



# ECE 250 Algorithms and Data Structure Project Three: Leftist heap

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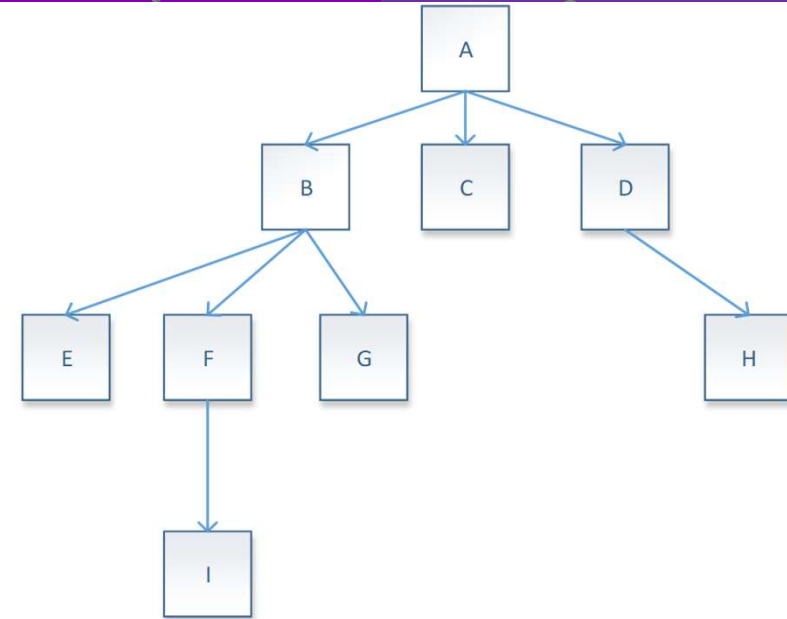
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# Tree Data Structure

- Several Data Structures:
  - Arrays:
    - Random Access
    - Easy to implement
  - Linked lists:
    - Ideal for frequent add, remove, and update
  - **Drawbacks:**
    - Time needed to search for an item:
      - Both arrays and linked lists are **linear structures**. The time needed to search a **linear list** is proportional to the size.
      - Imagine that you want to search for ***m*** items in an array with 1,000,000 elements. Even on a machine with 1 million comparison per second, this search takes ***m*** seconds which is not acceptable.
  - **Solution:**
    - Find more efficient data structures such as **tree**.



# Tree Data Structure



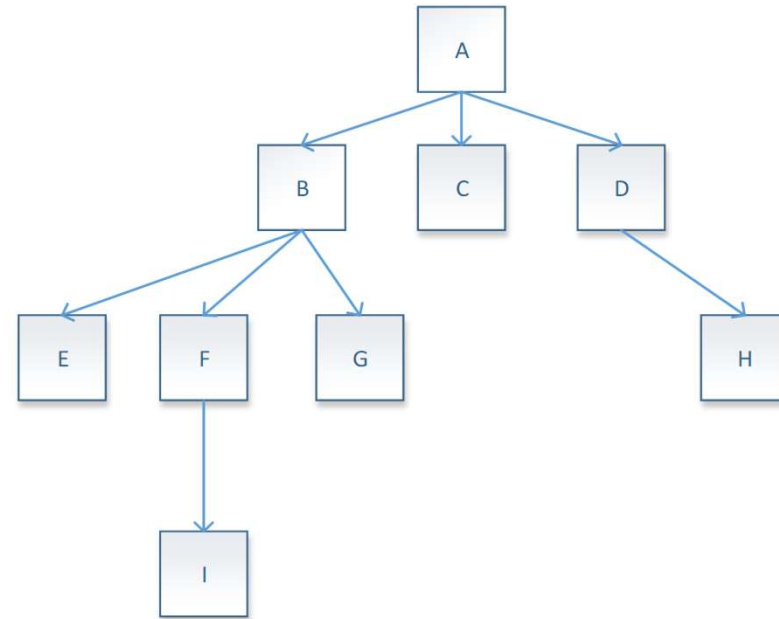
- Tree:
  - A collection of nodes connected by **edges**
  - Nonlinear (Hierarchical) compared to arrays, linked lists, stacks, and queues
  - Abstract Data Types (ADT)
  - Can be empty
  - One node is distinguished as a **root**
  - Every node (except root) is connected by a directed **edge** from **exactly** one other node



# Tree Data Structure

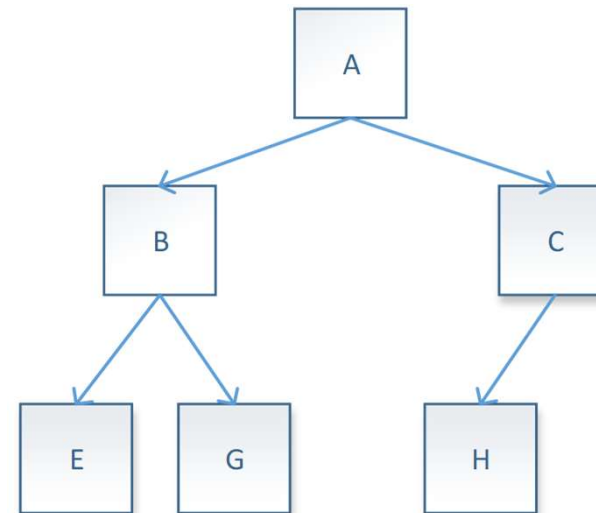
- Tree:

- A is the **parent** of B,C, and D
- B is a **child** of A
- B is the parent of E, F, and G
- B, C, and D are **siblings**
- Nodes with no children are called **leaves** or **external** nodes
- Nodes that are not leaves are called **internal** nodes
- The **depth** of a node is the number of edges from root to that node
- The **height** of a node is the number of edges from the node to the deepest leaf
- The **height** of a tree is the height of its root



# Tree Data Structure

- Binary Tree:
  - A tree where each node can have no more than **two children**:
    - The **left** child
    - The **right** child
  - E is the left child of B
  - G is the right child of B



