

「不二坊」再現排隊人龍 警加強拖吊



What will you do to process the long line?

Queues

Ch 3.3 & Ch 3.4

Queues

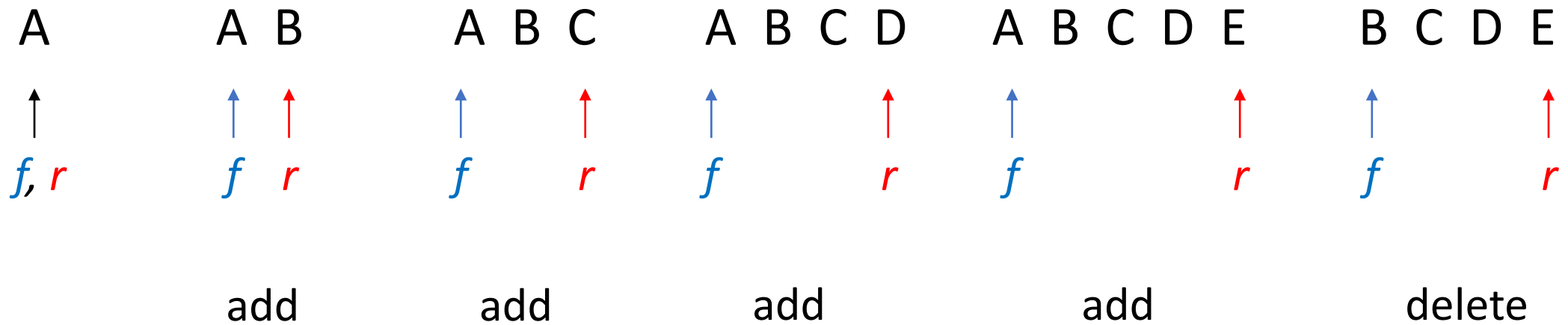
- An ordered list.
- Insertions and deletions take place at different ends.
- New elements are **added** at the **rear** end.
- Old elements are **deleted** at the **front** end.
- The first element inserted into a queue is the first element removed.

First-In-First-Out (FIFO)





Insertion and deletion of elements

- f : front of the queue
- r : rear of the queue



Operations of queues

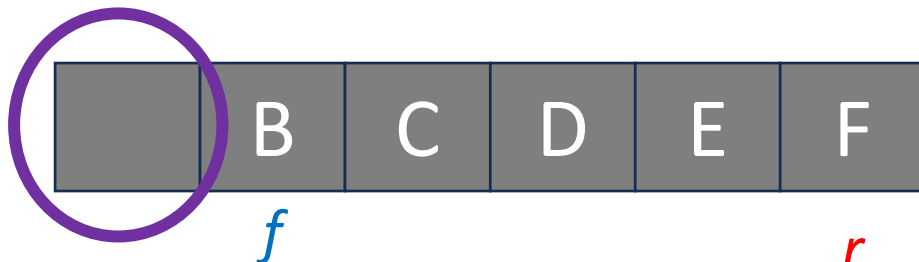
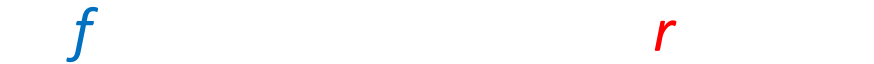
- **CreateQ**: create an empty queue and set up variables *rear* and *front* as -1.
- **IsFullQ**: return True if queue is full 

rear == MAX_QUEUE_SIZE - 1
- **IsEmptyQ**: return True if queue is empty 

front == *rear*
- **AddQ**: insert item at rear of queue
- **DeleteQ**: remove and return the item at front of queue

Representation of queues (1)

Sequential representation



Delete

Add

The queue is full because $r == \text{MAX_QUEUE_SIZE}$.

But the array is not full.

Some empty space remains.

When a customer at front of the line leaves, what will other customers do?

They move forward.

So, let's shift the array after each delete.



