

Assignment Cover Letter (Group Work)

Student Information: Surname Given Names Student ID Number

 1. Chandra
 Jason Christopher
 2101725033

 2. Vallenchee
 Jose
 2101718482

 2.101699672
 2101699672

3. Joyan Vincent

Course Code :COMP6340 Course Name : Analysis of Algorithm

Class Major :L3BC Name of Lecturer(s) :1. Ms. Maria

:Computer Science Seraphina

Title of Assignment:

(if any)

Type of Assignment :Final Report Paper

Submission Pattern

Due Date :15th January 2019 Submission Date :15th January 2019

The assignment should meet the below requirements.

- 1. Assignment (hard copy) is required to be submitted on clean paper, and (soft copy) as per lecturer's instructions.
- 2. Soft copy assignment also requires the signed (hardcopy) submission of this form, which automatically validates the softcopy submission.
- 3. The above information is complete and legible.
- 4. Compiled pages are firmly stapled.
- 5. Assignment has been copied (soft copy and hard copy) for each student ahead of the submission.

Plagiarism/Cheating

BiNus International seriously regards all forms of plagiarism, cheating and collusion as academic offenses which may result in severe penalties, including loss/drop of marks, course/class discontinuity and other possible penalties executed by the university. Please refer to the related course syllabus for further information.

Declaration of Originality

By signing this assignment, we understand, accept and consent to BiNus International terms and policy on plagiarism. Herewith we declare that the work contained in this assignment is our own work and has not been submitted for the use of assessment in another course or class, except where this has been notified and accepted in advance.

1. Introduction

Background

Nowadays, Everyone saves most of their documents in a digital way. This is good because it means that they don't have to carry physical documents with them and this means that they can view their documents anywhere in the world. But, as we all know digital files will take up a lot of memory space.

For example, in the 2017 Panama Papers leak, 13.4 million secret documents were leaked to the public. These documents took up to 2 terabytes of data therefore we believe that compressing these file could have saved up a lot of memory space.

So, For our analysis of algorithms final project, we decided to compare the results between 2 lossless compression algorithms which can be used to compress text files.

Problems

While a single text file might be small in file size, we all know that if you have more files, more space will be used.

Aims

To efficiently compress text files without losing any data...

Benefits

To reduce amount of space used in electronics for text files.

Features

- Allows the user to choose between 2 lossless algorithms
- Allows the user to choose to compress or decompress.
- Create output file after compression and decompression.

2. Related Work

7zip, winrar, winzip are a few programs that execute lossless compression

3. Implementation

• Formal description of the problem

As more people use text files, we wanted to reduce the amount of space that text files take up by finding the most efficient lossless compression algorithm.

Design of algorithm

In this project, we implemented 2 algorithms:

- Lempel-Ziv (LZ-77)
- Used to losslessly compress and decompress text files
- Lempel-Ziv-Welch
- Used to losslessly compress and decompress text files

Proof of correctness

Compare the output file from the compression and the output file from decompression with the original text file.

Compression Time Complexity in milliseconds

N (Total Words)	LZ-77	Lempel-Ziv-Welch
500	660	10
1000	1200	20
10000	12520	120

Compression Space Complexity in bytes

N (Total Words)	LZ-77	Lempel-Ziv-Welch
500	28646	41141
1000	54786	41154
10000	546972	327873

Decompression Time Complexity in milliseconds

N (Total Words)	LZ-77	Lempel-Ziv-Welch
500	10	10
1000	10	10
10000	140	50

Decompression Space Complexity in bytes

N (Total Words)	LZ-77	Lempel-Ziv-Welch
500	28982	7382
1000	55798	11907
10000	530824	55542

4. Evaluation

- Theoretical analysis of the algorithms
- Lempel-Ziv (LZ-77)
- Lempel-Ziv-Welch
- Implementation Details
- ➤ Lempel-Ziv-Welch analysis:

```
import sys
      import time
3
4
       def create dict(mode='c-'):
           if mode == 'c-':
5
6
                dictionary = dict()
7
                for code in range (256):
8
                    dictionary[chr(code)] = code
9
                return dictionary
           elif mode == 'd-':
10
11
               dictionary = list()
12
               for code in range (256):
13
                    dictionary.append(chr(code))
14
               return dictionary
```

