- [§1.1] 2. (a) This is not a declarative sentence so it is not a proposition.
 - (b) This is not a declarative sentence either; not a proposition.
 - (c) Google indicates black flies do exist in Maine, so this is a FALSE proposition.
 - (d) No value is assigned to x; not a proposition.
 - (e) The moon is made of rock, not green cheese; FALSE proposition.
 - 3. (a) "Mei does not have an MP3 player."
 - (b) "There is pollution in New Jersey."
 - (c) " $2 + 1 \neq 3$."
 - (d) "The summer in Maine is not hot or not sunny."

6.

Smartphone	RAM (MB)	ROM (GB)	Camera (MP)	
A	256	32	8	
В	288	64	4	
\overline{C}	128	32	5	

- (a) TRUE. 288 > 256 and 288 > 128.
- (b) TRUE. 32 > 64 but 5 > 4.
- (c) FALSE. 288 > 256 and 64 > 32, but $4 \ge 8$.
- (d) FALSE. "If p then q" is false:
 - p: "B has more RAM and more ROM than C" is true; $(288 > 128) \land (64 > 32)$.
 - q: "B has a higher resolution camera than C" is false; $4 \ge 5$.
- (e) FALSE. "p if and only if q" is false:
 - p: "A has more RAM than B" is false; $256 \ge 288$.
 - q: "B has more RAM than A" is true; 288 > 256.
- 10. (b) "The election is decided or the votes have been counted."
 - (d) "If the votes have been counted, then the election is decided."
 - (f) "If the election is not decided, then the votes have not been counted."
- 12. (a) "If you have the flu, then you miss the final examination."
 - (c) "If you miss the final examination, then you do not pass the course."
 - (e) "If you have the flu, then you do not pass the course, or if you miss the final examination, then you do not pass the course."
- 17. These are all of the form "if p then q," which is true when q is true or p is false.
 - (a) p is true and q is false, so the statement is FALSE.
 - (b) p is false, so the statement is TRUE.

- (c) p is false, so the statement is TRUE.
- (d) p is false (monkeys cannot fly) so the statement is TRUE.
- 22. (a) "If one gets promoted, then one has washed the boss's car."
 - (d) "If he cheats, then Willy gets caught."
- 32. (a)

p	$\neg p$	$p \to \neg p$
Т	F	F
F	Т	${ m T}$

(e)

١								
,	p	q	$\neg p$	$q \to \neg p$	$p \leftrightarrow q$	$(q \to \neg p) \leftrightarrow (p \leftrightarrow q)$		
	Т	Т	F	F	Т	F		
	Τ	F	F	Τ	F	F		
	F	Т	Т	Т	F	F		
	F	F	Γ	Т	T	Т		

37. (a)

\						
)	p	q	r	$\neg q$	$\neg q \lor r$	$p \to (\neg q \lor r)$
	Т	Т	Т	F	Т	T
	Т	Т	F	F	F	F
	Τ	F	Т	T	Т	${ m T}$
	Т	F	F	Γ	Т	${ m T}$
	F	Т	Т	F	Т	${ m T}$
	F	Т	F	F	F	${ m T}$
	\mathbf{F}	F	Т	Τ	Т	${ m T}$
	F	F	F	T	Т	${ m T}$

(d)

1)	p	q	r	$\neg p$	$p \rightarrow q$	$\neg p \rightarrow r$	$(p \to q) \land (\neg p \to r)$
	Т	Т	Т	F	Т	Т	T
	Τ	Т	F	F	T	Т	m T
	Т	F	Т	F	F	Т	F
	Т	F	F	F	F	Т	F
	F	Т	Т	Т	Т	Т	${ m T}$
	F	Т	F	Т	Т	F	F
	F	F	Т	T	Т	Т	m T
	F	F	F	T	Т	F	F

[§1.3] 1. (a)

p	$p \wedge \mathbf{T}$
Τ	T
F	F

(b)

p	$p \vee \mathbf{F}$
Τ	Т
F	F

(c)

)	p	$p \wedge \mathbf{F}$	\mathbf{F}
	Т	F	F
	F	F	F

(d)	p	$p \lor \mathbf{T}$	\mathbf{T}
	Т	Т	Т
	\mathbf{F}	Т	$\mid T \mid$

(0)		
(e)	p	$p \lor p$
	Т	Τ
	F	F

(f)		
(1)	p	$p \wedge p$
	Т	Т
	\mathbf{F}	F

6. $\neg (p \land q) \equiv \neg p \lor \neg q$ as shown in class on January 13:

p	q	$p \wedge q$	$\neg (p \land q)$	$\neg p$	$\neg q$	$\neg p \lor \neg q$
Τ	Т	${ m T}$	F	F	F	F
Τ	F	\mathbf{F}	T	F	Γ	Τ
\mathbf{F}	Τ	\mathbf{F}	T	T	F	Τ
\mathbf{F}	F	\mathbf{F}	T	T	T	${ m T}$

- 7. (a) $[\neg(p \land q) \equiv \neg p \lor \neg q]$ "Jan is not rich or not happy."
 - (b) $[\neg (p \vee q) \equiv \neg p \wedge \neg q]$ "Carlos will neither bicycle nor run tomorrow."
 - (c) $[\neg(p \lor q) \equiv \neg p \land \neg q]$ "Mei neither walks nor takes the bus to class."
 - (d) $[\neg (p \wedge q) \equiv \neg p \vee \neg q]$ "Ibrahim is not smart or not hard working."

Ω	(a)					
9.		p	q	$(p \wedge q)$	$(p \land q) \to p$	
		Т	Т	Т		
		T	F	F	al-ways T	
		F	Т	F	always T	
		E	E	l E		

(a)					
(c)	p	q	$\neg p$	$p \rightarrow q$	$\neg p \to (p \to q)$
	Т	Т	F	Т	
	Τ	F	F	F	1 00
	F	Т	Т	Т	always T
	F	F	Γ	T	

(0)							
(e)	p	q	$p \rightarrow q$	$\neg(p \to q)$	$\neg(p \to q) \to p$		
	Т	Т	Т	F			
	Τ	F	F	Γ	- 1 T		
	F	Т	Т	F	always T		
	\mathbf{F}	F	Т	F			

11. (a)

$$\begin{array}{ll} (p \wedge q) \rightarrow p \equiv p \vee \neg (p \wedge q) & \text{by logical equivalence involving conditionals (LEIC)} \\ \equiv p \vee (\neg p \vee \neg q) & \text{by the first De Morgan law} \\ \equiv (p \vee \neg p) \vee q & \text{by associativity} \\ \equiv \mathbf{T} \vee q & \text{by negation} \\ \equiv \mathbf{T} & \text{by the identity laws} \end{array}$$

(c)

(e)

17. $\neg(p \leftrightarrow q) \equiv p \leftrightarrow \neg q$ because both sides are true (or both sides are false) for all combinations of truth values of p and q:

or trath variate of p and q.							
	p	q	$\neg q$	$p \leftrightarrow q$	$\neg(p \leftrightarrow q)$	$p \leftrightarrow \neg q$	
	Τ	Т	F	Т	F	F	
	Τ	F	Т	F	T	Т	
	F	Т	F	F	$^{\mathrm{T}}$	T	
	\mathbf{F}	F	Т	m T	F	F	

32. Proof by counterexample. Suppose

$$p$$
: true q : false r : false

Then $(p \land q)$ is false, so $(p \land q) \rightarrow r$ is true.

Since $(p \to r)$ is false, $(p \to r) \land (q \to r)$ is false.

Therefore $(p \land q) \to r$ is not logically equivalent to $(p \to r) \land (q \to r)$.