

# 1 Representations

## 1.1 Source Language

statement ::= compound-statement  
| if ( expression ) compound-statement else compound-statement  
| while ( expression ) compound-statement  
| for ( statement ; statement ; statement ) compound-statement  
| qualifier type name ;  
| qualifier type name = expression ;  
| return expression ;  
| expression ;

compound-statement ::= { statement\* }

primary-expression ::= identifier  
| constant  
| string-literal  
| ( expression )

postfix-expression ::= primary-expression  
| postfix-expression [ expression ]  
| name ( argument-list )  
| postfix-expression ++  
| postfix-expression --

unary-expression ::= postfix-expression  
| ++ unary-expression  
| -- unary-expression  
| unary-operator postfix-expression

multiplicative-expression ::= unary-expression  
| multiplicative-expression \* unary-expression  
| multiplicative-expression / unary-expression  
| multiplicative-expression % unary-expression

additive-expression ::= multiplicative-expression  
| additive-expression + multiplicative-expression  
| additive-expression - multiplicative-expression

shift-expression ::= additive-expression  
| shift-expression << additive-expression  
| shift-expression >> additive-expression

relational-expression ::= shift-expression  
| relational-expression < shift-expression  
| relational-expression > shift-expression

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| relational-expression <= shift-expression
| relational-expression >= shift-expression

equality-expression ::= relational-expression
| equality-expression == relational-expression
| equality-expression != relational-expression

bitwise-and-expression ::= equality-expression
| and-expression & equality-expression

exclusive-or-expression ::= and-expression
| exclusive-or-expression | exclusive-or-expression

bitwise-or-expression ::= exclusive-or-expression
| inclusive-or-expression | exclusive-or-expression

logical-and-expression ::= bitwise-or-expression
| logical-and-expression && bitwise-or-expression

logical-or-expression ::= logical-and-expression
| logical-or-expression || logical-and-expression

conditional-expression ::= logical-or-expression
| conditional-expression ? expression : conditional-expression

assignment-expression ::= conditional-expression
| unary-expression assignment-operator assignment-expression

expression ::= assignment-expression

parameter-list ::= qualifier type name , argument-list
| type-qualifier type name

argument-list ::= expression , argument-list
| expression

assignment-operator ::= = | *= | /= | %= | += | -=

unary-operator ::= * | - | ! | ~

type-qualifier ::= const | volatile

type-specifier ::= int | unsigned | char

type ::= type-qualifier type-specifier

function-definition ::= type name ( parameter-list ) compound-expression

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## 2 Symbolic Execution

### 2.1 Metavariables

$e$  - source language expressions  
 $c$  - source language statements  
 $s$  - symbolic expressions  
 $x, \alpha$  - names  
 $v$  - symbolic values

### 2.2 Symbolic Expressions

$s ::= v \mid x$   
 $\mid s \text{ binop } s$   
 $\mid \text{unop } s$   
 $\mid \text{sel}(s, (s, \dots, s))$   
 $\mid \text{upd}(s, (s, \dots, s), s)$   
 $\mid x(s, \dots, s)$   
 $\mid s? : s : s$   
 $\mid \text{NewArr } \tau$

$\tau ::= \text{Arr } \tau$   
 $\mid \text{int8}$   
 $\mid \text{int32}$   
 $\mid \text{word32}$

### 2.3 Symbolic State

$S = \langle g; \rho; \mu \rangle$   
 $g$  - path condition  
 $\rho$  - mapping of names to stack memory  
 $\mu$  - mapping of names to heap memory

### 2.4 Expressions

$$\frac{}{\langle S; v \rangle \Downarrow \langle S; v \rangle} \text{LITERAL}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S'; s \rangle}{\langle S; -e \rangle \Downarrow \langle S'; -s \rangle} \text{NEGATE}$$

$$\frac{\langle S; e_1 \rangle \Downarrow \langle S_1; s_1 \rangle \quad \langle S_1; e_2 \rangle \Downarrow \langle S_2; s_2 \rangle}{\langle S; e_1 + e_2 \rangle \Downarrow \langle S_2; s_1 + s_2 \rangle} \text{ADD}$$

$$\frac{\forall i \in 1..n, \langle S_i; e_i \rangle \Downarrow \langle S_{i+1}; s_i \rangle}{\langle S_1; x(e_1, \dots, e_n) \rangle \Downarrow \langle S_{n+1}; x(s_1, \dots, s_n) \rangle} \text{FUNCALL}$$

## 2.5 Statements

$$\frac{\langle S; e \rangle \Downarrow \langle S'; s \rangle}{\langle S; e; \rangle \Downarrow \langle S'; 0 \rangle} \text{EXPRESSION}$$

$$\frac{\forall i \in \{1..n\}, \langle S_k; c_i \rangle \Downarrow \langle S_{i+1}; s_{i+1} \rangle}{\langle S_1; \{c_1..c_n\} \rangle \Downarrow \langle S_{n+1}[\rho \mapsto \rho(S_1)]; s_{n+1} \rangle} \text{COMPOUNDSTATEMENT}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S_1; g_1 \rangle \quad \langle S_1[g \mapsto g(S_1) \wedge g_1]; c_1 \rangle \Downarrow \langle S_2; s_2 \rangle}{\langle S; \text{if } e \text{ } c_1 \text{ else } c_2 \rangle \Downarrow \langle S_2; 0 \rangle} \text{IFTRUE}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S_1; g_1 \rangle \quad \langle S_1[g \mapsto g(S_1) \wedge \neg g_1]; c_2 \rangle \Downarrow \langle S_2; s_2 \rangle}{\langle S; \text{if } e \text{ } c_1 \text{ else } c_2 \rangle \Downarrow \langle S_2; 0 \rangle} \text{IFFALSE}$$

