

Problem Set #3

Due Tuesday, November 12th (before the class)

Problem 1 (10pts)

Consider a network with MPLS enabled routers as shown in Figure 1 below. We would like to perform traffic engineering using MPLS so that traffic from R1 to R6 will be routed as R1->R3->R5->R6->A and traffic from R2 to R6 will be routed as R2->R3->R4->C.

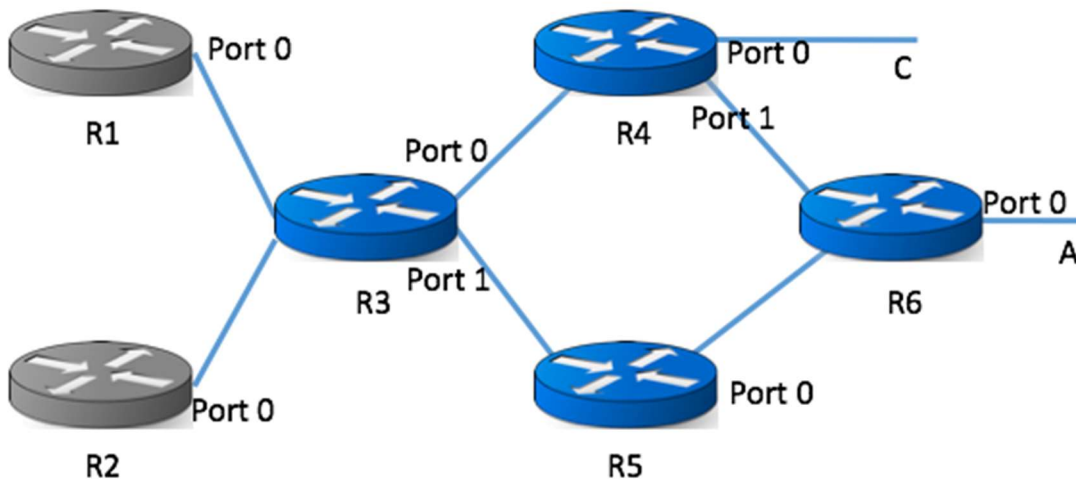


Figure 1. MPLS enabled network for Problem 1

Please fill in the following tables of MPLS entries for each router.

R1			
In label	Out label	Dst	Out interface
-	1	A	0

R2			
In Label	Out Label	Dst	Out interface
-	2	C	0

R3			
In label	Out label	Dst	Out interface
1	4	A	1
2	3	C	0

R4			
In label	Out label	Dst	Out interface
3	-	C	0
3	X1	A	1

R5			
In label	Out label	Dst	Out interface
4	5	A	0

R6			
In label	Out label	Dst	Out interface
5	-	A	0

Problem 2 (20pts)