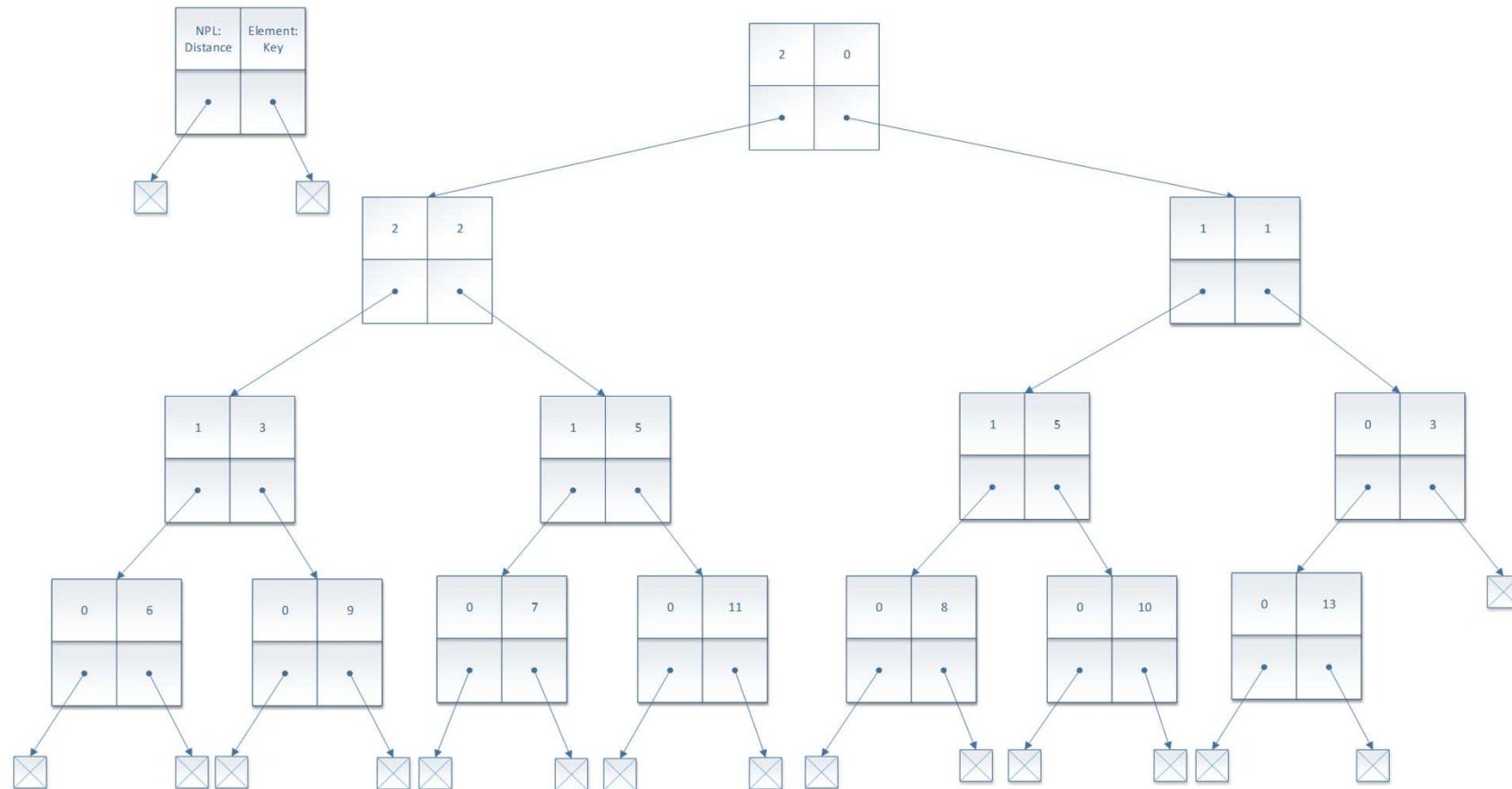


# Leftist Heap

- Leftist tree: is an implementation of a **mergeable heap**
  - Mergeable heap: is a **heap** that supports an additional operation (i.e., merging two heaps)
    - Heap: is a **tree-based** ADT that satisfies **heap property**
      - Heap property (Max heap): the key (value) of each parent node is greater than or equal to its children
      - Heap property (Min heap): the key (value) of each parent node is less than or equal to its children
- Usage: priority queues
- **Distance** (null\_path\_length): the number of edges in the shortest path from a node to its descendent external nodes
- Properties of leftist tree:
  - $\text{Key}(i) \geq \text{key}(\text{parent}(i))$
  - $\text{Distance}(\text{right}(i)) \leq \text{Distance}(\text{left}(i))$
- Note:
  - Each sub-tree of a leftist tree is also a leftist tree
  - $\text{Distance}(i) = 1 + \text{Distance}(\text{right}(i))$

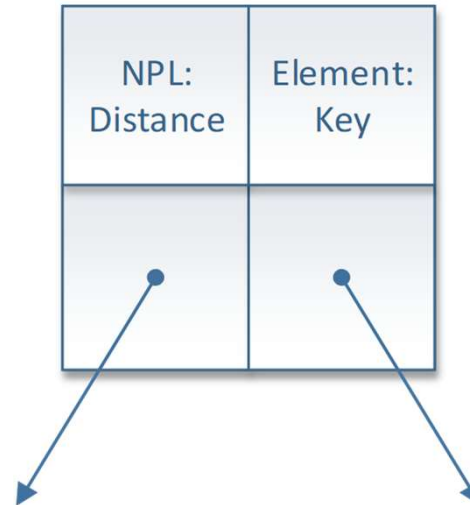


# Leftist Heap Example (i.e., int.in.txt)



# Leftist\_node

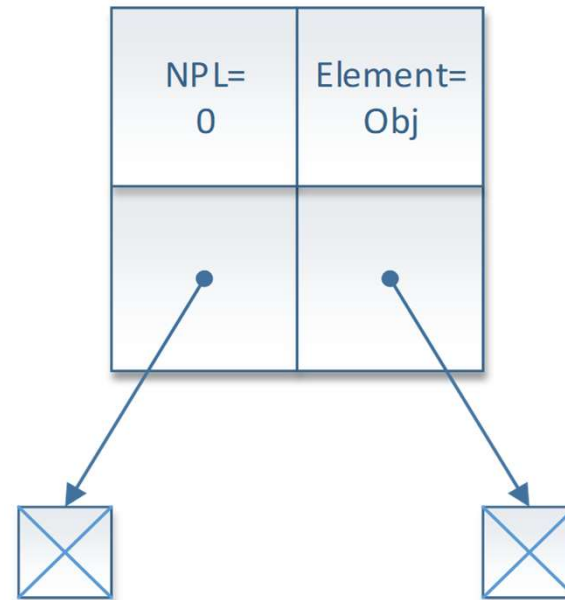
```
template <typename Type>
class Leftist_node {
private:
    Type element;
    Leftist_node *left_tree;
    Leftist_node *right_tree;
    int heap_null_path_length;
public:
    Leftist_node(Type const &);
    Type retrieve() const;
    bool empty() const;
    Leftist_node *left() const;
    Leftist_node *right() const;
    int count(Type const &) const;
    int null_path_length() const;
    void push(Leftist_node *, Leftist_node *&);
    void clear(); };
```





