

# 1 Source Language

$\langle \text{statement} \rangle ::= \langle \text{compound-statement} \rangle$   
| **if** (  $\langle \text{expression} \rangle$  )  $\langle \text{compound-statement} \rangle$  **else**  $\langle \text{compound-statement} \rangle$   
| **while** (  $\langle \text{expression} \rangle$  )  $\langle \text{compound-statement} \rangle$   
| **for** (  $\langle \text{statement} \rangle$  ;  $\langle \text{statement} \rangle$  ;  $\langle \text{statement} \rangle$  )  $\langle \text{compound-statement} \rangle$   
|  $\langle \text{qualifier} \rangle$   $\langle \text{type} \rangle$   $\langle \text{name} \rangle$  ;  
|  $\langle \text{qualifier} \rangle$   $\langle \text{type} \rangle$   $\langle \text{name} \rangle = \langle \text{expression} \rangle$  ;  
| **return**  $\langle \text{expression} \rangle$  ;  
|  $\langle \text{expression} \rangle$  ;

$\langle \text{compound-statement} \rangle ::= \{ \langle \text{statement} \rangle^* \}$

$\langle \text{primary-expression} \rangle ::= \langle \text{identifier} \rangle$   
|  $\langle \text{constant} \rangle$   
|  $\langle \text{string-literal} \rangle$   
| (  $\langle \text{expression} \rangle$  )

$\langle \text{postfix-expression} \rangle ::= \langle \text{primary-expression} \rangle$   
|  $\langle \text{postfix-expression} \rangle$  [  $\langle \text{expression} \rangle$  ]  
|  $\langle \text{name} \rangle$  (  $\langle \text{argument-list} \rangle$  )  
|  $\langle \text{name} \rangle ++$   
|  $\langle \text{name} \rangle --$

$\langle \text{unary-expression} \rangle ::= \langle \text{postfix-expression} \rangle$   
|  $++ \langle \text{name} \rangle$   
|  $-- \langle \text{name} \rangle$   
|  $\& \langle \text{name} \rangle$   
|  $\langle \text{unary-operator} \rangle \langle \text{cast-expression} \rangle$

$\langle \text{multiplicative-expression} \rangle ::= \langle \text{unary-expression} \rangle$   
|  $\langle \text{multiplicative-expression} \rangle * \langle \text{unary-expression} \rangle$   
|  $\langle \text{multiplicative-expression} \rangle / \langle \text{unary-expression} \rangle$   
|  $\langle \text{multiplicative-expression} \rangle \% \langle \text{unary-expression} \rangle$

$\langle \text{additive-expression} \rangle ::= \langle \text{multiplicative-expression} \rangle$   
|  $\langle \text{additive-expression} \rangle + \langle \text{multiplicative-expression} \rangle$   
|  $\langle \text{additive-expression} \rangle - \langle \text{multiplicative-expression} \rangle$

$\langle \text{shift-expression} \rangle ::= \langle \text{additive-expression} \rangle$   
|  $\langle \text{shift-expression} \rangle \ll \langle \text{additive-expression} \rangle$   
|  $\langle \text{shift-expression} \rangle \gg \langle \text{additive-expression} \rangle$

$\langle \text{relational-expression} \rangle ::= \langle \text{shift-expression} \rangle$   
|  $\langle \text{relational-expression} \rangle < \langle \text{shift-expression} \rangle$   
|  $\langle \text{relational-expression} \rangle > \langle \text{shift-expression} \rangle$   
|  $\langle \text{relational-expression} \rangle \leq \langle \text{shift-expression} \rangle$   
|  $\langle \text{relational-expression} \rangle \geq \langle \text{shift-expression} \rangle$

$\langle \text{equality-expression} \rangle ::= \langle \text{relational-expression} \rangle$   
 $\quad | \quad \langle \text{equality-expression} \rangle == \langle \text{relational-expression} \rangle$   
 $\quad | \quad \langle \text{equality-expression} \rangle != \langle \text{relational-expression} \rangle$

$\langle \text{bitwise-and-expression} \rangle ::= \langle \text{equality-expression} \rangle$   
 $\quad | \quad \langle \text{and-expression} \rangle \& \langle \text{equality-expression} \rangle$

$\langle \text{exclusive-or-expression} \rangle ::= \langle \text{and-expression} \rangle$   
 $\quad | \quad \langle \text{exclusive-or-expression} \rangle | \langle \text{exclusive-or-expression} \rangle$

