

Internet Security

Public-Key Encryption and PKI

Public-Key Cryptography: History and Concept

1969 James Ellis — { No-secret key : encryption
secret key : decryption

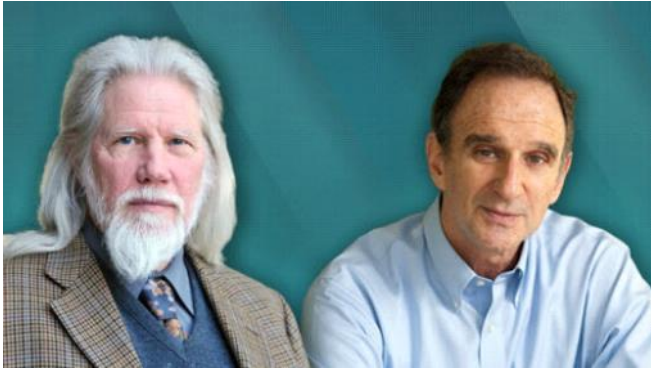
1976 { Whitfield Diffie
Martin Hellman (Stanford)

Diffie - Hellman Key Exchange ✓

1976 [Rivest
Shamir
Adleman] (MIT) \Rightarrow RSA

GCHQ Clifford Cocks 1973

Inventors of Public-Key Encryption



Whitfield Diffie

Martin Hellman

2015 Turing Award Winner



Leonard Adleman



Ron Rivest



Adi Shamir

2002 Turing Award Winner



Clifford Cocks

RSA Algorithm

public key (e, n)

private key (d, n)

3

$$\underline{65537} = 2^{16} + 1$$

Encryption: $\frac{M^e \bmod n}{\Rightarrow C}$

Decryption: $C^d \bmod n = M$

$$(M^e)^d \bmod n = \underline{M^{e \cdot d} \bmod n} = M$$

Euler theorem

For any $M < p$ and q (p, q are prime #s)

$$\Rightarrow M^{\underline{(p-1)(q-1)+1}} \bmod \underline{p \cdot q} = \underline{1} \cdot M$$

find p, q

$$\underline{n = p \cdot q}$$

$$e = \underline{65537}$$

$$\underline{e \cdot d = 1(p-1)(q-1) + 1}$$

$$e \cdot d = \underline{k(p-1)(q-1) + 1}$$

$$\underline{e \cdot d \bmod (p-1)(q-1) = 1}$$

$$M^{e \cdot d} \bmod n$$

$$= M^{k(p-1)(q-1)} \cdot M \bmod n$$

$$= (M^{(p-1)(q-1)})^k \cdot M \bmod n$$

$$= \cancel{(M^{(p-1)(q-1)})^k \bmod n} \cdot M \bmod n$$

find p, q ,

$$\underline{n = p \cdot q}$$

find e

find d , s.t.

$$\underline{e \cdot d \bmod (p-1)(q-1) = 1}$$

Extended

Euclidean Algorithm

$$a \cdot x \bmod n = 1$$
$$x = \underline{a^{-1} \bmod n}$$

$$n = 33$$
$$e = 17$$

$$\underline{M = 31}$$

Exercise Related to RSA

Let $n = 33$ and $e = 17$.

1. Find the private key d .
2. Encrypt the message $M = 31$.

Assume RSA is used.

For 2, you don't need to get the final numeric results; showing the expression is sufficient. ~~You do need to find the numeric value of the private key, though.~~

$$n = 33 = 3 \times 11$$

$$e \cdot d \mod (3-1)(11-1) = 1$$

$$e \cdot d \mod 20 = 1$$

$$\left\{ \begin{array}{l} M^{17} \mod 33 = C \\ C^{13} \mod 33 = M \end{array} \right.$$

$\nwarrow \frac{p-1}{q-1}$

$$17 \cdot d \mod 20 = 1$$

\downarrow
13

$$17 \cdot 13 = 221 \mod 20 = 1$$

Computing Using Big Numbers

$$(31^{17}, \text{mod } 33)$$

$$\underbrace{31 \cdot 31 \cdot 31 \cdot 31 \cdot 31 \dots 31}_{\text{mod } n}$$

$$\underbrace{(31^{17})}_{\text{mod } 33}$$

$$(31^2 \text{ mod } 33) \cdot 31^{15}$$

$$\underbrace{(1234567891011 \dots)}_{1024 \text{ bit}}$$

$$31^Z = (31^{32})^{\frac{Z}{32}} \cdot 31^{\frac{Z}{32} \cdot 32}$$

$$e^{2^{1024}} = (((e^2)^2)^2)^2$$

$$2^{1024} / 2^5$$

$$2^{1019}$$

$$(2^{1024})^2$$

$$(2^2)^2$$

Digital Signature

$$[Hash(M)]^d \pmod n = C$$

Signature

(e, n)

C^e

$$\{ \underline{M} \rightarrow \underline{C} \}$$

$$\underline{C}^e \pmod n \Rightarrow \underline{M}$$

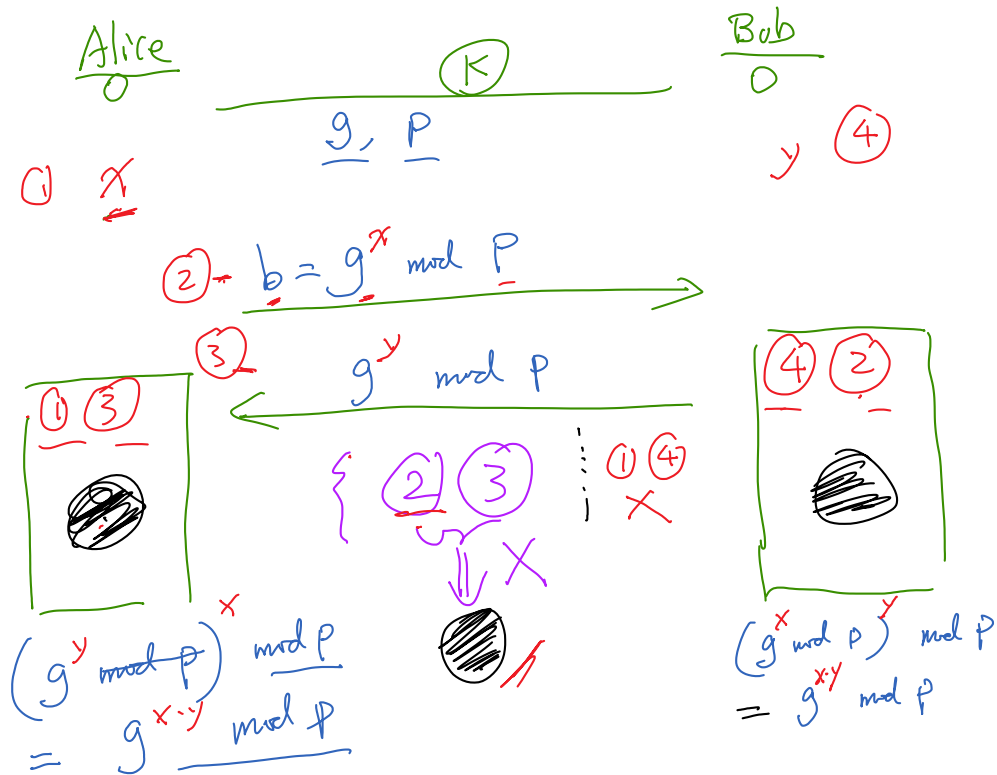
Letter: M

num

Signature
Verification

M, signature, (e, n)