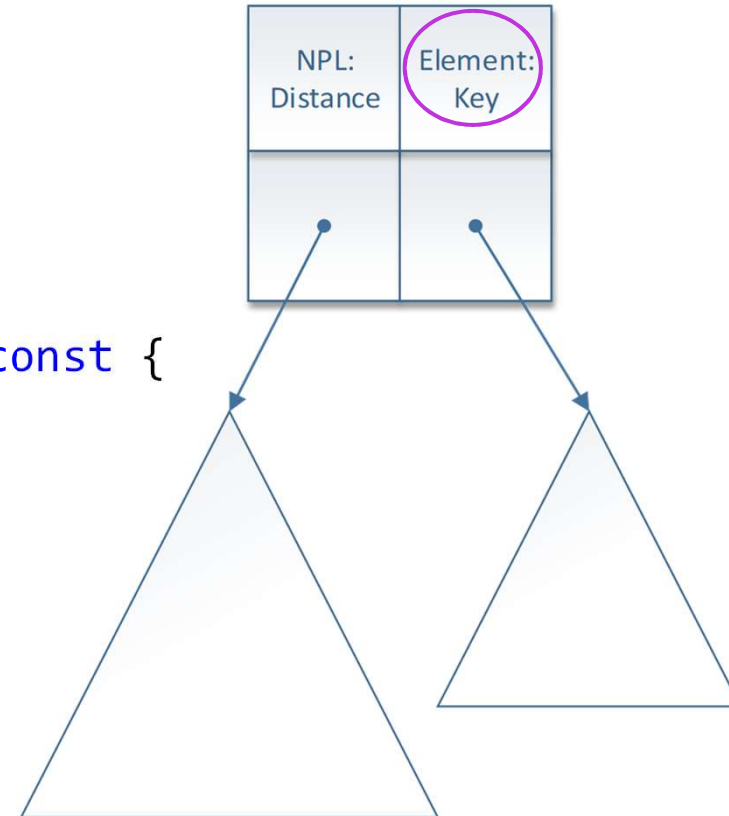
## Leftist\_node Constructor

```
template <typename Type>
Leftist_node<Type>::Leftist_node(Type const
&obj) :
element(obj),
left_tree(nullptr),
right_tree(nullptr),
heap_null_path_length(0) {
// does nothing
}
```



# Leftist\_node::retrieve()

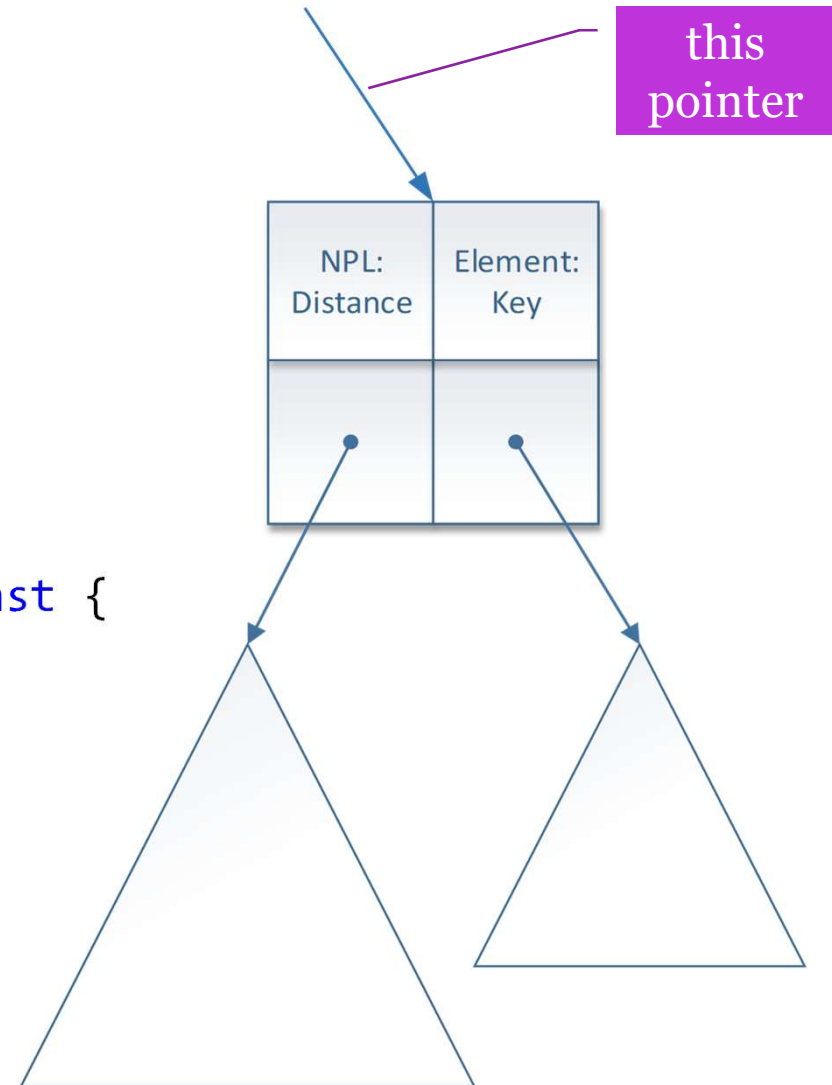
```
template <typename Type>
Type Leftist_node<Type>::retrieve() const {
    //Return the element.
}
```