## 2.6 Memory

$$\frac{x \in \text{dom } \rho(S) \quad \rho(S)[x] = s}{\langle S; x \rangle \Downarrow \langle S; s \rangle} \text{VAR}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S'; s \rangle \quad \text{toOffset}(s) = (\text{ptr } \alpha + s_2) \quad \rho(S')[\alpha] = s_1}{\langle S; *e \rangle \Downarrow \langle S'; \text{sel}(s_1, s_2) \rangle} \text{DEREFSTACK}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S'; s \rangle \quad \text{toOffset}(s) = (\text{ptr } \alpha + s_2) \quad \mu(S')[\alpha] = s_1}{\langle S; *e \rangle \Downarrow \langle S'; \text{sel}(s_1, s_2) \rangle} \text{DEREFHEAP}$$

$$\frac{\rho(S_1)[x] = \text{ptr } \alpha \quad \rho(S_1)[\alpha] = s \quad \forall i \in \{1..n\}, \langle S_i; e_i \rangle \Downarrow \langle S_{i+1}; s_i \rangle}{\langle S_1; x[e_1] \dots [e_n] \rangle \Downarrow \langle S_{i+1}; \text{Sel}(s, (s_1, \dots, s_n)) \rangle} \text{SELLOCAL}$$

$$\frac{\rho(S_1)[x] = \text{ptr } \alpha \quad \mu(S_1)[\alpha] = s \quad \forall i \in \{1..n\}, \langle S_i; e_i \rangle \Downarrow \langle S_{i+1}; s_i \rangle}{\langle S_1; x[e_1] \dots [e_n] \rangle \Downarrow \langle S_{i+1}; \text{Sel}(s, (s_1, \dots, s_n)) \rangle} \text{SELGLOBAL}$$

$$\frac{x \notin \text{dom } (\rho(S)) \quad \dim(\tau) = 0}{\langle S; \tau \ x; \rangle \Downarrow \langle S[\rho \mapsto \rho[x \mapsto 0]]; 0 \rangle} \text{DECLARELOCAL1}$$

$$\frac{x \notin \text{dom } (\rho(S)) \quad \dim(\tau) > 0 \quad \alpha \text{ is fresh}}{\langle S; \tau \ x; \rangle \Downarrow \langle S[\rho \mapsto \rho[x \mapsto \text{ptr } \alpha, \alpha \mapsto \text{NewArr } \tau]]; \text{ptr } \alpha \rangle} \text{DECLARELOCAL2}$$

$$\frac{x \in \text{dom } (\rho(S)) \quad \langle S; e \rangle \Downarrow \langle S'; s \rangle}{\langle S; x = e; \rangle \Downarrow \langle S'[\rho \mapsto \rho[x \mapsto s]]; s \rangle} \text{ASSIGNLOCAL}$$

$$\frac{\rho(S_1)[x] = \text{ptr } \alpha \quad \rho(S_1)[\alpha] = s \quad \forall i \in \{1..n\}. \langle S_i; e_i \rangle \Downarrow \langle S_{i+1}; s_i \rangle \quad \langle S_{n+1}; e_v \rangle \Downarrow \langle S'; s' \rangle}{\langle S_1; x[e_1] \dots [e_n] = e_v \rangle \Downarrow \langle S'[\rho \mapsto \rho(S')[\alpha \mapsto \text{upd}(s, (s_1, \dots, s_n), s')]]; s' \rangle} \text{UPDLOCAL}$$

$$\frac{\rho(S_1)[x] = \text{ptr } \alpha \quad \mu(S_1)[\alpha] = s \quad \forall i \in \{1..n\}. \langle S_i; e_i \rangle \Downarrow \langle S_{i+1}; s_i \rangle \quad \langle S_{n+1}; e_v \rangle \Downarrow \langle S'; s' \rangle}{\langle S_1; x[e_1] \dots [e_n] = e_v \rangle \Downarrow \langle S'[\mu \mapsto \mu(S')[\alpha \mapsto \text{upd}(s, (s_1, \dots, s_n), s')]]; s' \rangle} \text{UPDGLOBAL}$$

$$\frac{\rho(S)[x] = s}{\langle S; ++x \rangle \Downarrow \langle S[\rho \mapsto \rho(S)[x \mapsto s + 1]]; s + 1 \rangle} \text{INCPRELOCAL}$$

$$\frac{\rho(S)[x] = s}{\langle S; x++ \rangle \Downarrow \langle S[\rho \mapsto \rho(S)[x \mapsto s + 1]]; s \rangle} \text{INCPOSTLOCAL}$$