
MODULE *Assumes*

EXTENDS *Integers, Sequences, FiniteSets, Naturals*

You can run this as a model using “No behavior spec” mode
 Single line comment

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ASSUME
   $\wedge \text{TRUE} = \text{TRUE}$ 
   $\wedge \neg \text{FALSE} = \text{TRUE}$ 

Jason  $\triangleq$  “jason”
ASSUME
  Jason = “jason”

record  $\triangleq$  [name  $\mapsto$  “jason”, age  $\mapsto$  2]
ASSUME
   $\wedge \text{record.name} = \text{“jason”}$ 
   $\wedge \text{record.name} \neq \text{“foo”}$ 

ASSUME
   $\forall F \in \{\text{TRUE}\} : F = F$ 

ASSUME
   $\forall F \in \{\text{FALSE}\} : F = F$ 

ASSUME  $\Rightarrow$  means “implies”, as in  $A \Rightarrow B$  is “(not  $A$ ) OR  $B$ ”
   $\text{FALSE} \Rightarrow \text{TRUE} = \text{TRUE}$ 

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ASSUME
   $\text{TRUE} \equiv \text{TRUE}$ 

ASSUME
   $\text{FALSE} \equiv \text{FALSE}$ 

ASSUME
   $\forall F, G \in \{\text{TRUE}, \text{FALSE}\} : (F \Rightarrow G) \equiv \neg F \vee G$ 

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Sets

ASSUME

$$\{1, 2, 2, 2, 3\} = \{1, 2, 3\}$$

ASSUME

$$\{1, 2, 3, 3, 4, 4\} \setminus \{4\} = \{1, 2, 3\}$$

ASSUME

$$\{1, 2, 3\} \cup \{4, 5, 6\} = \{1, 2, 3, 4, 5, 6\}$$

ASSUME

$$\exists x \in \{3, 4, 5\} : x = 5$$

ASSUME

$$\{1, 3\} \subseteq \{3, 2, 1\}$$

ASSUME

$$(\forall i \in \{2, 4, 8\} : i \% 2 = 0) = \text{TRUE}$$

ASSUME

$$(\{1, 2\} \in \text{SUBSET } \{1, 2, 3\}) = \text{TRUE}$$

ASSUME

$$(\{1, 2\} \in \text{SUBSET } (\{1, 3\} \cup \{4, 2\})) = \text{TRUE}$$

ASSUME

$$\begin{aligned} &\wedge \{ \text{"one"}, \text{"two"} \} \neq \{ \} \\ &\wedge \{ \text{"one"}, \text{"two"} \} \neq \{ \} \\ &\wedge \{ \text{"one"}, \text{"two"} \} \setminus \{ \text{"one"} \} = \{ \text{"two"} \} \end{aligned}$$

ASSUME

$$\begin{aligned} &\wedge \text{IsFiniteSet}(\{1, 2, 3\}) \\ &\wedge \neg \text{IsFiniteSet}(\text{Nat}) \end{aligned}$$

ASSUME

$$\begin{aligned} &\wedge \text{Cardinality}(\{3, 4, 1\}) = 3 \\ &\wedge \text{Cardinality}(\{ \}) = 0 \end{aligned}$$

ASSUME

$$\{x \in 1 \dots 8 : x \% 2 = 1\} = \{1, 3, 5, 7\}$$

ASSUME

$$\{x \in 1 \dots 8 : x \% 2 = 1 \wedge \neg(x \% 5 = 0)\} = \{1, 3, 7\}$$

ASSUME
 $\{\langle x, y \rangle \in \{\langle 1, 2 \rangle, \langle 4, 2 \rangle\} : x > y\} = \{\langle 4, 2 \rangle\}$

IF STATEMENTS

ASSUME
 $\wedge (\text{IF } 1 < 3 \text{ THEN } 1 \text{ ELSE } 0) = 1$
 $\wedge (\text{IF } 1 < 3 \text{ THEN IF } 2 > 1 \text{ THEN } 6 \text{ ELSE } 4 \text{ ELSE } 7) = 6$

LET STATEMENTS

ASSUME
 $\wedge \text{LET } x \triangleq 6 \text{ IN } x \in \{6, 7\}$

For all y in S such that y is not prime or y is less than or equal to x

$IsPrime(x) \triangleq x > 1 \wedge \neg \exists d \in 2 \dots (x-1) : x \% d = 0$