Creating a function to make a dictionary and appending 255 ASCII characters inside it.

```
16
        def compress(input file, output file):
17
            dictionary = create dict()
            temp string = str()
18
19
            code = int()
20
            temp list = list()
21
            temp char = str()
            space var = int()
22
23
            encoded file = open(output file, 'wb')
24
25
            with open (input file, 'r') as data file:
26
                temp char = data file.read(1)
27
28
                while temp char:
29
                     temp string += temp char
30
31
                     if len(dictionary) >= 65536:
32
                         del dictionary
33
                         dictionary = create dict()
34
35
                     if temp string not in dictionary:
                         code = len(dictionary)
36
                         dictionary[temp string] = code
37
38
                         temp list = list(temp string)
                         del temp list[-1]
39
                         temp string = "".join(temp list)
40
```

Creating the compress function to declare variables (even though it is not needed just to prevent confusion), then open the file and reads it.

Using while loop to read the character one by one until and add it to the string variable. Length of the dictionary is checked first to make sure it doesn't reach 65536 entries. 65536 is the max value of the string's code because it uses 2 bytes / 16 bit only. There is also if condition to check if the string is not in the dictionary.

```
if temp_string not in dictionary:

code = len(dictionary)
dictionary(temp_string] = code
temp_list = list(temp_string)
del temp_list[-1]
temp_string = "".join(temp_list)
temp_string = "".join(temp_list)
encoded_file.write(dictionary(temp_string).to_bytes(2, byteorder='big', signed=False))

temp_string = temp_char
temp_char = data_file.read(1)

if temp_string:
encoded_file.write(dictionary(temp_string).to_bytes(2, byteorder='big', signed=False))

encoded_file.write(dictionary(temp_string).to_bytes(2, byteorder='big', signed=False))

encoded_file.close()
space_var += (sys.getsizeof(dictionary) + sys.getsizeof(temp_string) + sys.getsizeof(code) + sys.getsizeof(temp_list) + sys.getsizeof(temp_char))

return space_var
```

If it's not, it will append the string and its code to the dictionary. Then remove the last character from the string first and after that, write it to the file in bytes. The next string will then use the last read character and read a new one from the source file

If there is any string left after the source file's content has been read all, it will be written to the file in bytes. After the file is closed, the space used is calculated and returned.

```
def decompress(input_file, output_file):
53
           dictionary = create_dict('d-')
54
           temp string = str()
55
           code = int()
56
           space var = int()
57
           decoded_file = open(output_file, 'w')
58
59
           with open(input file, 'rb') as encoded file:
60
               code = int.from_bytes(encoded_file.read(2), byteorder='big', signed=False)
61
62
               while code:
63
                    if code > len(dictionary):
64
                       raise IndexError("Out of range index, invalid code")
65
                   if len(dictionary) >= 65536:
67
                        del dictionary
68
                        dictionary = create dict('d-')
69
                    if code == len(dictionary):
                        dictionary.append(temp_string + temp_string[0])
71
72
                    elif temp_string:
73
                        dictionary.append(temp_string + dictionary[code][0])
74
75
                    decoded file.write(dictionary[code])
                    temp_string = dictionary[code]
76
        compress() > with open(input... > while temp_char > if len(dictiona...
```

It works first by creating the dictionary with 256 ascii pre-initialized first, same as the compress function. It will then read the bytes, 2 at each time one by one. It will then convert the bytes to int. If the code (int) is bigger than the current length of the dictionary. That means, the dictionary is corrupted. That is because the dictionary is built incrementally (relative with the dictionary's length) along with reading the file. If the code is the same as the dictionary's length (not yet exists), it will be added to the dictionary (the string and its first character) firsthand. If the code is below the current dictionary's length (already exists), the string combined with the first character of the string which the code refers to will be appended to the dictionary.

```
74
                    decoded file.write(dictionary[code])
76
                    temp string = dictionary[code]
                    code = int.from bytes(encoded file.read(2), byteorder='big', signed=False)
78
79
           decoded_file.close()
80
           space_var += (sys.getsizeof(dictionary) + sys.getsizeof(temp_string) + sys.getsizeof(code))
81
           return space_var
84
       start time = time.time()
       space_var = int()
86
87
      while True:
           action = input("'c-' for compression\n'd-' for decompression\n'e-' for exit\n> ")
89
90
           if action == 'e-':
91
             sys.exit()
           input_file = input("Type the input file\n> ")
93
94
           output_file = input("Type the output file\n> ")
95
96
           if action == 'c-':
               space var = compress(input file, output file)
        decompress() > with open(input... > while code > if len(dictiona...
```

