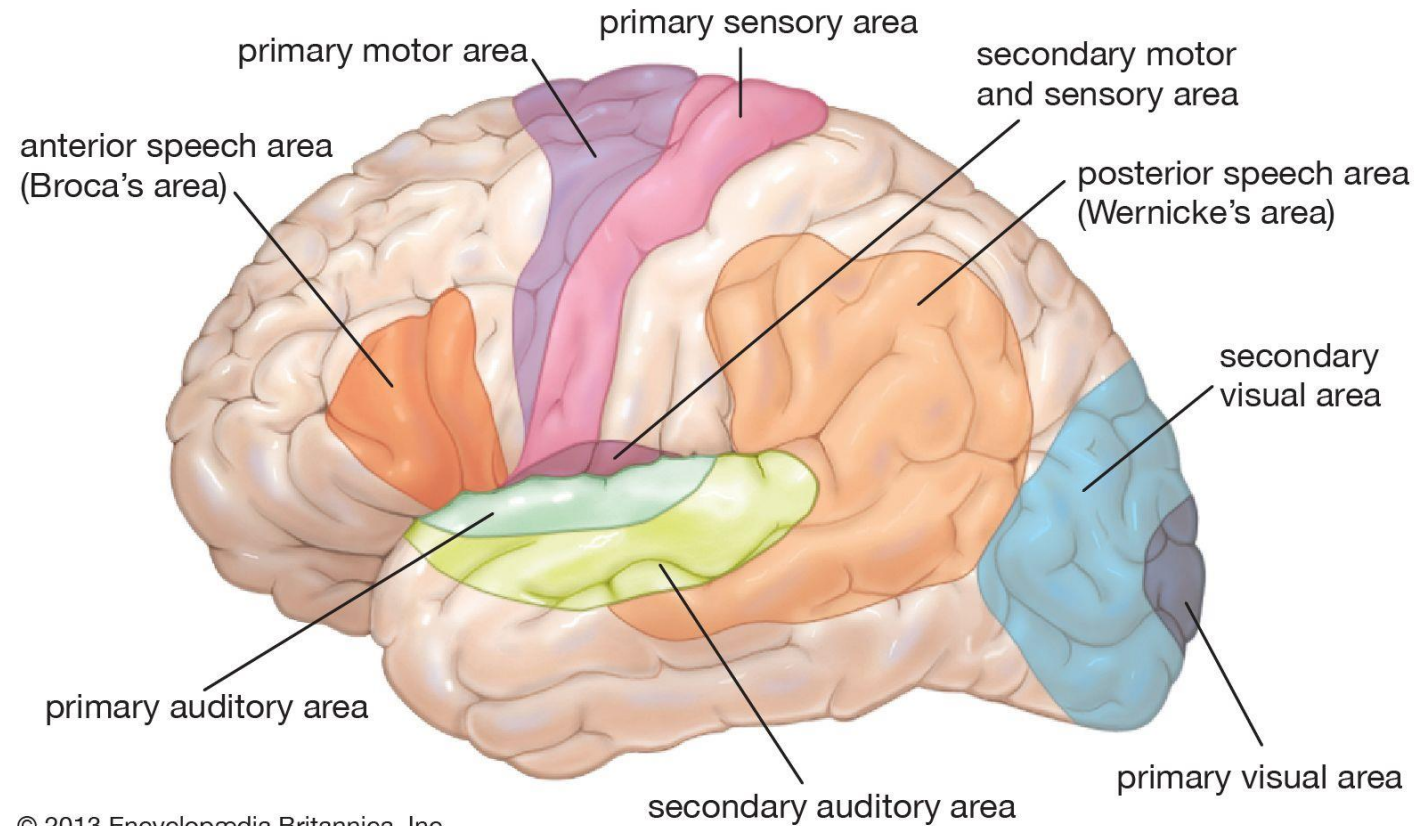


# Graphs

Ch. 6

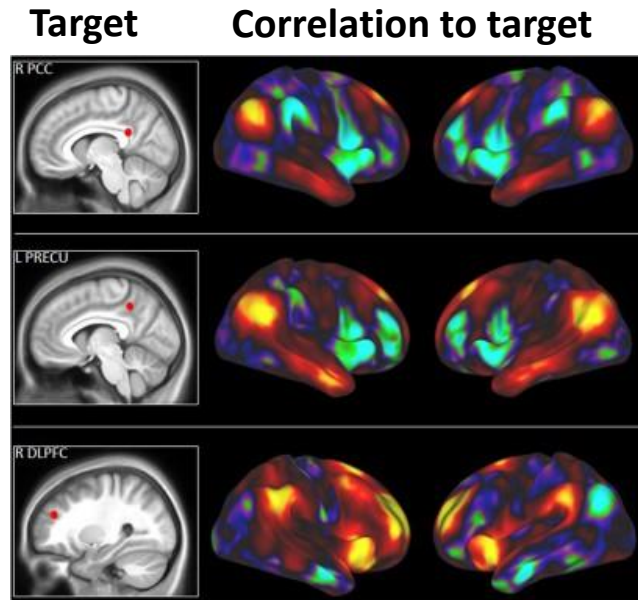
# Regions in the brain

- Each region is associated with different functions.

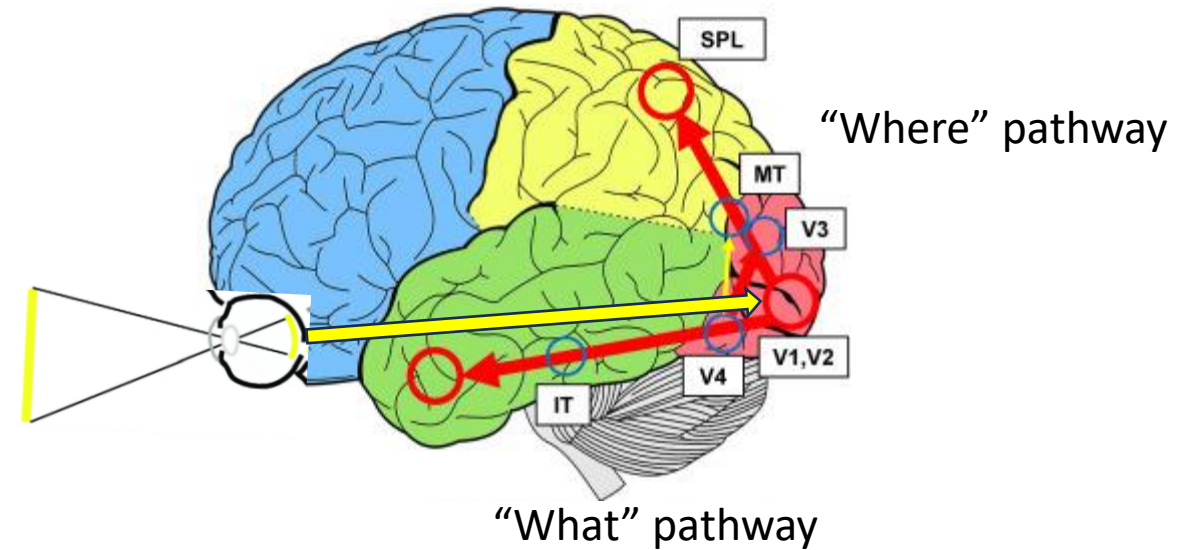


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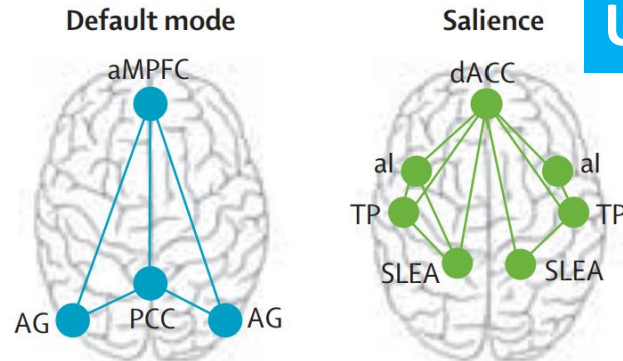
# Connections between regions



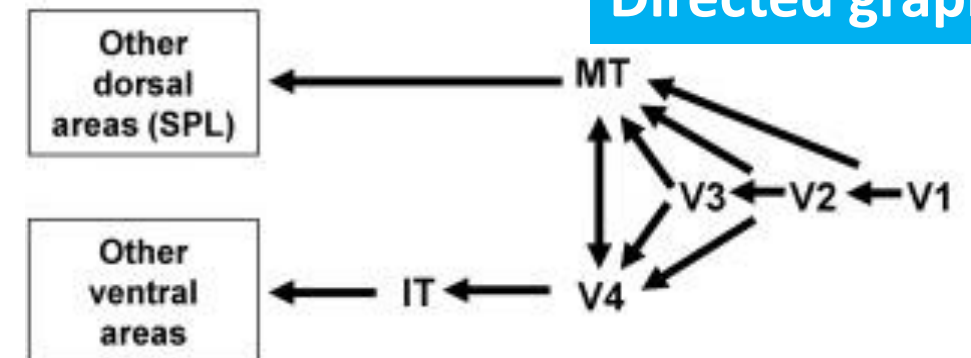
[Beaty et al., Scientific Reports 2015; 5: 10964]



- Using graphs to represent the connections for visualization and further analysis.



**Undirected graph**



**Directed graph**

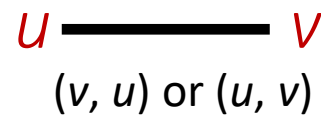
[Leanne M Williams, Lancet Psychiatry 2016; 3: 472–80]

[Choi et al., NeuroImage 2020; 220: 117145]

# Graphs

$$G = (V, E)$$

- $V$ : Set of vertices
  - Vertices are also called nodes or points.
- $E$ : Set of edges
  - Each edge connects two different vertices.
  - Edges are called arcs and lines.
  - In **undirected** graph,  $(u, v)$  and  $(v, u)$  represent the same edge.