Diffie-Hellman Key Exchange



✓

$$g^x = b$$

$$\log_2 g^x = \log_2 b$$

$$x \log_2 g = \log_2 b$$

$$x = \frac{\log_2 b}{\log_2 g}$$

Discrete logarithm problem

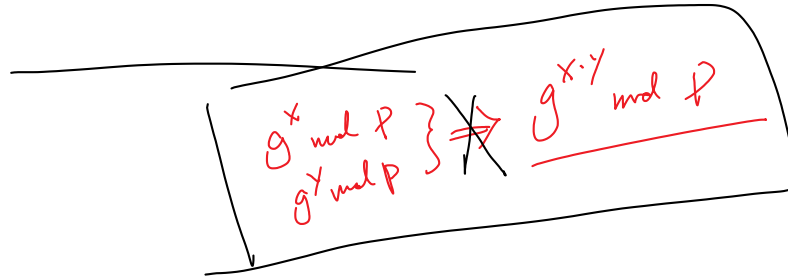
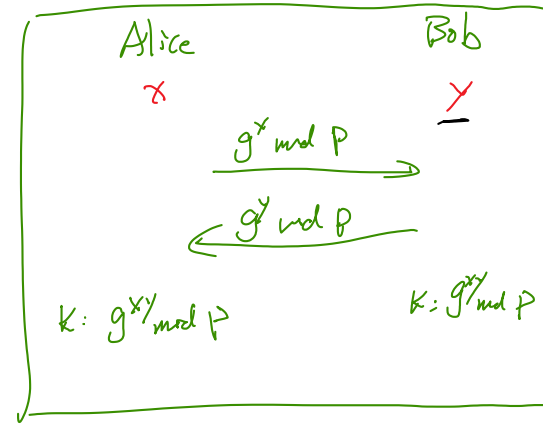
Given $a^x \text{ mod } n = b$
 solve x

Turn DH to Public-Key Encryption

Alice - public key: $g^x \bmod P$ ✓ ?
 private key: x ?

Bob.
 Encryption (M): $K \cdot M \bmod P$ $K = g^{x \cdot y} \bmod P$
~~AES (M)~~ $g^y \bmod P$ ✓

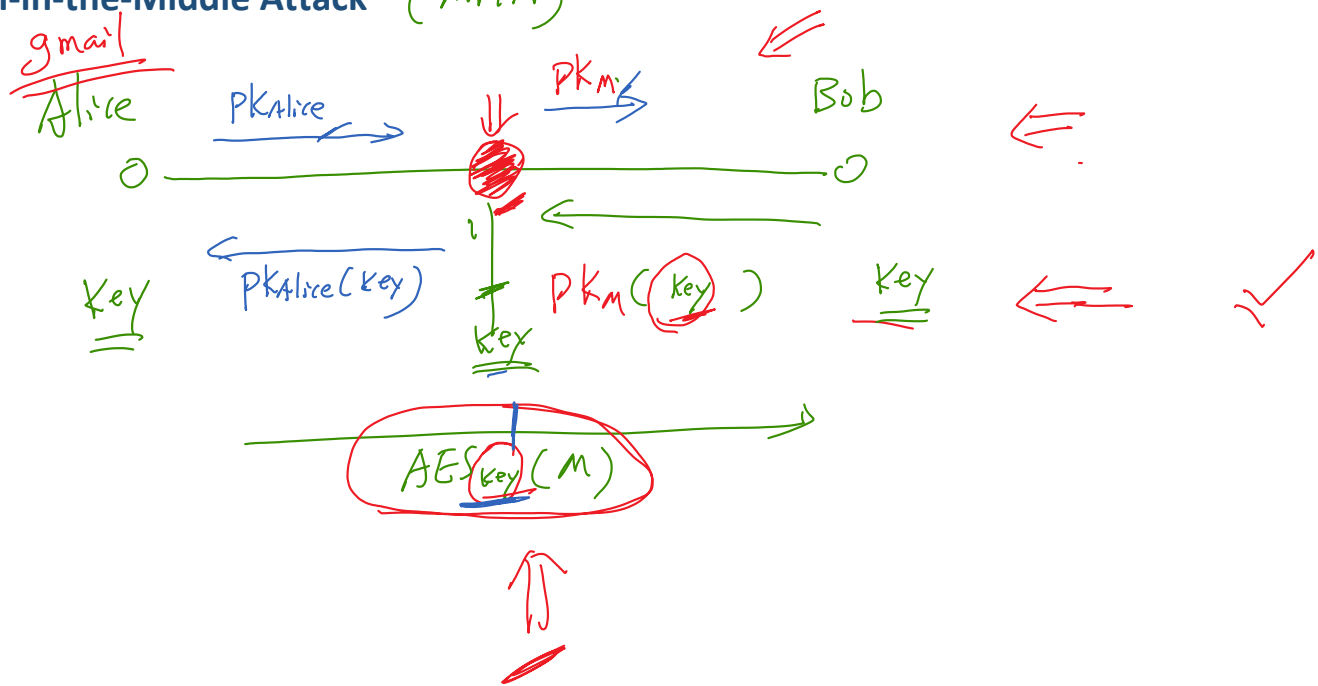
Alice.
 Decryption



$AES_K(M)$ $RSA(K)$

PUBLIC-KEY INFRASTRUCTURE (PKI)

Man-in-the-Middle Attack (MITM)



c key

15

Defeating the Man-in-the-Middle Attack Using Digital Signature

❖ Digital Signatures

owner. (M) sig Anybody can verify

RSA

$M^e \bmod n \Rightarrow$ encryption

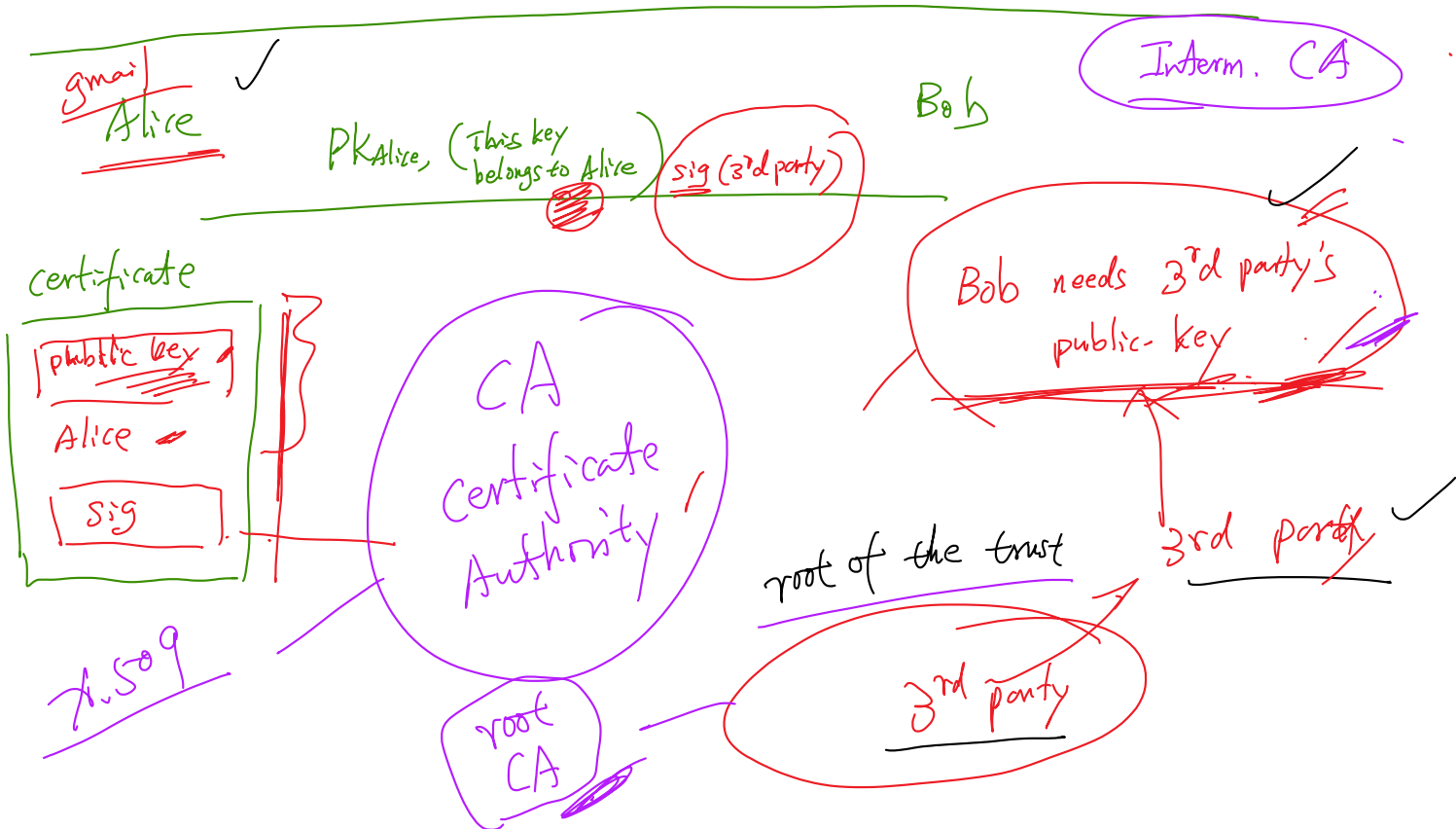
\Rightarrow everybody can do it
owner can decrypt

