

Homework 4 - Logic Circuits and Number Representations

Problem 4.1

Solution:

A	B	C	$A \oplus B$	$B \odot C$	X
0	0	0	0	1	0
0	0	1	0	0	0
0	1	0	1	0	0
0	1	1	1	1	1
1	0	0	1	1	0
1	0	1	1	0	0
1	1	0	0	0	0
1	1	1	0	1	0

As shown in the truth table above, we can see that the input condition needed to produce X as 1 is A as 0, B as 1, and C as 1. After computing the XOR and XNOR gates, we just compare those two columns with the C values and use the AND gate to produce the results.

Problem 4.2

Solution:

a.

A	B	C	$A \wedge B$	$B \vee C$	$(B \vee C) \wedge \bar{A}$	$\overline{(A \wedge B)} \vee C$	Y
0	0	0	0	0	0	1	1
0	0	1	0	1	1	0	1
0	1	0	0	1	1	1	0
0	1	1	0	1	1	0	1
1	0	0	0	0	0	1	1
1	0	1	0	1	0	0	0
1	1	0	1	1	0	0	0
1	1	1	1	1	0	0	0

$$Y = ((B \vee C) \wedge \bar{A}) \oplus (\overline{(A \wedge B)} \vee C)$$

b.

$$\begin{aligned} Y &= \overline{ABC} \vee \overline{ABC} \vee \overline{ABC} \vee \overline{ABC} \\ &= \overline{AB} \wedge (\overline{C} \vee C) \vee \overline{ABC} \vee \overline{ABC} \\ &= \overline{AB} \vee 1 \vee \overline{ABC} \vee \overline{ABC} \end{aligned}$$

Problem 4.3

Solution:

a. $+27_{10}$

$$\begin{aligned} 27 \div 2 &= 13 + 1 \\ 13 \div 2 &= 6 + 1 \\ 6 \div 2 &= 3 + 0 \\ 3 \div 2 &= 1 + 1 \\ 1 \div 2 &= 0 + 1 \end{aligned}$$

Due to the fact that we always have to use 8 bits, the answer is 00011011_2

b. $+66_{10}$

$$\begin{aligned} 66 \div 2 &= 33 + 0 \\ 33 \div 2 &= 16 + 1 \\ 16 \div 2 &= 8 + 0 \\ 8 \div 2 &= 4 + 0 \\ 4 \div 2 &= 2 + 0 \\ 2 \div 2 &= 1 + 0 \\ 1 \div 2 &= 0 + 1 \end{aligned}$$

The answer is: 01000010_2

c. -18_{10}

$$18 \div 2 = 9 + 0$$

$$9 \div 2 = 4 + 1$$

$$4 \div 2 = 2 + 0$$

$$2 \div 2 = 1 + 0$$

$$1 \div 2 = 0 + 1$$

$$+18_{10} = 00010010_2$$

$$-18_{10} = 11101101_2 \text{ (One's Complement)}$$

$$-18_{10} = 11101110_2 \text{ (Two's Complement)}$$

d. 127_{10}

$$127 \div 2 = 63 + 1$$

$$63 \div 2 = 31 + 1$$

$$31 \div 2 = 15 + 1$$

$$15 \div 2 = 7 + 1$$

$$7 \div 2 = 3 + 1$$

$$3 \div 2 = 1 + 1$$

$$1 \div 2 = 0 + 1$$

The answer is: 01111111_2

e. -127_{10}

$$+127_{10} = 01111111_2$$

$$-127_{10} = 10000000_2 \text{ (One's Complement)}$$

$$-127_{10} = 10000001_2 \text{ (Two's Complement)}$$

f. $+131_{10}$

$$131 \div 2 = 65 + 1$$

$$65 \div 2 = 32 + 1$$

$$32 \div 2 = 16 + 0$$

$$16 \div 2 = 8 + 0$$

$$8 \div 2 = 4 + 0$$

$$4 \div 2 = 2 + 0$$

$$2 \div 2 = 1 + 0$$

$$1 \div 2 = 0 + 1$$

The answer is: 10000011_2

g. -7_{10}

$$7 \div 2 = 3 + 1$$

$$3 \div 2 = 1 + 1$$

$$1 \div 2 = 0 + 1$$

$$7_{10} = 00000111_2$$

$$-7_{10} = 11111000_2 \text{ (One's Complement)}$$

$$-7_{10} = 11111001_2 \text{ (Two's Complement)}$$

Problem 4.4

Solution:

a. 00011000_2

$$(0 \times 2^7) + (0 \times 2^6) + (0 \times 2^5) + (1 \times 2^4) + (1 \times 2^3) + (0 \times 2^2) + (0 \times 2^1) + (0 \times 2^0)$$