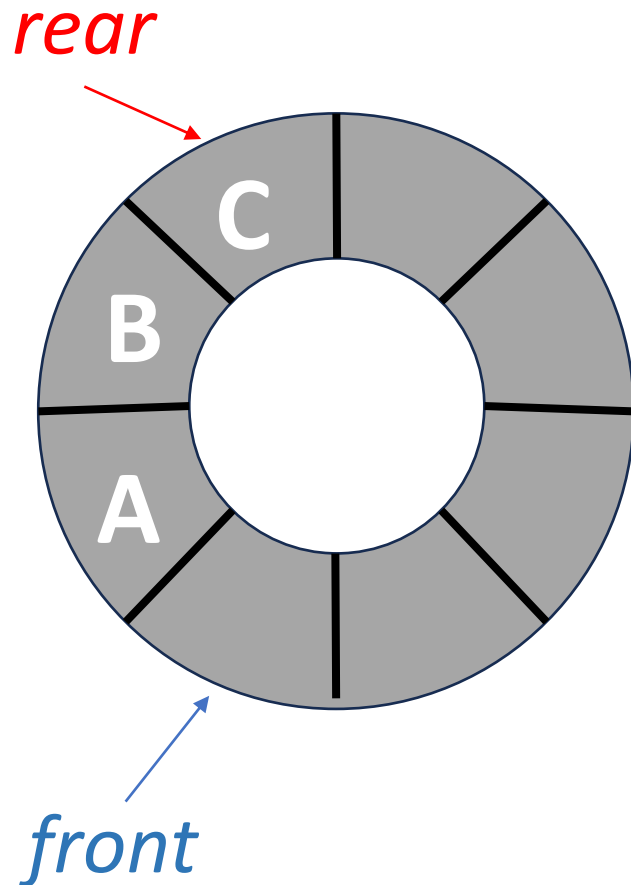
Delete

Shift

Time consuming

Representation of queues (2)

Circular queue



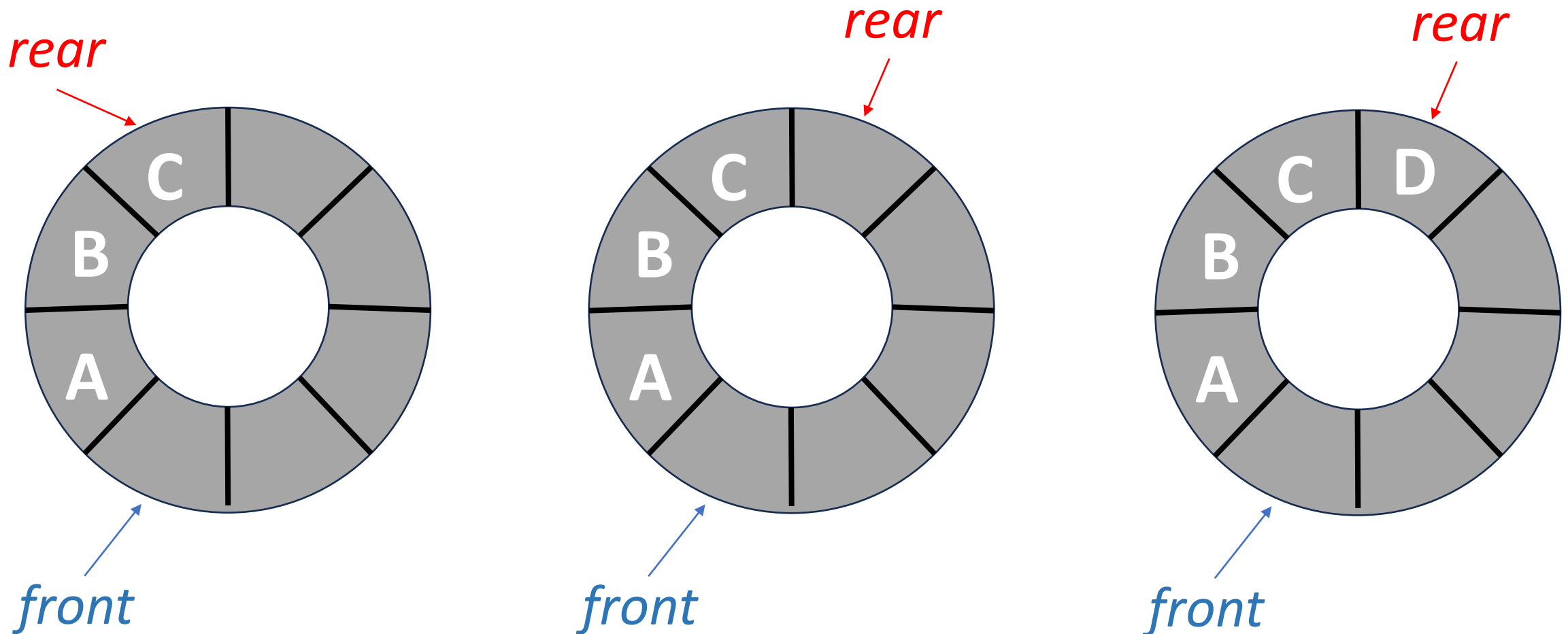
- *front*: One position counterclockwise from the position of the front element.
- *rear*: The position of the rear element.

