一次mysql死锁的排查

背景

pop购药上线后解冻操作经常发生死锁,报错日志如下:



死锁sql语句

session1	session2
DELETE from order_pay_status where id in (select b.id from (select id from order_pay_status where id > 3 order by id limit 500) b)	<pre>update order_pay_status set curr_status = 4 where id = 9;</pre>

死锁日志

```
*** (1) TRANSACTION:
TRANSACTION 24836, ACTIVE 9 sec starting index read
mysql tables in use 1, locked 1
LOCK WAIT 6 lock struct(s), heap size 1136, 3 row lock(s), undo log entries 2
MySQL thread id 374, OS thread handle 140602328172288, query id 6711 10.0.56.104 root updating
UDDATE order_pay_status
SET curr_status = 4,
modified = now()
WHERE
id = 9
*** (1) WAITING FOR THIS LOCK TO BE GRANTED:
RECORD LOCKS space id 251 page no 3 n bits 72 index PRIMARY of table `med_settle_purse`.`order_pay_status` trx id 24836 lock_mode X locks rec but not gap waiting
```

```
*** (2) TRANSACTION:
TRANSACTION 24842, ACTIVE 9 sec starting index read
{\it mysq1 tables in use 2, locked 2}
5 lock struct(s), heap size 1136, 4 row lock(s)
MySQL thread id 373, OS thread handle 140602736256768, query id 6717 10.0.56.104 root updating
DELETE from order_pay_status
where id in (
select b. id from (
select id from order_pay_status
where id > 0
AND DATE_FORMAT(created, '%Y-%m-%d') < DATE_FORMAT('2019-05-02 09:20:48.976', '%Y-%m-%d')
order by id
limit 500
) b
*** (2) HOLDS THE LOCK(S):
RECORD LOCKS space id 251 page no 3 n bits 72 index PRIMARY of table `med_settle_purse`.`order_pay_status` trx id 24842 lock mode S
Record lock, heap no 1 PHYSICAL RECORD: n_fields 1; compact format; info bits 0
0: len 8; hex 73757072656d756d; asc supremum;;
*** (2) WAITING FOR THIS LOCK TO BE GRANTED:
RECORD LOCKS space id 251 page no 3 n bits 72 index PRIMARY of table `med_settle_purse`.`order_pay_status` trx id 24842 lock_mode X
waiting
```

死锁特征

- 1: lock_mode X locks rec but not gap waiting
- 2: hold lock mode S, lock mode X waiting
- 3: 隔离级别: RR
- 4: 索引: 主键索引

四种类型的行锁

- 1: 记录锁 (LOCK_REC_NOT_GAP) : lock_mode X locks rec but not gap
- 2: 间隙锁 (LOCK_GAP) : lock_mode X locks gap before rec
- 3: Next-key 锁 (LOCK ORNIDARY) : lock mode X
- 4: 插入意向锁 (LOCK_INSERT_INTENTION) : lock_mode X locks gap before rec insert intention

重现步骤

session1	session2
DELETE from order_pay_status where id in (select b.id from (select id from order_pay_status where id > 3 order by id limit 500) b)	
首先子查询对符合条件的记录逐行加S锁)(包括id=9的记录)	
	update order_pay_status set curr_status = 4 where id = 9;
	其次更新语句对id=9的记录加排他锁,等待
DELETE from order_pay_status where id in (select b.id from (select id from order_pay_status where id > 3 order by id limit 500) b)	
执行delete操作对所有记录加排他锁,等待	

deadLock deadLock

分析

- 1: mysql的锁机制是公平锁,即有个锁队列,先到先得。
- 2: 不同事物下S锁与X锁互斥
- 3: 锁的范围是锁整个事务, 事务不结束, 锁不释放。

清楚上面的概念后就不难得出结论了:

- 事物A对id=9的记录加S锁,处于锁队列的第一个位置,并加锁成功,。 事物B对id=9的记录加X锁,处于锁队列的第二个位置,因为第一个位置的S锁不是自己事物内的所以互斥加锁失败,处于等待。 事物A对id=9的记录加X锁,处于锁队列的第三个位置,因为锁队列是先到先得的,第二位置的X锁不是自己事物内的锁所以互斥,处于等待。

这样就出现了锁等待,发生死锁。

解决

经过上面的分析,不难看出问题出现在了S锁上。如果没有S锁就不会发生死锁。那么S锁是怎么加的呢?因为删除语句做了子查询,为了防止读出数据不一致,所以读的时候加S锁,然后再执行删除。所以解决办法就是删除子查询,把session1的语句拆成两部分,第一部分普通查询(普通查询不会加S锁实际上是快照查),第一部分的结果作为第二部分的入参执行真正的删除操作。这样就不会加S锁,也不会出现死锁了。

资料

mysql发生死锁的种类有很多,这只是一个案例,更多案例https://github.com/aneasystone/mysql-deadlocks

参考博客:解决mysql死锁之路