*Edge without orientation*

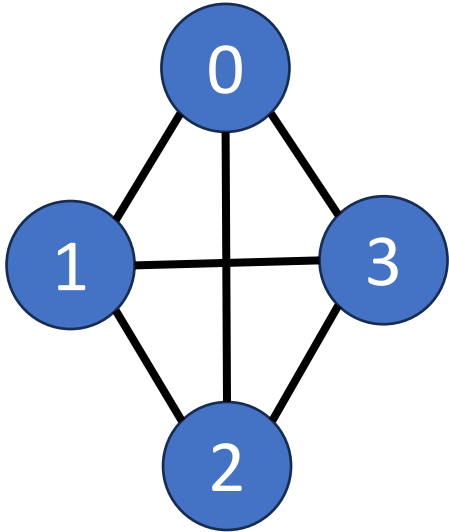
- In **directed** graph,  $\langle u, v \rangle$  and  $\langle v, u \rangle$  represent different edges.



*Edge with orientation*

# Sample graphs

Undirected graph

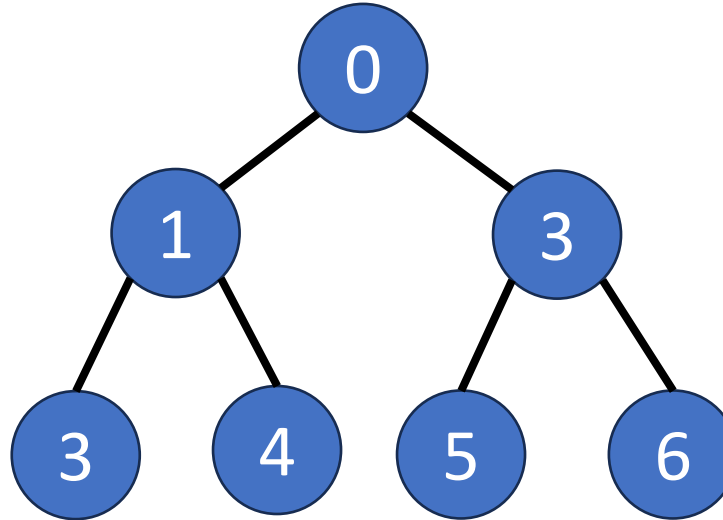


$G_1$

**Set representation:**

$$V(G_1) = \{0, 1, 2, 3\}$$

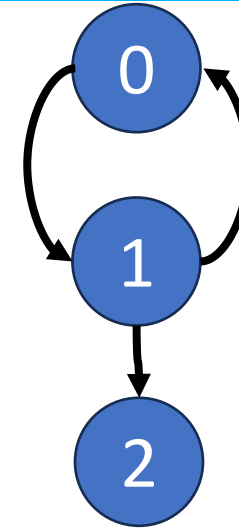
$$E(G_1) = \{(0, 1), (1, 2), (2, 3), (0, 3), (0, 2), (1, 3)\}$$



$G_2$

Tree is a type of graph.

Directed graph  
(Digraph)



$G_3$

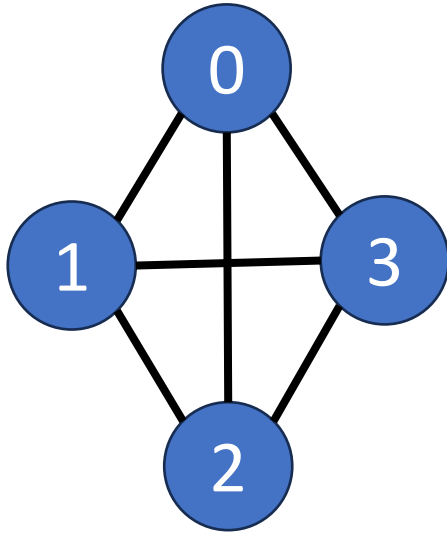
**Set representation:**

$$V(G_3) = \{0, 1, 2\}$$

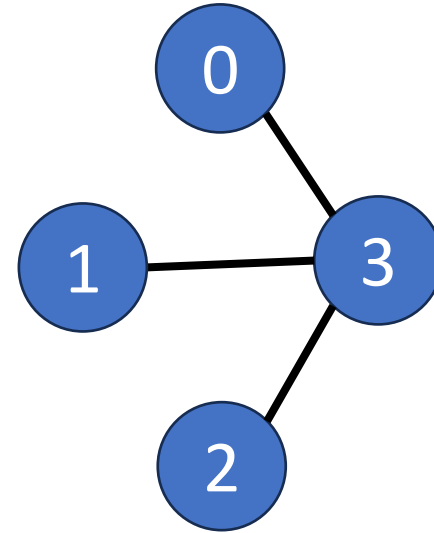
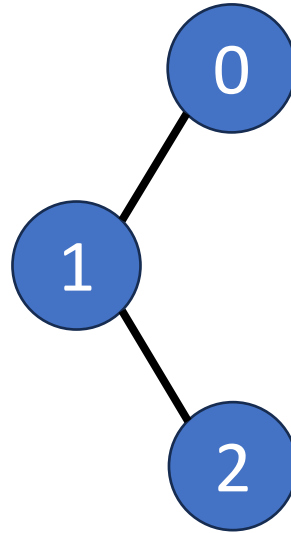
$$E(G_3) = \{<0, 1>, <1, 0>, <1, 2>\}$$

# Terminology: subgraph

- **Subgraph** of  $G$ : a graph is composed of a subset of vertices  $G$  and a subset of edges  $E$ .



$G_1$



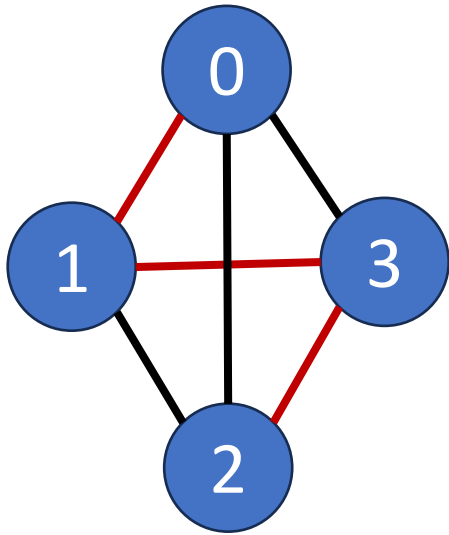
Some subgraphs of  $G_1$

# Terminology: path

- Path

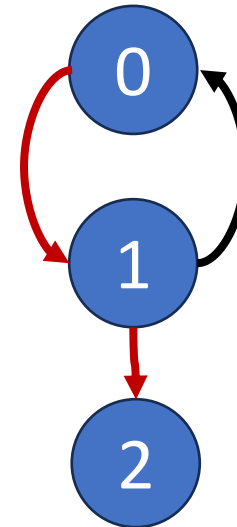
- Length of a path is the number of edges on it.

$(0,1),(1,3),(3,2)$  or  $0,1,3,2$  is a path from 0 to 2 in  $G_1$ .  
The length of path is 3.



$G_1$

$\langle 0,1 \rangle, \langle 1,2 \rangle$  or  $0,1,2$  is a path from 0 to 2 in  $G_3$ .  
The length of path is 2.



$G_3$

# Terminology: simple path and cycle

- **Simple path**: No repeating vertices (except for the case that the 1<sup>st</sup> and the last are the same).

$(0,1),(1,3),(3,2)$  or  $0,1,3,2$ .

Simple path

$(0,1),(1,3),(3,1)$  or  $0,\underline{1},3,\underline{1}$ .

Repeated vertices

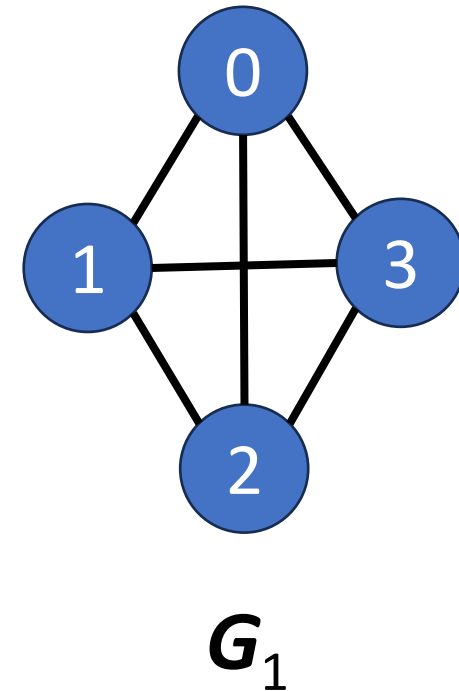
Not simple path

$(0,1),(1,2),(2,0)$  or  $\underline{0},1,2,\underline{0}$ .

