# Trees and binary trees

Ch. 5

### Tree

- A finite set of one or more nodes.
  - A node called *root*.
  - Remaining nodes are partitioned into disjoint sets, called subtrees.

A recursive definition

**Degree of a node**: number of subtrees **LEVEL** root **Leaf**: nodes have degree zeros Internal node: not leaf and not root Subtree  $T_1$ degree=3 1 Subtree T<sub>3</sub> Subtree T<sub>2</sub> degree=2 B 2 degree=1 degree=0 3 leaf leaf leaf leaf degree=0 4 leaf 1eaf

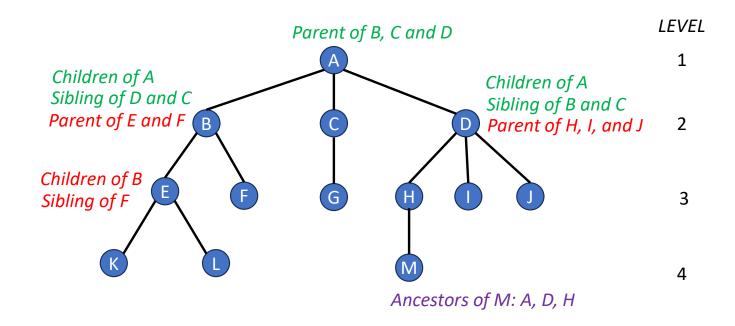
### Tree

Children: roots of subtrees of a node X

Parent: X is parent

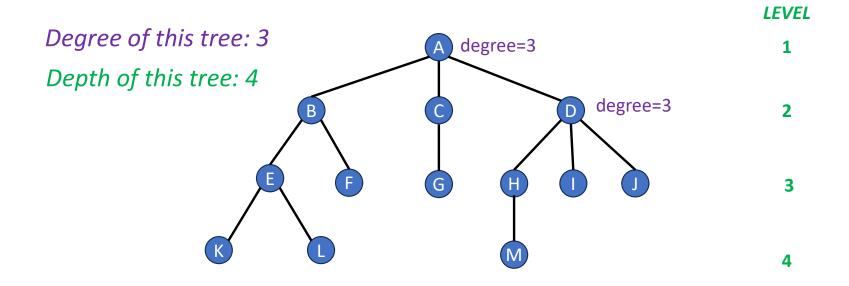
• Sibling: children of the same parent

Ancestors: all the nodes along the path from the root to the node.



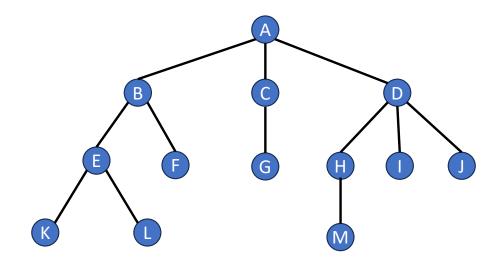
### Tree

- Degree of a tree: maximum of the degree of the nodes in the tree
- Height (depth) of a tree: maximum level of any node in the tree



## List representation

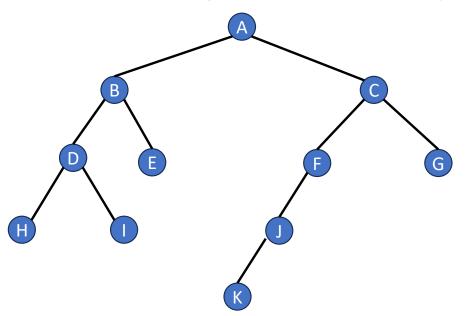
- Root node comes first
- Followed by a list of subtrees of that node



(A (B (E (K, L), F), C (G), D (H (M), I, J)))

### Binary tree

- Degree-two tree
- Recursive definition:
  - Each node has left subtree and/or right subtree.
  - Left subtree (or right subtree) is a binary tree.