Add an element into a circular queue

1. Move *rear* one clockwise.

$\text{rear} = (\text{rear} + 1) \% \text{MAX_QUEUE_SIZE}$

2. Put new element into $\text{queue}[\text{rear}]$.

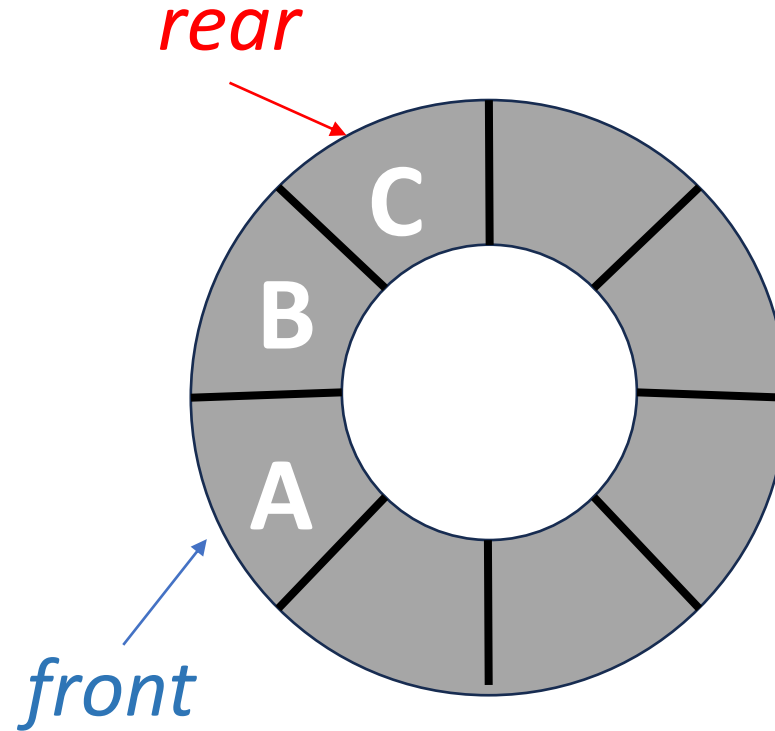
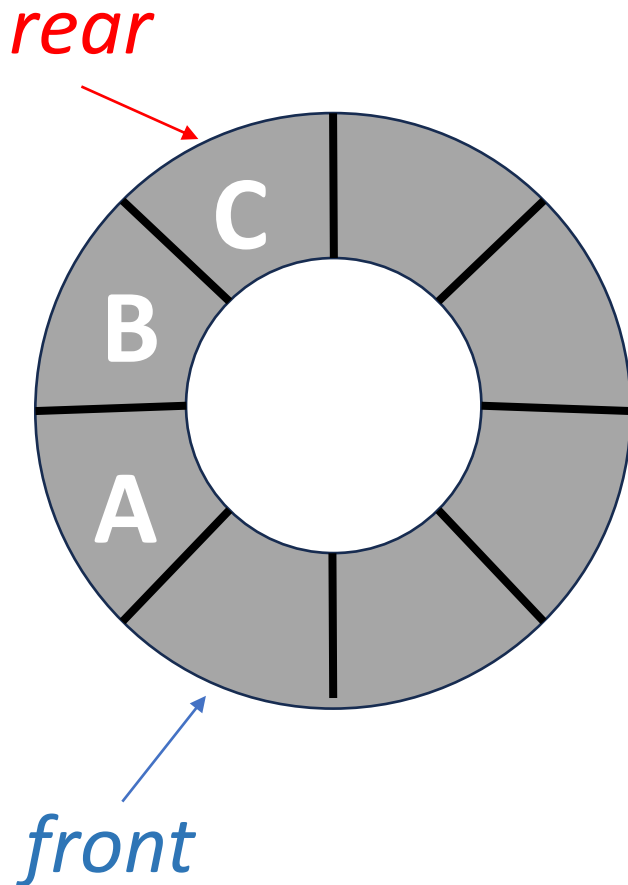


