			C tone	
	INSERTION SE	PR		
	Algorithm: Insertion soot (A)			
	1 for j =	o to Atlength		
			of the second se	
	4) wh	ile (120 f ALIZ	rkey)	
	5	ile (170 f ALIZ DLI+1] = ALI	T	
	VV	1 = 1=1 1 = 1=1 = end while		
		CTI+1] = EXY	CUN 1E	
		For		
	L	and the state of t	and the second s	
1	Cost of line	No:	of times Run	
	C ₁		n	
	Cg	and the second s	n=1	
	Cy	W.	n=1 (=,(t1)	
	54			
	Ce	2	1=1(t1-1)	
		No	(((till)	
	the state of the s			
	C		N=1	
	$T(n) = c_1 + n + (c_2 + c_3) \times h - D + \sum_{j=1}^{n-1} (j) \times c_4$ + $(c_4 + c_6) \times \sum_{j=1}^{n-1} (j+1) + c_8 + (n-1)$			
		to the first of the Miles	101/18/	
		lein or had	-n) (-n) O(n)	
The same				

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	Optimizing further		
)	we can apprimize to the searching by win		
/	we can aptimiz to the searching by wing binary search.		
<i>/</i>			
/	from O(n) to olalogn) for one		
<i>-</i>	element and to no oligan as		
1	which will insproving the searching complexity from O(n) to dealogn) for one of element and to not o(rog n) as o(n) for nelements.		
	But since it will take O(n) for one		
<i>y</i>	element to be placed at the correct		
<u> </u>	position, nelements will take not o(n)		
	or o(n2) time for being placed at their		
	remains o(n)		
	link & list.		
<u> </u>	link solist.		
-	Instead of Array that will improve the		
	complexity of swepping from o(n) to Edi) as we can insert an element in a linked		
<u></u>	list by changing habitery (will at 11/4/2)		
4	rest of element) But line the moderal		
1	rest of clament) But since the completely be search remains o(n) as are		
1	Cannot use binary search in link 1114		
4	Lannot use binary learth in link 11st Hence overall complexity remain some		
To.			