Repeated vertices located at 1<sup>st</sup> and last positions

Simple path

Cycle



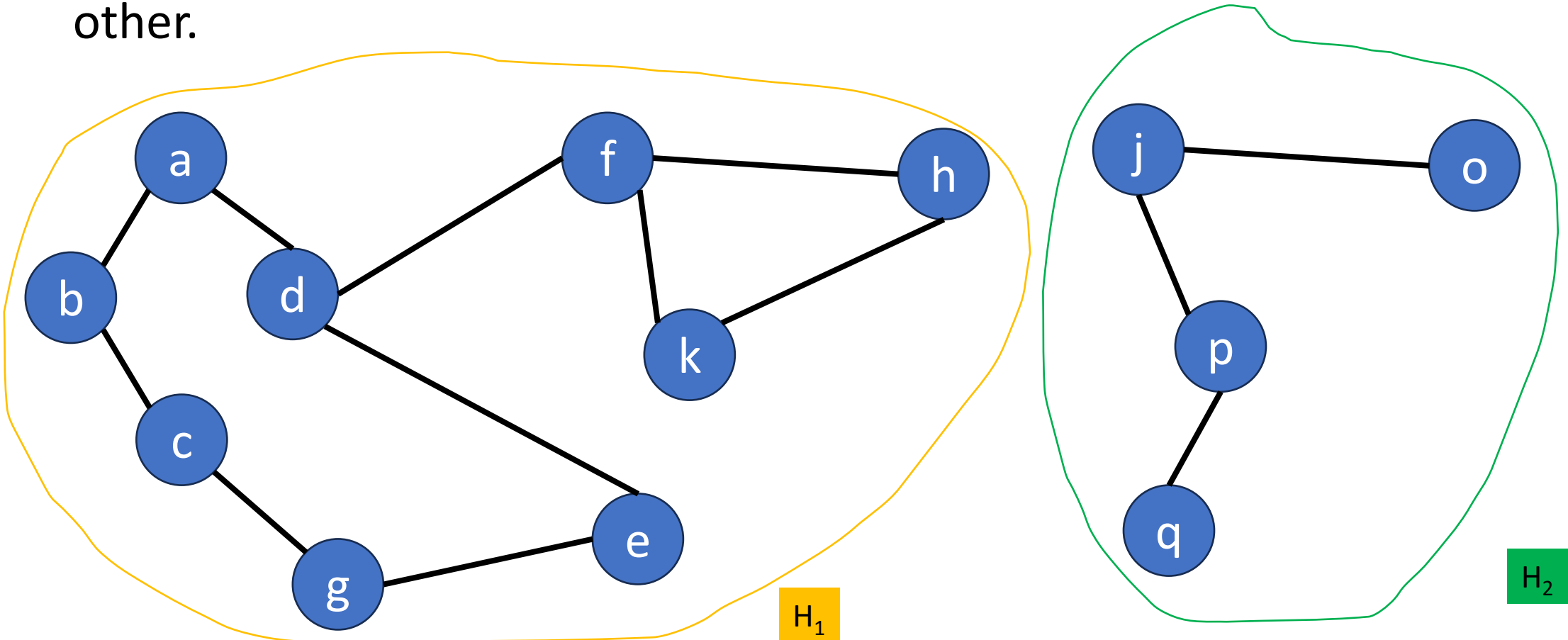
- **Cycle**:

- Simple path
- 1<sup>st</sup> and the last vertices are the same.



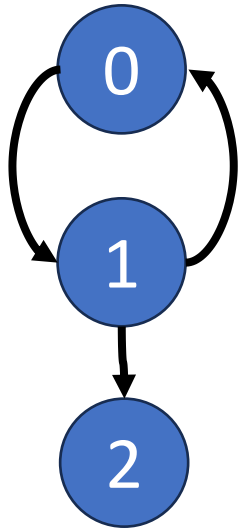
# More applications of graphs

- Find connected components:
  - Two vertices  $u$  and  $v$  are **connected** iff there is a path from  $u$  to  $v$ .
  - **Connected component**: a maximum subgraph that are connected to each other.

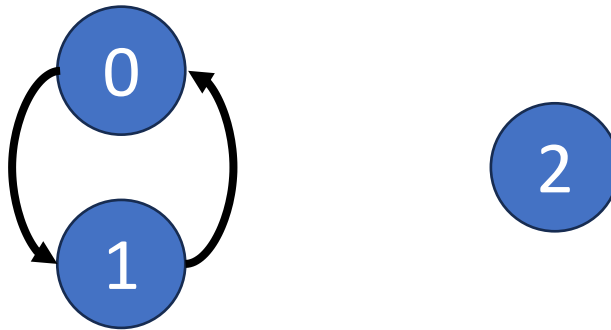
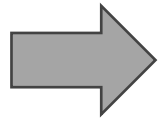


# Strongly connected components

- In a **directed** graph, every pair of vertices  $u$  and  $v$  has a directed path from  $u$  to  $v$  and also from  $v$  to  $u$ .



$G_3$



Two strongly connected components

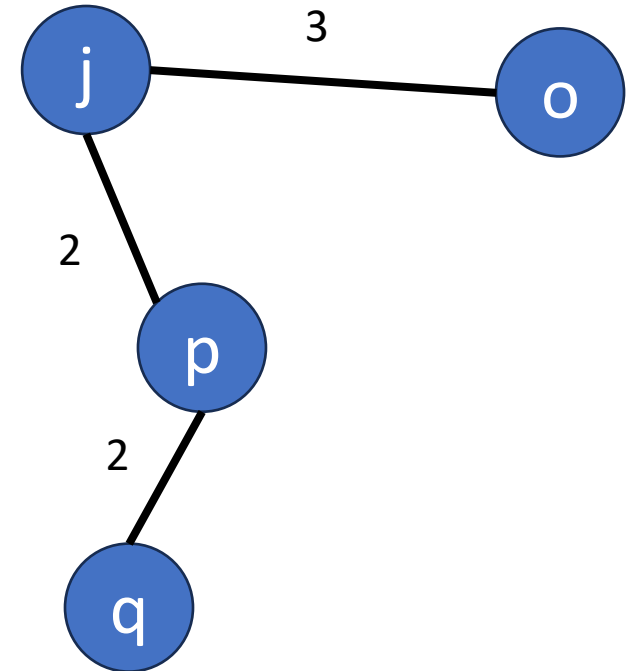
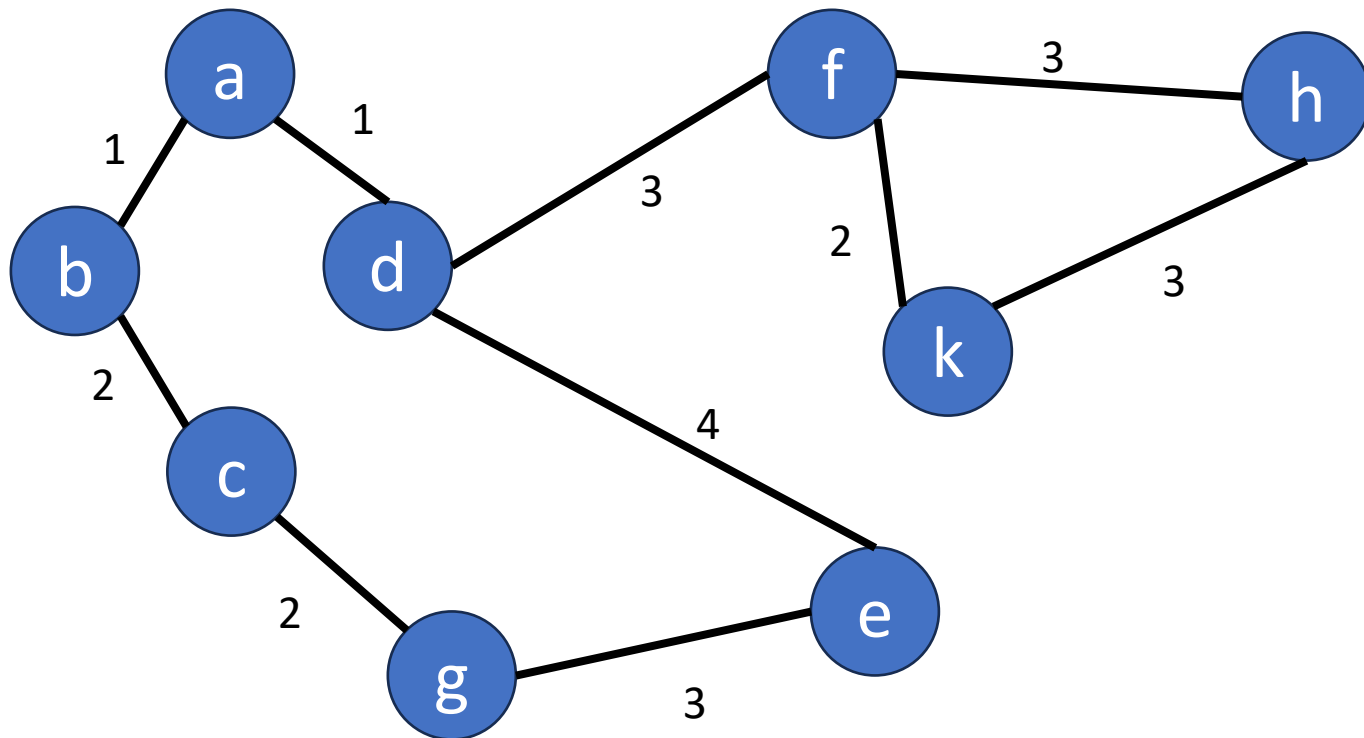
# More applications of graphs

- Plan the route from city *b* to city *k*.

*Vertex = city*

*Edge = roads*

*Edge weight = distance or time*



# Complete graph

- Having the maximum number of edges.

$n = 1$

$n = 2$

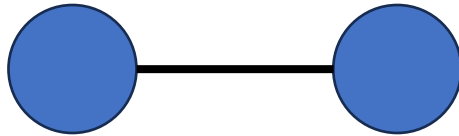
$n = 3$

$n = 4$

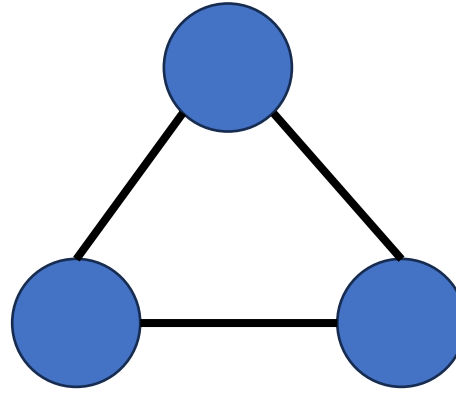
Undirected  
graph



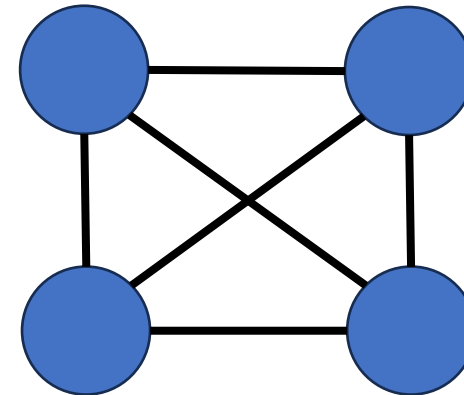
0



1



3



6

Edge  $(u,v)$  is the  
same as edge  $(v,u)$ .