Delete an element from a circular queue

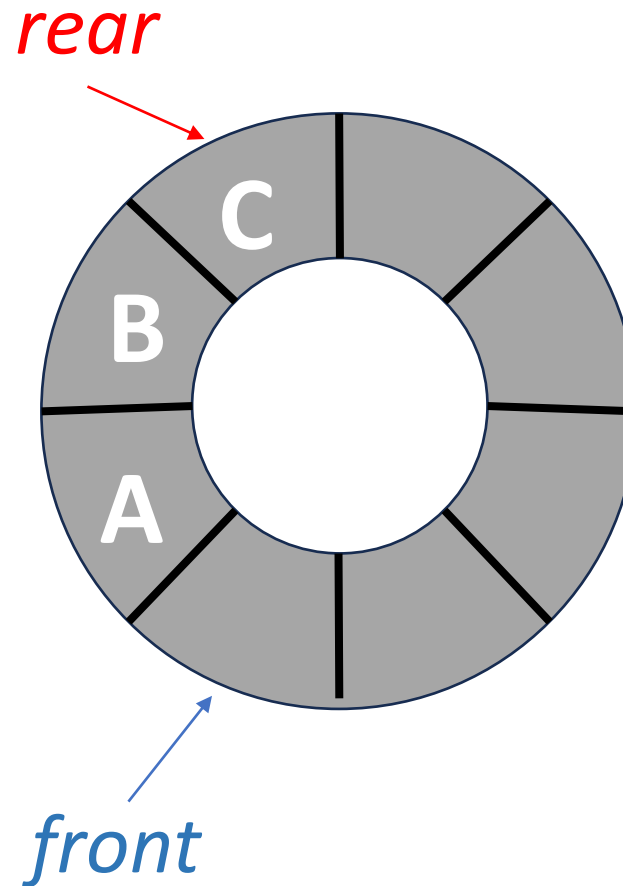
1. Move *front* one position clockwise.

$\text{front} = (\text{front} + 1) \% \text{MAX_QUEUE_SIZE}$

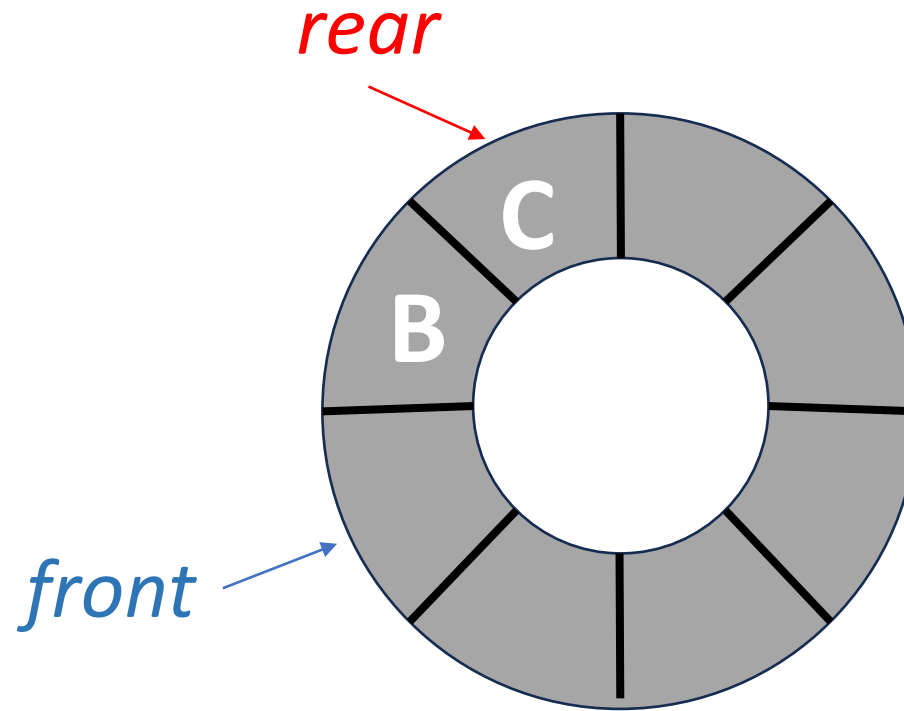
2. Return value of $\text{queue}[\text{front}]$.



