

MODULE *Assumes*
EXTENDS *Integers, Sequences*

You can run this as a model using “No behavior spec” mode
Single line comment

ASSUME
 $\wedge \text{TRUE} = \text{TRUE}$
 $\wedge \neg \text{FALSE} = \text{TRUE}$

$Jason \triangleq \text{“jason”}$
ASSUME
 $Jason = \text{“jason”}$

$record \triangleq [name \mapsto \text{“jason”}, age \mapsto 2]$
ASSUME
 $\wedge record.name = \text{“jason”}$
 $\wedge record.name \neq \text{“foo”}$

ASSUME
 $\forall F \in \{\text{TRUE}\} : F = F$

ASSUME
 $\forall F \in \{\text{FALSE}\} : F = F$

ASSUME \Rightarrow means “implies”, as in $A \Rightarrow B$ is “(not A) OR B”
 $\text{FALSE} \Rightarrow \text{TRUE} = \text{TRUE}$

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 $\text{FALSE} \Rightarrow \text{FALSE} = \text{TRUE}$

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 $\text{TRUE} \Rightarrow \text{FALSE} = \text{FALSE}$

ASSUME
 $\text{TRUE} \equiv \text{TRUE}$

ASSUME
 $\text{FALSE} \equiv \text{FALSE}$

ASSUME
 $\forall F, G \in \{\text{TRUE}, \text{FALSE}\} : (F \Rightarrow G) \equiv \neg F \vee G$

Sets

ASSUME

$$\{1, 2, 2, 2, 3\} = \{1, 2, 3\}$$

ASSUME

$$\{1, 2, 3, 3, 4, 4\} \setminus \{4\} = \{1, 2, 3\}$$

ASSUME

$$\{1, 2, 3\} \cup \{4, 5, 6\} = \{1, 2, 3, 4, 5, 6\}$$

ASSUME

$$\exists x \in \{3, 4, 5\} : x = 5$$

ASSUME

$$\{1, 3\} \subseteq \{3, 2, 1\}$$

Filtering a set

ASSUME

$$\{x \in 1 \dots 8 : x \% 2 = 1\} = \{1, 3, 5, 7\}$$

ASSUME

$$\{x \in 1 \dots 8 : x \% 2 = 1 \wedge \neg(x \% 5 = 0)\} = \{1, 3, 7\}$$

ASSUME

$$\{\langle x, y \rangle \in \{\langle 1, 2 \rangle, \langle 4, 2 \rangle\} : x > y\} = \{\langle 4, 2 \rangle\}$$

$$IsPrime(x) \triangleq x > 1 \wedge \neg \exists d \in 2 \dots (x - 1) : x \% d = 0$$

For all y in S such that y is not prime or y is less than or equal to x

$$\begin{aligned} LargestPrime(S) &\triangleq \text{CHOOSE } x \in S : \\ &\quad \wedge IsPrime(x) \\ &\quad \wedge \forall y \in S : \\ &\quad \quad IsPrime(y) \Rightarrow y \leq x \\ &\quad \text{or } y > x \Rightarrow \neg IsPrime(y) \end{aligned}$$

ASSUME

$$LargestPrime(\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}) = 7$$

$$IsEven(x) \triangleq x \% 2 = 0$$

$$\begin{aligned} LargetEven(S) &\triangleq \text{CHOOSE } x \in S : \\ &\quad \wedge IsEven(x) \\ &\quad \wedge \forall y \in S : \\ &\quad \quad IsEven(y) \Rightarrow y \leq x \end{aligned}$$

ASSUME
 $TargetEven(\{1, 2, 3, 4, 5, 5, 5\}) = 4$

ASSUME
 $\forall x \in \{\} : \text{FALSE}$

ASSUME
 $\forall x \in \{\} : \text{TRUE}$

ASSUME
 $\forall x \in \{\} : 7$

ASSUME
 $\forall x \in \{\text{FALSE}\} : \text{TRUE}$

ASSUME
 $\forall x \in \{\text{TRUE}\} : \text{TRUE}$

ASSUME
 $(\forall x \in \{\text{FALSE}\} : \text{FALSE}) = \text{FALSE}$

$IsCommutative(Op(-, -), S) \triangleq \forall x \in S :$
 $\forall y \in S : Op(x, y) = Op(y, x)$

$Add(x, y) \triangleq x + y$
 $Divide(x, y) \triangleq x \div y$

ASSUME
 $IsCommutative(Add, \{1, 2, 3\})$

ASSUME
 $IsCommutative(Divide, \{1, 2, 3\}) = \text{FALSE}$

ASSUME
 $IsCommutative(Divide, \{1, 2, 3\}) \Rightarrow \text{FALSE}$

ASSUME
 $IsCommutative(Divide, \{1, 2, 3\}) \Rightarrow \text{TRUE}$

ASSUME
 $\neg IsCommutative(Divide, \{1, 2, 3\})$

ASSUME
 $\neg \exists x \in \{1, 3, 5\} : IsEven(x)$

$Pick(S) \triangleq \text{CHOOSE } s \in S : \text{TRUE}$
 RECURSIVE $SetReduce(-, -, -)$
 $SetReduce(Op(-, -), S, value) \triangleq \text{IF } S = \{\} \text{ THEN } value$