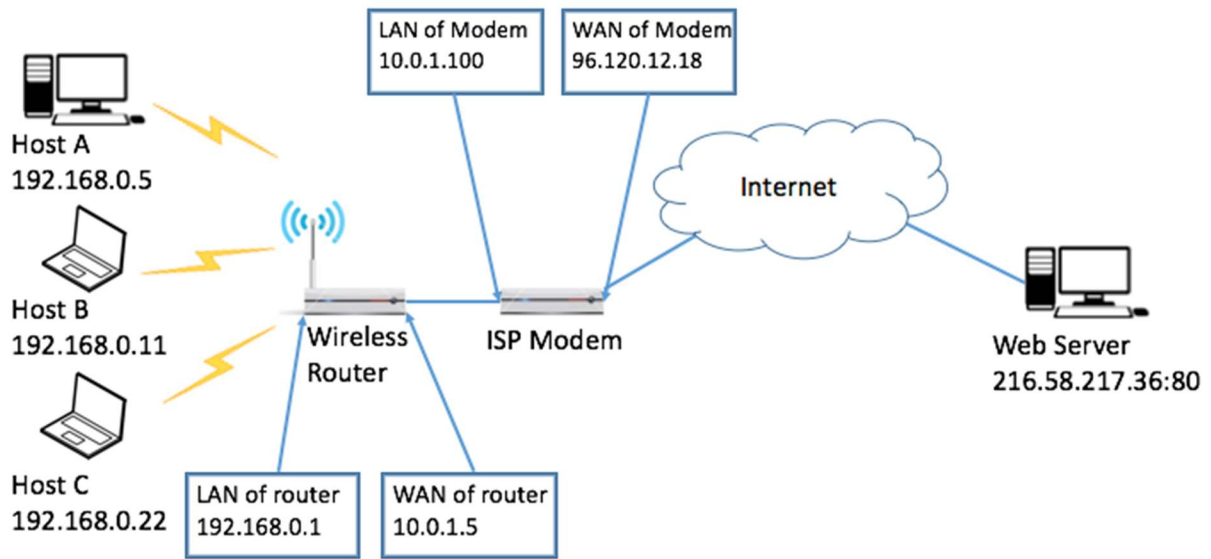


Figure 2. Network setup for Problem 2.

The Figure 2. above is a typical home network setup. An ISP Modem provides internet service; a wireless router is connected to the ISP Modem via Ethernet. Hosts A, B and C are connected to the wireless router to access the Internet.

- (a) In order for the hosts A, B, and C to access the Web Server, Network Address Translation (NAT) with random port mapping needs to be enabled for both the Wireless Router and the ISP Modem. Assume Hosts will pick a random port number between 8000 and 9000, the Wireless Router can choose a random port number between 2000 and 2500, and the ISP Modem can choose a random port number between 3000 and 4000. Please fill in the NAT table for the Wireless Router and the ISP Modem below.

NAT Table of Wireless Router	
LAN side	WAN side
192.168.0.22:8001	10.0.1.5:2001
192.168.0.11:8002	10.0.1.5:2002
192.168.0.5:8003	10.0.1.5:2003

NAT Table of ISP Modem	
LAN side	WAN side
10.0.1.5:2001	96.120.12.18:3001
10.0.1.5:2002	96.120.12.18:3002
10.0.1.5:2003	96.120.12.18:3003

- (b) Now we look into the details about how packets are exchanged between Host B and Web Server. Assume Host B sends a HTTP request packet to Web Server. And Web Server then sends HTTP content back to Host B. Please fill in the tables below to show how the packet's IP header changed along the route. (Please formulate your answer based on your answers for (a).)

HTTP request Before entering Router	
Src IP	192.168.0.11
Src Port	8002
Dst IP	216.58.217.36
Dst Port	80

HTTP request After exiting Router	
Src IP	10.0.15
Src Port	2002
Dst IP	216.58.217.36
Dst Port	80

HTTP request After exiting Modem	
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Src IP	96.120.12.18
Src Port	3002
Dst IP	216.58.217.36
Dst Port	80

HTTP response Before entering Modem	
Src IP	216.58.217.36
Src Port	80
Dst IP	96.120.12.18
Dst Port	3002

HTTP response After exiting Modem	
Src IP	216.58.217.36
Src Port	80
Dst IP	10.0.1.5
Dst Port	2002

HTTP response After exiting Router	
Src IP	216.58.217.36
Src Port	80
Dst IP	192.168.0.11
Dst Port	8002

- (c) Suppose now Host A also runs a webserver on port 8888, it is attached to a domain name <http://www.mylocalhomeserver.com>, explain **what NAT entries**

should be added so that people from the internet can assess this webserver via URL. You can assume that the above domain name is registered properly.

The NAT table of the wireless Router and the Wireless Modem need to be updated.

NEW NAT Table of Wireless Router	
LAN side	WAN side
192.168.0.22:8001	10.0.1.5:2001
192.168.0.11:8002	10.0.1.5:2002
192.168.0.5:8003	10.0.1.5:2003
192.168.0.5:8008	10.0.1.5:8888

NEW NAT Table of ISP Modem	
LAN side	WAN side
10.0.1.5:2001	96.120.12.18:3001
10.0.1.5:2002	96.120.12.18:3002
10.0.1.5:2003	96.120.12.18:3003
10.0.1.5:8888	96.120.12.18:8888

- (d) The wireless link at the last mile is very error prone and you would like to improve the performance. What would you do in this case?

Move your servers one mile closer.

Problem 3 (10pts)

Suppose a router has three input flows and one output flow. It receives the packets listed in the Table 1. below, all at about the same time, in the order listed, during a period in which the output port is busy but all queues are otherwise empty. Give the order in which the packets are transmitted, assuming:

- (a) Fair queuing

Packet	Size	Flow	F
1	200	1	200

2	200	1	400
3	160	2	160
4	200	2	360
5	160	2	520
6	210	3	210
7	120	3	330
8	90	3	420

P3, P1, P6, P7, P4, P2, P8, P5

- (b) Weighted fair queuing with flow 2 having twice as much share as flow 1, and flow 3 having 1.5 times as much share as flow 1. Note that ties are to be solved in the order of flow1, flow2 and flow3.