# Leftist\_node::empty()

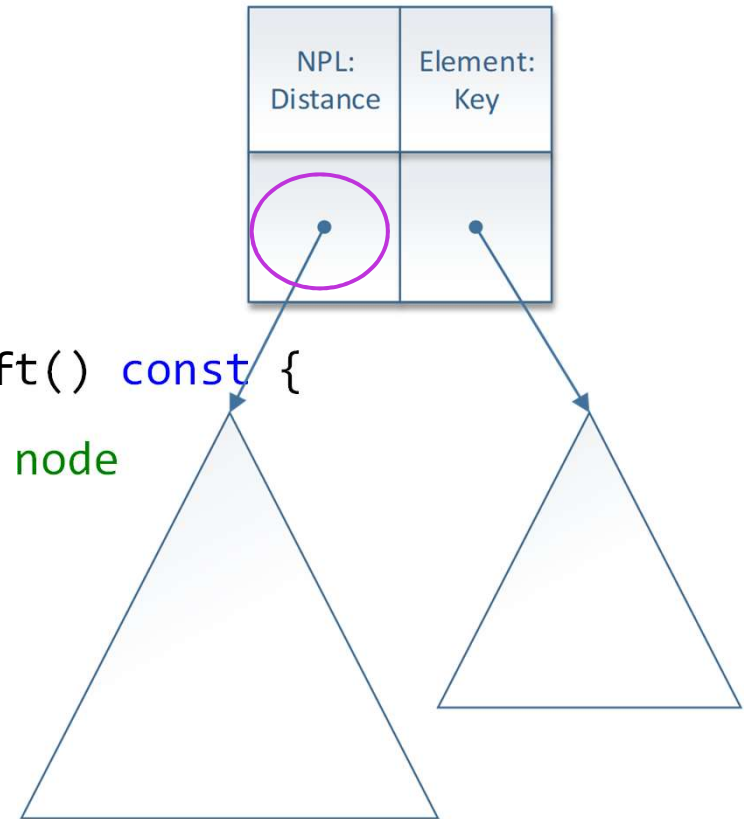
```
template <typename Type>
bool Leftist_node<Type>::empty() const {
    //Check if this equals to nullptr.
}

```



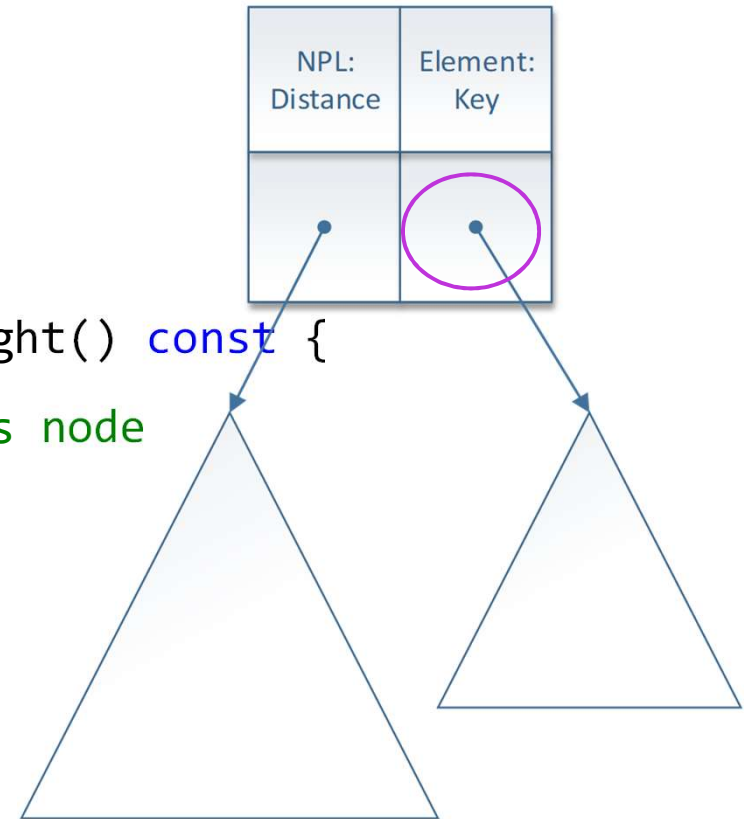
# Leftist\_node::left()

```
template <class Type>
Leftist_node<Type> *Leftist_node<Type>::left() const {
// Return the address of left tree of this node
}
```



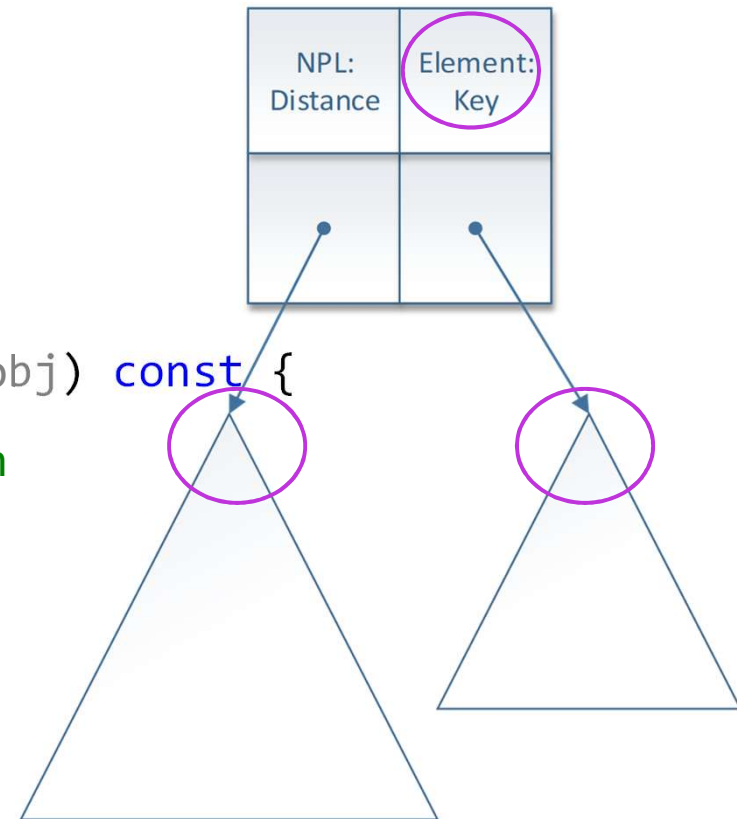
# Leftist\_node::right()

```
template <class Type>
Leftist_node<Type> *Leftist_node<Type>::right() const {
// Return the address of right tree of this node
}
```



