru: A
  def fls: A

def not(value: A): A

def and(lhs: A, rhs: A): A
  def or(lhs: A, rhs: A): A
}
```

```
sealed abstract class FreeBool[+A]
1
2
     case object Tru extends FreeBool[Nothing]
3
     case object Fls extends FreeBool[Nothing]
4
5
     case class Not[A](value: FreeBool[A])
6
          extends FreeBool[A]
7
     case class And[A](lhs: FreeBool[A], rhs: FreeBool[A])
8
          extends FreeBool[A]
9
     case class Or[A](lhs: FreeBool[A], rhs: FreeBool[A])
10
          extends FreeBool[A]
11
     case class Inject[A](value: A) extends FreeBool[A]
12
```



```
def runFreeBool[A, B](fb: FreeBool[A])(f: A \Rightarrow B)(
1
          implicit B: BoolAlgebra[B]): B = {
2
        fb match {
3
          case Tru
                      => B.tru
4
          case Fls
                       => B.fls
5
          case Inject(v) \Rightarrow f(v)
6
          case Not(v) => B.not(runFreeBool(v)(f))
7
          case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
          case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A \Rightarrow B)(
1
          implicit B: BoolAlgebra[B]): B = {
2
        fb match {
3
          case Tru
                       => B.tru
4
          case Fls
                       => B.fls
5
          case Inject(v) \Rightarrow f(v)
6
          case Not(v) => B.not(runFreeBool(v)(f))
7
          case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
          case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A \Rightarrow B)(
1
          implicit B: BoolAlgebra[B]): B = {
2
        fb match {
3
          case Tru
                       => B.tru
4
          case Fls
                       => B.fls
5
          case Inject(v) \Rightarrow f(v)
6
          case Not(v) => B.not(runFreeBool(v)(f))
7
          case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
          case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

Free Boolean Algebra

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A \Rightarrow B)(
1
          implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru
                      => B.tru
4
         case Fls
                       => B.fls
5
          case Inject(v) => f(v)
6
         case Not(v) => B.not(runFreeBool(v)(f))
7
          case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
          case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

roduction Free Monad Free Applicative Free Functor Free Monoid **Free Boolean Algebra** Conclusio

Using Free Bool

- that was simple (though boilerplate-y)
- what can we do with our new discovered structure
- reminder: boolean operators
 - true, false
 - and, or
 - · xor, implies, nand, nor, nxor



```
sealed trait Search
case class Term(t: String) extends Search
case class After(date: Date) extends Search
case class InText(t: String) extends Search
case class InUrl(url: String) extends Search
// and the usual smart ctors
```

roduction Free Monad Free Applicative Free Functor Free Monoid Free Boolean Algebra Conclusion

Free Bool Example: Search

assuming some implicits we can write:

```
val search = term("FP") &
after("20180101") &
!(term("Java") | inText("spring")) &
inUrl("lambdaconf")
```



```
def evalSearch(pred: FreeBool[Search])(
1
         site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
         case Term(t) => site.terms.contains(t)
4
         case After(d)
                                 => site.indexedAt > d
5
         case InText(t: String) => site.text.contains(t)
6
         case InUrl(w)
                                 => site.url.contains(w)
8
9
       runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

```
def evalSearch(pred: FreeBool[Search])(
1
          site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
                                  => site.terms.contains(t)
          case Term(t)
4
         case After(d)
                                  => site.indexedAt > d
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But Wait There's More

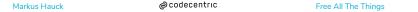
- short circuiting and other optimization
- what if you don't have all the information?
 - partially evaluate predicates
 - if evaluates successfully, done
 - else, send it on
- core language vs extension
 - · Chris also demonstrates extension
 - translate a rich language to base instructions
 - · with all the advantages

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We Freed Boolean Algebras

- good example of underused free structure
- partial evaluation
- serialize the AST (JSON, Protobuf, Avro, ...)
- exercise: minimize AST representation



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Resources

- Free Boolean Algebra by Chris Stucchio https://engineering.wingify.com/posts/Free-objects/
- Source Code: https://github.com/markus1189/ free-all-the-things/tree/lambdaconf



Go And Free All The Things!

Markus Hauck (@markus1189)



Introduction

Free Monad

Free Applicative

Free Functor

Free Monoid

Free Boolean Algebra

Conclusion

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Bonus

- try to encode normal form for boolean algebra
- remove Inject cases from Monad and Applicative
- free Magmas
- define free X using alternative minimal set of ops of the typeclass