sig $M^d \bmod n \Rightarrow$ signature

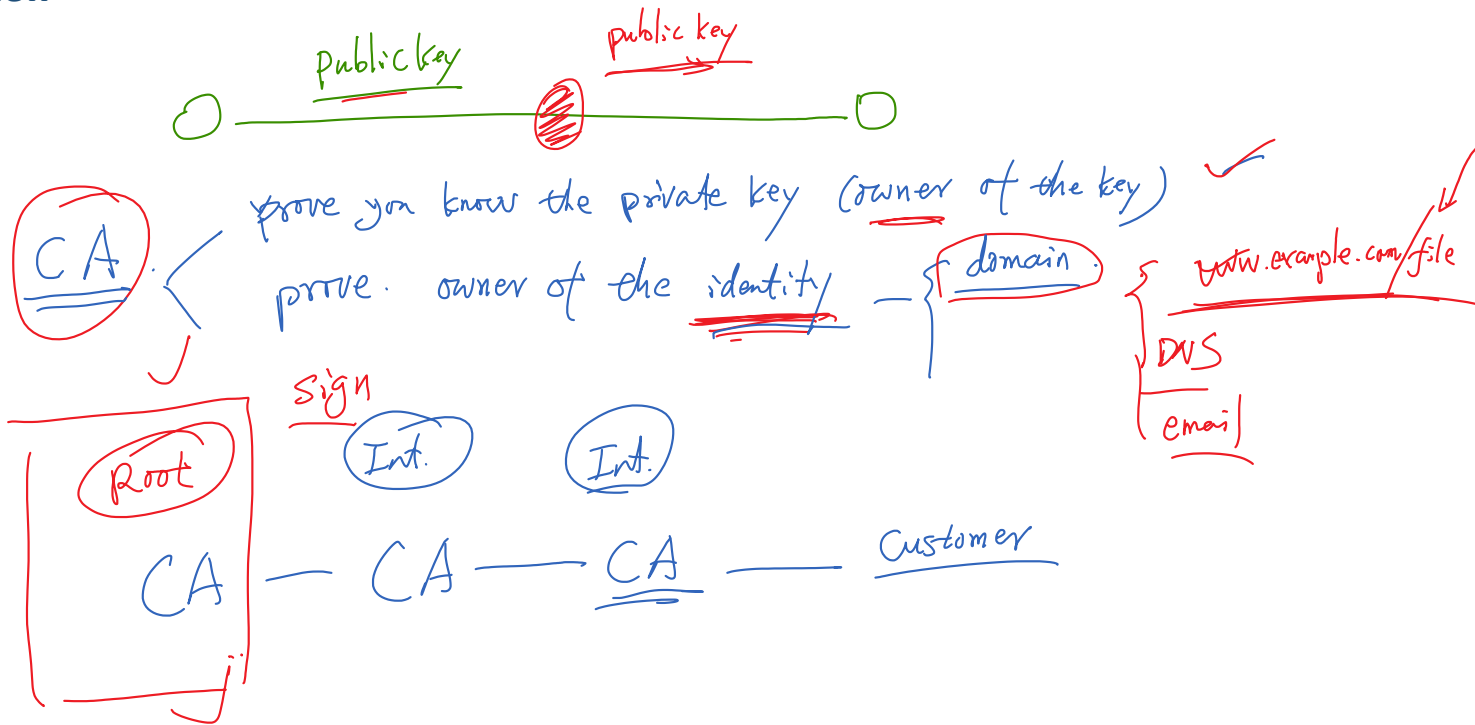
\Rightarrow only owner can do it
everybody can verify

(M, sig)

$$H(M) := (sig)^e \bmod n$$



Review



X.509 Certificate (paypal)

```
$ openssl x509 -in certificate.crt -text -noout
```

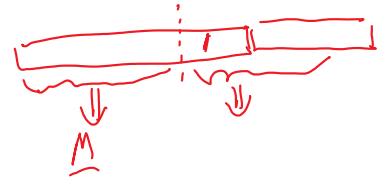
Certificate:
Data:
Version: 3 (0x2)
Serial Number:
2c:d1:95:10:54:37:d0:de:4a:39:20:05:6a:f6:c2:7f
Signature Algorithm: sha256WithRSAEncryption
Issuer: C=US, O=Symantec Corporation, OU=Symantec Trust Network,
CN=Symantec Class 3 EV SSL CA - G3
Validity
Not Before: Feb 2 00:00:00 2016 GMT
Not After : Oct 30 23:59:59 2017 GMT
Subject: 1.3.6.1.4.1.311.60.2.1.3=US/
1.3.6.1.4.1.311.60.2.1.2=Delaware/
businessCategory=Private Organization/
serialNumber=3014267, C=US/
postalCode=95131-2021, ST=California, L=San Jose/
street=2211 N 1st St, O=PayPal, Inc., OU=CDN Support,
CN=www.paypal.com
Subject Public Key Info:
Public Key Algorithm: rsaEncryption
Public-Key: (2048 bit)
Modulus:
00:da:43:c8:b3:a6:33:5d:83:c0:63:14:47:fd:6b:
22:bd:bf:4e:a7:43:11:55:eb:20:8b:e4:61:13:ee:
.....
00:c5:01:69:b5:10:16:a5:85:f8:fd:07:84:9a:c9:
14:91
Exponent: 65537 (0x10001)
Signature Algorithm: sha256WithRSAEncryption
4b:a9:64:20:cc:77:0b:30:ab:69:50:d3:7f:de:dc:7c:e2:fb:
93:84:fd:78:a7:06:e8:14:03:99:c0:e4:4a:ef:c3:5d:15:2a:
...
7d:6a:de:cb:9f:ff:ef:8c:65:35:e4:22:b5:88:b2:48:32:1e:
a4:71:a7:9e