# Leftist\_node::count()

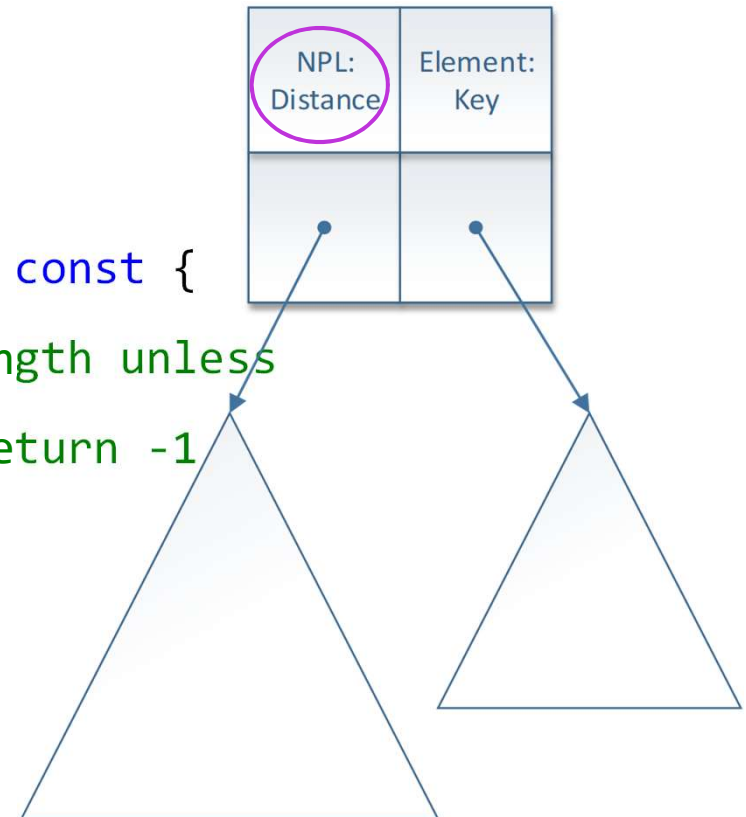
```
template <typename Type>
int Leftist_node<Type>::count(Type const &obj) const {
    // Return the number of instances of obj in
    // this sub-tree.
    // You can do it recursively
}
```





# Leftist\_node:: null\_path\_length()

```
template <typename Type>
int Leftist_node<Type>::null_path_length() const {
    // Return the member variable null-path length unless
    // this is the null pointer, in which case, return -1
}
```



## Leftist\_node:: push

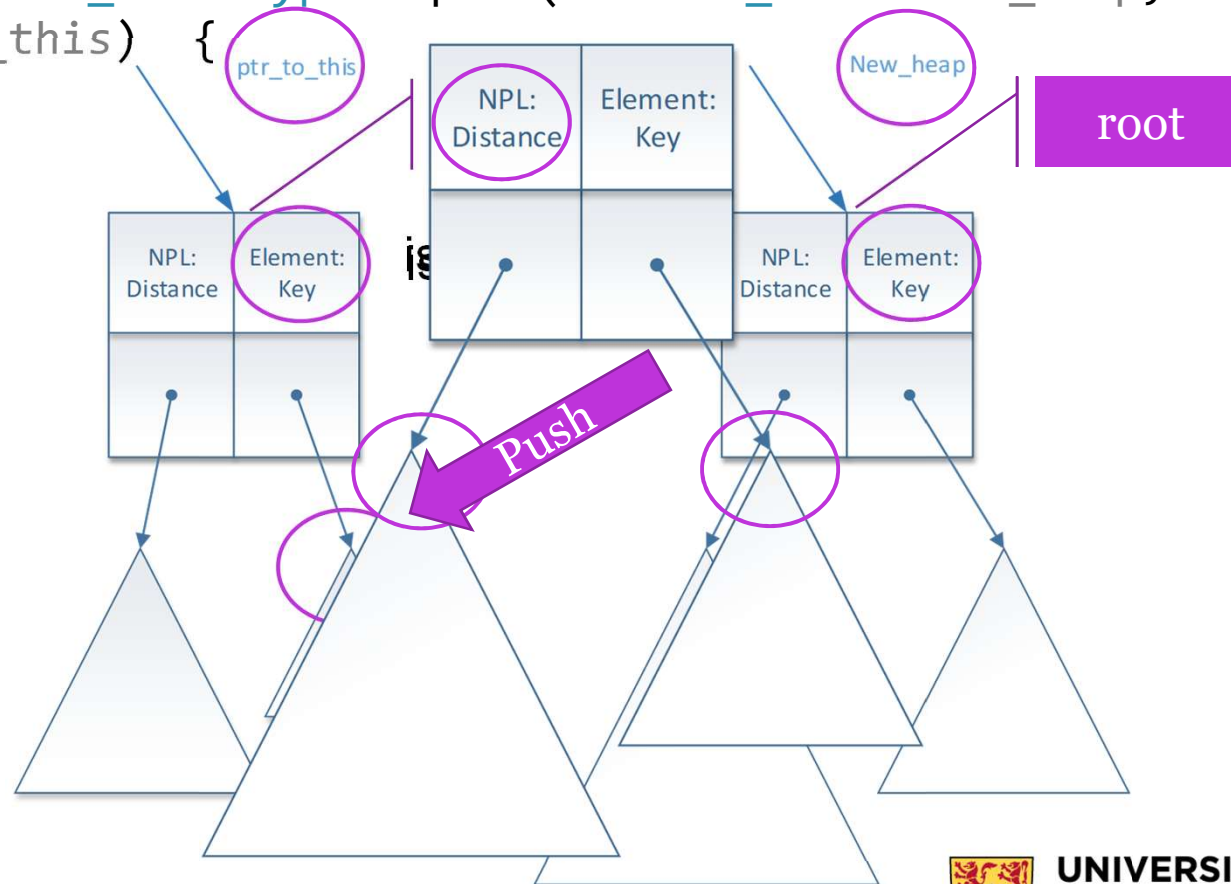
```
template <typename Type>
void Leftist_node<Type>::push(Leftist_node *new_heap, Leftist_node *&ptr_to_this) {
    //If the new_heap is null return
    //if this is null, set the pointer to this to be the new heap and return
    // If the element (value) of current node > new_heap's element, set the pointer to
    this to be the new heap and push this node to the new heap
    //If the element of this node  $\leq$  new_heap's element, push the node into the right
    subtree.
    // Update the null_path length
    // if the left sub-tree has a smaller null_path_length than the right sub-tree, swap
    the two sub-trees
}
```