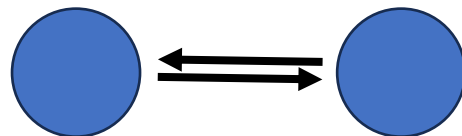
$$\frac{n(n-1)}{2}$$



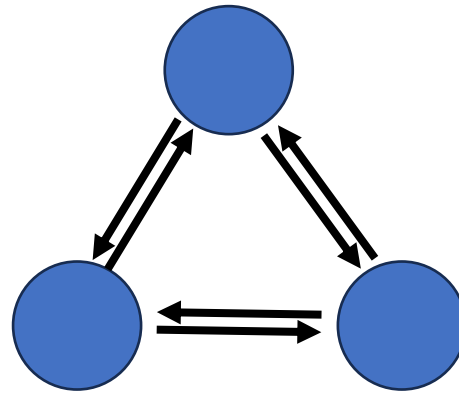
Digraph



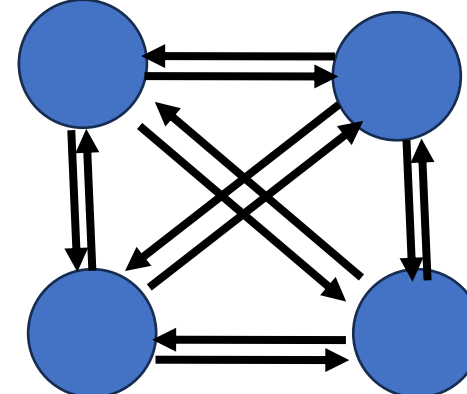
0



2



6



12

Edge  $(u,v)$  is not the  
same as edge  $(v,u)$ .

$$n(n-1)$$



Number of edges

Number of edges

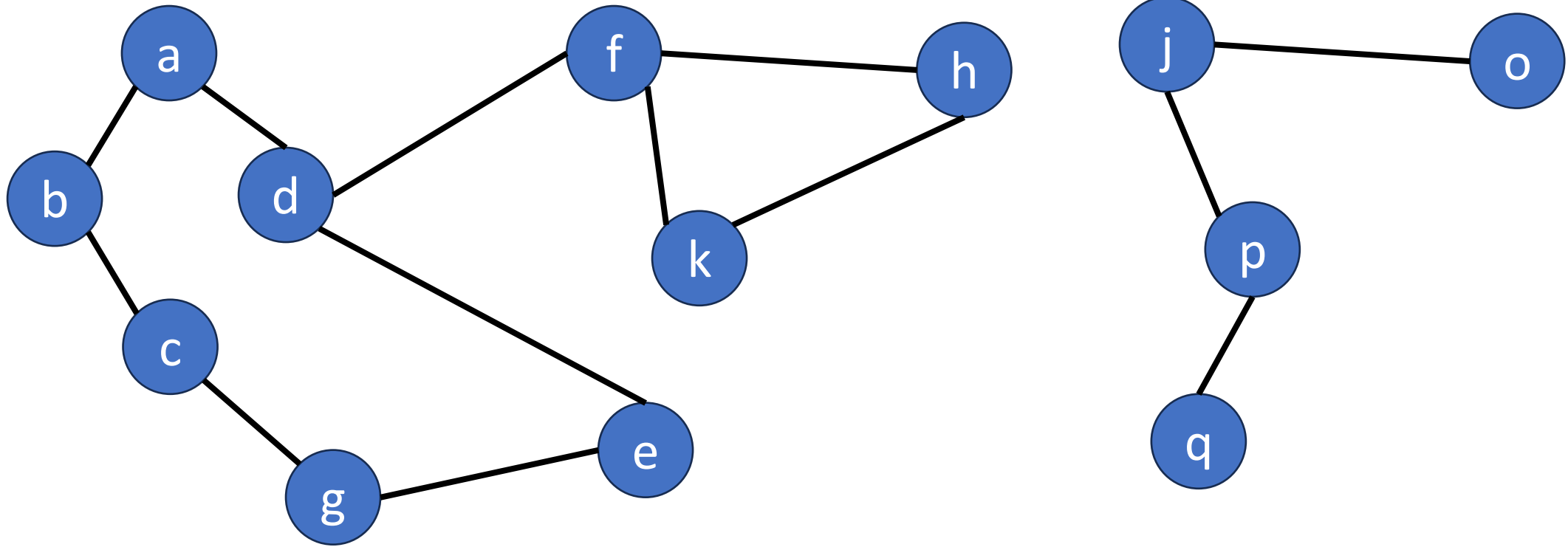
# Property: Number of edges

$n$ : number of vertices

- Number of edges in an undirected graph is  $\leq n(n-1)/2$ .
- Number of edges in a directed graph is  $\leq n(n-1)$ .

# Property: Vertex degree

- Number of edges incident to vertex.

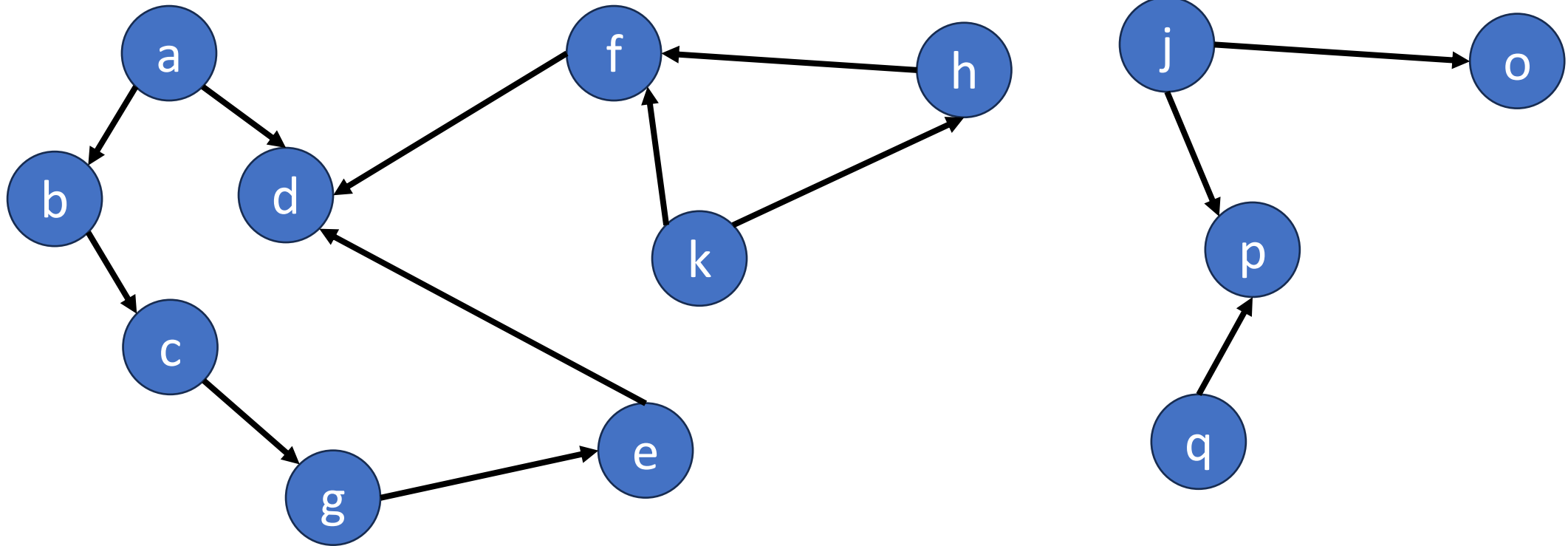


Example:

$$\text{degree}(d) = 3, \text{degree}(f) = 3, \text{degree}(h) = 2, \text{degree}(q) = 1$$

# Property: In-degree of a vertex

- Number of incoming edges.

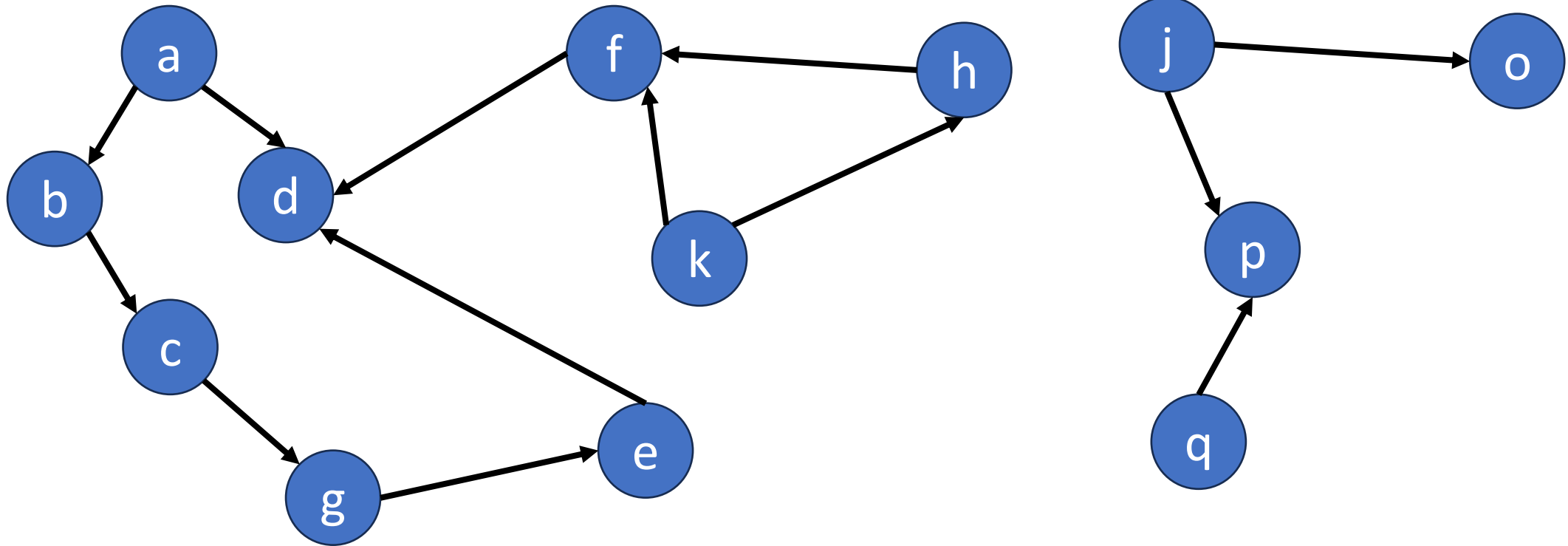


Example:

$\text{in-degree}(d) = 3$ ,  $\text{in-degree}(f) = 2$ ,  $\text{in-degree}(h) = 1$ ,  $\text{in-degree}(q) = 0$

# Property: Out-degree of a vertex

- Number of outbound edges.