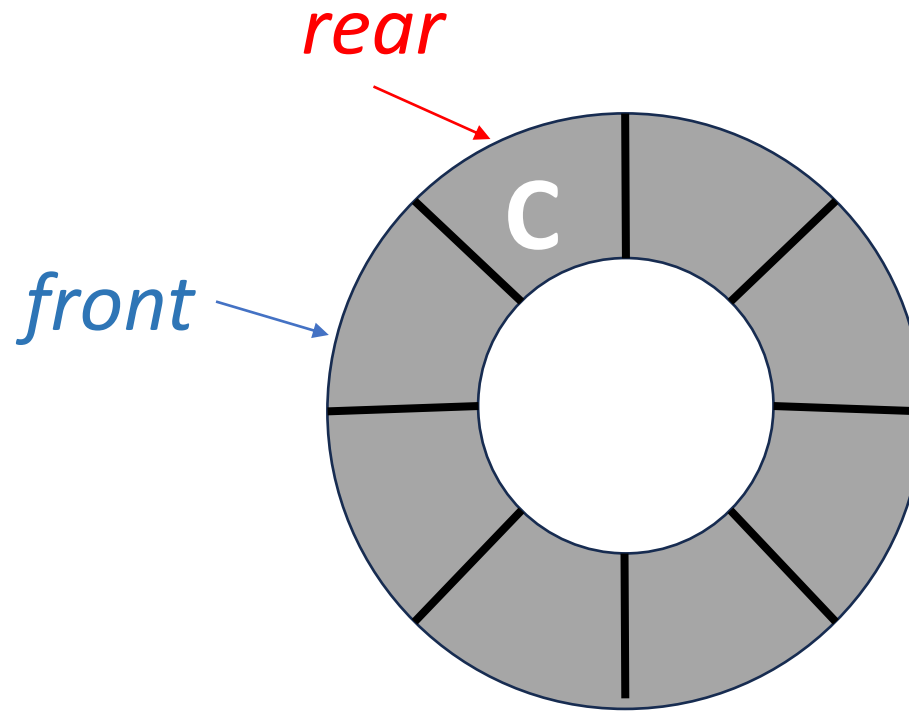
Emptying the circular queue



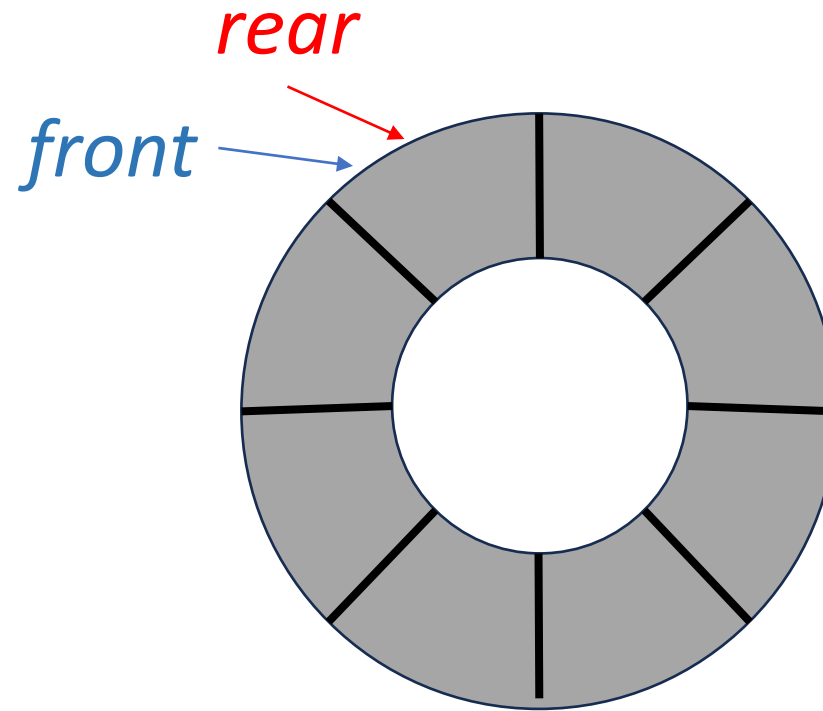
Emptying the circular queue



Emptying the circular queue