-----BEGIN CERTIFICATE-----
MIIHWTCCBkGgAwIBAgIOLNGVEFQ30N5K0SAFavbCfzANBgkqhkiG9w0BAQsFADB3
MQswCQYDVQQGEwJVUzEdMBsGA1UEChMUU3ltYw50ZWNgQ29ycG9yYXRpb24xHzAd
BgNVBAsTF1N5bWVudGVjIFRydXN0IE5ldHdvcmxKDAwBgNVBAMTH1N5bWVudGVj
IENsYXNzIDMgRVYyU1NMIENBIC0gRzMwHhcNMTYwMjAyMDAwMDAwWhcNMTcxMDMw
MjM1OTUwUjCCAQkxZARBgsrBgEEAYI3PAIBAxMCMVVMxGTAXBgsrBgEEAYI3PAIB
.....
w3NLCcoN9KcCVKsPx70KwIgeYKaNe98YBdY9b4nw+KcJRzjZZIFJVIu7R53cf01
wv4AdQBo9pj4H2ScvjQM7rkoHUz8cVfDZ5PURNEKZ6y7T0/7xAAAaVKKVnLXAAAE
AwBGMEQCIHQpjXQ06Mf0V9DjzEnQm2CLPnu18P/LLyZrM6sEZvCNAiAziN0uyunX
wsaILVE7FMjg96sY02A0dsW/mGVPps7LJDANBgkqhkiG9w0BAQsFAA0CAQEAS6lk
IMx3CzCraVDTf97cf0L7k4T9eKcG6BQDmcDkSu/DXRUqgaG5/9w6r82A8HyPjh1X
BWLw0Zr6JZ87V8IxdYV/UQWQLRnnEp9yaRT/4f/fbS90bsQH3YmMbLds2I2zAIB
ZdZuwa0v/PAR29XusH8FY//HNR2I2wTXGg8Ztpad6KT9gIqFfHvfSZ8VDSU9IdjN
fDlUABWAm1B+nDxoZWlyvHHmmOgw6m4wm5ANFul1hjAWear/TLWd20lj7iXUt+dw
GN/QMQ3a55rjwNQnA3s2WwUHGPaE/jMG17iiL20/hUdIvLE9+wA+fWrey5//74x1
NeQitYiySDIepHGngg==
-----END CERTIFICATE-----

Binary form

Base64

256



CA's X.509 Certificate

```
Certificate:
Data:
  Version: 3 (0x2)
  Serial Number:
    7e:e1:4a:6f:ef:f2:d3:7f:3f:ad:65:4d:3a:da:b4
  Signature Algorithm: sha256WithRSAEncryption
  Issuer: C=US, O=VeriSign, Inc., OU=VeriSign Trust Network,
  OU=(c) 2006 VeriSign, Inc. - For authorized use only,
  CN=VeriSign Class 3 Public Primary Certification Authority - G5
  Validity
    Not Before: Oct 31 00:00:00 2013 GMT
    Not After : Oct 30 23:59:59 2023 GMT
  Subject: C=US, O=Symantec Corporation, OU=Symantec Trust Network,
  CN=Symantec Class 3 EV SSL CA - G3
  Subject Public Key Info:
    Public Key Algorithm: rsaEncryption
    Public-Key: (2048 bit)
    Modulus:
      00:d8:a1:65:74:23:e8:2b:64:e2:32:d7:33:37:3d:
      ...
      66:80:af:b3:2f:29:1d:23:b8:8a:e1:a1:70:07:0c:
      34:0f
    Exponent: 65537 (0x10001)
  Signature Algorithm: sha256WithRSAEncryption
    42:01:55:7b:d0:16:1a:5d:58:e8:bb:9b:a8:4d:d7:f3:d7:eb:
    ...
    86:4b:29:4c:e1:dc:b5:e1:e0:33:9d:b3:cb:36:91:4b:fe:a1:
    b4:ee:f0:f9
```

CA

Int. CA

Root CA's X.509 Certificate

```
Certificate:
Data:
  Version: 3 (0x2)
  Serial Number:
    18:da:d1:9e:26:7d:e8:bb:4a:21:58:cd:cc:6b:3b:4a
  Signature Algorithm: sha1WithRSAEncryption
  Issuer: C=US, O=VeriSign, Inc., OU=VeriSign Trust Network,
  OU=(c) 2006 VeriSign, Inc. - For authorized use only,
  CN=VeriSign Class 3 Public Primary Certification Authority - G5
  Validity
    Not Before: Nov  8 00:00:00 2006 GMT
    Not After : Jul 16 23:59:59 2036 GMT
  Subject: C=US, O=VeriSign, Inc., OU=VeriSign Trust Network,
  OU=(c) 2006 VeriSign, Inc. - For authorized use only,
  CN=VeriSign Class 3 Public Primary Certification Authority - G5
  Subject Public Key Info:
    Public Key Algorithm: rsaEncryption
    Public-Key: (2048 bit)
    Modulus:
      00:af:24:08:08:29:7a:35:9e:60:0c:aa:e7:4b:3b:
      ...
      9f:73:b8:33:0a:cf:5d:3f:34:87:96:8a:ee:53:e8:
      25:15
    Exponent: 65537 (0x10001)
  Signature Algorithm: sha1WithRSAEncryption
    93:24:4a:30:5f:62:cf:d8:1a:98:2f:3d:ea:dc:99:2d:bd:77:
    ...
    3f:68:5c:f2:42:4a:85:38:54:83:5f:d1:e8:2c:f2:ac:11:d6:
    a8:ed:63:6a
```

self-signed

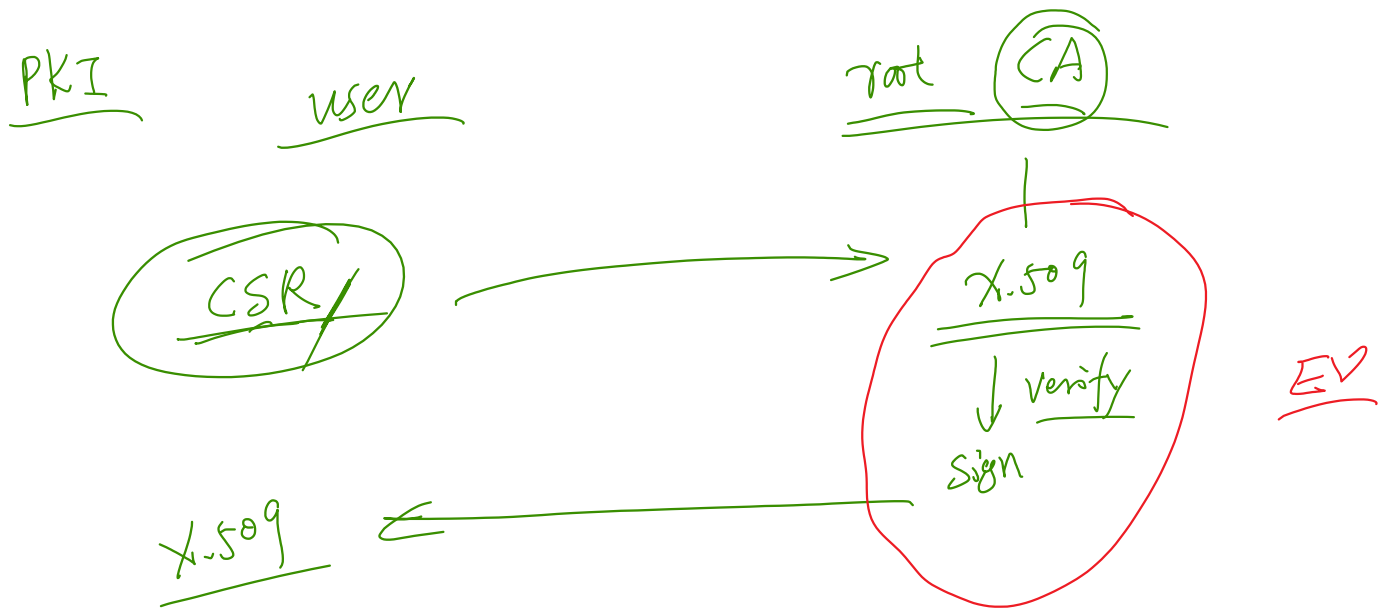
root - CA

Root Certificate Authority (CA)