Packet	Size	Flow
1	200	1
2	200	1
3	160	2
4	200	2
5	160	2
6	210	3
7	120	3
8	90	3

Table 1.

Packet	Size	Flow	F
1	200	1	100
2	200	1	200
3	160	2	40
4	200	2	90
5	160	2	130
6	210	3	70
7	120	3	110
8	90	3	132.5

P3, P6, P4, P1, P7, P5, P8, P2

Problem 4 (15pts)

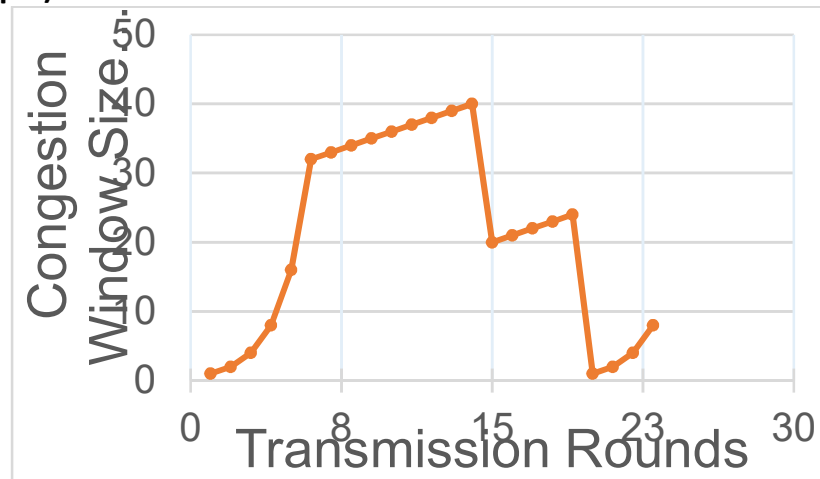


Figure 3. Congestion Window Size

Assuming TCP Reno is the protocol experiencing the behavior shown above, answer the following questions:

- (a) Identify the RTT rounds when TCP runs Slow Start (e.g., from the 1th round to which round?)

Slow start for first 6 RTT and last 4 RTT

- (b) Identify the RTT rounds when TCP runs Congestion Avoidance

Congestion Avoidance runs from the 6th to the 14th transmission

- (c) After the 14th RTT round, is segment loss detected by a triple duplicate ACK or by a timeout and why?

Triple duplicate ACK. Timeout puts the window to 0.

- (d) During which RTT round the 170th segment is sent?

$1 (1) + 2 (3) + 4 (7) + 8 (15) + 16 (31) + 32 (63) + 33 (96) + 34 + 35 + 36 (165) = 201$; 10th transmission

- (e) Assuming a packet loss is detected after the 23th RTT round by the receipt of triple duplicate ACKs, what will be the value of the congestion window?

The congestion window halves, so it is 4.

Problem 5 (15pts)

Figure 4. below shows how 2 disconnected LAN are connected by IP tunnel (the dash line). For each interface the IP and MAC addresses are shown in the figure. (HW1-HW14 are used to represent hardware addresses)

