

Assignment 3 : CSL 201
Topic : Red-black Trees

The following describes the outline of a template class implementing red-black trees:

```
template <class T>
class TreeNode {
friend class RedBlackTree<T>;
T x;
TreeNode<T> *left, *right, *parent;
int leftSize; // number of nodes in left subtree
int r; // the number of black nodes on any path
        // from the current node to a leaf in its subtree
bool color;
}

template <class T>
class RedBlackTree {
private:
TreeNode<T> *root;
}
```

Your objective is to add the following functions to the template class for red-black trees. Each function should run in $O(\log n)$ time, where n is the number of nodes in the red-black tree. The only exception is **ThreeWayJoin** which should run in $O(1 + |r(T_1) - r(T_2)|)$ time, where $r(T_1)$ and $r(T_2)$ denote the number of black nodes on any root to leaf path in trees T_1 and T_2 , respectively.

1. **Search.** Searches for a key in the red-black tree. Returns a pointer to the key if it is present, otherwise returns **NULL**.

```
T* search(T x) {};
```

2. **Insert.** Inserts a key into the red-black tree. If the key is already present, prints a “Key already present” message.

```
void insert(T x) {};
```

3. **Delete.** Deletes the item with a particular key from the red-black tree. If the key is not present, prints a “Key not present” message.

```
void delete(T x) {};
```

4. **k-th Smallest.** Given an integer k , returns the k th smallest element in the red-black tree.

```
T kthSmallest(int k) {};
```

Note that to implement this function in $O(\log n)$ time, we need to store the number of nodes in the left subtree at each node (field **leftSize** in class `TreeNode`).

5. **Three Way Join.** Given a key k and two red-black trees T_1 and T_2 such that all keys in tree T_1 are $< k$ and all keys in tree T_2 are $> k$, returns a red-black tree T containing all keys in $T_1 \cup T_2 \cup \{k\}$.

```
RedBlackTree& threeWayJoin(T k, RedBlackTree& t1, RedBlackTree& t2) {};
```

Note that to achieve this operation in $O(1 + |r(T_1) - r(T_2)|)$ time, you need to maintain the rank of each node i.e., the number of black nodes on any path from the given node to a leaf in the node's subtree (field **r** in `TreeNode`).

6. **Split.** Given a key k and a red-black tree T on which the function is called, makes two red-black trees T_1 and T_2 such that T_1 contains all keys in tree T which are $< k$, and T_2 contains all keys in T which are $> k$.

```
void split(T k, RedBlackTree& t1, RedBlackTree& t2) {};
```

Use field **r** in `TreeNode` and the function **threeWayJoin** to achieve this operation in $O(\log n)$ time.

For now, complete the above task and test it on your own inputs. The procedure for grading your assignment will be described later.

1 Input Format

Your algorithm will maintain a sequence of trees T_0, T_1, T_2, \dots and so on. Initially, all the trees are empty. You can assume that there will never be more than 10000 trees in the program.

The first line of the input will contain a single number n , which denotes the number of tree operations. Each line after that will contain a single instruction which can be an insert, a delete, a split, a join, or finding the k -th smallest element.

An insert instruction consists of the keyword "insert" followed by the number i of the tree T_i in which insertion is to be made. After this is a single number n , which denotes the number of integers to be inserted. This is followed by a sequence of n integers. Note that we allow batch insertions, i.e., a single insert instruction can be used to insert any number of items into a tree.

For example, to insert 10 numbers 45, 13, 11, 77, 2, 9, 1, 66, 5, 22 into tree number 4, we will give the instruction:

```
insert 4 \\ insert operation on tree number 4
10 \\ number of integers to be inserted
45 13 11 77 2 9 1 66 5 22 \\ the list of integers to be inserted
```

A deletion instruction is similar to an insert instruction, except that we use the keyword “delete”. For example, the following instruction deletes the set of numbers 7, 32, 41, 15, 16, 19, 22 from tree 10:

```
delete 10 \\ delete operation on tree number 10
7 \\ number of integers to be deleted
7 32 41 15 16 19 22 \\ list of integers to be deleted
```

The instruction of a three-way join, consists of the keyword “join” followed by the number x of the first tree, the key k , and the number y of the second tree. This is followed by the number z of the tree which is obtained by joining T_x, k , and T_y . After the join, trees x and y will be empty, and tree z will contain all the keys in $T_x \cup \{k\} \cup T_y$. You can assume that all keys in the first tree are less than k , and all keys in the second tree are greater than k .

For example, the instruction to join trees 5 and 8 with intermediate key 77 into tree 11 is:

```
join 5 77 8 11
```

The instruction for a split, consists of the keyword “split” followed by the number x of the tree to be split, the splitting key k , and the numbers y and z of the target trees. After the split, tree y consists of all keys in tree x that were less than k , tree z consists of all keys in tree x which were greater than k , and tree x itself becomes empty.

Thus, the instruction to split tree 5 about key 61 with target trees 11 and 17 is:

```
split 5 61 11 17
```

The instruction for k -th smallest consists of the keyword “select” followed by the tree number and the rank of the key to be returned. To find the 11th smallest key in tree 7, we will write the following:

```
select 7 11
```

The output of “select” instruction is a single number on a line by itself, which is the value of the k -th smallest key in the corresponding red-black tree.

Note that all other instructions except “select” have no output.

We are giving six input test files $in0, in1, in2, \dots, in5$. The description of each test case along with the correct output is given in the file “test_cases”.

2 Example

Suppose your program is given the following input:

```
7
insert 0 5 1 3 5 7 9
insert 1 5 11 14 15 16 17
delete 1 2 14 16
join 0 10 1 2
split 2 7 3 4
select 3 2
select 4 4
```

The first number is equal to 7, and tells us that there are 7 instructions in the input. The first instruction inserts 5 numbers into tree 0. After this, $T_0 = (1\ 3\ 5\ 7\ 9)$. The second instruction inserts 5 numbers into tree 1. After this, $T_1 = (11\ 14\ 15\ 16\ 17)$. The third instruction deletes two numbers 14 and 16 from tree 1. After this, $T_1 = (11\ 15\ 17)$. The fourth instruction joins trees 0 and 1 along with intermediate key 10 into tree 2. After this instruction trees 0 and 1 are empty, and tree 2 contains the keys $(1\ 3\ 5\ 7\ 9\ 10\ 11\ 15\ 17)$. The fifth instruction splits tree 2 about key 7 into trees 3 and 4. After this instruction, trees 0, 1, and 2 are empty, and $T_3 = (1\ 3\ 5)$ and $T_4 = (9\ 10\ 11\ 15\ 17)$. The sixth instruction asks for the 2nd smallest key in tree 3. The output of this instruction is the number 3 on a line by itself. The seventh instruction asks for the 4th smallest key in tree 4. The output of this instruction is the number 15 on a line by itself.

The final output of the program consists of just two numbers, which are the outcome of the two “select” instructions:

```
3
15
```

Note that all instructions except “select” have no output.