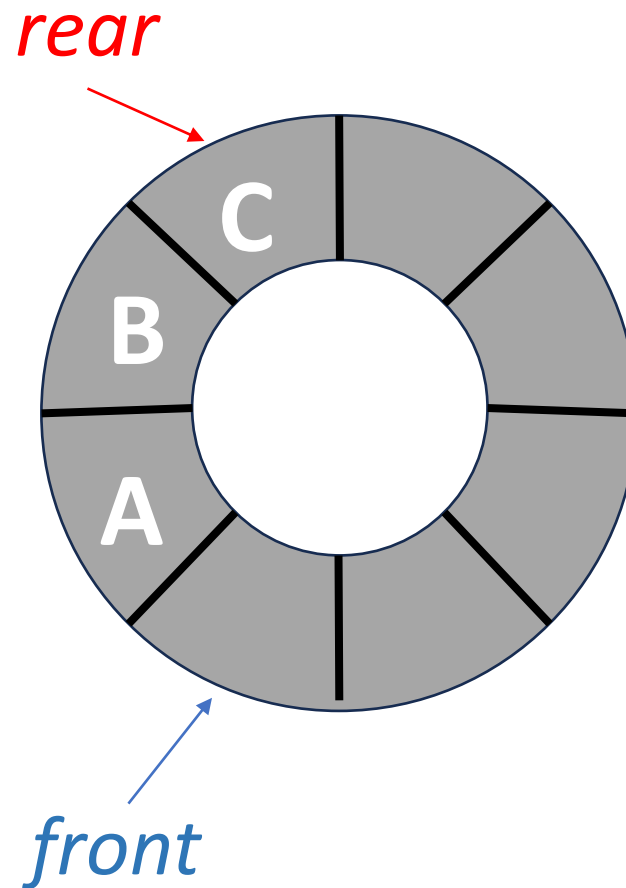


Emptying the circular queue

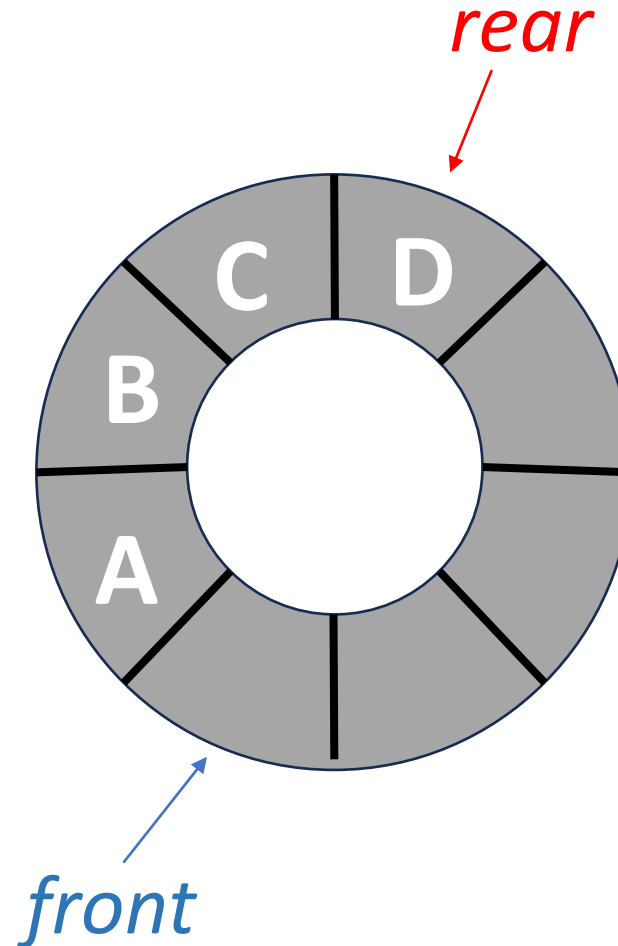


If (*front*==*rear*)
return queueEmpty();