$LargestPrime(S) \triangleq \text{CHOOSE } x \in S :$
 $\quad \wedge IsPrime(x)$
 $\quad \wedge \forall y \in S :$
 $\quad \quad IsPrime(y) \Rightarrow y \leq x$
 $\quad \text{or } y > x \Rightarrow \neg IsPrime(y)$

ASSUME
 $LargestPrime(\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}) = 7$

$IsEven(x) \triangleq x \% 2 = 0$

$LargetEven(S) \triangleq \text{CHOOSE } x \in S :$
 $\quad \wedge IsEven(x)$
 $\quad \wedge \forall y \in S :$
 $\quad \quad IsEven(y) \Rightarrow y \leq x$

ASSUME
 $LargetEven(\{1, 2, 3, 4, 5, 5, 5\}) = 4$

ASSUME
 $\forall x \in \{\} : \text{FALSE}$

ASSUME
 $\forall x \in \{\} : \text{TRUE}$

ASSUME

$\forall x \in \{\} : 7$

ASSUME

$\forall x \in \{\text{FALSE}\} : \text{TRUE}$

ASSUME

$\forall x \in \{\text{TRUE}\} : \text{TRUE}$

ASSUME

$(\forall x \in \{\text{FALSE}\} : \text{FALSE}) = \text{FALSE}$

$IsCommutative(Op(-, -), S) \triangleq \forall x \in S : \forall y \in S : Op(x, y) = Op(y, x)$

$Add(x, y) \triangleq x + y$

$Divide(x, y) \triangleq x \div y$

ASSUME

$IsCommutative(Add, \{1, 2, 3\})$

ASSUME

$IsCommutative(Divide, \{1, 2, 3\}) = \text{FALSE}$

ASSUME

$IsCommutative(Divide, \{1, 2, 3\}) \Rightarrow \text{FALSE}$

ASSUME

$IsCommutative(Divide, \{1, 2, 3\}) \Rightarrow \text{TRUE}$

ASSUME

$\neg IsCommutative(Divide, \{1, 2, 3\})$

ASSUME

$\neg \exists x \in \{1, 3, 5\} : IsEven(x)$

$Pick(S) \triangleq \text{CHOOSE } s \in S : \text{TRUE}$

RECURSIVE $SetReduce(-, -, -)$

$SetReduce(Op(-, -), S, value) \triangleq \text{IF } S = \{\} \text{ THEN } value$
 $\text{ELSE LET } s \triangleq Pick(S)$
 $\text{IN } SetReduce(Op, S \setminus \{s\}, Op(s, value))$

$Sum(S) \triangleq \text{LET } _op(a, b) \triangleq a + b$
 $\text{IN } SetReduce(_op, S, 0)$

ASSUME

$Sum(\{1, 2, 3\}) = 6$

$Min(S) \triangleq \text{CHOOSE } x \in S : \forall y \in S : x \leq y$

ASSUME

$$\text{Min}(\{5, 3, 7, 10, 2, 9\}) = 2$$

$$\text{Max}(S) \triangleq \text{CHOOSE } x \in S : \forall y \in S : x \geq y$$

ASSUME

$$\text{Max}(\{4, 6, 1, 2, 9, 3, 5\}) = 9$$

SEQUENCES

$\text{Seq}(S)$ is the set of all finite sequences of set S .

ASSUME

$$\wedge \langle \rangle \in \text{Seq}(\{1, 0\})$$

$$\wedge \langle 1 \rangle \in \text{Seq}(\{1, 0\})$$

$$\wedge \langle 0 \rangle \in \text{Seq}(\{1, 0\})$$

$$\wedge \langle 0, 0 \rangle \in \text{Seq}(\{1, 0\})$$

$$\wedge \langle 0, 1 \rangle \in \text{Seq}(\{1, 0\})$$

$$\wedge \langle 1, 0 \rangle \in \text{Seq}(\{1, 0\})$$

$$\wedge \langle 1, 1 \rangle \in \text{Seq}(\{1, 0\})$$

ASSUME

$$\wedge \{\langle 0 \rangle\} \subseteq \text{Seq}(\{0, 1\})$$

ASSUME

$$\langle 1, 2, 3 \rangle \in \text{Seq}(\{1, 2, 3\})$$

ASSUME

$$\langle 4 \rangle \notin \text{Seq}(\{1, 2, 3\})$$

ASSUME

$$\langle 1, 2, 3, 4 \rangle \notin \text{Seq}(\{1, 2, 3\})$$

ASSUME

$$\wedge \langle 1, 2 \rangle \circ \langle 3, 4 \rangle = \langle 1, 2, 3, 4 \rangle$$

