

LIFE OF A HELLO PROGRAM

% CSAPP-CHAPTER 1

TRANSLATION

SOURCE FILE

- interpretation of source file
 - sequence of bits (as file)
 - text characters (as text file)
 - program of C statements, syntax (as C source file)
- encoding: context decides representation

TRANSLATION

```
gcc hello.c
```

1. from: C statement @source file
2. to: machine-language instructions @executable object file

COMPILATION SYSTEMS

Four phases

1. Preprocessor (cpp): directives (#), .c, .h -> .i (text file)
2. compiler (cc1): .i -> .s (text file)
 - .s is assembly-lang. program, each statement is 1-to-1 mapped to a machine-lang. instruction
3. assembler (as): .s -> .o
 - relocatable object program, binary file: encode instructions not characters
4. linker (ld): .o, .o -> .o.
 - merge multiple relocatable objects to a single executable object

EXECUTION

in shell: command-line interpreter

`./a.out`

How does my screen know what to do, when I only tell my keyboard an object file name?

This is a process of three steps.

1. typing string `a.out`: from keyboard to main memory
2. load executable file `a.out`: from disk to main memory
3. execute object `a.out`: OS and HW

HW OVERVIEW

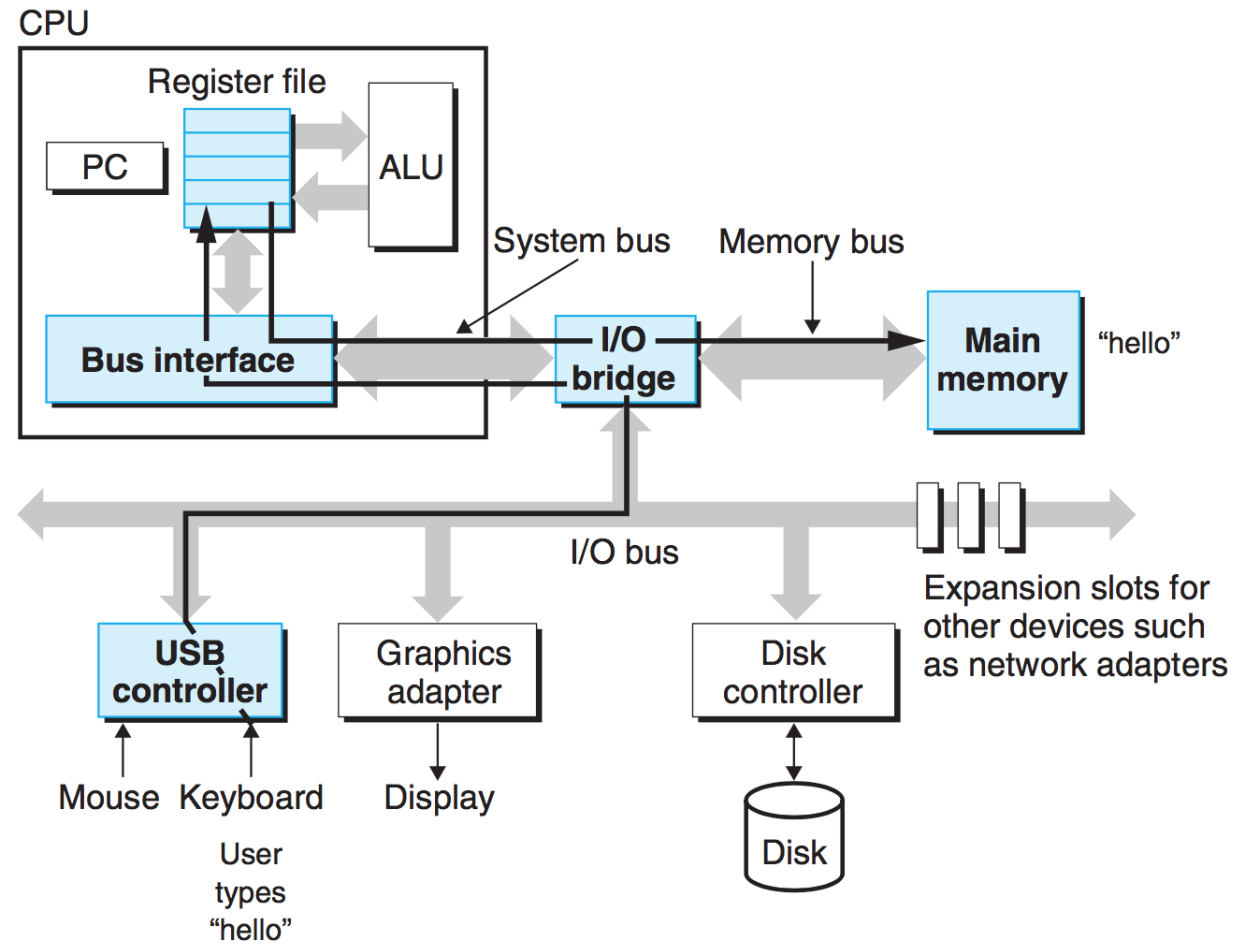
- Bus: words between HW components
- IO devices: keyboard, mouse, display, disk, network,
 - controller: transfer info. btwn IO bus and devices
 - DMA: copy data in disk device to memory without CPU
- Main memory (not virtual memory!):
 - DRAM chips (physically)
 - linear array of bytes (logically)
 - stores: 1. machine instructions, 2. C program variables
- Processor:
 - register: word-size storage
 - register PC: where to load instruction