Example:

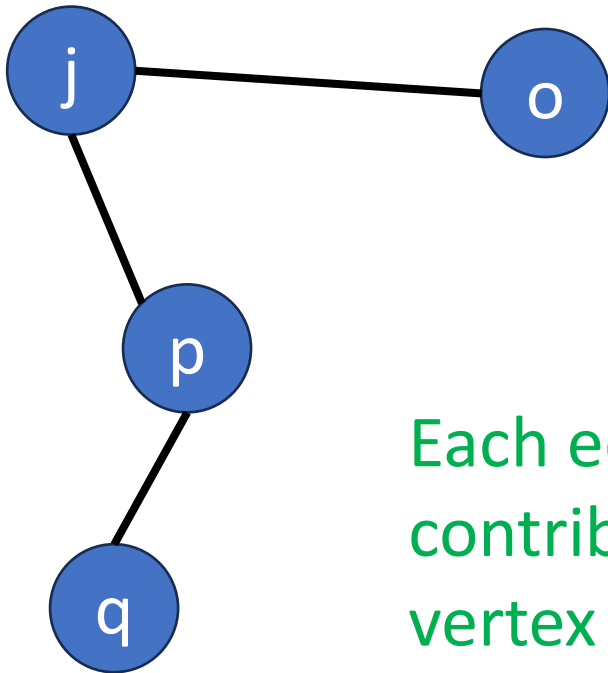
$\text{out-degree}(d) = 0$ ,  $\text{out-degree}(f) = 1$ ,  $\text{out-degree}(h) = 1$ ,  $\text{out-degree}(q) = 1$



# Sum of degree

- $e$ : number of edges

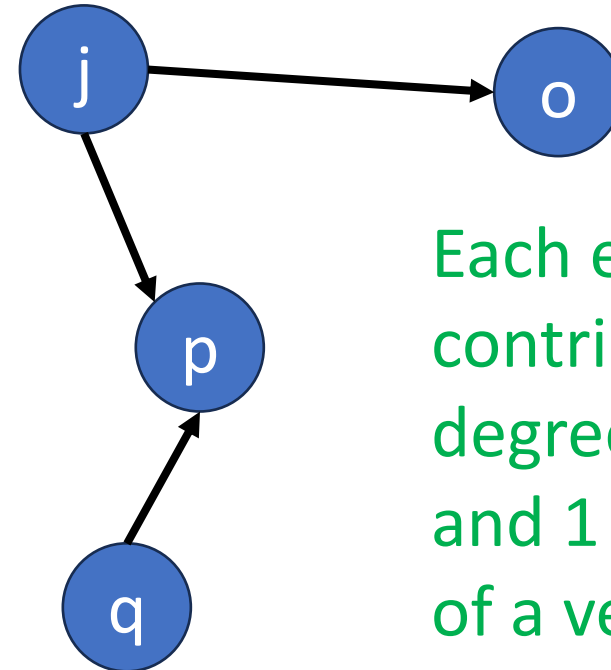
Undirected graph



Each edge  
contributes 2 to  
vertex degree.

Sum of vertex degrees =  $2e$

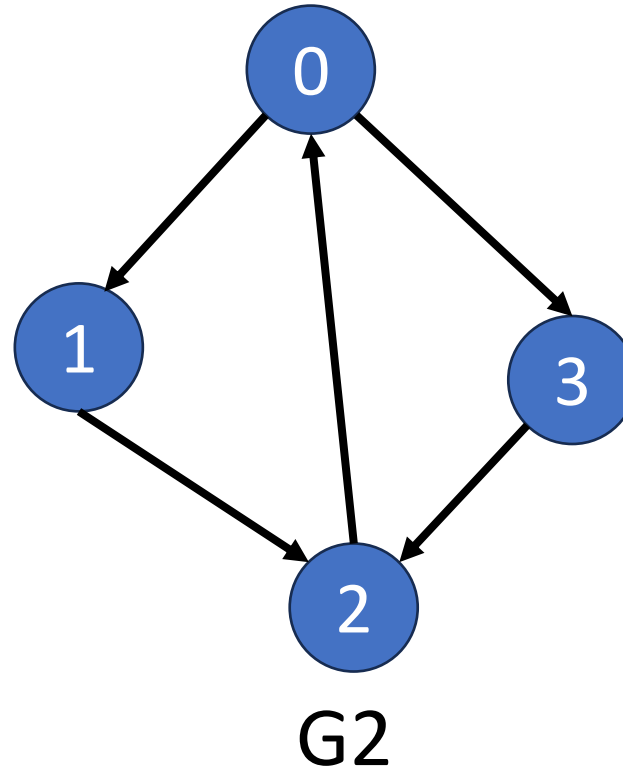
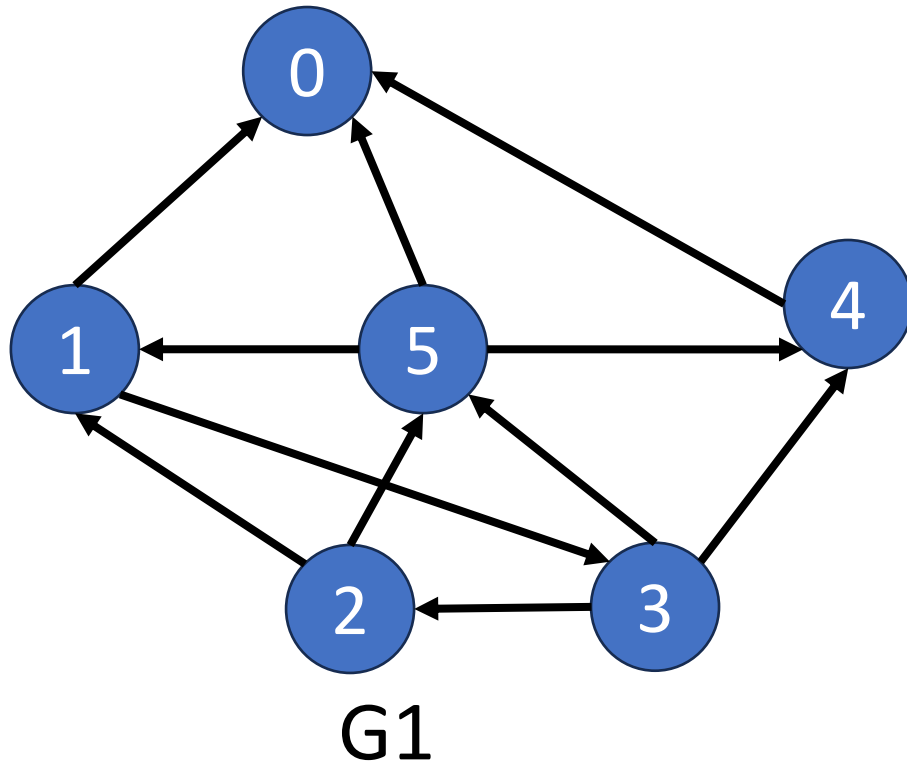
Digraph



Each edge  
contributes 1 to in-  
degree of a vertex  
and 1 to out-degree  
of a vertex.

sum of in-degrees  
= sum of out-degrees  
=  $e$

# Exercise



Please reply your answers of Q6 – Q8 via the following link:

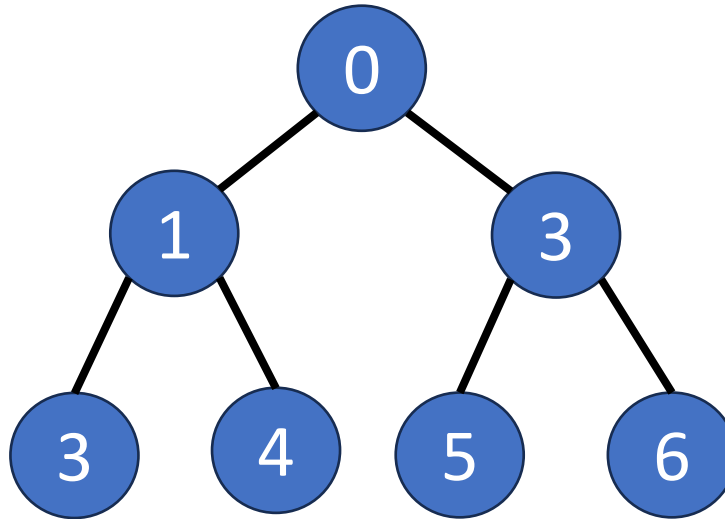


Group members: 1~3 people

- Q6: Write out the in-degree of vertices 0, 1, ..., 5 in G1.
- Q7: Write out the out-degree of vertices 0, 1, ..., 5 in G1.
- Q8: Is the directed graph G2 strongly connected? Why?

# Tree

- **Acyclic** graph: a tree has no cycles.
- **Connected** graph:
  - all pairs of nodes are connected.
  - **n vertices** connected graph with **n-1 edges**.

