System Programming HW3 report b07902129 何政勳

Stack frame	Base pointer	Stack pointer
Main	0x7fffffffdc50	0x7fffffffda60
dummy	0x7fffffffda50	0x7fffffff3df0
funct_1	0x7ffffff3de0	0x7ffffff3dc0
dummy	0x7ffffff3db0	0x7ffffffea150
funct_2	0x7ffffffea140	0x7ffffffea120
dummy	0x7ffffffea110	0x7ffffffe04b0
funct_3	0x7ffffffe04a0	0x7ffffffe0480
dummy	0x7ffffffe0470	0x7ffffffd6810
funct_4	0x7fffffd6800	0x7ffffffd67e0
	Main dummy funct_1 dummy funct_2 dummy funct_3 dummy	Main 0x7ffffffdc50 dummy 0x7fffffffda50 funct_1 0x7fffffff3de0 dummy 0x7fffffff3db0 funct_2 0x7ffffffea140 dummy 0x7ffffffea110 funct_3 0x7ffffffe04a0 dummy 0x7ffffffe0470

use gdb to find the address of each function

- 1. gcc -g hw3.c -o hw3
- 2. gdb hw3
- 3. b 95 //set breakpoint just before the long jump to main
- 4. r 5 3 1 0 //run hw3 with arguments 5 3 1 0
- 5. frame [0-8] //select a frame indexing from 0 to 8
- 6. info frame // show the information of the selected frame
- 7. info reg // show the contents of registers of the selected frame

B. Yes, setjmp saves the stack pointer of the environment, when longjmp is called, it can restore (rewind) the stack to the previous status. If a variable is saved in stack memory and the stack memory isn't damaged by other function, then the variable can be restored.

C. Dummy function is used to separate all the others function to prevent damage on the memory of stack frames. For example, if there wasn't a dummy function, after calling the scheduler function at the first time in main function, the stack frame of funct 1 would be damaged.

D. No, it can't. The stack frames of dummy, funct_3, dummy ... aren't restored. In gdb I set up two break points, the first one is right before funct_4's return, the second one is right before dummy's return. The funct_4 ones showed up correctly, however, before the breakpoint in dummy showed up, an error "stack smashing detected" came up. Also the process received a SIGABRT.

E. I finished my program mainly by checking up course slides, Linux programmer's manual and discussion with b07902125. One of the hardest part is understanding the task 3 since there are much more things to do than task 1 and 2. The queue in task 3 is kind of confusing, which takes me some time to figure out what it is. Also sigaction is used since signal is unsafe and system depending.