Q6: The degree of the tree

Q7: The depth of the tree

Q8: List the leaf nodes

Q9: The sibling of node E

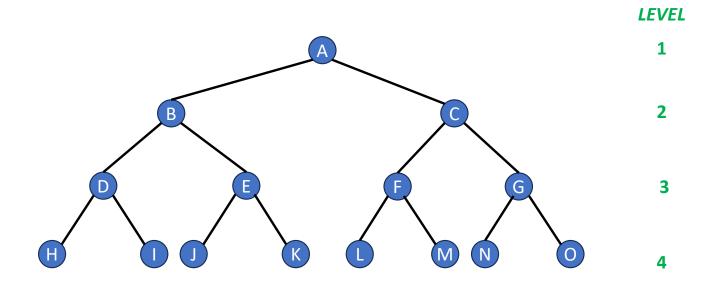
Q10: The children of node C

Q11: The level of node K

## Full binary tree

Maximum number of nodes =  $1 + 2 + 4 + 8 + ... + 2^{h-1}$ =  $2^h - 1$ 

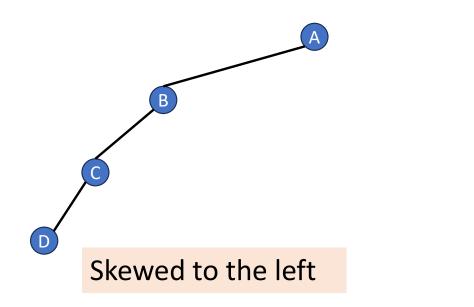
• A full binary tree of a given height h has  $2^h-1$  nodes.

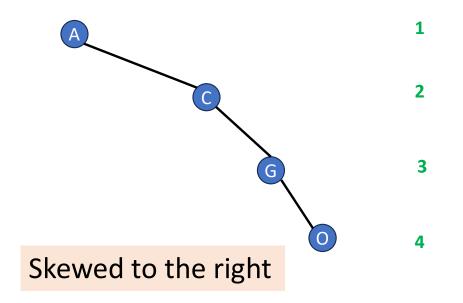


Height 4 full binary tree.

### Minimum number of nodes

- At least one node at each of first h levels
- Minimum number of nodes for a height h tree: h
- Special case of binary tree: skewed tree





**LEVEL** 

### Number of nodes and height

Let n be the number of nodes in a binary tree whose height is h.

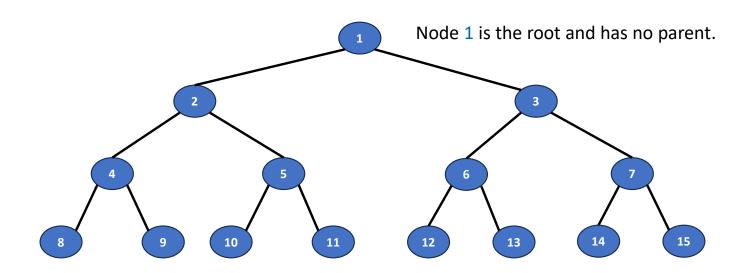
Minimum number of nodes 
$$h \le n \le 2^{h}-1$$
 Maximum number of nodes

$$Log_2(n+1) \le h$$

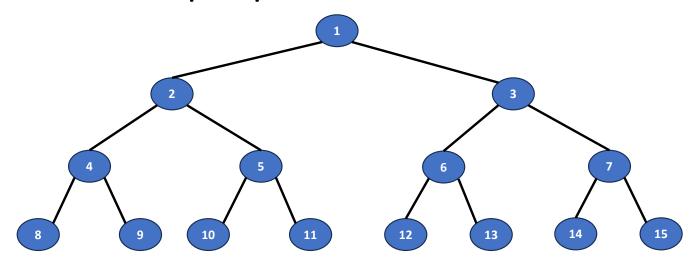
The height h of a binary tree is at least  $log_2(n+1)$ .

### Numbering nodes in a full binary tree

- Sequentially numbering the nodes 1 through  $2^h-1$ .
  - Number by levels from top to bottom.
  - Within a level number from left to right.



## Node number properties

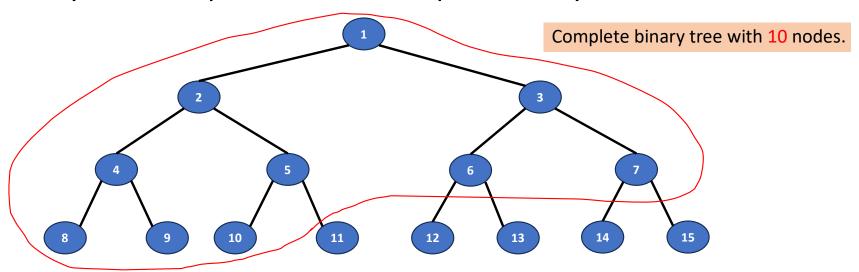


- Let *n* be the number of nodes in a full binary tree
  - Parent of node i is node floor(i / 2)
  - Left child of node *i* is node 2i
  - Right child of node *i* is node 2i+1