After that, it will write the string that the code refers to to the new file and use it for the next string. Last, it will calculate all the spaces that are used and return it.

```
action = input("'c-' for compression\n'd-' for decompression\n> ")
input file = input("Type the input file\n> ")
output file = input("Type the output file\n> ")
if action == 'c-':
   start time = time.time()
   space var, normal size, compressed size = lzw.compress(input file, output file)
   print('File successfully compressed')
   print('Compression ratio: %.2f' % float(100 - (compressed_size / normal_size * 100)))
   print("Runtime: %.2f seconds" % (time.time() - start_time))
   print("Space Complexity: %d bytes" % space_var)
elif action == 'd-':
   start time = time.time()
    space var = lzw.decompress(input file, output file)
   print('File successfully decompressed')
   print("Runtime: %.2f seconds" % (time.time() - start time))
   print("Space Complexity: %d bytes" % space_var)
```

Creating a while loop for the input and output file when the file is compressed and decompressed, also printing out the runtime, compression ratio and space complexity.

➤ Lempel-Ziv (LZ-77):

```
4
       class LZ77:
5
           def __init__(self):
6
               pass
7
8
           def filename_validation(self, filename):
             if '.' in filename:
10
                   return True
12
               return False
13
14
           def append filename (self, filename, string):
15
               if '.' in filename:
16
                   i = filename.index('.')
17
                   filename = filename[:i] + string + filename[i:]
18
               else:
19
                  filename += string
20
21
               return filename
22
```

Filename validation checks if the file name has a "." file extension. If it does, it returns true, if it doesn;t it returns false, while append filename is used to rewrite the output file.

```
def file_to_list(self, filename):
    # Read data from input file, store it in a list and then close

    data = []
    file_input = open(filename, 'rb')

    while True:
        byte = file_input.read(1)
        if not byte:
            break
        data.append(byte)

    file_input.close()
    return data
```

Reads data from input file, checks if it is byte and closes it. First,we open the file as read in binary mode, then we used a while loop to read the input file one byte at a time and checks if it is in the byte data type. If it isn't the code breaks. If it is, the byte from the file is stored into the list. After the whole file has been read, the file is closed and the list containing the byte data is returned.

```
def copy(self, data, start, end):
    length = len(data) - start
    ret_data = []

for i in range(0, end - start):
        ret_data.append(data[start:len(data)][i % length])

return ret_data

return ret_data
```

```
def output_write(self, filename, data):
    # Writes every element from a list of bytes to a file
    file_output = open(filename, 'wb')

for byte in data:
        file_output.write(byte)
    file_output.close()
```