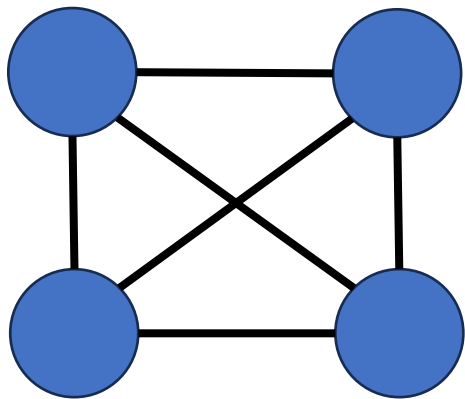


$G_2$

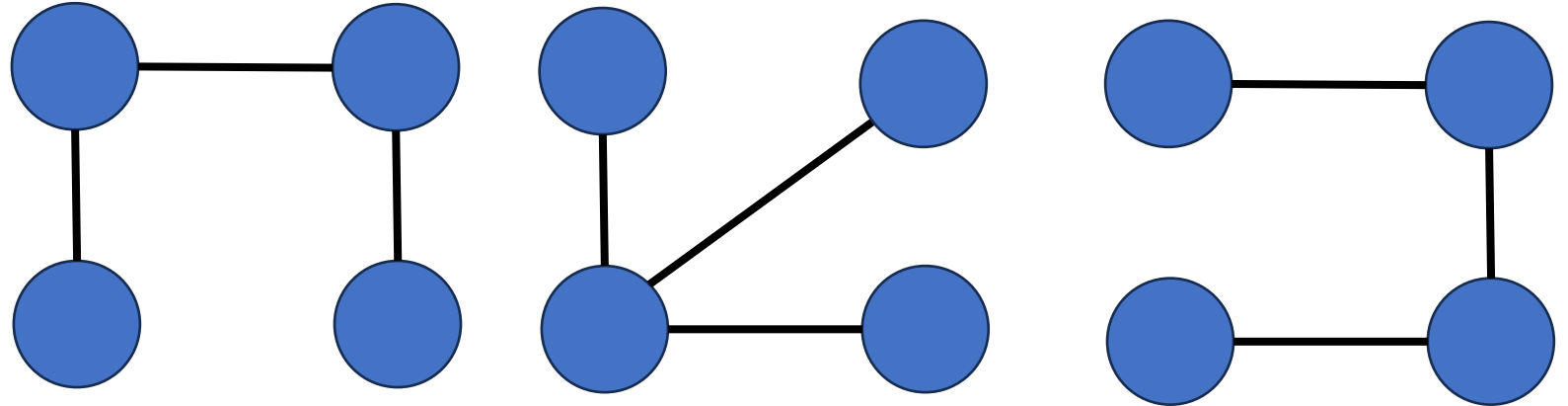
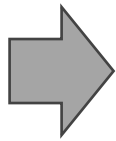
# Spanning tree

- A subgraph that includes **all vertices** of the original graph.
- A tree.

If the original graph has  $n$  vertices, the spanning tree has  $n$  vertices and  $n-1$  edges.



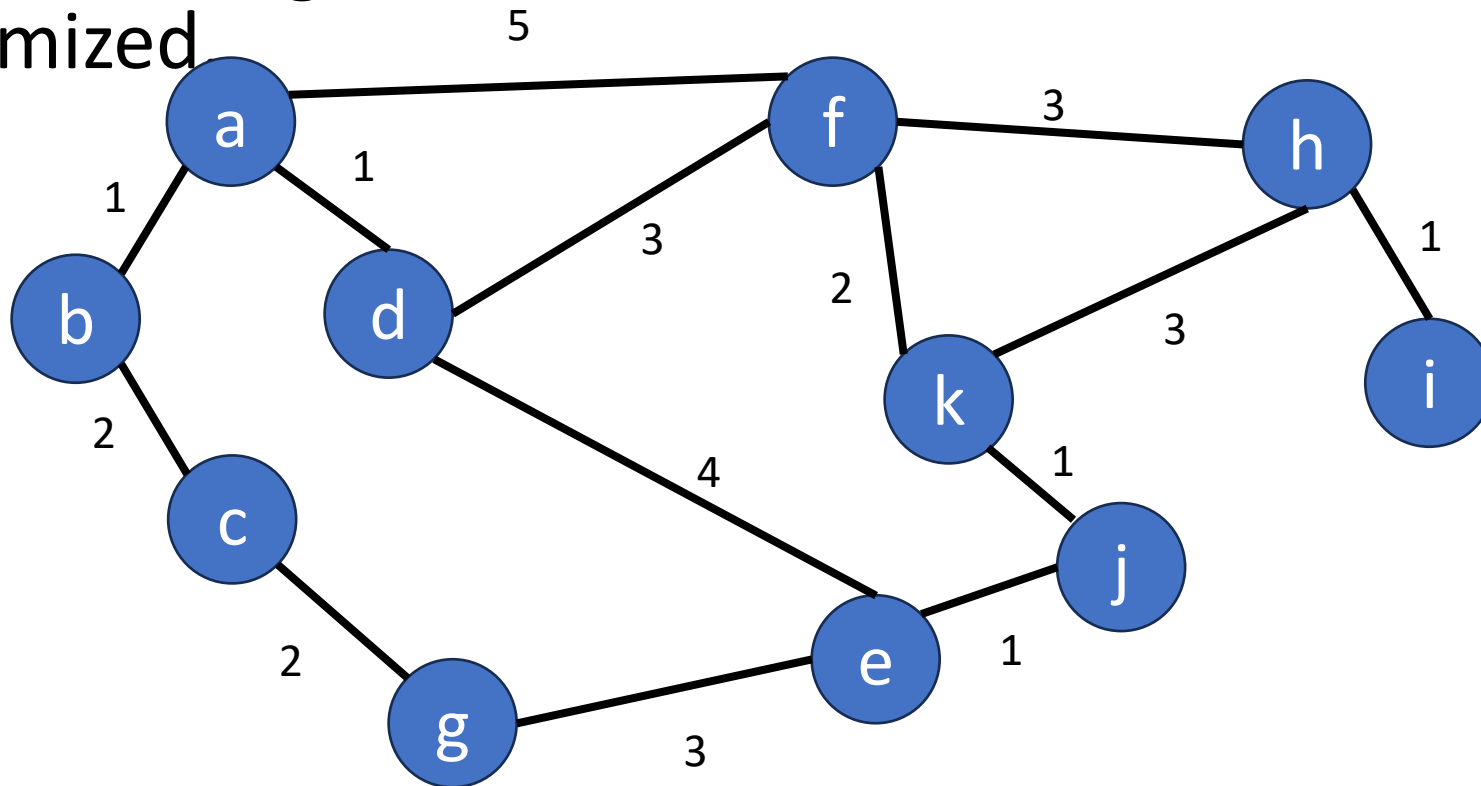
Original graph



Three examples of spanning trees

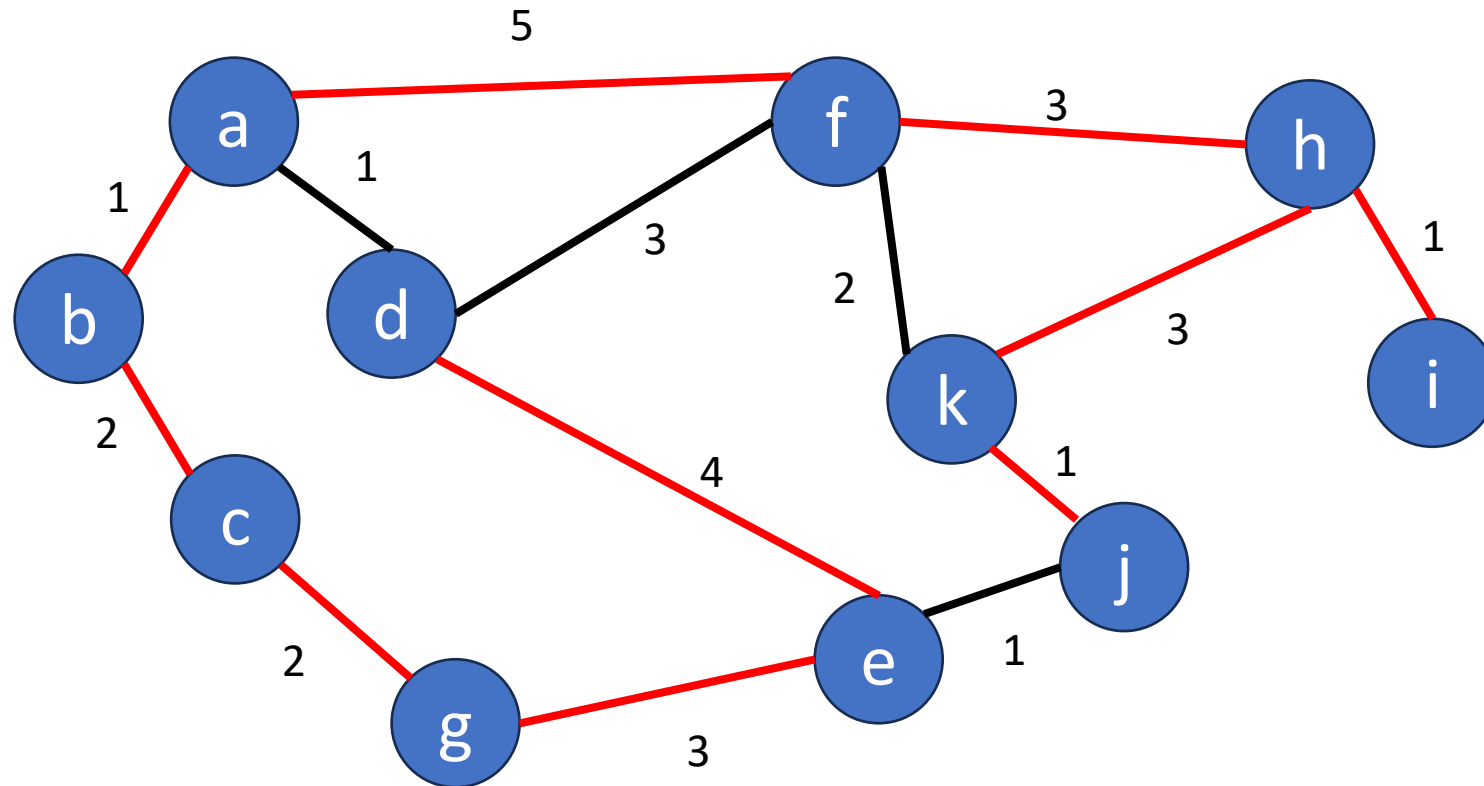
# Minimum cost spanning tree

- A spanning tree with least cost
  - **Tree cost:** sum of edge weights/costs.
- Application: Building communication links for all cities while the cost is minimized