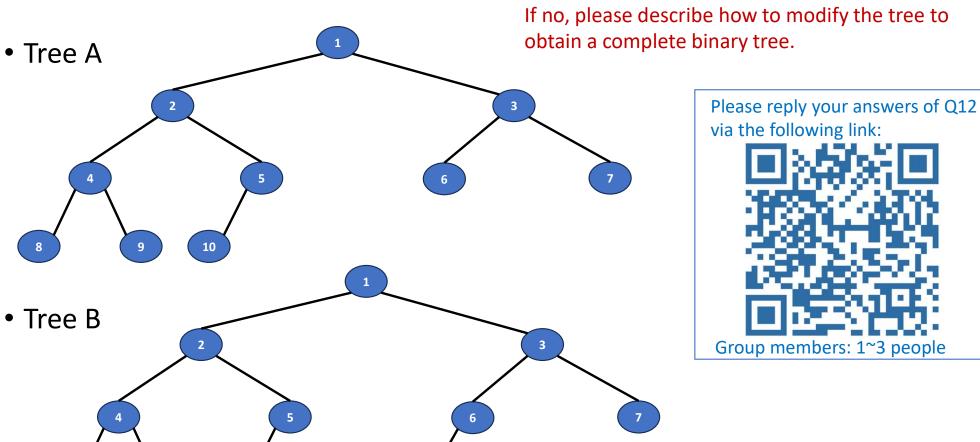
## Complete binary tree with *n* nodes

#### • Steps:

- Create a full binary tree has at least n nodes.
- 2. Number the nodes sequentially.
- 3. The binary tree defined by the node numbered 1 through n is the n-node complete binary tree.
- Full binary tree is a special case of complete binary tree.



