

# 1 Source Language

$\langle \text{statement} \rangle ::= \langle \text{compound-statement} \rangle$   
|  $\text{'if' } (\langle \text{expression} \rangle) \langle \text{compound-statement} \rangle \text{'else' } \langle \text{compound-statement} \rangle$   
|  $\text{'while' } (\langle \text{expression} \rangle) \langle \text{compound-statement} \rangle$   
|  $\text{'for' } (\langle \text{statement} \rangle \text{';' } \langle \text{statement} \rangle \text{';' } \langle \text{statement} \rangle) \langle \text{compound-statement} \rangle$   
|  $\langle \text{qualifier} \rangle \langle \text{type} \rangle \langle \text{name} \rangle \text{';'}$   
|  $\langle \text{qualifier} \rangle \langle \text{type} \rangle \langle \text{name} \rangle \text{'=' } \langle \text{expression} \rangle \text{';'}$   
|  $\text{'return' } \langle \text{expression} \rangle \text{';'}$   
|  $\langle \text{expression} \rangle \text{';'}$

$\langle \text{compound-statement} \rangle ::= \text{'{' } \langle \text{statement} \rangle^* \text{'}'}$

$\langle \text{expression} \rangle ::= \langle \text{name} \rangle \mid \langle \text{literal} \rangle$   
|  $(\langle \text{expression} \rangle)$   
|  $\langle \text{expression} \rangle \langle \text{binop} \rangle \langle \text{expression} \rangle$   
|  $\langle \text{unop} \rangle \langle \text{expression} \rangle$   
|  $\langle \text{name} \rangle (\langle \text{parameter-list} \rangle)$   
|  $\text{'\&' } \langle \text{name} \rangle$   
|  $\langle \text{expression} \rangle \text{'++'}$   
|  $\langle \text{expression} \rangle \text{'--'}$

$\langle \text{parameter-list} \rangle ::= \langle \text{qualifier} \rangle \langle \text{type} \rangle \langle \text{name} \rangle \text{' ,' } \langle \text{parameter-list} \rangle$   
|  $\langle \text{type-qualifier} \rangle \langle \text{type} \rangle \langle \text{name} \rangle$

$\langle \text{binop} \rangle ::= \text{'+'} \mid \text{'-'} \mid \text{'*'} \mid \text{'/'} \mid \text{'%'}$   
|  $\text{'\&\&'}$  |  $\text{'||'}$   
|  $\text{'\&'}$  |  $\text{'|'}$  |  $\text{'<<'}$  |  $\text{'>>'}$   
|  $\text{'=='}$  |  $\text{'<='}$  |  $\text{'>='}$  |  $\text{'<'}$  |  $\text{'>'}$

$\langle \text{unop} \rangle ::= \text{'-'} \mid \text{'!'}$  |  $\text{'\sim'}$  |  $\text{'*'} \mid \text{'++'}$  |  $\text{'--'}$

$\langle \text{type-qualifier} \rangle ::= \text{'const'}$  |  $\text{'volatile'}$

$\langle \text{type} \rangle ::= \langle \text{type-qualifier} \rangle \langle \text{type-specifier} \rangle$

## 2 Symbolic Execution

### 2.1 Expressions

$$\frac{}{\langle S; v \rangle \Downarrow \langle S; v \rangle} \text{LITERAL}$$

### 2.2 Statements

$$\frac{\langle S; e \rangle \Downarrow \langle S'; s \rangle}{\langle S; e; \rangle \Downarrow \langle S'; \emptyset \rangle} \text{EXPRESSION}$$

$$\frac{\forall i \in 1..n, \langle S_i; c_i \rangle \Downarrow \langle S_{i+1}; s_{i+1} \rangle}{\langle S_1; \{c_1..c_n\} \rangle \Downarrow \langle S_{n+1}; s_{n+1} \rangle} \text{COMPOUNDSTATEMENT}$$

$$\frac{\begin{array}{c} \langle S; e \rangle \Downarrow \langle S_1; g_1 \rangle \quad g(S) \not\Rightarrow g_1 \quad g(S) \not\Rightarrow \neg g_1 \\ \langle S_1[g \mapsto g(S_1) \wedge g_1]; c_1 \rangle \Downarrow \langle S_2; s_2 \rangle \\ \langle S_1[g \mapsto g(S_1) \wedge \neg g_1]; c_1 \rangle \Downarrow \langle S_3; s_3 \rangle \\ S' = \langle (g_1 ? g(S_2) : g(S_3)); (g_1 ? \rho(S_2) : \rho(S_3)); (g_1 ? \mu(S_2) : \mu(S_3)) \rangle \end{array}}{\langle S; \text{if } e \text{ } c_1 \text{ else } c_2 \rangle \Downarrow \langle S'; \emptyset \rangle} \text{IFELSE}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S_1; g_1 \rangle \quad g(S) \Rightarrow g_1 \quad \langle S_1; c_1 \rangle \Downarrow \langle S_2; s \rangle}{\langle S; \text{if } e \text{ } c_1 \text{ else } c_2 \rangle \Downarrow \langle S_2; \emptyset \rangle} \text{IFTRUE}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S_1; g_1 \rangle \quad g(S) \Rightarrow \neg g_1 \quad \langle S_1; c_2 \rangle \Downarrow \langle S_2; s \rangle}{\langle S; \text{if } e \text{ } c_1 \text{ else } c_2 \rangle \Downarrow \langle S_2; \emptyset \rangle} \text{IFFALSE}$$

## 2.3 Memory

$$\frac{\rho(S)[x] = s}{\langle S; x \rangle \Downarrow \langle S; s \rangle} \text{VAR}$$

$$\frac{x \notin \text{dom } \rho(S)}{\langle S; \tau \ x; \rangle \Downarrow \langle S[\rho \mapsto (\rho(S), (x \rightarrow \emptyset))]; s \rangle} \text{DECLARELOCAL}$$

$$\frac{x \notin \text{dom } \rho(S) \quad \langle S; e \rangle \Downarrow \langle S'; s \rangle}{\langle S; \tau \ x = e; \rangle \Downarrow \langle S'[\rho \mapsto (\rho(S'), (x \rightarrow s))]; s \rangle} \text{DECLAREASSIGNLOCAL}$$

$$\frac{\langle S; e_1 \rangle \Downarrow \langle S_1; \text{ptr } x \rangle \quad x \in \text{dom } \rho(S_1) \quad \langle S_1; e_2 \rangle \Downarrow \langle S_2; s \rangle}{\langle S; *e_1 = e_2 \rangle \Downarrow \langle S_2[\rho \mapsto (\rho(S_2), (x \rightarrow s))]; s \rangle} \text{UPDLOCAL}$$

$$\frac{\langle S; e_1 \rangle \Downarrow \langle S_1; s_1 \rangle \quad s_1 \neq \text{ptr } x \quad \langle S_1; e_2 \rangle \Downarrow \langle S_2; s_2 \rangle}{\langle S; *e_1 = e_2 \rangle \Downarrow \langle S_2[\mu \mapsto (\mu(S_2), (s_1 \rightarrow s_2))]; s_2 \rangle} \text{UPDGLOBAL}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S'; \text{ptr } x \rangle \quad \rho(S')[x] = s}{\langle S; *e \rangle \Downarrow \langle S'; s \rangle} \text{SELLOCAL}$$

$$\frac{\langle S; e \rangle \Downarrow \langle S'; \text{ptr } x \rangle \quad \rho(S')[x] = s}{\langle S; *e \rangle \Downarrow \langle S'; s \rangle} \text{SELGLOBAL}$$