Survey result on April 2016:

- Comodo Group: 40.6%
- Symantec: 26.0% market share
- GoDaddy: 11.8%
- GlobalSign: 9.7%

Getting X.509 Certificate from CA



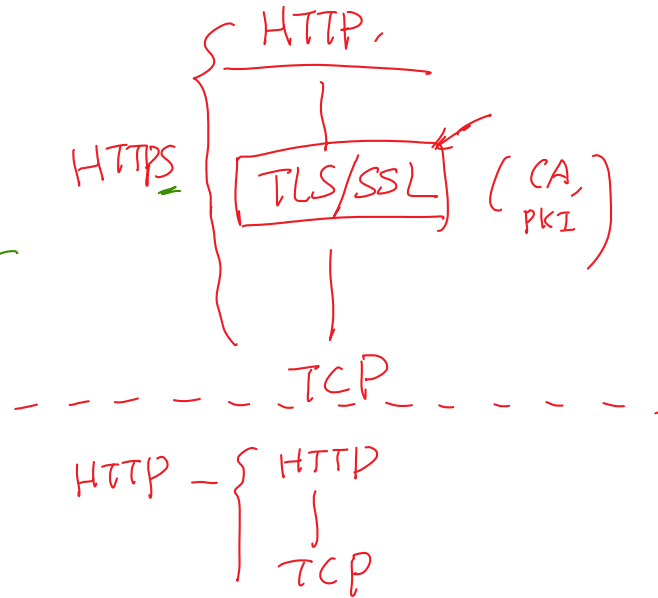
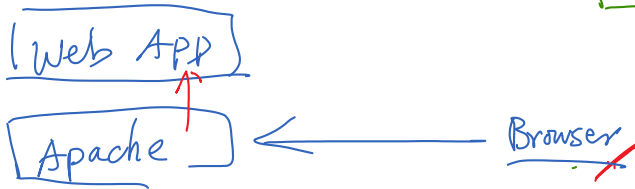
Set up HTTPS Web Server using X.509 Certificate

❖ Apache configuration file

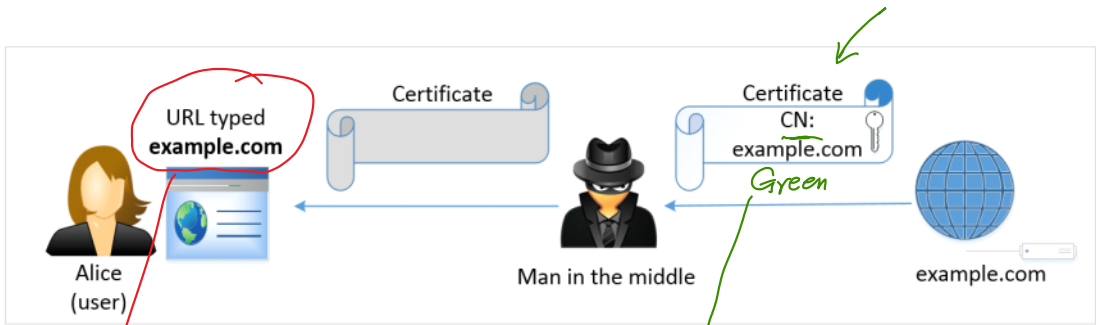
12.04

```
<VirtualHost *:443>
  ServerName example.com
  DocumentRoot /var/www/Example
  DirectoryIndex index.html

  SSLEngine On
  SSLCertificateFile      /etc/apache2/ssl/bank_cert.pem ①
  SSLCertificateKeyFile   /etc/apache2/ssl/bank_key.pem ②
</VirtualHost>
```



How PKI Defeats MITM Attacks



Host name \neq CN

certificate check

Attacker's public key

CN: attacker.com

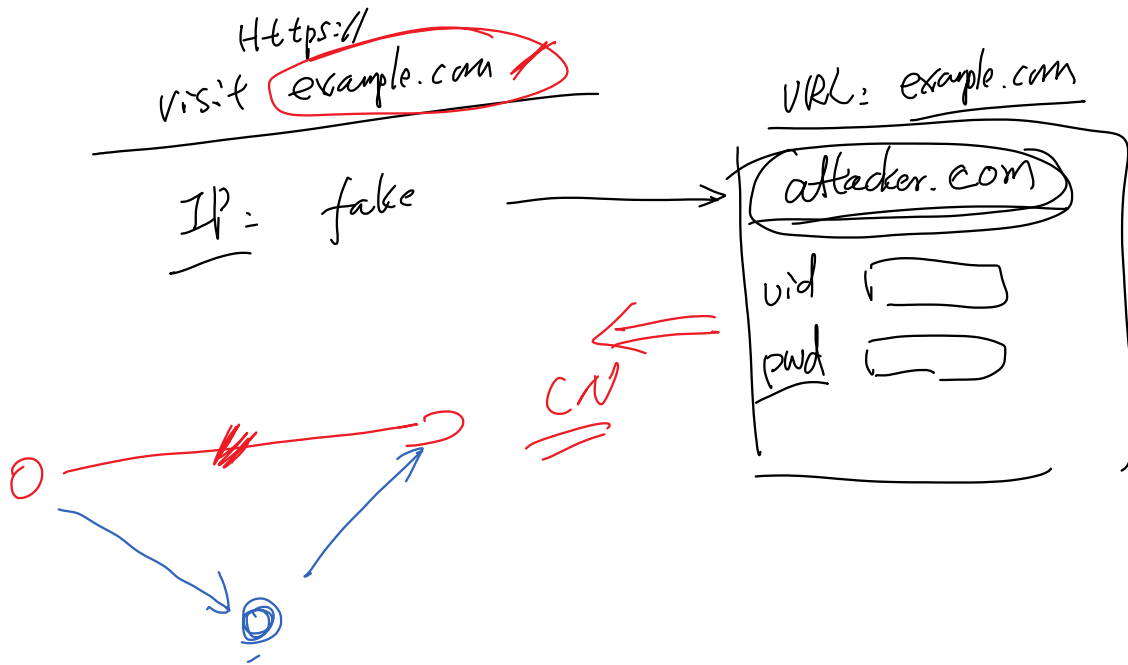
Sig

SV

UPN \leftarrow cert.