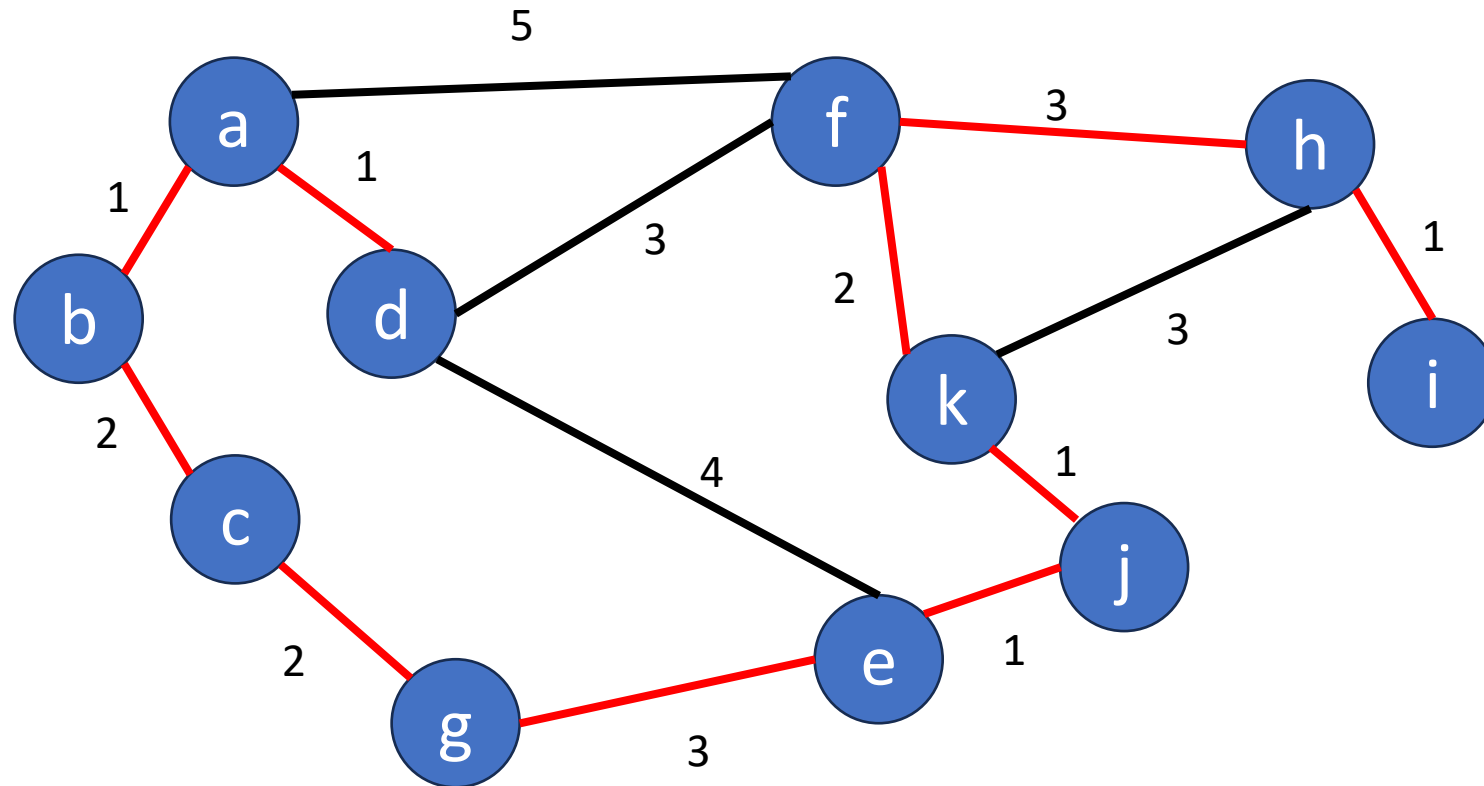
# Example of spanning tree

- Spanning tree cost =  $1 + 2 + 2 + 3 + 4 + 5 + 3 + 1 + 3 + 1 = 25$



# Example of spanning tree

- Spanning tree cost =  $1 + 1 + 2 + 2 + 3 + 1 + 1 + 2 + 3 + 1 = 17$



# Graph representations

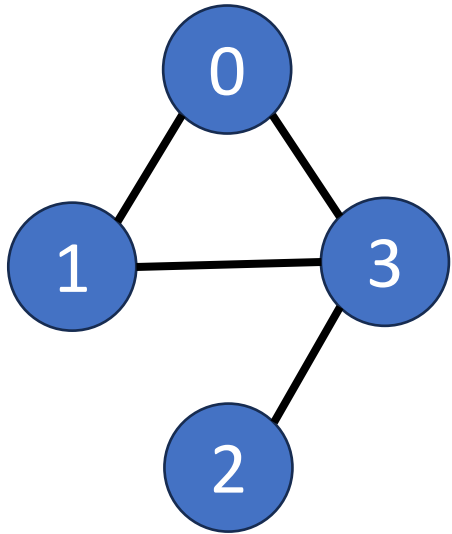
- Adjacency matrix
- Adjacency lists
  - Linked adjacency lists
  - Array adjacency lists



# Adjacency matrix

- Using 2D  $n$ -by- $n$  matrix  $A$ 
  - $n$ : number of vertices
  - $A[i][j] = 1$ : an edge between vertices  $i$  and  $j$
  - Diagonal entries are **zeros**.

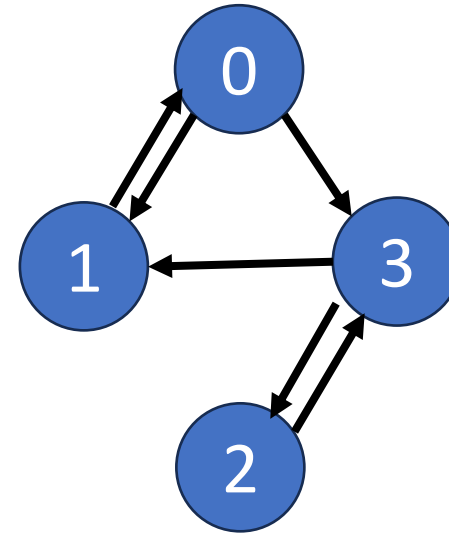
## Undirected graph



	0	1	2	3
0	0	1	0	1
1	1	0	0	1
2	0	0	0	1
3	1	1	1	0

Adjacency matrix of an undirected graph is symmetric.

## Digraph

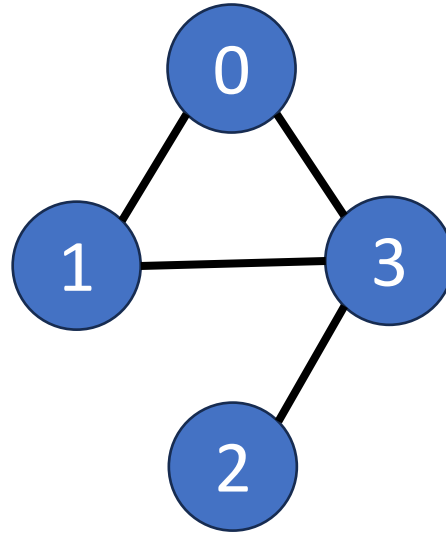


	0	1	2	3
0	0	1	0	1
1	1	0	0	0
2	0	0	0	1
3	0	1	1	0

Adjacency matrix of a directed graph need not be symmetric.

# Count number of edges in a graph

- Searching the matrix time complexity:  $O(n^2)$
- What if the graph is sparse?



	0	1	2	3
0	0	1	0	1
1	1	0	0	1
2	0	0	0	1
3	1	1	1	0

$$e \ll n^2$$

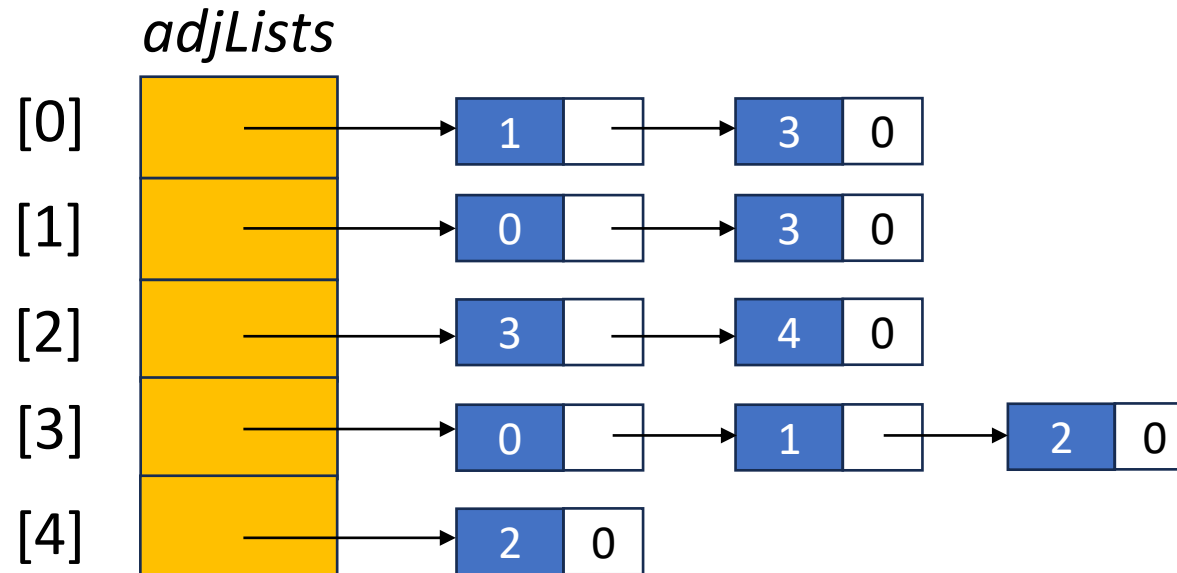
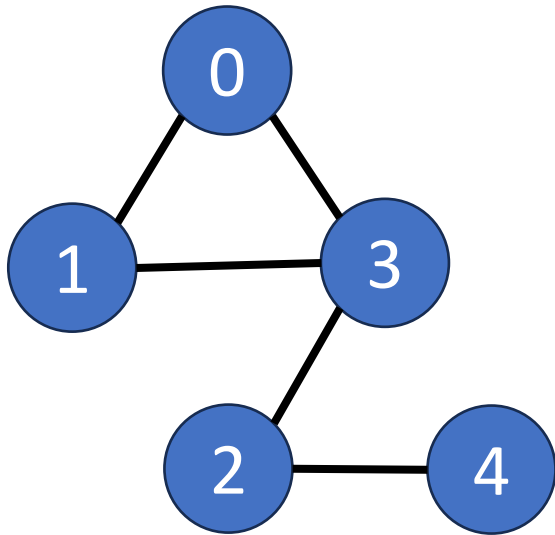
*Most elements in the adjacency matrix are zeros.*