$$\begin{aligned} & \text{ELSE LET } s \triangleq \text{Pick}(S) \\ & \text{IN } \text{SetReduce}(Op, S \setminus \{s\}, Op(s, \text{value})) \\ \\ Sum(S) & \triangleq \text{LET } _op(a, b) \triangleq a + b \\ & \text{IN } \text{SetReduce}(_op, S, 0) \\ \\ \text{ASSUME} \\ & Sum(\{1, 2, 3\}) = 6 \\ \\ Min(S) & \triangleq \text{CHOOSE } x \in S : \forall y \in S : x \leq y \\ \\ \text{ASSUME} \\ & Min(\{5, 3, 7, 10, 2, 9\}) = 2 \\ \\ Max(S) & \triangleq \text{CHOOSE } x \in S : \forall y \in S : x \geq y \\ \\ \text{ASSUME} \\ & Max(\{4, 6, 1, 2, 9, 3, 5\}) = 9 \\ \\ \text{ASSUME} \\ & \langle 1, 2, 3 \rangle \in Seq(\{1, 2, 3\}) \\ \\ \text{ASSUME} \\ & \langle 4 \rangle \notin Seq(\{1, 2, 3\}) \\ \\ \text{ASSUME} \\ & \langle 1, 2, 3, 4 \rangle \notin Seq(\{1, 2, 3\}) \end{aligned}$$

Sets of tuples.

$$\begin{aligned} chessboard_squares & \triangleq \{ \text{"a"}, \text{"b"}, \text{"c"}, \text{"d"}, \text{"e"}, \text{"f"}, \text{"g"}, \text{"h"} \} \times (1 \dots 8) \\ \\ \text{ASSUME} \\ & \wedge \langle \text{"a"}, 1 \rangle \in chessboard_squares \\ & \wedge \langle \text{"a"}, 2 \rangle \in chessboard_squares \\ & \wedge \langle \text{"a"}, 3 \rangle \in chessboard_squares \\ & \wedge \langle \text{"a"}, 4 \rangle \in chessboard_squares \\ \\ jason & \triangleq (1 \dots 2) \times \{ \text{"Jason"}, \text{"DeBolt"} \} \\ \\ \text{ASSUME} \\ & \wedge \langle 1, \text{"Jason"} \rangle \in jason \\ & \wedge \langle 2, \text{"Jason"} \rangle \in jason \\ & \wedge \langle 1, \text{"DeBolt"} \rangle \in jason \\ & \wedge \langle 2, \text{"DeBolt"} \rangle \in jason \end{aligned}$$

$$digits \triangleq \{ \text{"one"}, \text{"three"} \} \times \{ \text{"two"}, \text{"four"} \}$$

ASSUME

$$\begin{aligned} &\wedge \langle \text{"one"}, \text{"two"} \rangle \in digits \\ &\wedge \langle \text{"three"}, \text{"four"} \rangle \in digits \end{aligned}$$

$$A \triangleq \{1\}$$

$$B \triangleq \{2\}$$

$$C \triangleq \{3\}$$

ASSUME

$$\begin{aligned} &\wedge \langle 1, 2, 3 \rangle \in A \times B \times C \\ &\wedge \langle 1, \langle 2, 3 \rangle \rangle \in A \times (B \times C) \\ &\wedge \langle \langle 1, 2 \rangle, 3 \rangle \in (A \times B) \times C \end{aligned}$$

Structures.

Structures are hashes. They have keys and values. You specify them as $[\text{key} \mapsto \text{value}]$ and query them with either $[\text{"key"}]$ or $.key$. Both are legal and valid.

$$SomeHash \triangleq [x \mapsto 1, y \mapsto \{2, 3\}]$$

ASSUME

$$\begin{aligned} &\wedge SomeHash.x = 1 \\ &\wedge SomeHash[\text{"x"}] = 1 \\ &\wedge SomeHash.y = \{2, 3\} \\ &\wedge SomeHash[\text{"y"}] = \{2, 3\} \\ &\wedge \text{DOMAIN } SomeHash = \{ \text{"x"}, \text{"y"} \} \end{aligned}$$

Aside from that, there's one extra trick structures have. Instead of $\text{key} \mapsto \text{value}$, you can do $\text{key} : \text{set}$. In that case, instead of a structure you get the set of all structures which have, for each given key, a value in the set.

$$SetOfStructures \triangleq [x : \{1\}, y : \{2, 3, 4\}]$$

If you use $:$ syntax and any of the values are not sets, then the entire construct is invalid. In other words, while $[a: \{1\}, b: \{2, 3\}]$ is the above set, $[a: 1, b: \{2, 3\}]$ will throw an error if you try to use it.

ASSUME

$$\begin{aligned} &\wedge [x \mapsto 1, y \mapsto 2] \in SetOfStructures \\ &\wedge [x \mapsto 1, y \mapsto 3] \in SetOfStructures \\ &\wedge [x \mapsto 1, y \mapsto 4] \in SetOfStructures \end{aligned}$$

Type Composition

Any type can be squeezed inside any other type.

$crazy \triangleq [a \mapsto \{\langle \rangle, \langle 1, 2, 3 \rangle, \langle 3, 2, 1 \rangle\}, b \mapsto \langle [a \mapsto 0] \rangle]$

A function of keys mapping to sets of tuples or of keys mapping to tuples of functions.

ASSUME

$crazy.b[1].a = 0$ Remember that tuples are 1 indexed.

$blah \triangleq [name \mapsto \text{"jason"}, hobbies \mapsto [outdoor \mapsto \langle \text{"cycling"}, \text{"hiking"} \rangle, indoor \mapsto \langle \text{"reading"}, \text{"watching tv"} \rangle]]$

ASSUME

$\wedge blah.name = \text{"jason"}$

$\wedge blah.hobbies.outdoor[1] = \text{"cycling"}$

$sing \triangleq \langle \langle 4, 5, 6 \rangle, \langle \rangle, \langle \rangle \rangle$

ASSUME

DOMAIN $sing = \{1, 2, 3\}$

\ * Modification History
 \ * Last modified Sun Apr 21 21:47:34 PDT 2019 by jasondebolt
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