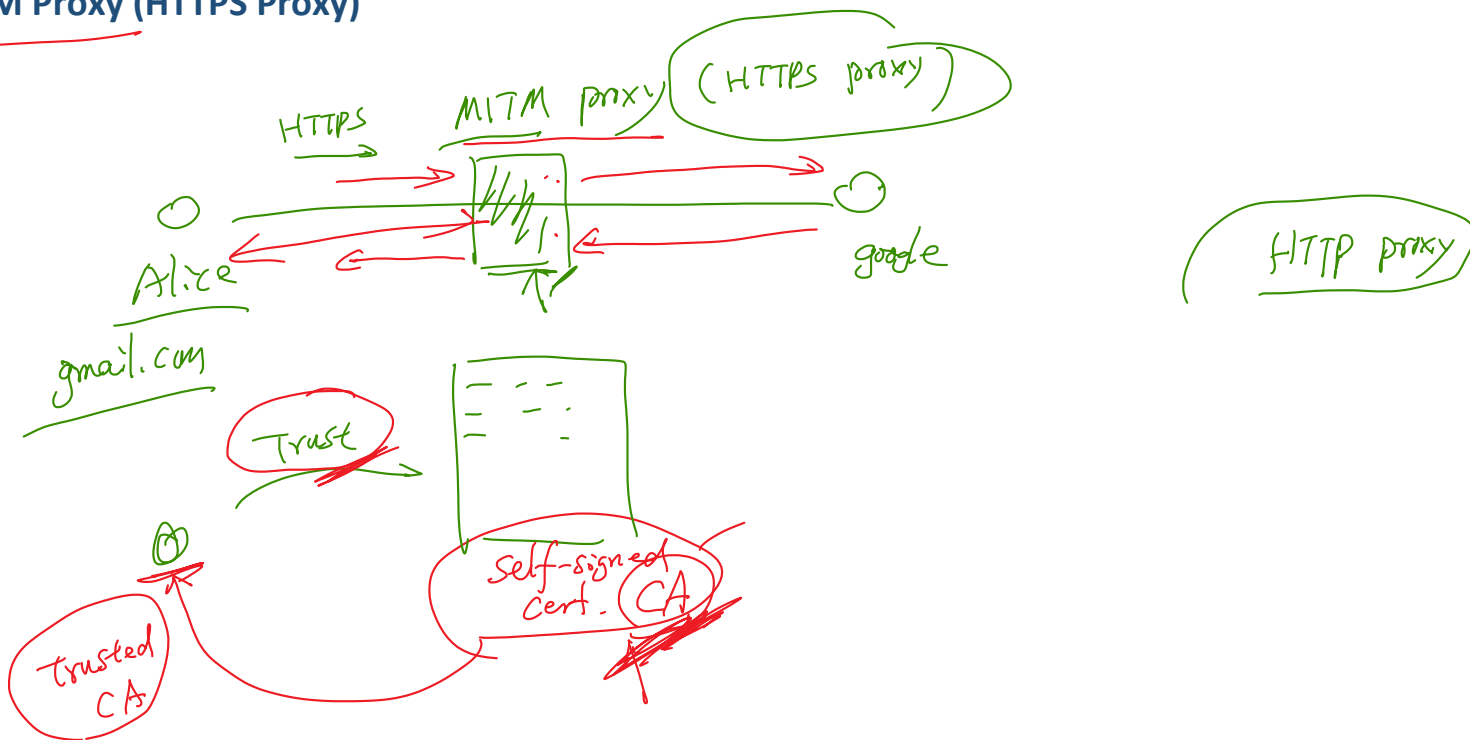
Question: DNS

For the DNS cache-poisoning attack (i.e., provide a fake IP address for a banking site), if the banking site uses HTTPS, can the attack still work?

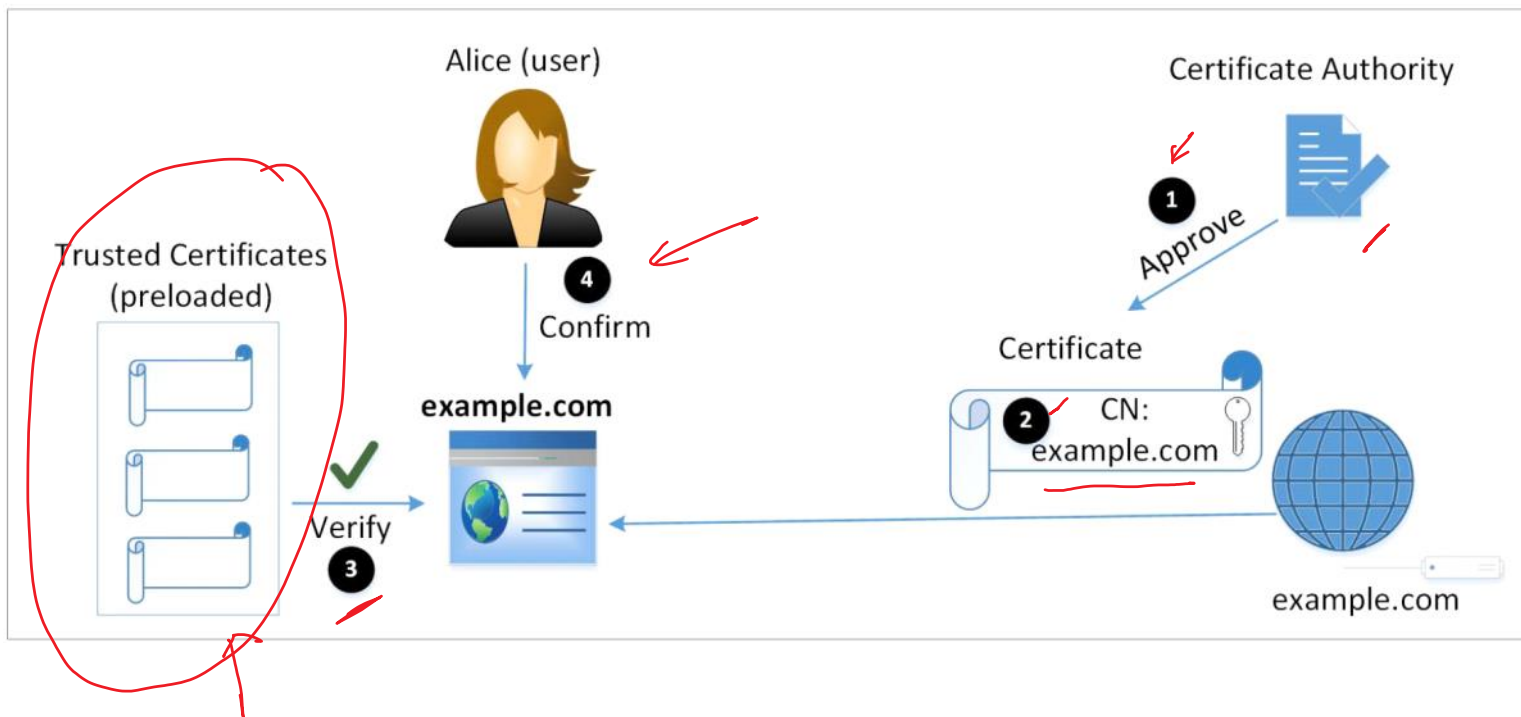


DASSEC

MITM Proxy (HTTPS Proxy)



Attacks on PKI ①②③④



Attack on CA ①

DigiNotar Case Study

DigiNotar B.V.³ was a Certificate Authority that provided digital certificate services. The digital certificates were used to secure Internet traffic, to issue (qualified) electronic signatures and to provide data encryption. DigiNotar also issued government accredited PKIoverheid certificates. During the months of June and July of 2011, the security of DigiNotar was breached and rogue certificates were issued. One of these certificates, a rogue Google certificate, was abused on a large scale in August of 2011 targeting primarily Iranian Internet users. At the end of August the intrusion became public knowledge and set into motion a chain of events that eventually led to the removal of all the Certificate Authorities that were hosted by DigiNotar from trust lists and ultimately the bankruptcy of the company.

"Using this [Gmail authentication] cookie, the hacker is able to log in directly to the Gmail mailbox of the victim and also read the stored emails," said Fox-IT. The hackers could also use the same credentials to log onto other Google services, including Google Docs and Google Latitude -- in the latter case, to identify the exact location of the victim -- and hijack [Facebook](#) and [Twitter](#) accounts.

Fox-IT said that approximately 300,000 IP addresses, each representing at least one computer and so at least one user, had accessed sites displaying a fake certificate for *google.com* between July 27 and Aug. 29. Nearly all -- Fox-IT said 99% -- of those IP addresses originated in Iran.