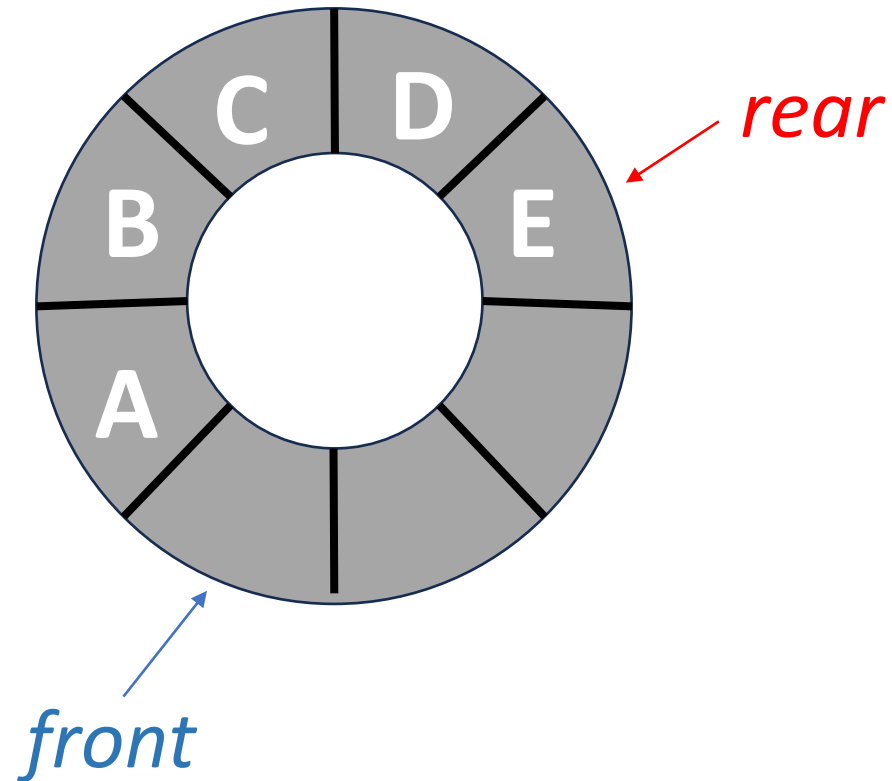
Addition until the circular queue is full



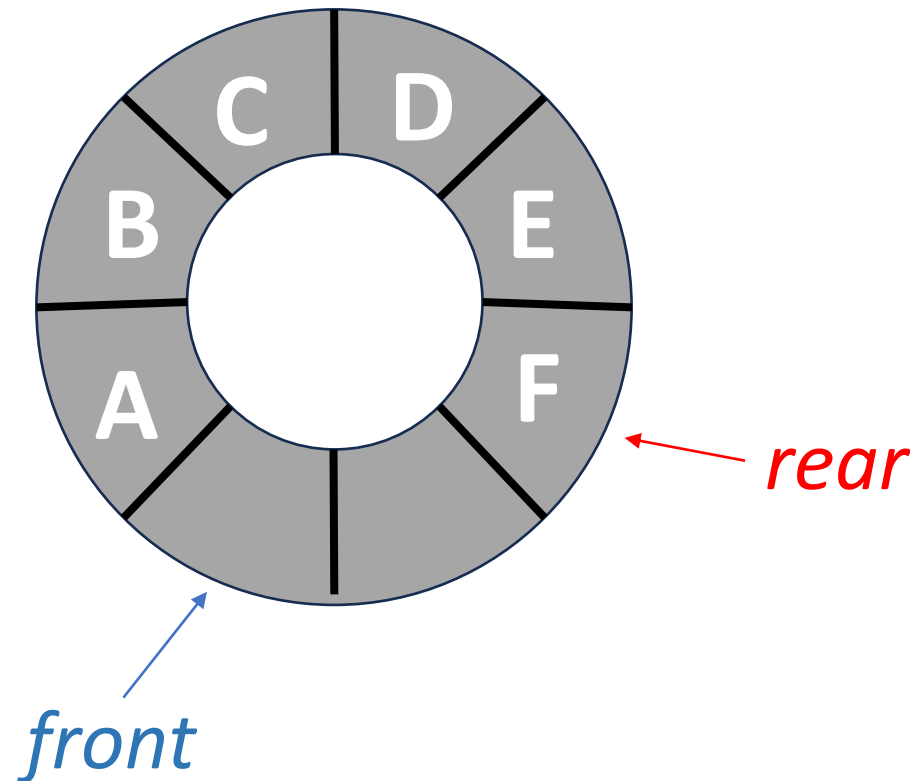
Addition until the circular queue is full