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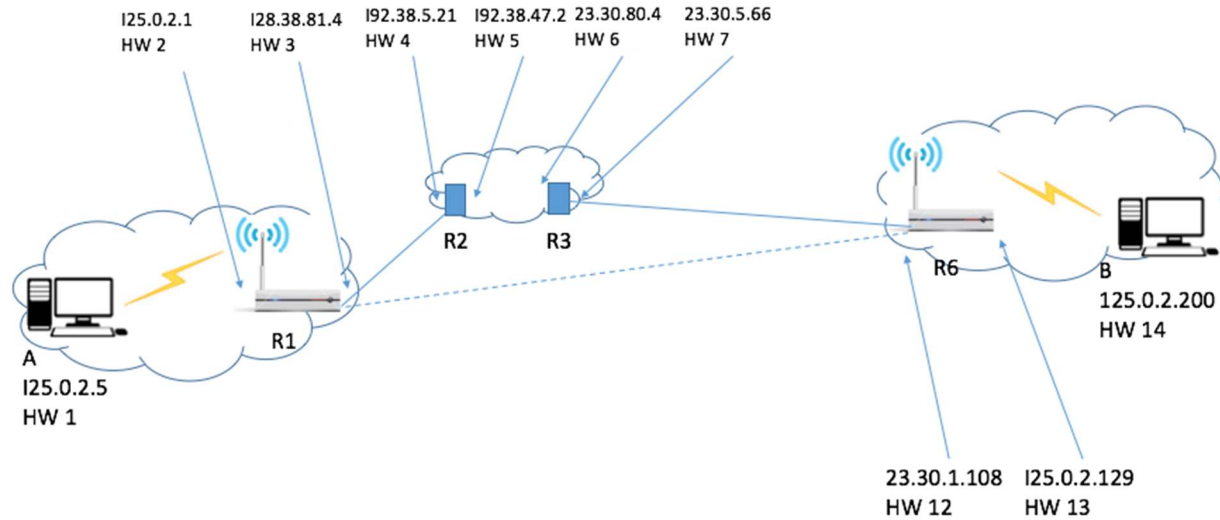


Figure 4. Network setup for Problem 5

Now Host B sends a packet to Host A. Please show how the packet travels along the route, please describe header information along the route.

B -> R6	
Src IP	125.0.2.200
Src Mac	HW14
Dest IP	125.0.2.5
Dest Mac	HW13

R6 -> R3	
Src IP	23.30.1.108
Src Mac	HW12
Dest IP	128.38.81.4
Dest Mac	HW7

R3 -> R2	
Src IP	23.30.1.108
Src Mac	HW6
Dest IP	128.38.81.4
Dest Mac	HW5

R2 -> R1	
Src IP	23.30.1.108

Src Mac	HW4
Dest IP	128.38.81.4
Dest Mac	HW3

R1 -> A	
Src IP	125.0.2.200
Src Mac	HW2
Dest IP	125.0.2.5
Dest Mac	HW1

Problem 6 (20pts)

Derive the expected throughput of the following TCP congestion control algorithm: The additive increment factor is α . Multiplicative decrease factor β , which means after loss, the windows size will change from W to $(1-\beta)W$. Please order the throughput for each flow. AIMD(a,b) means the cwnd increases a per each round trip time and the cwnd set to $(1-b)W$ from W when the loss happens.

Flow1: AIMD(a=1,b=0.5), RTT=10ms, loss rate = 10^{-6}
Flow2: AIMD(a=2,b=0.2), RTT=100ms, loss rate = 10^{-8}
Flow3: AIMD(a=5,b=0.8), RTT=300ms, loss rate = 10^{-9}
Flow4: AIMD(a=8,b=0.4), RTT=1000ms, loss rate = 10^{-4}
Flow5: AIMD(a=6,b=0.5), RTT=100ms, loss rate = 10^{-10}

$$Throughput = \frac{\sqrt{2-\beta} \times \sqrt{\alpha}}{RTT \times \sqrt{2\beta\sqrt{p}}}$$

Flow1: 122.5 – pkts/ms
Flow2: 300 – pkts/ms
Flow3: 204.1 – pkts/ms
Flow4: 0.4 – pkts/ms
Flow5: 3000 – pkts/ms

Problem 7 (10pts)

Suppose that TCP uses the combination of quick acknowledgements (quick ack) and delayed acknowledgements (delayed ack). The quick ack only triggers up to 16 packets starting from 1 packet during slow start. The maximum capacity of the link is 5000 KBps, the RTT is 10ms, and 1MSS = 1KB. Note that KBps is KB per second).

(a) About what is cwnd at the time of first packet loss?

The Packet needs to be sent between the 12th RTT = 4096 and the 13th RTT = 8192

(b) About how long until sender discovers first loss?

The sender sends the packet and its lost. The loss of a packet triggers the triple duplicate ACK response. The sender sends three more packets and gets the ack from the lost packet.

$$d_{\text{total}} = 3 * d_{\text{Trans}} + d_{\text{prop}} = 3 * L/R + (1/2 \text{ RTT})$$

$$d_{\text{total}} = 3 * 1\text{KB}/5000\text{-KB/s} + 10\text{-ms} / 2 = 5.6\text{-ms}$$

A timeout loss is when the sender receives nothing from the receiver. This is normally 2RTT.