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Jason Seminara
js4308@nyu.edu
N12517906
Homework #3
2015-04-01
Problem #1
== Conversion into FOL
A. \forall s: \neg G(s, Blues) \lor \neg G(s, Opera)
B. \forall s \ W(Amy,s) \Rightarrow S(Amy,s)
C. \exists s \ W(Amy,s) \ \land \ G(s,Blues)
D. \forall s [\exists p [F(p,Joe) \land S(p,s)] \Rightarrow L(Joe,s)]
E. F(Amy, Joe)
F. \exists s: L(Joe,s) \land \neg G(s,0pera)
G. \forall s [ L(Joe,s) \Rightarrow \exists p [ F(p,Joe) \land S(s,p) ] ]
H. \forall s [\exists p [F(p,Joe) \land S(p,s)] \Rightarrow W(Ted,s)]
I. \forall s [ W(Ted,s) \Rightarrow \neg L(Joe,s) ]
==Problem #2 (prove F is a consequence of A-E )==
=== CNF
A. \forall s: \neg G(s, Blues)
                              v \neg G(s, 0pera)
B. ∀s: ¬W(Amy,s)
                              v S(Amy,s)
C. ∃s: W(Amy,s)
                              Λ G(s,Blues)
                             v \neg S(p,s) v L(Joe,s)
D. ∀p,s: ¬F(p,Joe)
E. F(Amy, Joe)
F. ∃s: L(Joe,s) ∧ G(s,Opera)
=== remove Quantifiers & break up top-level conjunctions
A. \neg G(s, Blues)
                        v ¬G(s,Opera)
B. \neg W(Amy,s)
                         v S(Amy,s)
C1. W(Amy, SK1)
C2. G(SK1,Blues)
D. \neg F(p, Joe)
                   v \neg S(p,s)
                                 v L(Joe,s)
E. F(Amy, Joe)
=== negate
F' = \neg L(Joe,s) \lor G(s,Opera)
=== Resolution
                                        (from E + D)
G. \neg S(Amy,s)
                        v L(Joe,s)
H. \neg G(SK1, Opera)
                                        (from A + C2)
                                        (from B + C1)
I. S(Amy, SK1)
J. L(Joe, SK1)
                                        (from G + I)
                                        (from H + K)
K. G(SK1,Opera)
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==Problem #3 (prove I is a consequence of G+H )==
G. \forall s [ L(Joe,s) \Rightarrow \exists p [ F(p,Joe) \land S(s,p) ] ]
    \forall s[ \sim L(Joe,s) \lor \exists p[ F(p,Joe) \land S(s,p) ] ]
    \forall s [ \sim L(Joe, s) \lor [F(SK0(s), Joe) \land S(s, SK0(s))] ]
    [ \simL(Joe,s) v [F(SK0(s),Joe) \wedge S(s,SK0(s))] ]
    (\sim L(Joe,s) \lor F(SKO(s),Joe)) \land (\sim L(Joe,s) \lor S(s,SKO(s)))
H. \forall s[ \exists p[ F(p,Joe) \land S(p,s) ] \Rightarrow \sim W(Ted,s) ]
    \forall s[ \sim \exists p[ F(p,Joe) \land S(p,s) ] \lor \sim W(Ted,s) ]
    \forall s,p \ [ \sim F(p,Joe) \ v \sim S(p,s) \ v \sim W(Ted,s) ]
I. ~∀s[ W(Ted,s) => ¬L(Joe,s) ]
     \exists s \sim (W(Ted,s) \Rightarrow \neg L(Joe,s))
     ∃s W(Ted,SK2) ∧ L(Joe,SK2)
     W(Ted,SK2) \wedge L(Joe,SK2)
===== CNF
G1. \simL(Joe,s) v F(SK0(s),Joe)
G2. \simL(Joe,s) v S(SK0(s),s)
H. \sim F(p, Joe) \vee \sim S(p, s) \vee \sim W(Ted, s)
I1. W(Ted,SK2)
I2. L(Joe,SK2)
=== Resolution
     F(SK0(SK2), Joe)
                               (from G1 + I2)
J.
K. S(SK0(SK2),SK2)
                               (from G2 + I2)
    ~F(p,Joe) v ~S(p,SK2) (from H + I1)
Μ.
    \simS(SK0(SK2),SK2) (from J + L)
N. \varnothing (from K + M)
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