$$\text{LessThanThree}(x) \triangleq x < 3$$

ASSUME

$$\wedge \text{Head}(\langle 2, 3, 4 \rangle) = 2$$

$$\wedge \text{Tail}(\langle 2, 3, 4 \rangle) = \langle 3, 4 \rangle$$

$$\wedge \text{Append}(\langle 1, 2 \rangle, 3) = \langle 1, 2, 3 \rangle$$

$$\wedge \text{Len}(\langle 5, 2, 1 \rangle) = 3$$

$$\wedge \text{SubSeq}(\langle 9, 3, 5, 6 \rangle, 1, 3) = \langle 9, 3, 5 \rangle$$

$$\wedge \text{SelectSeq}(\langle 5, 2, 9 \rangle, \text{LessThanThree}) = \langle 2 \rangle$$

TUPLES

$chessboard_squares \triangleq \{“a”, “b”, “c”, “d”, “e”, “f”, “g”, “h”\} \times (1 .. 8)$

ASSUME

$\wedge \langle “a”, 1 \rangle \in chessboard_squares$
 $\wedge \langle “a”, 2 \rangle \in chessboard_squares$
 $\wedge \langle “a”, 3 \rangle \in chessboard_squares$
 $\wedge \langle “a”, 4 \rangle \in chessboard_squares$

$jason \triangleq (1 .. 2) \times \{“Jason”, “DeBolt”\}$

ASSUME

$\wedge \langle 1, “Jason” \rangle \in jason$
 $\wedge \langle 2, “Jason” \rangle \in jason$
 $\wedge \langle 1, “DeBolt” \rangle \in jason$
 $\wedge \langle 2, “DeBolt” \rangle \in jason$

$digits \triangleq \{“one”, “three”\} \times \{“two”, “four”\}$

ASSUME

$\wedge \langle “one”, “two” \rangle \in digits$
 $\wedge \langle “three”, “four” \rangle \in digits$

ASSUME

$\wedge \langle “one”, “two” \rangle \circ \langle “three” \rangle = \langle “one”, “two”, “three” \rangle$
 $\wedge \langle “one”, “two” \rangle \circ \langle “three” \rangle = \langle “one”, “two”, “three” \rangle$

$A \triangleq \{1\}$

$B \triangleq \{2\}$

$C \triangleq \{3\}$

ASSUME

$\wedge \langle 1, 2, 3 \rangle \in A \times B \times C$
 $\wedge \langle 1, \langle 2, 3 \rangle \rangle \in A \times (B \times C)$
 $\wedge \langle \langle 1, 2 \rangle, 3 \rangle \in (A \times B) \times C$

Structures.

Structures are hashes. They have keys and values. You specify them as $[key \mapsto value]$ and query them with either $[“key”]$ or $.key$. Both are legal and valid.

$SomeHash \triangleq [x \mapsto 1, y \mapsto \{2, 3\}]$

ASSUME

$\wedge SomeHash.x = 1$

$$\begin{aligned}
&\wedge \text{SomeHash}["x"] = 1 \\
&\wedge \text{SomeHash}.y = \{2, 3\} \\
&\wedge \text{SomeHash}["y"] = \{2, 3\} \\
&\wedge \text{DOMAIN } \text{SomeHash} = \{ "x", "y" \}
\end{aligned}$$

$$\begin{aligned}
\text{SomeHash2} &\triangleq [x \mapsto 1, y \mapsto \{2, 3\}] \\
\text{SomeHash3} &\triangleq [\text{SomeHash2} \text{ EXCEPT } !["x"] = 6]
\end{aligned}$$

ASSUME

$$\wedge \text{SomeHash3}.x = 6$$

Aside from that, there's one extra trick structures have. Instead of $\text{key} \mapsto \text{value}$, you can do $\text{key} : \text{set}$. In that case, instead of a structure you get the set of all structures which have, for each given key, a value in the set.

$$\text{SetOfStructures} \triangleq [x : \{1\}, y : \{2, 3, 4\}]$$

If you use $:$ syntax and any of the values are not sets, then the entire construct is invalid. In other words, while $[a: \{1\}, b: \{2, 3\}]$ is the above set, $[a: 1, b: \{2, 3\}]$ will throw an error if you try to use it.