## Q12: Are Tree A and Tree B complete binary trees?

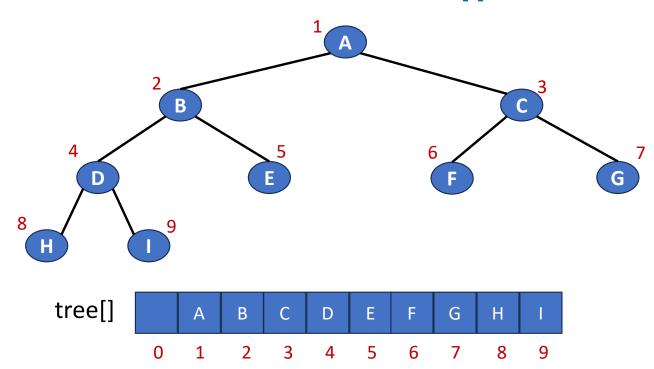


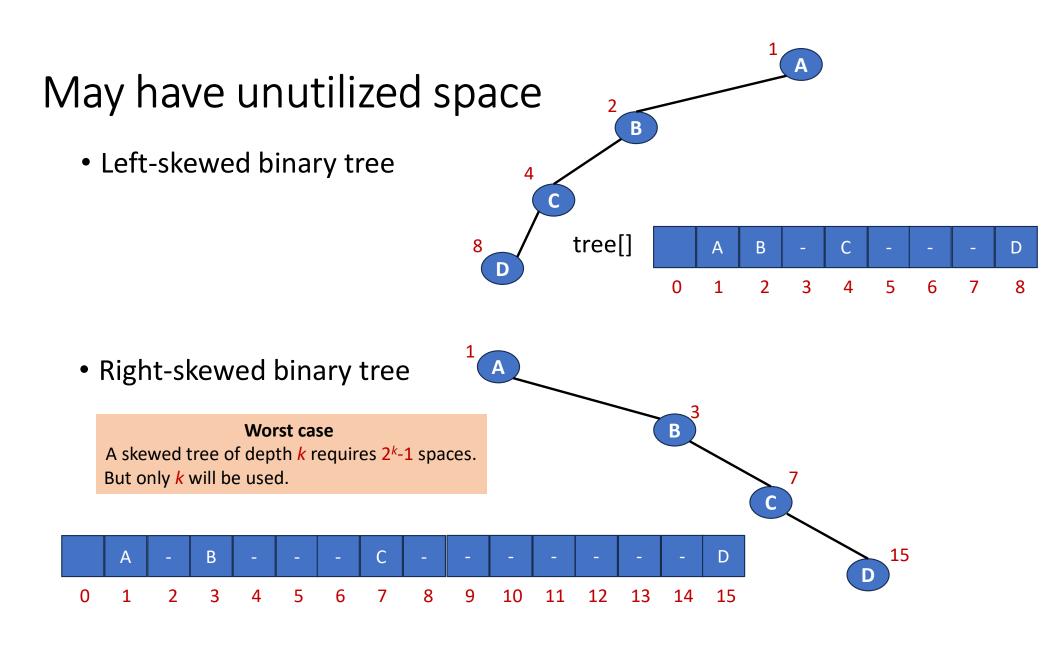
## Binary tree representations

- Array representation
- Linked representation

## Array representation

• Number the nodes using the numbering scheme for a full binary tree. The node that is numbered *i* is stored in tree[*i*].





### Linked representation

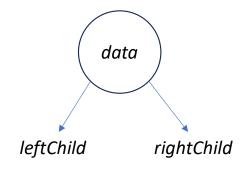
• Each binary tree node is represented as an object whose data type is TreeNode.

The space required by an n node binary tree is
 n \* (space required by one node).

## Node representations

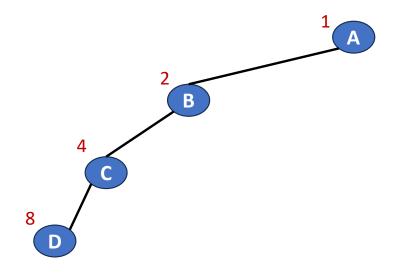
```
typedef struct node *treePointer;
typedef struct node{
    char data;
    treePointer leftChild, rightChild;
};
```

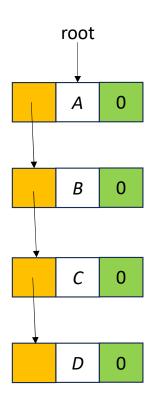
leftChild	data	rightChild

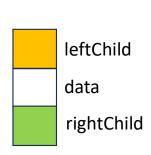


# Linked representation for binary tree

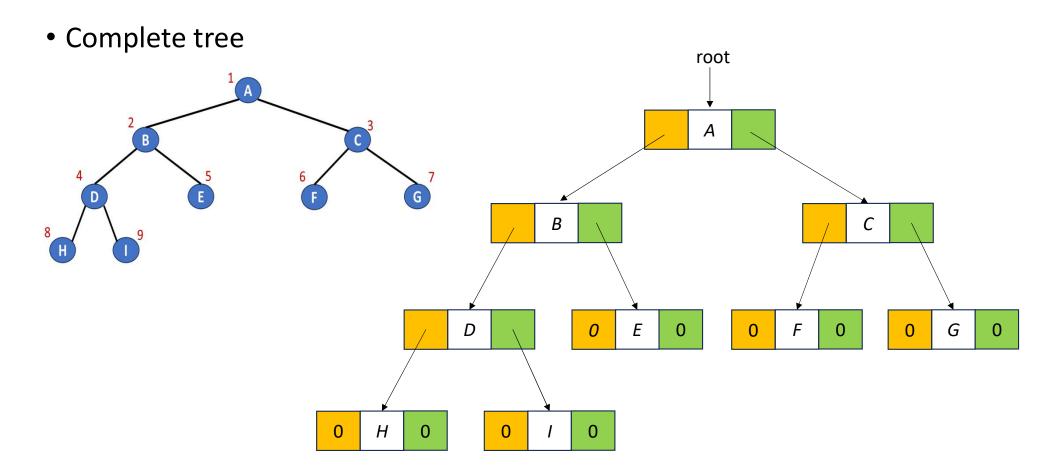
Skewed binary tree





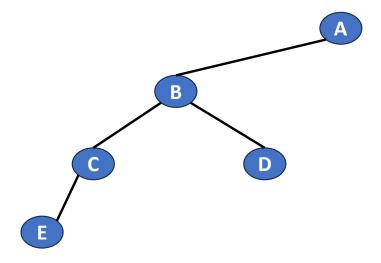


# Linked representation for binary tree



### Exercise

- Draw the internal memory representation of the binary tree
  - 1. Using sequential (array) representation
  - 2. Using linked representation



### Exercise

• Draw the internal memory representation of the binary tree