Used to write every element from a list of bytes to a file. First, it opens the output file as wb which means that it is writing to a file in binary mode, then we used a for loop to write every element from the list into the output file. After that,we close the output file.

```
def CommonSubseq_length(self, sequence_1, sequence_2):
    i = 0
    while i < len(sequence_1) and i < len(sequence_2):
        if not sequence_1[i] == sequence_2[i]:
            break
        i += 1</pre>
return i
```

```
def longest_prefix(self, data, position, win_size, pv_size):
    start = 0 + (position > win_size) * (position - win_size)
   longest start = start
   longest length = 0
    if position + pv_size < len(data):</pre>
        end = position + pv size
    else:
       end = len(data)
    index = start
    while index < position:
       commonseq_length = self.CommonSubseq_length(data[index:end], data[position:end])
        if commonseq length > longest length:
            longest_length = commonseq_length
            longest start = abs(position - index)
        index += 1
    return longest_start, longest_length
```

```
def compress(self, data):
   compressed data = []
   literal mode = {'status': False, 'position': 0}
   # The index of this loop indicates the position of the sliding window. Everything before
   # the index is already compressed and everything after is uncompressed.
   index = 0
   while index < len(data):
       longest_start, longest_length = self.longest_prefix(data, index, 127, 10)
       # When the length of the longest sequence of data that can be compressed is less than
       # three bytes, it's not worth it to compress this sequence since compression adds two
       # bytes of overhead. In this case, the better option is to turn on 'literal mode' and
       # keep adding the next byte to the compressed data until a sequence is found that is
       # worth compressing.
if longest_length < 3:
    if literal_mode['status'] == True:
        # If literal mode is already on, get the current length of the literal sequence
        # and add one since one byte is about to be added
        literal length = int.from bytes(compressed data[literal mode['position']], byteorder='big')
        literal length += 1
        compressed_data[literal_mode['position']] = literal_length.to_bytes(1, byteorder='big')
        compressed_data.append(data[index])
        # Exit literal mode if the maximum length has been reached. If the next byte
        # is still supposed to be part of a literal sequence, a new literal sequence
        # will have to be started
        if compressed data[literal mode['position']] == 255:
           literal mode['status'] = False
```

```
else:
           # When literal mode is first turned on, one byte is added to the beginning of the
           # sequence. The highest order bit indicates that literal mode is on, and the
           # remaining seven bits encode the length of the literal bytes to follow
           compressed_data.append((129).to_bytes(1, byteorder='big'))
           compressed data.append(data[index])
           literal mode['status'] = True
           literal mode['position'] = len(compressed data) - 2
       # When literal mode is on, only one byte will be added to the compressed data
       # on each iteration
       index += 1
   else:
       # Set the next byte variable equal to the next byte in the data
       # list that comes after the data to be added to the dictionary.
       # If the index of this byte is past the end of the list, set
       # next byte equal to the end character instead. The end character
       # is '$' which has the decimal value 36 in UTF-8 encoding.
       if index + longest_length < len(data):
           next byte = data[index + longest length]
        # If the index of this byte is past the end of the list, set
        # next byte equal to the end character instead. The end character
        # is '$' which has the decimal value 36 in UTF-8 encoding.
        if index + longest_length < len(data):
            next byte = data[index + longest length]
        else:
            next_byte = (36).to_bytes(1, byteorder='big')
        # Append the compressed data to the comp data list
        compressed_data.append(longest_start.to_bytes(1, byteorder='big'))
        compressed data.append(longest length.to bytes(1, byteorder='big'))
        compressed data.append(next byte)
        # If literal mode was on, turn it off
        literal_mode['status'] = False
        # Since longest len+1 bytes were just added to the compressed data,
        # advance the list index ahead by that many bytes
        index += longest length + 1
return compressed data
```

```
def decompress(self, data):
    decompressed data = []
   literal mode = {'status': False, 'length': 0}
   index = 0
   while index < len(data):
        if literal mode['status'] == True:
            # If literal mode is on, simply append the next byte onto the
            # decompressed data list. Decrement the count of remaining
            # literal bytes in this sequence by one
            decompressed data += [data[index]]
            literal mode['length'] -= 1
            index += 1
            # If there are no remaining literal bytes in this sequence, exit
            # literal mode
            if literal_mode['length'] == 0:
               literal mode['status'] = False
```

```
else:
   parameter_start = int.from_bytes(data[index], byteorder='big')
   if parameter start > 128:
        # If the first bit of the parameter start is on, the other
       # seven bits encode the length of the literal bytes to follow
       literal mode['status'] = True
       literal mode['length'] = parameter start - 128
       index += 1
   else:
       parameter length copy = int.from bytes(data[index + 1], byteorder='big')
       start = len(decompressed_data) - parameter_start
       end = start + parameter_length_copy
       # Copy data onto the end of the decomp list from the dictionary, which is
       # made up of data that was previously decompressed. Also add the next
        # byte onto the end of the list as long as it's not the end character ('$').
        # If one or more '$' characters were part of the actual compressed data,
        # distinguish them from the end character by checking the value of the
```

```
* copy data onto the end of the decomp fist from the dictionary, which is 
# made up of data that was previously decompressed. Also add the next
              \sharp byte onto the end of the list as long as it's not the end character ('$').
              # If one or more '$' characters were part of the actual compressed data,
              # distinguish them from the end character by checking the value of the # loop index to see if it's at the end of the compressed data or not
              decompressed data += self.copy(decompressed data, start, end) + [((not data[index + 2] == b'$') or (index < len(data) - 3)) * data[index + 2]]
              index += 3
    return decompressed_data
def 1z77 main(1z77):
     action = input("c- for compression\nd- for decompression\n> ")
     input_file = input("Type your input file\n> ")
     output_file = input("Type your output file\n> ")
if action == "c-":
           if output_file == '':
               output_file = lz77.append_filename(input_file, ' comp')
           start_time = time.time()
          data = 1z77.file to list(input file)
           # Pass data to be compressed to the compress function. Store returned compressed data in the
           # compressed data variable
           compressed data = 1z77.compress(data)
           # Write compressed data to output file
           1z77.output write(output file, compressed data)
           print('File successfully compressed')
           print('Compression ratio is: {0:1.2f}'.format(len(compressed data) / len(data)))
           print("Time Complexity: %s seconds" % (time.time() - start time))
```