ASSUME

$$\begin{aligned}
&\wedge [x \mapsto 1, y \mapsto 2] \in \text{SetOfStructures} \\
&\wedge [x \mapsto 1, y \mapsto 3] \in \text{SetOfStructures} \\
&\wedge [x \mapsto 1, y \mapsto 4] \in \text{SetOfStructures}
\end{aligned}$$

Functions

ASSUME

$$\begin{aligned}
&\wedge [\{1, 2, 3\} \rightarrow \{\text{"done"}\}] = \{\langle \text{"done"}, \text{"done"}, \text{"done"} \rangle\} \quad \text{Turns a function into a set of tuples} \\
&\wedge [\{\text{"a"}, \text{"b"}\} \rightarrow \{\text{"done"}\}] = \{[a \mapsto \text{"done"}, b \mapsto \text{"done"}]\} \quad \text{Turns a function into a set of structs} \\
&\wedge [\{\text{"a"}, \text{"b"}\} \rightarrow \{\text{"done"}, \text{"pc"}\}] = \{[a \mapsto \text{"done"}, b \mapsto \text{"done"}], \quad \text{Turns a function into a set of structs} \\
&\quad [a \mapsto \text{"done"}, b \mapsto \text{"pc"}], \\
&\quad [a \mapsto \text{"pc"}, b \mapsto \text{"done"}], \\
&\quad [a \mapsto \text{"pc"}, b \mapsto \text{"pc"}]\} \\
&\wedge [\{\text{"p1"}, \text{"p2"}\} \rightarrow \{\text{"a"}, \text{"b"}, \text{"c"}, \text{"done"}\}] = \{[p1 \mapsto \text{"a"}, p2 \mapsto \text{"a"}], \\
&\quad [p1 \mapsto \text{"a"}, p2 \mapsto \text{"b"}], \\
&\quad [p1 \mapsto \text{"a"}, p2 \mapsto \text{"c"}], \\
&\quad [p1 \mapsto \text{"a"}, p2 \mapsto \text{"done"}], \\
&\quad [p1 \mapsto \text{"b"}, p2 \mapsto \text{"a"}], \\
&\quad [p1 \mapsto \text{"b"}, p2 \mapsto \text{"b"}], \\
&\quad [p1 \mapsto \text{"b"}, p2 \mapsto \text{"c"}], \\
&\quad [p1 \mapsto \text{"b"}, p2 \mapsto \text{"done"}], \\
&\quad [p1 \mapsto \text{"c"}, p2 \mapsto \text{"a"}], \\
&\quad [p1 \mapsto \text{"c"}, p2 \mapsto \text{"b"}], \\
&\quad [p1 \mapsto \text{"c"}, p2 \mapsto \text{"c"}], \\
&\quad [p1 \mapsto \text{"c"}, p2 \mapsto \text{"done"}],
\end{aligned}$$