Investigators assumed that the *google.com* certificate was used primarily to spy on Iranians' Gmail accounts.

CNNIC Case Study

Google to drop China's CNNIC Root Certificate Authority after trust breach

Last month, a Chinese certificate authority issued valid security certificates for a number of domains, including Google's, without their permission, which resulted in a major trust breach in the crypto chain.



OWEN WILLIAMS

12 days ago



CNNIC had delegated its authority to Egyptian intermediary MCS Holdings to issue the certificates in question and the company installed it in a man-in-the-middle proxy internally.

Protecting CA



Comodo Case Study (March 2011)

Countermeasure: Certificate Pinning

Attack on Algorithm ②

The Importance of Collision Resistance

Certificate:

Data:

Version: 3 (0x2)

Serial Number:

2c:d1:95:10:54:37:d0:de:4a:39:20:05:6a:f6:c2:7f

Signature Algorithm: sha256WithRSAEncryption

Issuer: C=US, O=Symantec Corporation, OU=Symantec Trust Network,
CN=Symantec Class 3 EV SSL CA - G3

Validity

Not Before: Feb 2 00:00:00 2016 GMT

Not After : Oct 30 23:59:59 2017 GMT

Subject: 1.3.6.1.4.1.311.60.2.1.3=US/
1.3.6.1.4.1.311.60.2.1.2=Delaware/
businessCategory=Private Organization/
serialNumber=3014267, C=US/
postalCode=95131-2021, ST=California, L=San Jose/
street=2211 N 1st St, O=PayPal, Inc., OU=CDN Support,
CN=www.paypal.com

Subject Public Key Info:

Public Key Algorithm: rsaEncryption

Public-Key: (2048 bit)

Modulus:

00:da:43:c8:b3:a6:33:5d:83:c0:63:14:47:fd:6b:
22:bd:bf:4e:a7:43:11:55:eb:20:8b:e4:61:13:ee:
.....
00:c5:01:69:b5:10:16:a5:85:f8:fd:07:84:9a:c9:
14:91

Exponent: 65537 (0x10001)

Signature Algorithm: sha256WithRSAEncryption

4b:a9:64:20:cc:77:0b:30:ab:69:50:d3:7f:de:dc:7c:e2:fb:
93:84:fd:78:a7:06:e8:14:03:99:c0:e4:4a:ef:c3:5d:15:2a:
...
7d:6a:de:cb:9f:ff:ef:8c:65:35:e4:22:b5:88:b2:48:32:1e:
a4:71:a7:9e

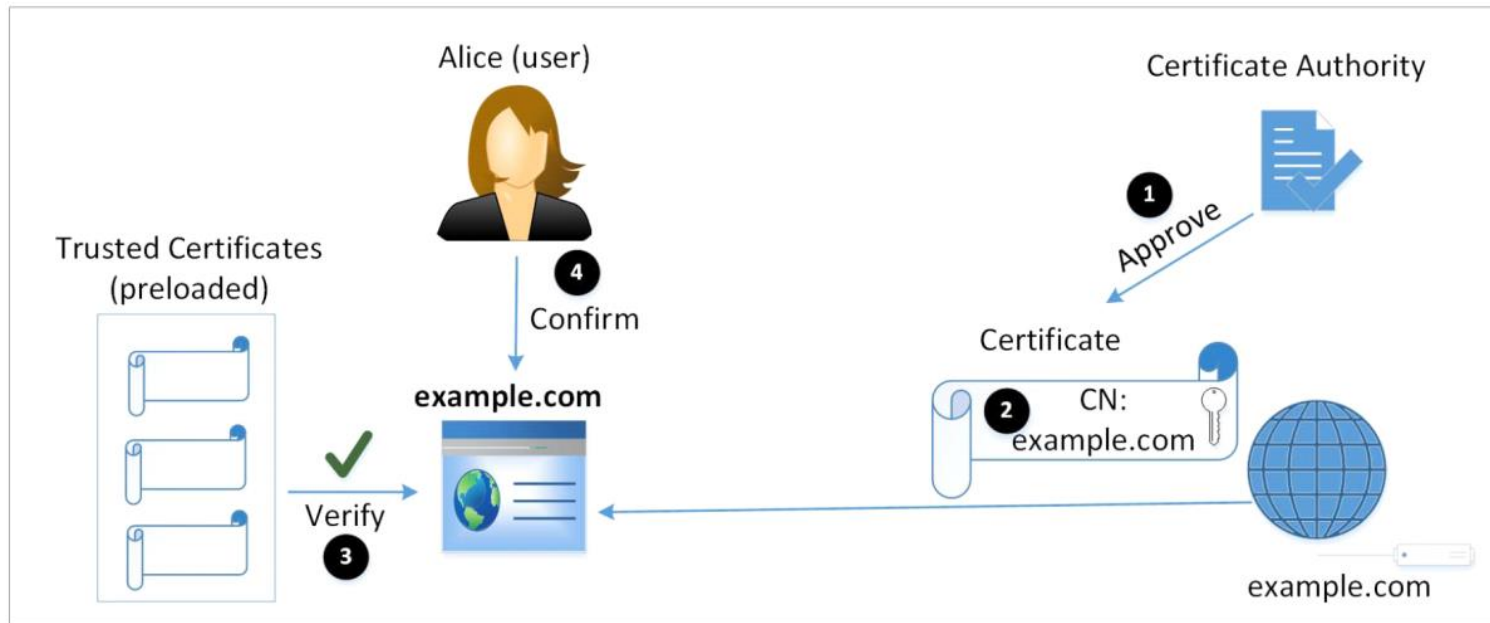
Question: Root CA's certificate and SHA1

Question: I notice that VeriSign's G5 certificate (self-signed) uses sha1, which is proven not to be collision resistant in February 2017, should VeriSign immediately revoke this certificate? Why or Why not?

```
Certificate:
  Data:
    Version: 3 (0x2)
    Serial Number:
      18:da:d1:9e:26:7d:e8:bb:4a:21:58:cd:cc:6b:3b:4a
    Signature Algorithm: sha1WithRSAEncryption
    Issuer: C=US, O=VeriSign, Inc., OU=VeriSign Trust Network,
      OU=(c) 2006 VeriSign, Inc. - For authorized use only,
      CN=VeriSign Class 3 Public Primary Certification Authority - G5
    Validity
      Not Before: Nov  8 00:00:00 2006 GMT
      Not After : Jul 16 23:59:59 2036 GMT
    Subject: C=US, O=VeriSign, Inc., OU=VeriSign Trust Network,
      OU=(c) 2006 VeriSign, Inc. - For authorized use only,
      CN=VeriSign Class 3 Public Primary Certification Authority - G5
    Subject Public Key Info:
      Public Key Algorithm: rsaEncryption
      Public-Key: (2048 bit)
      Modulus:
        00:af:24:08:08:29:7a:35:9e:60:0c:aa:e7:4b:3b:
        ...
        9f:73:b8:33:0a:cf:5d:3f:34:87:96:8a:ee:53:e8:
        25:15
      Exponent: 65537 (0x10001)
    Signature Algorithm: sha1WithRSAEncryption
      93:24:4a:30:5f:62:cf:d8:1a:98:2f:3d:ea:dc:99:2d:bd:77:
      ...
      3f:68:5c:f2:42:4a:85:38:54:83:5f:d1:e8:2c:f2:ac:11:d6:
      a8:ed:63:6a
```