The main of the Iz77, here the user will have to type the input and output file. Then the file will be appended inside the Iz77 function and therefore will be compressed, and the time complexity along with the compression ratio will be printed out.

```
# Write compressed data to output file
   1z77.output write(output file, compressed data)
   print('File successfully compressed')
   print('Compression ratio is: {0:1.2f}'.format(len(compressed_data) / len(data)))
   print("Time Complexity: %s seconds" % (time.time() - start time))
elif action == "d-":
   if output_file == '':
      output_file = lz77.append_filename(input_file, ' decomp')
   start time = time.time()
   data = lz77.file_to_list(input_file)
   # Pass data to be decompressed to the decompress function. Store returned decompressed data in the
   # decompressed data variable
   decompressed data = 1z77.decompress(data)
   # Write decompressed data to output file
   lz77.output_write(output_file, decompressed_data)
   print('File successfully decompressed')
   print("Time Complexity: %s seconds" % (time.time() - start_time))
```

Else if function is used to produce the decompressed file along with its time complexity.

```
algorithm = str()
1z77 = L277()
1zw = LZW()

while True:
    algorithm = input("Which algorithm do you want to use?\n1. LZ77\n2. LZW\n3. Exit\n> ")
    if algorithm == '1':
        lz77_main(lz77)
    elif algorithm == '2':
        lzw_main(lzw)
    elif algorithm == '3':
        sys.exit()
```

The Iz77 and Izw function are called here inside the while loop where the user will have to choose to run either the compressor or exit.

Explanations of each algorithms:

- 1. Lempel-Ziv (LZ-77):
 - LZ77 algorithms achieve compression by replacing repeated occurrences
 of data with references to a single copy of that data existing earlier in the
 uncompressed data stream. A match is encoded by a pair of numbers
 called a *length-distance pair*, which is equivalent to the statement "each of
 the next *length* characters is equal to the characters exactly *distance*characters behind it in the uncompressed stream".

 LZ77 maintains a sliding window during compression (the encoder must keep track of some amount of the most recent data, such as the last 2 kB, 4 kB, or 32 kB. The structure in which this data is held is called a sliding window, which is why LZ77 is sometimes called sliding-window compression).

2. Lempel-Ziv-Welch:

• Encodes sequences of 8-bit data as fixed-length 12-bit codes. The codes from 0 to 255 represent 1-character sequences consisting of the corresponding 8-bit character, and the codes 256 through 4095 are created in a dictionary for sequences encountered in the data as it is encoded. At each stage in compression, input bytes are gathered into a sequence until the next character would make a sequence for which there is no code yet in the dictionary. The code for the sequence (without that character) is added to the output, and a new code (for the sequence with that character) is added to the dictionary.

Encoding:

- o 1. Initialize the dictionary to contain all strings of length one.
- 2. Find the longest string W in the dictionary that matches the current input.
- 3. Emit the dictionary index for W to output and remove W from the input.
- 4. Add W followed by the next symbol in the input to the dictionary.
- 5. Go to step 2.

Decoding :

The decoding algorithm works by reading a value from the encoded input and outputting the corresponding string from the initialized dictionary. In order to rebuild the dictionary in the same way as it was built during encoding, it also obtains the next value from the input and adds to the dictionary the concatenation of the current string and the first character of the string obtained by decoding the next input value, or the first character of the string just output if the next value can not be decoded (If the next value is unknown to the decoder, then it must be the value that will be added to the dictionary this iteration, and so its first character must be the same as the first character of the current string being sent to decoded

- output). The decoder then proceeds to the next input value (which was already read in as the "next value" in the previous pass) and repeats the process until there is no more input, at which point the final input value is decoded without any more additions to the dictionary.
- In this way the decoder builds up a dictionary which is identical to that used by the encoder, and uses it to decode subsequent input values. Thus the full dictionary does not need to be sent with the encoded data; just the initial dictionary containing the single-character strings is sufficient (and is typically defined beforehand within the encoder and decoder rather than being explicitly sent with the encoded data).