$$\begin{aligned}
& [p1 \mapsto \text{"done"}, p2 \mapsto \text{"a"}], \\
& [p1 \mapsto \text{"done"}, p2 \mapsto \text{"b"}], \\
& [p1 \mapsto \text{"done"}, p2 \mapsto \text{"c"}], \\
& [p1 \mapsto \text{"done"}, p2 \mapsto \text{"done"}] \} \\
\wedge [\{ \text{"p1"}, \text{"p2"}, \text{"p3"} \} \rightarrow \{ \text{"a"}, \text{"b"}, \text{"c"} \}] = & \{ [p1 \mapsto \text{"a"}, p2 \mapsto \text{"a"}, p3 \mapsto \text{"a"}], \\
& [p1 \mapsto \text{"a"}, p2 \mapsto \text{"a"}, p3 \mapsto \text{"b"}], \\
& [p1 \mapsto \text{"a"}, p2 \mapsto \text{"a"}, p3 \mapsto \text{"c"}], \\
& [p1 \mapsto \text{"a"}, p2 \mapsto \text{"b"}, p3 \mapsto \text{"a"}], \\
& [p1 \mapsto \text{"a"}, p2 \mapsto \text{"b"}, p3 \mapsto \text{"b"}], \\
& [p1 \mapsto \text{"a"}, p2 \mapsto \text{"b"}, p3 \mapsto \text{"c"}], \\
& [p1 \mapsto \text{"a"}, p2 \mapsto \text{"c"}, p3 \mapsto \text{"a"}], \\
& [p1 \mapsto \text{"a"}, p2 \mapsto \text{"c"}, p3 \mapsto \text{"b"}], \\
& [p1 \mapsto \text{"a"}, p2 \mapsto \text{"c"}, p3 \mapsto \text{"c"}], \\
& [p1 \mapsto \text{"b"}, p2 \mapsto \text{"a"}, p3 \mapsto \text{"a"}], \\
& [p1 \mapsto \text{"b"}, p2 \mapsto \text{"a"}, p3 \mapsto \text{"b"}], \\
& [p1 \mapsto \text{"b"}, p2 \mapsto \text{"a"}, p3 \mapsto \text{"c"}], \\
& [p1 \mapsto \text{"b"}, p2 \mapsto \text{"b"}, p3 \mapsto \text{"a"}], \\
& [p1 \mapsto \text{"b"}, p2 \mapsto \text{"b"}, p3 \mapsto \text{"b"}], \\
& [p1 \mapsto \text{"b"}, p2 \mapsto \text{"b"}, p3 \mapsto \text{"c"}], \\
& [p1 \mapsto \text{"b"}, p2 \mapsto \text{"c"}, p3 \mapsto \text{"a"}], \\
& [p1 \mapsto \text{"b"}, p2 \mapsto \text{"c"}, p3 \mapsto \text{"b"}], \\
& [p1 \mapsto \text{"b"}, p2 \mapsto \text{"c"}, p3 \mapsto \text{"c"}], \\
& [p1 \mapsto \text{"c"}, p2 \mapsto \text{"a"}, p3 \mapsto \text{"a"}], \\
& [p1 \mapsto \text{"c"}, p2 \mapsto \text{"a"}, p3 \mapsto \text{"b"}], \\
& [p1 \mapsto \text{"c"}, p2 \mapsto \text{"a"}, p3 \mapsto \text{"c"}], \\
& [p1 \mapsto \text{"c"}, p2 \mapsto \text{"b"}, p3 \mapsto \text{"a"}], \\
& [p1 \mapsto \text{"c"}, p2 \mapsto \text{"b"}, p3 \mapsto \text{"b"}], \\
& [p1 \mapsto \text{"c"}, p2 \mapsto \text{"b"}, p3 \mapsto \text{"c"}], \\
& [p1 \mapsto \text{"c"}, p2 \mapsto \text{"c"}, p3 \mapsto \text{"a"}], \\
& [p1 \mapsto \text{"c"}, p2 \mapsto \text{"c"}, p3 \mapsto \text{"b"}], \\
& [p1 \mapsto \text{"c"}, p2 \mapsto \text{"c"}, p3 \mapsto \text{"c"}] \}
\end{aligned}$$

Type Composition

Any type can be squeezed inside any other type.

$$crazy \triangleq [a \mapsto \{ \langle \rangle, \langle 1, 2, 3 \rangle, \langle 3, 2, 1 \rangle \}, b \mapsto \langle [a \mapsto 0] \rangle]$$

A function of keys mapping to sets of tuples or of keys mapping to tuples of functions.

ASSUME

$$crazy.b[1].a = 0 \quad \text{Remember that tuples are 1 indexed.}$$

$$blah \triangleq [name \mapsto \text{"jason"}, hobbies \mapsto [outdoor \mapsto \langle \text{"cycling"}, \text{"hiking"} \rangle, indoor \mapsto \langle \text{"reading"}, \text{"watching tv"} \rangle]]$$

ASSUME

$\wedge \text{blah.name} = \text{"jason"}$
 $\wedge \text{blah.hobbies.outdoor}[1] = \text{"cycling"}$

$\text{sing} \triangleq \langle \langle 4, 5, 6 \rangle, \langle \rangle, \langle \rangle \rangle$

ASSUME

DOMAIN $\text{sing} = \{1, 2, 3\}$

\ * Modification History
\ * Last modified Sun May 12 10:42:59 *PDT* 2019 by *jasondebolt*
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