# Linked adjacency list

- Each adjacency list is a chain. One chain for each vertex.
  - Array length =  $n$
  - In undirected graph, total number of chain nodes =  $2e$
  - In directed graph, total number of chain nodes =  $e$

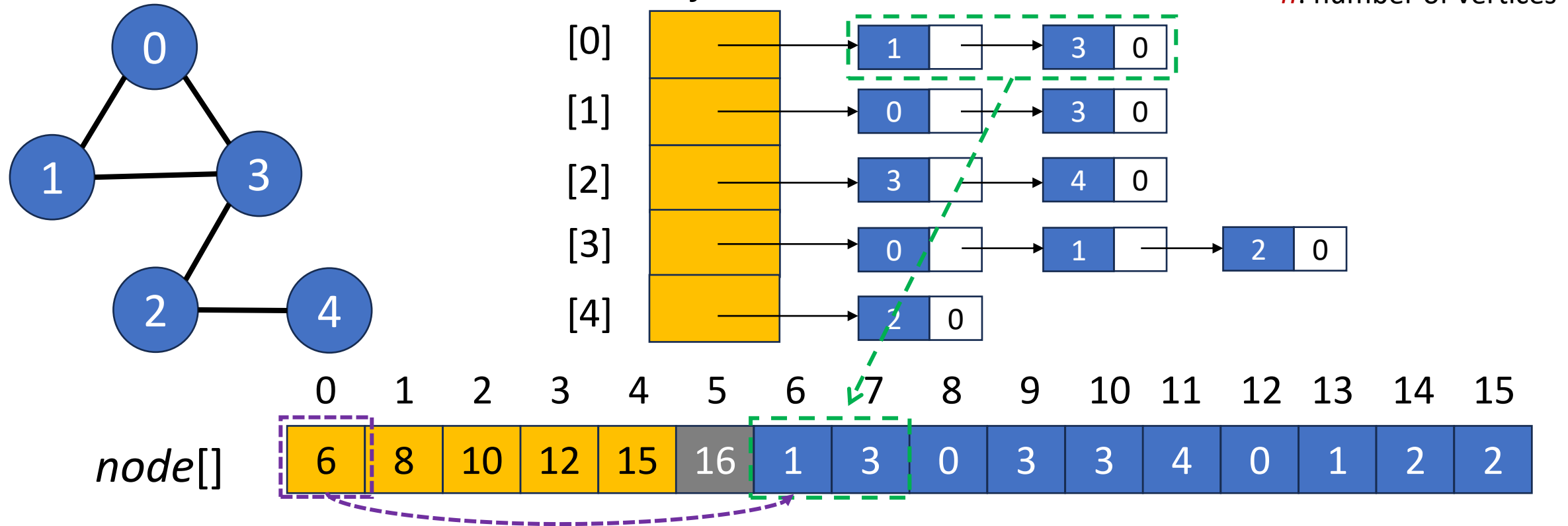
$e$ : number of edges

$n$ : number of vertices

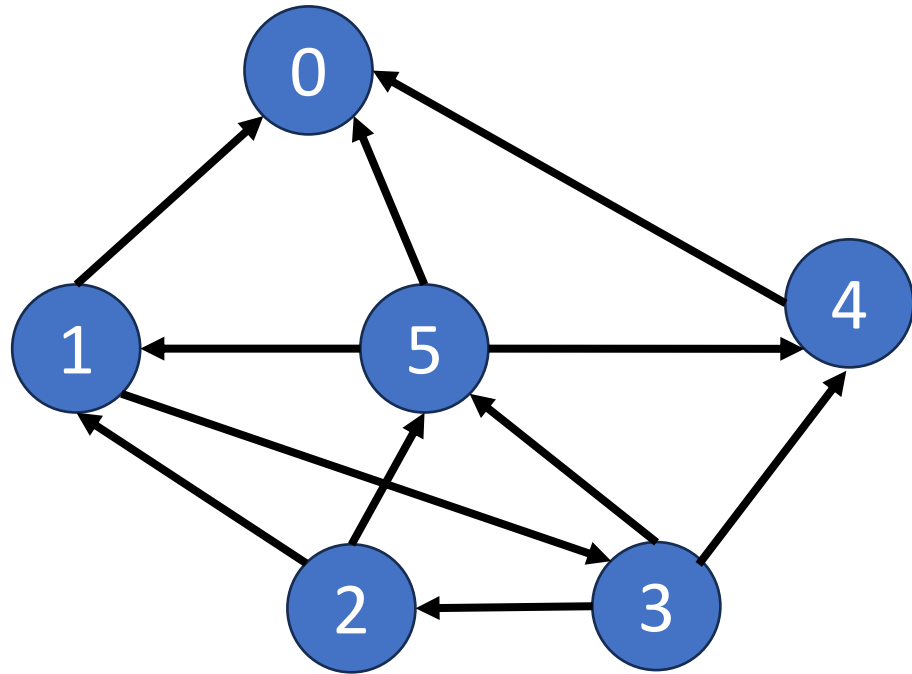


# Array adjacency list

- Using an integer array  $node[]$  to store all adjacency lists.
  - array length =  $n + 2e + 1$
  - $i$ -th element in  $node[0, 1, \dots, n-1]$ : starting point of the list for vertex  $i$ .
  - $node[n]$ : set to  $n + 2e + 1$



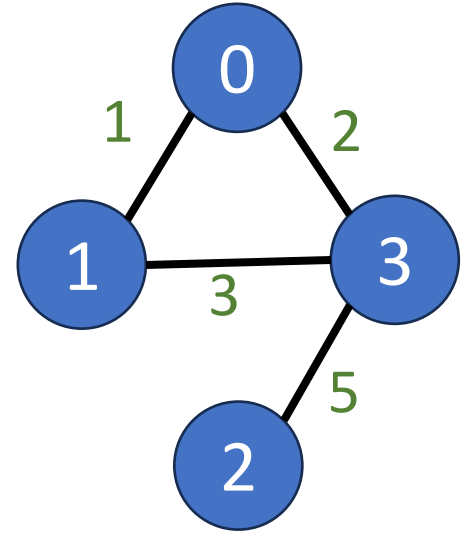
# Exercise



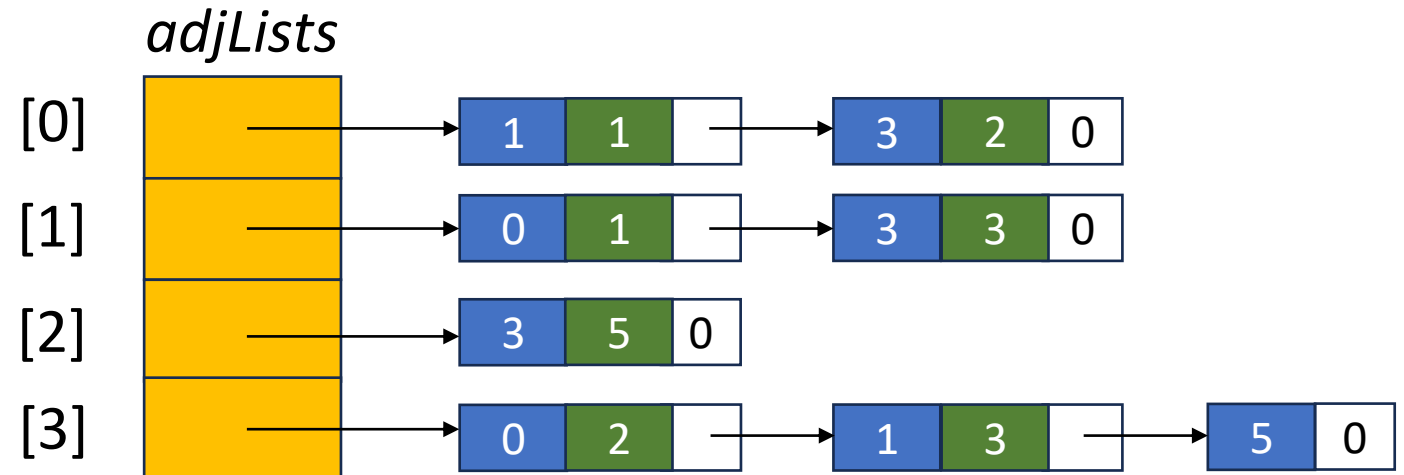
- Write out the adjacency matrix.
- Write out the linked adjacency-list representation.

# Weighted graphs

- Cost adjacency matrix:  $A[i][j]$  = cost of edge( $i, j$ )
- Adjacency lists: Each list element is a pair (adjacent vertex, edge weight)
- A graph with weighted edges is called a network.



	0	1	2	3
0	0	1	0	2
1	1	0	0	3
2	0	0	0	5
3	2	3	5	0



# Summary

- What is directed graph and undirected graph?
- Terminology of graph
  - graph, subgraph, path, simple path, cycle, connected components, degree
- Spanning tree
- Graph representation
  - Adjacency matrix
  - Adjacency lists