5. Discussion:

- We chose to use the Lempel-Ziv (LZ-77) because of its efficiency in data compression which is very suitable for our situation.
- We chose to use the Lempel-Ziv-Welch because it is a universal lossless data compression algorithm which is simple to implement and has the potential for very high throughput in hardware implementations.
- As our purpose is to compare 2 algorithms, these 2 algorithms is very suitable to be compared as the Lempel-Ziv-Welch is the improved implementation of LZ-78 (which is pretty similar to LZ-77 but the differences are LZ-78 is faster than LZ-77 but doesn't always achieve as high a compression ratio as LZ-77. The biggest advantage LZ-78 has over the LZ-77 algorithm is the reduced number of string comparisons in each encoding step. The only similarity between LZ-77 & LZ-78 is that they have fast decompression but slow compression).

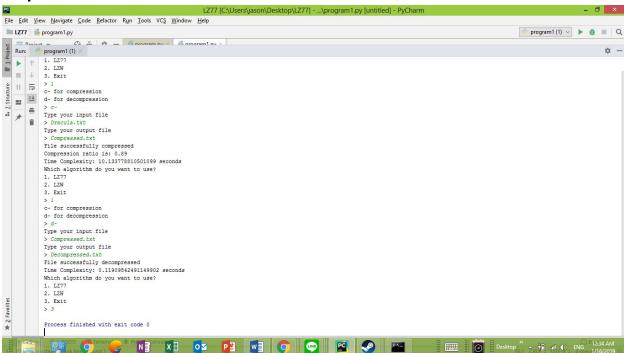
 We will be comparing these 2 codes in terms of their time & space complexity and their better usage on which conditions.

6. Conclusion & Recommendation:

★ As expected, the Lempel-Ziv-Welch algorithm is better at both compression and decompression in terms of both time complexity and space complexity in almost every situation. So, it is more recommended to use the Lempel-Ziv-Welch algorithm compared to the LZ-77.

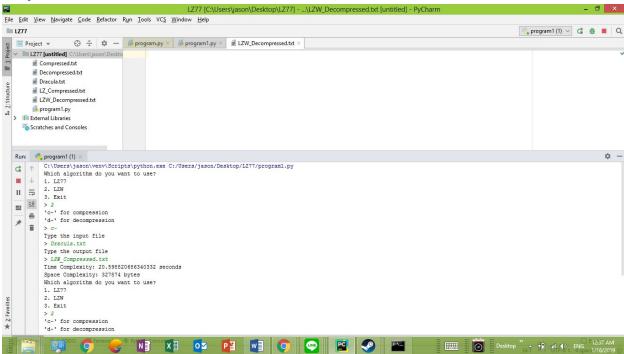
7. Program's Manual:

Step 1:



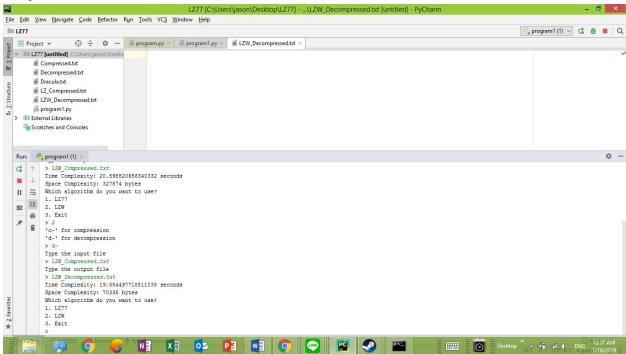
There are 3 choices to be made: 1), 2) & 3). If the user press 1, then the encoding for Lempel-Ziv (LZ-77) will run. If the user press 2, then the encoding for Lempel-Ziv-Welch will run. If the user press 3, then the system will stop running.

Step 2:



If the user press either 1) or 2) then the user will have to type either "c-" or "d-". In which "c-" is for compression and "d-" is for decompression.

Step 3:



The user will have to type in the input and output file to be compressed/decompressed in order to run the program.

After the program is done running, it will shot the time and space complexity. Then the program will go back to **Step 1** where the user has to make the first choices again.

*Notes : The text file must be in the same folder as the program in order to run.