$$= 8 + 16$$

$$= 24_{10}$$

b. 11110101_2

We must now subtract one from the binary number since we are going backwards

$$= 11110100_2 \text{ (One's Complement)}$$

$$\begin{aligned}
&= 00001011_2 \text{ (Invert)} \\
&= (0 \times 2^7) + (0 \times 2^6) + (0 \times 2^5) + (0 \times 2^4) + (1 \times 2^3) + (0 \times 2^2) + (1 \times 2^1) + (1 \times 2^0) \\
&= 8 + 2 + 1 \\
&= -11_{10}
\end{aligned}$$

$$\begin{aligned}
\text{c. } &01011011_2 \\
&(0 \times 2^7) + (1 \times 2^6) + (0 \times 2^5) + (1 \times 2^4) + (1 \times 2^3) + (0 \times 2^2) + (1 \times 2^1) + (1 \times 2^0) \\
&= 64 + 16 + 8 + 2 + 1 \\
&= 91_{10}
\end{aligned}$$

$$\begin{aligned}
\text{d. } &10110110_2 \\
&= 10110101_2 \text{ (One's Complement)} \\
&= 01001010_2 \text{ (Invert)} \\
&= (0 \times 2^7) + (1 \times 2^6) + (0 \times 2^5) + (0 \times 2^4) + (1 \times 2^3) + (0 \times 2^2) + (1 \times 2^1) + (0 \times 2^0) \\
&= 64 + 8 + 2 \\
&= -74_{10}
\end{aligned}$$

$$\begin{aligned}
\text{e. } &11111111_2 \\
&= 11111110_2 \text{ (One's Complement)} \\
&= 00000001_2 \text{ (Invert)} \\
&= (0 \times 2^7) + (0 \times 2^6) + (0 \times 2^5) + (0 \times 2^4) + (0 \times 2^3) + (0 \times 2^2) + (0 \times 2^1) + (1 \times 2^0) \\
&= -1_{10}
\end{aligned}$$

$$\begin{aligned}
\text{f. } &01101111_2 \\
&(0 \times 2^7) + (1 \times 2^6) + (1 \times 2^5) + (0 \times 2^4) + (1 \times 2^3) + (1 \times 2^2) + (1 \times 2^1) + (1 \times 2^0) \\
&= 64 + 32 + 8 + 4 + 2 + 1 \\
&= 111_{10}
\end{aligned}$$

$$\begin{aligned}
\text{g. } &10000001_2 \\
&= 10000000_2 \text{ (One's Complement)} \\
&= 01111111_2 \text{ (Invert)} \\
&= (0 \times 2^7) + (1 \times 2^6) + (1 \times 2^5) + (1 \times 2^4) + (1 \times 2^3) + (1 \times 2^2) + (1 \times 2^1) + (1 \times 2^0) \\
&= 64 + 32 + 16 + 8 + 4 + 2 + 1 \\
&= -127_{10}
\end{aligned}$$

$$\begin{aligned}
\text{h. } &10000000_2 \\
&= 01111111_2 \text{ (One's Complement)} \\
&= 10000000_2 \text{ (Invert)} \\
&= (1 \times 2^7) + (0 \times 2^6) + (0 \times 2^5) + (0 \times 2^4) + (0 \times 2^3) + (0 \times 2^2) + (0 \times 2^1) + (0 \times 2^0) \\
&= -128_{10}
\end{aligned}$$

Problem 4.5

Solution:

$$\begin{aligned}
\text{a. } &(27 + 36)_{\text{BCD}} \\
&= 0010\ 0111 + 0011\ 0110 \\
&= 0110\ 1101 \text{ (The binary on the right creates a number greater than 9, which is why 6 (or 0110) is going to be added)} \\
&= 0110\ (1101 + 0110) \\
&= 0110\ 0100_{\text{BCD}}
\end{aligned}$$

$$\begin{aligned}
\text{b. } &(73 + 29)_{\text{BCD}} \\
&= 0111\ 0011 + 0010\ 1001 \\
&= 1001\ 1100 \text{ (Right side creates a number greater than 9, so 0110 is added)} \\
&= 1001\ (1100 + 0110) \\
&= 1010\ 0010 \text{ (However, now the left side is greater than 9 after 0110 was added to the right side, therefore now 0110 has to be added to the left side as well)} \\
&= 0001\ 0000\ 0010_{\text{BCD}}
\end{aligned}$$

Problem 4.6

Solution:

- The range of unsigned decimal numbers that can be represented by using 8 bits is between 0 and 255.
- The range of signed decimal numbers that can be represented by using 8 bits is between -128 and 127.
- The range of unsigned decimal numbers that can be represented by using 11 bits is between 0 and 2407.
- The range of signed decimal numbers that can be represented by using 11 bits is between -1024 and 1023.
- The range of signed decimal numbers that can be represented by using 16 bits is between -32768 and 32767.