- Processor (continued)
 - instruction execution model
 - appear to execute in sequence (actually pipelined, out-of-order)
 - instruction read by PC
 - Turing machine
 - internal:
 - ALU: arithmetic/logic unit
 - register file: a bunch of named registers
 - ISA
 1. Load/Store: a word from main memory to a register
 2. Operate: read register content to ALU, arithmetic op on two words, store result to a register
 3. Jump: overwriting PC

- Processor (continued 2)
 - Cache
 - process-memory gap: register >100 faster than memory
 - L1 size: 10KB
 - L2 size: 1MB
 - SRAM chip
 - Locality: program to access code/data in localized regions

Figure 1.5

Reading the hello command from the keyboard.

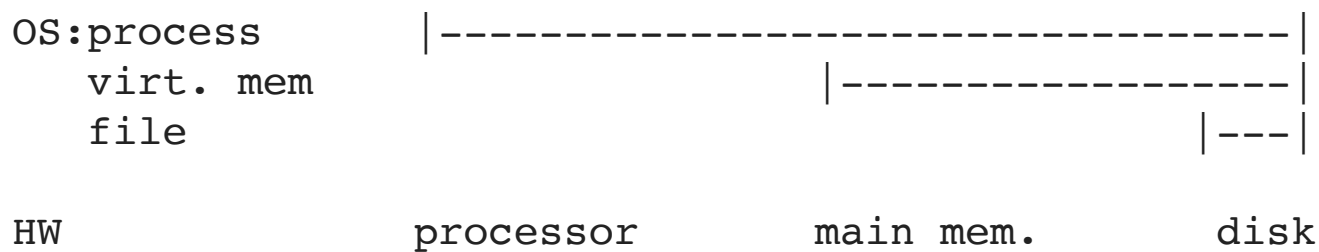


EXECUTION IN STEP 1,2 FROM HW PERSPECTIVE

1. Step 1: the shell read through keyboard IO characters (`./a.out`) into register, and then store them to main memory
2. Step 2: the shell read character `ENTER`, starts to load `a.out`
 - `load.o`: load executable file `a.out` by copy program code/data from file to main memory.
3. Step 3: What happens when executing `a.out`?
 - `a.out` is executed in a **process**, has its own address space (in **virtual memory**), and can print out string "helloworld" on screen by writing them to a **file**.

SYSTEMS/OS OVERVIEW

- OS goal:
 1. protecting HW from misuse by runaway app
 2. providing app easy way to manipulate HW
- fundamental abstractions: **process, virtual memory, file**



1. Process

- look like exclusive use of HW (no interrupt, only obj in memory)
- actually, run concurrently
 - instructions of different process interleaved
 - context switch: OS as mediator when switching from one process to another

2. Virtual memory

- virtual address space
 1. top-most region: kernel
 2. lower region: user

3. file: sequence of bytes

- read/write file through syscall (Unix IO)

+-----+	
kernel	
+-----+	+
user stack	
+-----+	
shared lib	dynamic sized
+-----+	
heap	
+-----+	+
program	
code/data	fixed size
+-----+	+

EXECUTION FROM SYSTEMS PERSPECTIVE

- Shell and `./hello` are two processes, run concurrently

SUMMARY (IN-CLASS EVALUATION)