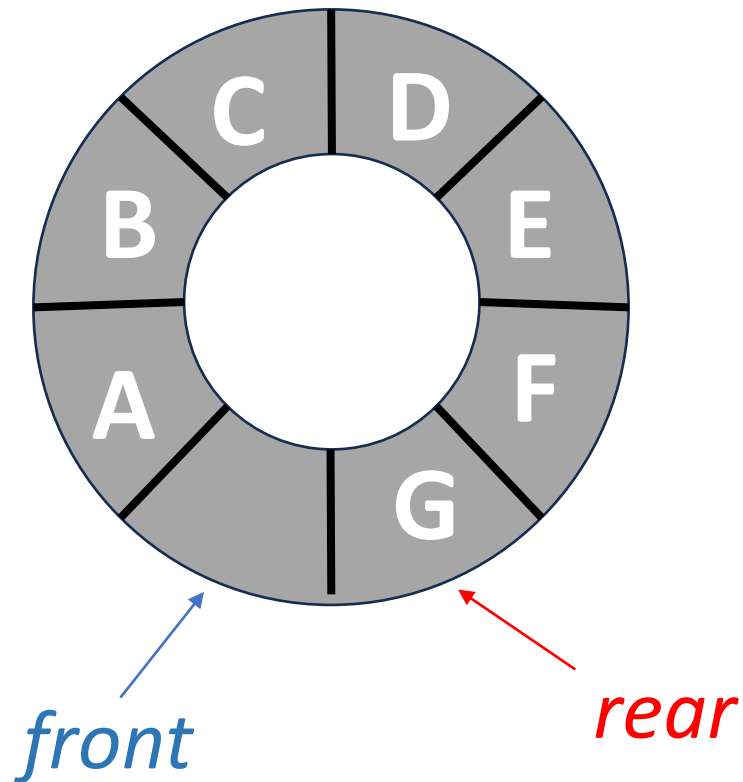
Addition until the circular queue is full



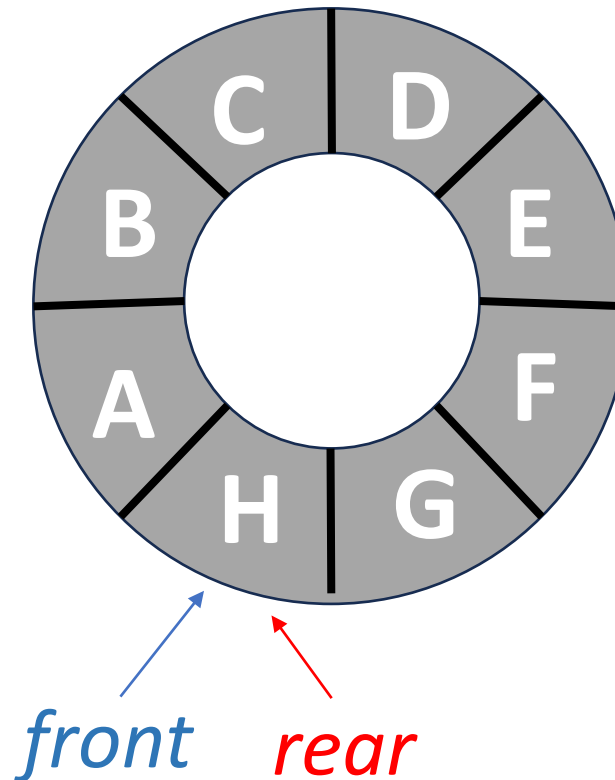
Addition until the circular queue is full



Addition until the circular queue is full



Addition until the circular queue is full



If (**front**==**rear**)
return queueFull();

Edge conditions should be considered.

- In DeleteQ, test for an empty queue is *front* == *rear*.
- In AddQ, test for a full queue is also *front* == *rear*.

We cannot distinguish between the case of empty and full circular queues.

Q6: What can we do to avoid this?

Please reply your answers via the following link:



Group members: 1~3 people

Hint: No standard answer. There are several ways.

Possible solutions

1. Don't let the queue get full.
 - When the addition of an element will signal *queueFull()*, increase array size.
2. Define a Boolean variable *lastOperationIsAddQ*.
 - Following each *AddQ*, set this variable to *True*.
 - Following each *DeleteQ*, set this variable to *False*.
 - Queue is *empty* if $front == rear$ and $\neg lastOperationIsAddQ$
 - Queue is *full* if $front == rear$ and $lastOperationIsAddQ$

Possible solutions

3. Define an integer variable **size**.
 - Following each **AddQ** do **size++**.
 - Following each **DeleteQ** do **size--**.
 - Queue is **empty** if (**size == 0**)
 - Queue is **full** if (**size == arrayLength**)