$\langle \text{bitwise-or-expression} \rangle ::= \langle \text{exclusive-or-expression} \rangle$   
 $\quad | \quad \langle \text{inclusive-or-expression} \rangle | \langle \text{exclusive-or-expression} \rangle$

$\langle \text{logical-and-expression} \rangle ::= \langle \text{bitwise-or-expression} \rangle$   
 $\quad | \quad \langle \text{logical-and-expression} \rangle \&\& \langle \text{bitwise-or-expression} \rangle$

$\langle \text{logical-or-expression} \rangle ::= \langle \text{logical-and-expression} \rangle$   
 $\quad | \quad \langle \text{logical-or-expression} \rangle || \langle \text{logical-and-expression} \rangle$

$\langle \text{conditional-expression} \rangle ::= \langle \text{logical-or-expression} \rangle$   
 $\quad | \quad \langle \text{conditional-expression} \rangle ? \langle \text{expression} \rangle : \langle \text{conditional-expression} \rangle$

$\langle \text{assignment-expression} \rangle ::= \langle \text{conditional-expression} \rangle$   
 $\quad | \quad \langle \text{unary-expression} \rangle \langle \text{assignment-operator} \rangle \langle \text{assignment-expression} \rangle$

$\langle \text{expression} \rangle ::= \langle \text{assignment-expression} \rangle$

$\langle \text{parameter-list} \rangle ::= \langle \text{qualifier} \rangle \langle \text{type} \rangle \langle \text{name} \rangle , \langle \text{argument-list} \rangle$   
 $\quad | \quad \langle \text{type-qualifier} \rangle \langle \text{type} \rangle \langle \text{name} \rangle$

$\langle \text{argument-list} \rangle ::= \langle \text{expression} \rangle , \langle \text{argument-list} \rangle$   
 $\quad | \quad \langle \text{expression} \rangle$

$\langle \text{assignment-operator} \rangle ::= = | * = | / = | \% = | + = | - =$

$\langle \text{unary-operator} \rangle ::= * | - | ! | \sim$

$\langle \text{type-qualifier} \rangle ::= \text{const} | \text{volatile}$

$\langle \text{type-specifier} \rangle ::= \text{int} | \text{char} | \text{float}$

$\langle \text{type} \rangle ::= \langle \text{type-qualifier} \rangle \langle \text{type-specifier} \rangle$

$\langle \text{function-definition} \rangle ::= \langle \text{type} \rangle \langle \text{name} \rangle ( \langle \text{parameter-list} \rangle ) \langle \text{compound-expression} \rangle$

## 2 Symbolic Execution

$S = \langle g; \rho; \mu \rangle$

## 2.1 Expressions

$$\frac{}{\langle S; v \rangle \Downarrow \langle S; v \rangle} \text{ LITERAL}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S'; s \rangle}{\langle S; -e \rangle \Downarrow \langle S'; -s \rangle} \text{ NEGATE}$$

$$\frac{\langle S; e_1 \rangle \Downarrow \langle S_1; s_1 \rangle \quad \langle S_1; e_2 \rangle \Downarrow \langle S_2; s_2 \rangle}{\langle S; e_1 + e_2 \rangle \Downarrow \langle S_2; e_1 + e_2 \rangle} \text{ ADD}$$

$$\frac{\langle S; e_1 = e_1 + e_2 \rangle \Downarrow \langle S'; s \rangle}{\langle S; e_1 += e_2 \rangle \Downarrow \langle S'; s \rangle} \text{ ASSIGNADD}$$

$$\frac{\forall i \in 1..n, \langle S_i; e_i \rangle \Downarrow \langle S_{i+1}; s_i \rangle}{\langle S_1; \mathbf{x}(e_1, \dots, e_n) \rangle \Downarrow \langle S_{n+1}; x(s_1, \dots, s_n) \rangle} \text{ FUNCALL}$$