# Leftist\_node:: push

```
template <typename Type>
```

```
void Leftist_node<Type>::push(Leftist_node *new_heap, Leftist_node
*&ptr_to_this) {
```



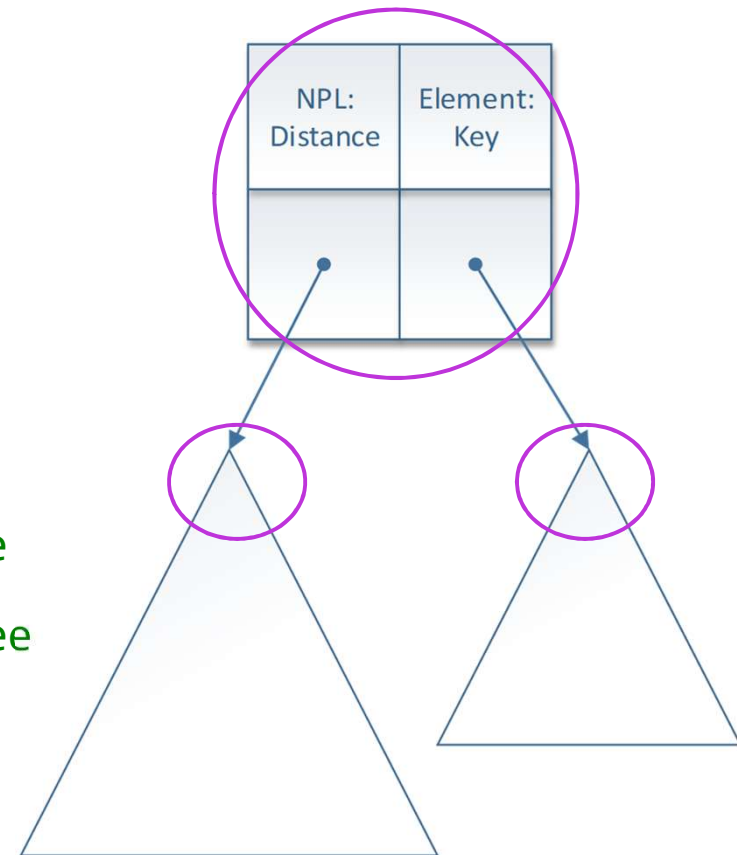
## Why \*&ptr\_to\_this?

- Why don't we use "this pointer"?
  - "this pointer" is a prvalue
  - Similar to e.g., 55 : 55=x ?!!
  - You can not reassign "this pointer"
- So we need to pass the address of current node:
  - Similar to passing by reference or passing by value in C++:
    - Passing by **reference**: the caller and the callee use the same variable
    - Passing by value: the caller and the callee have two different variables



# Leftist\_node::clear()

```
template<typename Type>
void Leftist_node<Type>::clear() {
    //If new heap is null return
    // Call clear function on the left sub-tree
    // Call clear function on the right sub-tree
    // Delete this node
}
```



# Leftist\_heap

```

class Leftist_heap {
private:
    Leftist_node<Type> *root_node;
    int heap_size;
public:
    Leftist_heap();
    ~Leftist_heap();
    void swap(Leftist_heap &heap);
    bool empty() const;
    int size() const;
    int null_path_length() const;
    Type top() const;
    int count(Type const &) const;
    void push(Type const &);
    Type pop();
    void clear();

// Friends
template <typename T>
friend std::ostream &operator<<(std::ostream &, Leftist_heap<T> const &);
};

```





# Leftist\_heap Constructor

```
template <typename Type>
Leftist_heap<Type>::Leftist_heap() :
root_node(nullptr),
heap_size(0) {
// does nothing
}
```



## Leftist\_heap destructor

```
template <typename Type>
Leftist_heap<Type>::~~Leftist_heap() {
clear(); // might as well use it...
}
```



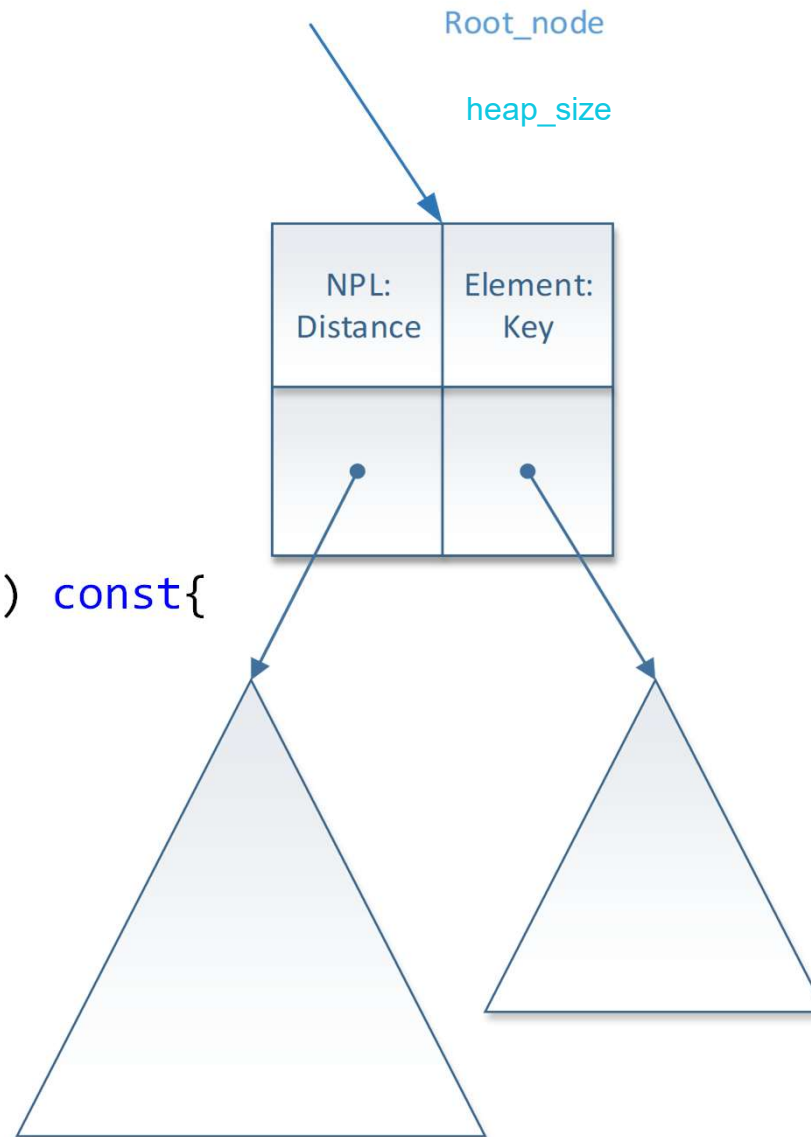
## Leftist\_heap::swap()

```
template <typename Type>
void Leftist_heap<Type>::swap(Leftist_heap<Type> &heap) {
    std::swap(root_node, heap.root_node);
    std::swap(heap_size, heap.heap_size);
}
```