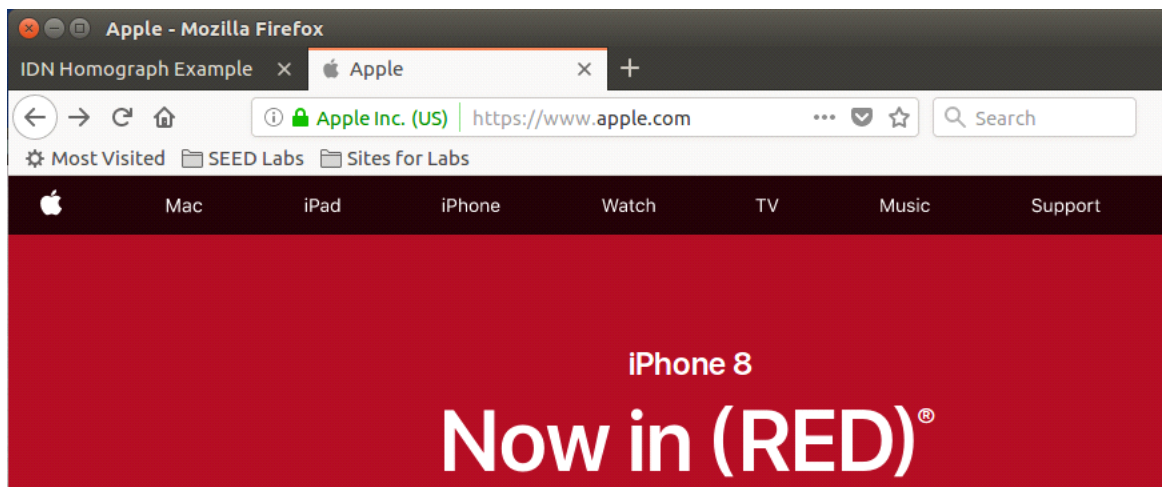
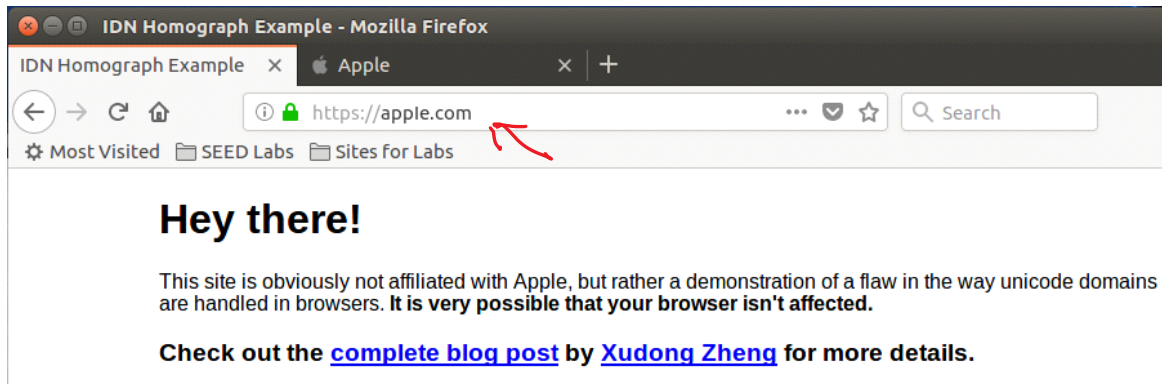
Attack on Trusted Certificates ③

Attack on User Confirmation ④

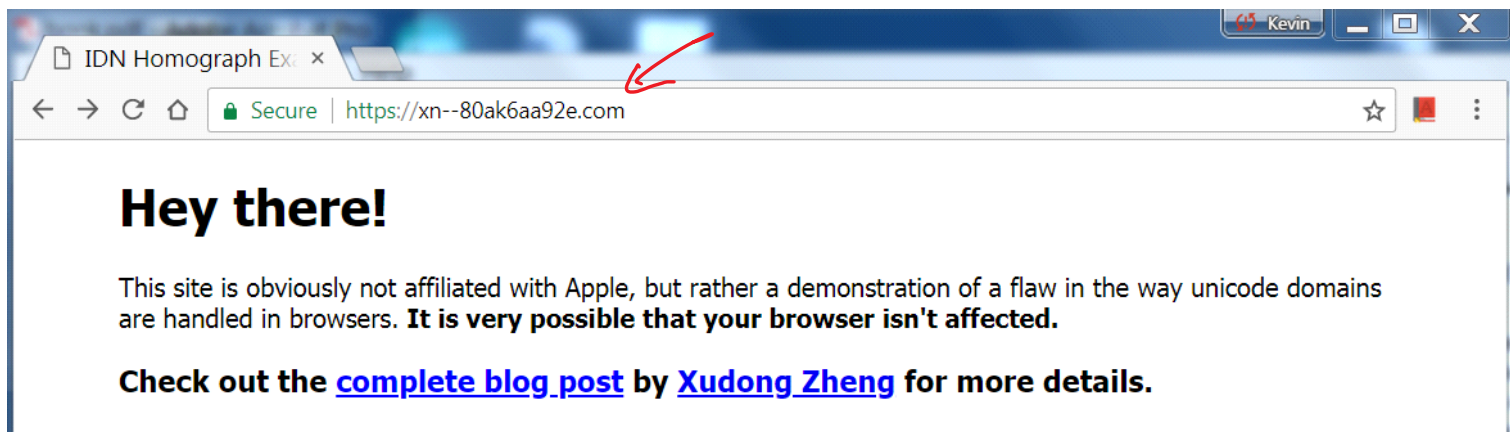


IDN Homograph Attack

❖ What you see (from Firefox)



❖ The actual name (from Chrome)



Types of X.509 Certificate

- DV: Domain Validated
- OV: Organization Validated
- EV: Extended Validated

 JPMorgan Chase and Co. (US) | <https://www.chase.com>

Issued To

Common Name (CN)	www.chase.com
Organization (O)	JPMorgan Chase and Co.
Organizational Unit (OU)	GTI GNS
Serial Number	62:5A:65:43:01:7A:7E:D1:2E:E4:46:20:39:0E:02:7C

Issued By

Common Name (CN)	Symantec Class 3 EV SSL CA - G3
Organization (O)	Symantec Corporation
Organizational Unit (OU)	Symantec Trust Network

 <https://www.amazon.com>

Issued To

Common Name (CN)	www.amazon.com
Organization (O)	Amazon.com, Inc.
Organizational Unit (OU)	<Not Part Of Certificate>
Serial Number	1D:4A:BD:AA:78:D0:9A:FE:79:9D:41:BC:EB:7A:76:62

Issued By




Common Name (CN)	Symantec Class 3 Secure Server CA - G4
Organization (O)	Symantec Corporation
Organizational Unit (OU)	Symantec Trust Network

Types of X.509 Certificates

Chrome browser

Cannot be verified	 Not secure https:// test-sspev.verisign.com:2443/test-SSPEV-revoked-verisign.html
DV/OV Certificate	 Secure https://www.microsoft.com/en-us/
EV Certificate	 PayPal, Inc. [US] https://www.paypal.com/us/home

Firefox browser

Cannot be verified	 https://test-sspev.verisign.com:2443/test-SSPEV-revoked-verisign.html
DV/OV Certificate	 https://www.microsoft.com/en-us/
EV Certificate	 PayPal, Inc. (US) https://www.paypal.com/us/home

Certificate Revocation List (CRL)

Summary

- ❖ Public key encryption concept
- ❖ Diffie-Hellman key exchange protocol
- ❖ RSA algorithm
- ❖ Man-in-the-middle attack
- ❖ Digital signature, X.509 certificate, and CA
- ❖ How PKI defeats MITM
- ❖ Attacks on PKI