## 2.2 Statements

$$\frac{\langle S; e \rangle \Downarrow \langle S'; s \rangle}{\langle S; e; \rangle \Downarrow \langle S'; \emptyset \rangle} \text{ EXPRESSION}$$

$$\frac{\forall i \in 1..n, \langle S_i; c_i \rangle \Downarrow \langle S_{i+1}; s_{i+1} \rangle}{\langle S_1; \{c_1..c_n\} \rangle \Downarrow \langle S_{n+1}; s_{n+1} \rangle} \text{ COMPOUNDSTATEMENT}$$

$$\frac{\begin{array}{l} \langle S; e \rangle \Downarrow \langle S_1; g_1 \rangle \quad g(S) \not\Rightarrow g_1 \quad g(S) \not\Rightarrow \neg g_1 \\ \langle S_1[g \mapsto g(S_1) \wedge g_1]; c_1 \rangle \Downarrow \langle S_2; s_2 \rangle \\ \langle S_1[g \mapsto g(S_1) \wedge \neg g_1]; c_1 \rangle \Downarrow \langle S_3; s_3 \rangle \\ S' = \langle (g_1 ? g(S_2) : g(S_3)); (g_1 ? \rho(S_2) : \rho(S_3)); (g_1 ? \mu(S_2) : \mu(S_3)) \rangle \end{array}}{\langle S; \mathbf{if } e \text{ } c_1 \text{ } \mathbf{else } c_2 \rangle \Downarrow \langle S'; \emptyset \rangle} \text{ IFELSE}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S_1; g_1 \rangle \quad g(S) \Rightarrow g_1 \quad \langle S_1; c_1 \rangle \Downarrow \langle S_2; s \rangle}{\langle S; \mathbf{if } e \text{ } c_1 \text{ } \mathbf{else } c_2 \rangle \Downarrow \langle S_2; \emptyset \rangle} \text{ IFTRUE}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S_1; g_1 \rangle \quad g(S) \Rightarrow \neg g_1 \quad \langle S_1; c_2 \rangle \Downarrow \langle S_2; s \rangle}{\langle S; \mathbf{if } e \text{ } c_1 \text{ } \mathbf{else } c_2 \rangle \Downarrow \langle S_2; \emptyset \rangle} \text{ IFFALSE}$$

## 2.3 Memory

$$\frac{\rho(S)[x] = s}{\langle S; x \rangle \Downarrow \langle S; s \rangle} \text{VAR}$$

$$\frac{x \notin \text{dom } \rho(S)}{\langle S; \tau \ x; \rangle \Downarrow \langle S[\rho \mapsto (\rho(S), (x \rightarrow \emptyset))]; s \rangle} \text{DECLARELOCAL}$$

$$\frac{x \notin \text{dom } \rho(S) \quad \langle S; e \rangle \Downarrow \langle S'; s \rangle}{\langle S; \tau \ x = e; \rangle \Downarrow \langle S'[\rho \mapsto (\rho(S'), (x \rightarrow s))]; s \rangle} \text{DECLAREASSIGNLOCAL}$$

$$\frac{\langle S; e_1 \rangle \Downarrow \langle S_1; \text{ptr } x \rangle \quad x \in \text{dom } \rho(S_1) \quad \langle S_1; e_2 \rangle \Downarrow \langle S_2; s \rangle}{\langle S; *e_1 = e_2 \rangle \Downarrow \langle S_2[\rho \mapsto (\rho(S_2), (x \rightarrow s))]; s \rangle} \text{UPDLOCAL}$$

$$\frac{\langle S; e_1 \rangle \Downarrow \langle S_1; s_1 \rangle \quad s_1 \neq \text{ptr } x \quad \langle S_1; e_2 \rangle \Downarrow \langle S_2; s_2 \rangle}{\langle S; *e_1 = e_2 \rangle \Downarrow \langle S_2[\mu \mapsto (\mu(S_2), (s_1 \rightarrow s_2))]; s_2 \rangle} \text{UPDGLOBAL}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S'; \text{ptr } x \rangle \quad \rho(S')[x] = s}{\langle S; *e \rangle \Downarrow \langle S'; s \rangle} \text{SELLOCAL}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S'; \text{ptr } x \rangle \quad \rho(S')[x] = s}{\langle S; *e \rangle \Downarrow \langle S'; s \rangle} \text{SELGLOBAL}$$

$$\frac{\rho(S)[x] = s}{\langle S; ++x \rangle \Downarrow \langle S[\rho \mapsto (\rho(S), x \rightarrow s + 1)]; s + 1 \rangle} \text{INCPRE}$$

$$\frac{\rho(S)[x] = s}{\langle S; x++ \rangle \Downarrow \langle S[\rho \mapsto (\rho(S), x \rightarrow s + 1)]; s \rangle} \text{INCPST}$$