# Leftist\_heap::empty()

```
template <typename Type>
bool Leftist_heap<Type>::empty() const{
    // Check if the heap is empty
}
```



## Leftist\_heap::size()

```
template<typename Type>
int Leftist_heap<Type>::size() const{
// Return the number of nodes in the heap
}
```



## Leftist\_heap::null\_path\_length()

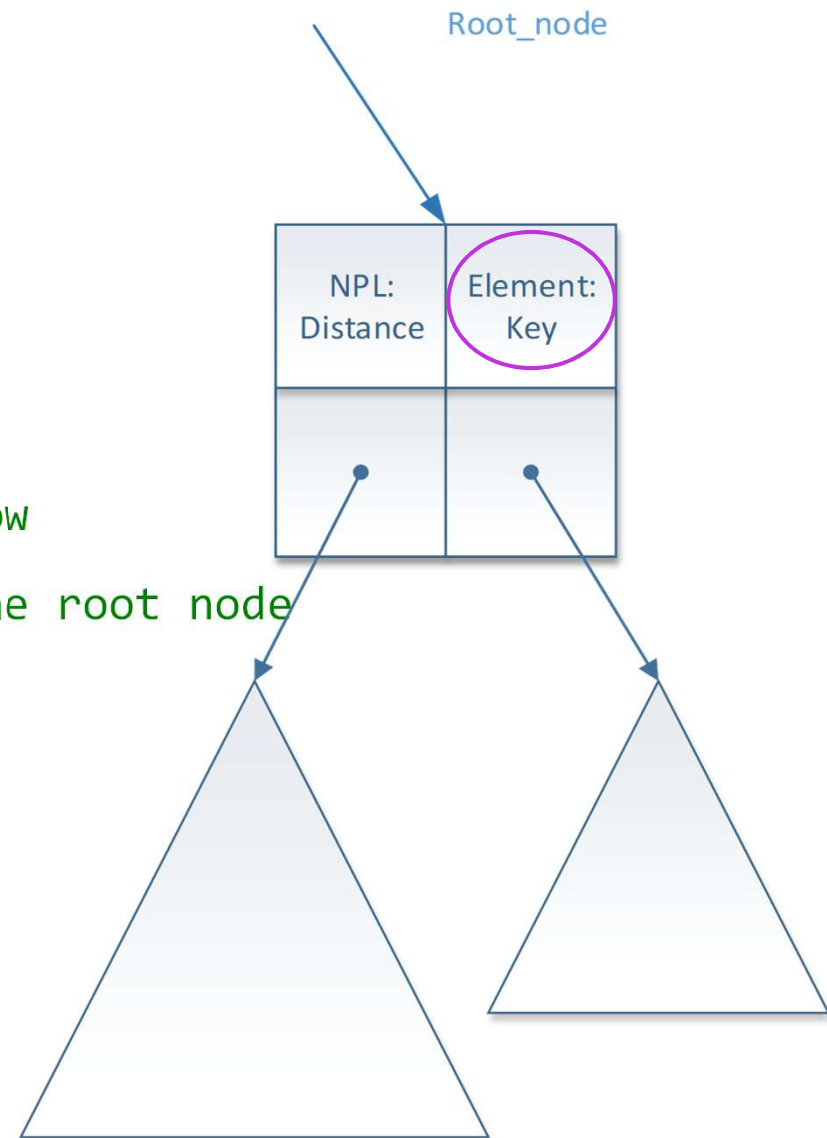
```
template<typename Type>
int Leftist_heap<Type>::null_path_length() const {
// Return the heap_null_path_length of the root node
}
```





# Leftist\_heap::top()

```
template<typename Type>
Type Leftist_heap<Type>::top() const{
// If the heap is empty throw underflow
// Otherwise, return the element of the root node
}
```



## Leftist\_heap::count()

```
template<typename Type>
int Leftist_heap<Type>::count(Type const &obj) const {
// Return the number of instances of obj in the heap
}
```



## Leftist\_heap::clear()

```
template<typename Type>
void Leftist_heap<Type>::clear() {
    // Call clear on the root node
    // Reset the root node
    // Reset the heap size
}
```



# Leftist\_heap::push()

```
template<typename Type>
```

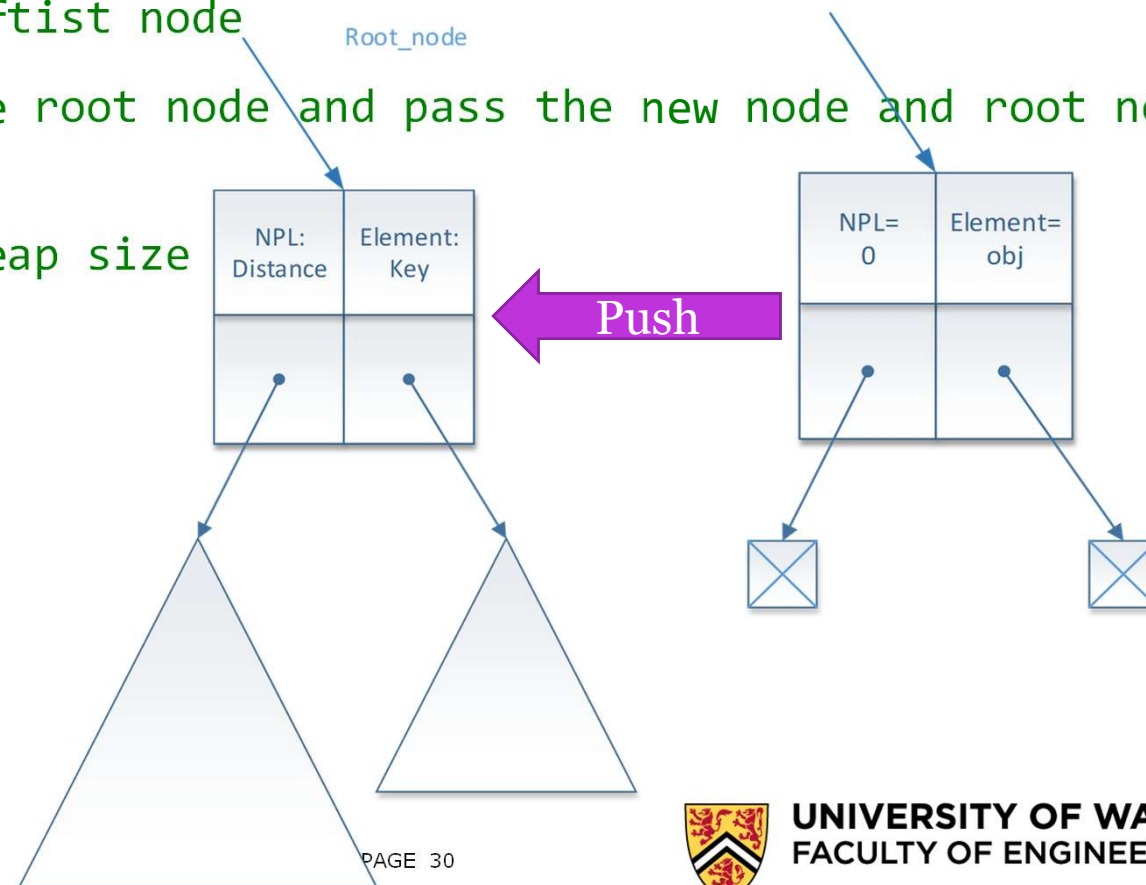
```
void Leftist_heap<Type>::push(Type const &obj) {
```

```
// Create a new leftist node
```

```
// Call push on the root node and pass the new node and root node as  
the arguments
```

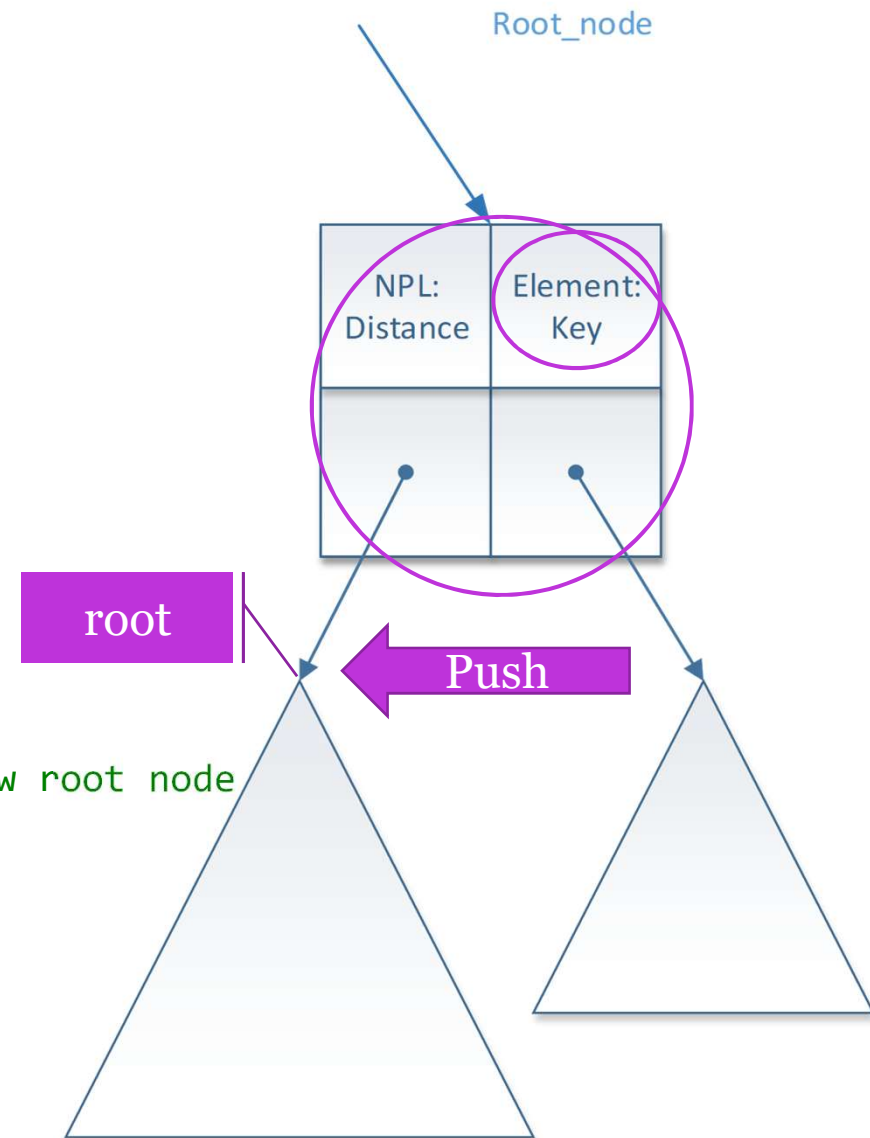
```
//Increment the heap size
```

```
}
```



# Leftist\_heap::pop()

```
template<typename Type>
Type Leftist_heap<Type>::pop() {
    // If the heap is empty throw underflow
    // Pop the last element and delete its node
    // The left sub-tree becomes the root node
    // The right sub-tree is pushed into the new root node
    // Decrement the heap size
    // Return the element of the popped node
}
```



# Leftist Heap

- Visualization:
  - <https://www.cs.usfca.edu/~galles/visualization/LeftistHeap.html>
- References
  - [https://ece.uwaterloo.ca/~dwharder/aads/Algorithms/Leftist\\_heaps/](https://ece.uwaterloo.ca/~dwharder/aads/Algorithms/Leftist_heaps/)
  - <http://www.cs.cmu.edu/~clo/www/CMU/DataStructures/>
  - <http://www.dgp.toronto.edu/people/JamesStewart/378notes/>
  - [https://en.wikipedia.org/wiki/Leftist\\_tree](https://en.wikipedia.org/wiki/Leftist_tree)
  - Weiss, Mark A. *Data structures & algorithm analysis in C++*. Pearson Education, 2012.

